Back Track

演算法	Time Complexity
DFS \ BFS	Adjacency Matrix : O(n2)
	Adjacency List : $O(E + V)$

Minimal Spanning Tree(DS 多談 Adj Matrix、演算法多談 Adj List)

演算法	Time Complexity	策略	
Kruskal'	Adjacency Matrix : O(V ²)	Greedy	Adjacency Matrix
s	Adjacency List(Binary Heap 製作): O(E log E)		Min-Heap
	Adjacency List(Fibonacci Heap 製作):O(E log V)		Disjoint Sets Tree
Prim's	Adjacency Matrix : O(V²)		Adjacency Matrix
	Adjacency List(Binary Heap 製作): O(E log V)		Min-Heap
	Adjacency List(Fibonacci Heap 製作):O(V log V + E)		Array

Shortest Path

演算法	Time Complexity	策略	負邊	負環
Dijkstra's	Adjacency Matrix : O(V2)	Greedy	X	X
	Adjacency List: O(V2)			
	Binary Heap : $O(E + V \log V)$			
	Fibonacci Heap : O(E + V log V)			
Bellman-Ford	Adjacency Matrix : O(V ³)	Dynamic Programming	O	X
	Adjacency List: O(V×E)			
Floyd-Warshall	Adjacency Matrix : O(V ³)	Dynamic Programming	О	X

Comparison of Various Structures

*	Array		Linked-List	AVL Tree
Insert	O(1)		O(n)	O(log n)
Delete	O(1)		O(n)	O(log n)
Search	O(log n)	//Binary Search	O(n)	O(log n)
Search kth item	O(1)		O(k)	O(log n)
Delete kth item	O(n-k)	//後面 n-k 個都要往前移一格	O(k)	O(log n)
Output in order	O(n)		O(n)	O(n)

[Thormas C 演算法版]

Operation	Binary Heap(Worst, DS 特愛考)	Binomial Heap(Worst)	Fibonacci Heap(攤)
Create-Heap	O(1)	O(1)	O(1)
Insert	O(log n)	O(log n) //[1]	O(1) //Lazy Merge
Delete	O(log n) //[6]	O(log n)	O(log n)
Find-Min	O(1)	O(log n) //[2]	O(1) //[3]
Extract-Min	O(log n)	O(log n)	O(log n)
Union	O(n)	O(log n)	O(1) //[4]
Decrease key	O(log n)	O(log n) //[5]	O(1)

[1]Binomial Heap 的 Worst \mathcal{F} ,Insert $O(\log n)$,但分攤 \mathcal{F} 為 O(1)

[2]因為Binomial Heap 最差下需比較 log n 棵樹的 Root 找 min

[3]Fibonacci Heap 有設定指標指向 Min

[4] 因為 Fibonacci Heap 採用 Lazy Merge

[5]例:看高度,因此O(log n)

[6]用最差情況 Delete Min 來考量: O(log n)

Sort

		Time Complexity			Space Complexity	Stable/Unstable
		Best	Average	Worst		
初等	Insertion	O(n)	$O(n^2)$	$O(n^2)$	O(1)	Stable
	Selection	$O(n^2)$	O(n ²)	O(n ²)	O(1)	Unstable
	Bubble	O(n)	O(n ²)	$O(n^2)$	O(1)	Stable
	Shell	$O(n^{3/2})$	O(n ²)	O(n ²)	O(1)	Unstable
高等	Quick	O(nlog n)	O(nlog n)	O(n ²)	$O(\log n) \sim O(n)$	Unstable
	Merge	O(nlog n)	O(nlog n)	O(nlog n)	O(n)	Stable
	Неар	O(nlog n)	O(nlog n)	O(nlog n)	O(1)	Unstable
線性	Radix(LSD)	$O(d^*(n+r)) \Longrightarrow O(n)$			O(r*n)	Stable
	Bucket(MSD)	$O(n+r) \Longrightarrow O(n) // \angle Radix \not$			O(r*n)	Unstable
	Counting	$O(n+k) \Longrightarrow O(n)$			O(n+k)	Stable

		Best	Average	Worst
初	Insertion	T(n)=1	$T(n)=T(n-1)+\Theta(n)$	T(n)=T(n-1)+(n-1)
等	Selection	T(n)=T(n-1)+1	$T(n)=T(n-1)+\Theta(n)$	T(n)=T(n-1)+(n-1)
	Bubble	T(n)=1	$T(n)=T(n-1)+\Theta(n)$	T(n)=T(n-1)+(n-1)
高	Quick	$T(n)=2T(n/2)+\Theta(n)$	$T(n)=1/n\sum[T(j-1)+T(n-j)+\Theta(n)]$	$T(n)=T(n-1)+\Theta(n)$
等	Merge	$T(n)=2T(n/2)+\Theta(n)$	$T(n)=2T(n/2)+\Theta(n)$	$T(n)=2T(n/2)+\Theta(n)$
	Heap	$T(n)=\Theta(n)+\Theta(n\log n)$	$T(n)=\Theta(n)+\Theta(n\log n)$	$T(n)=\Theta(n)+\Theta(n\log n)$
)

Greedy

Huffman	O(n log n)
Fraction Knapsack Problem	O(n log n)
Kruskal's	Adjacency List : O(E log E)
Prim's	Adjacency Matrix : O(V ²)
Sollin's	$O(V^2)$
Dijkstra's	Adjacency Matrix : O(V2)
Convex Hull	O(n log n)

Dynamic Programming

<u> </u>	
LCS / LIS	$O(mn) / O(n^2)$
0/1 Knapsack Problem	O(nW)
Chain Matrix	$O(n^3)$
Bellman-Ford	Adjacency Matrix : O(V ³)
Floyd-Warshell	Adjacency Matrix : O(V3)
OBST	O(n ³)
TSP(Traveling Salesman Problem)	$O(n^22^n)$

Divide-and-Conquer

Tower of Hanoi	O(2 ⁿ)
Binary Search	O(log n)
Quick Sort	O(n log n)
Merge Sort	$O(n^2)$
Bucket Sort(MSD)	O(n+r)
Closet Pair	O(n log n)
Strassen's Matrix	$O(n^{\log 7})$

Multiplication	

演算法解題技巧

1/5 1/5 1/5 V	
Branch and Bound 解 KP	O(nW) //Worst Case 亦為 NPC
Prune and Search 選第 k 小之數	$T(n)=T(n/5)+T(3n/4)+\Theta(n)=\Theta(n)$
陣列合併	$\Theta(\text{klogk}) + \Theta(\text{k}) = \Theta(\text{klogk})$
列出子集	$\Theta(2^{\rm n})$
找 1-1 函數	$\Theta(n)$
名人問題	$\Theta(n^2)$
平面上極大點	$\Theta(\text{nlogn})$
點之Rank	$T(n)=2T(n/2)+\Theta(n)=\Theta(n\log n)$
最大連續整數和	$\Theta(n)$

	Unsorted Single	Sorted Single	Unsorted Double	Sorted Double
	Linked List	Linked List	Linked List	Linked List
Search(L, k)	O(n)	O(1)	O(n)	O(logn)
Insert(L, p)	O(1)	O(n)	O(1)	O(n)
Delete(L, p)	O(n)	O(n)	O(1)	O(1)
Successor(L, p)	O(1)	O(1)	O(1)	O(1)
Predecessor(L, p)	O(n)	O(n)	O(1)	O(1)
Minimum(L)	O(n)	O(1)	O(n)	O(1)
Maximum(L)	O(n)	O(1)	O(n)	O(1)