



Reusable Enclaves for Confidential Serverless Computing

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The old way

- Your own server/VM
- Libs, OS, updates...
- Too heavy for a small app





Serverless

- Platform does them!
- In JS, Python...
- Commercialised



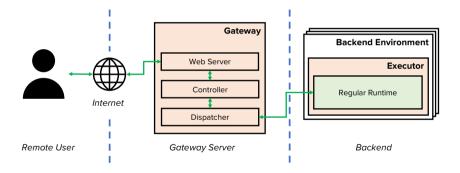




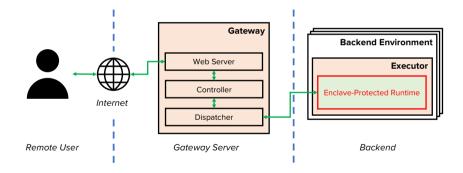




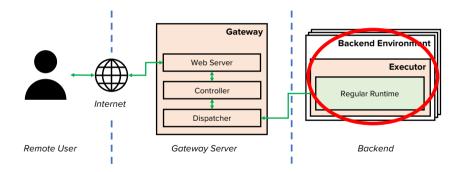
Serverless

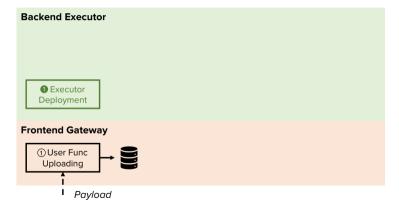


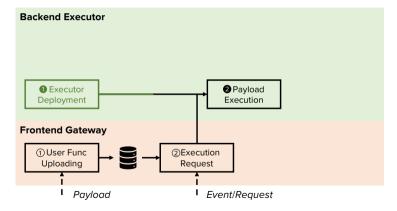
Confidential serverless

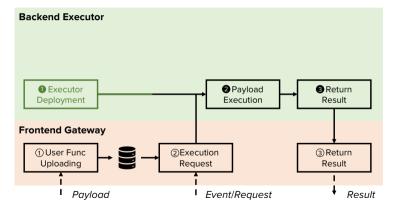


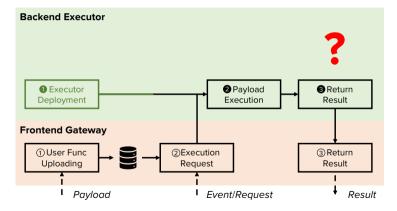
Root cause

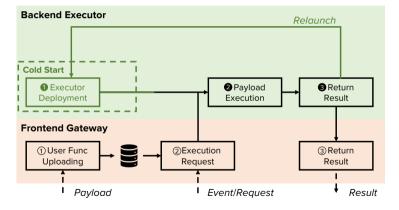










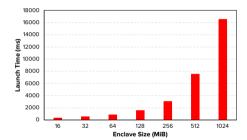


Penalty of Cold Start

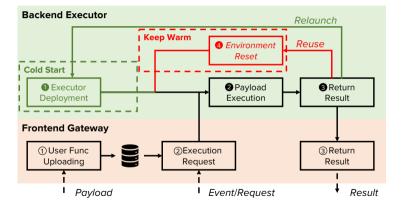
Facts [1]

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- Industry median memory
 - 170 MiB
- 50% workloads ends within
 - 1s



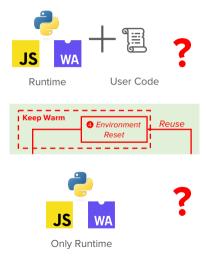
The solution?



The solution?

Keep warm?

- Keep user environment + runtime?
- Reset runtime?



The solution?

Keep warm?

- Keep user environment + runtime?
 Precious enclave memory!
- Reset runtime? Buggy Runtime!



Not good enough for confidential serverless!



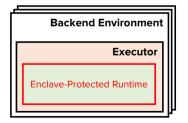


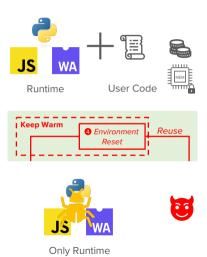




Question

Can we reset the enclave?





Challenges

How to reset?

How to prove the reset?

How to secure the reset?

Challenges

How to reset?

Enclave snapshot & rewinding

How to prove the reset?

Nested attestation

How to secure the reset?

Multi-Layer Intra-Enclave Compartmentalisation (MLIEC)

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Multi-Layer Intra-Enclave Compartmentalisation (MLIEC)

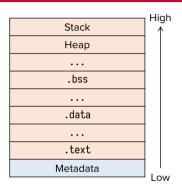
Generic architecture-independent method!

Enclave snapshot & rewinding

- Reset: Bring the enclave back to a known good state
- Take a snapshot and rewind

What a snapshot needs?

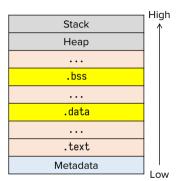
- Small memory footprint
- Fast to rewind



Enclave snapshot & rewinding

Initially...

- Stack, heap: Empty (zeros)
- .text: Read only (for now)
- .data, .bss



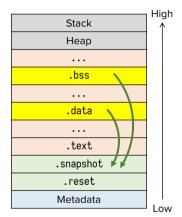
Enclave snapshot & rewinding

Initially...

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New reset module!

- Snapshot = copy
- Rewinding = copy back + zeroing



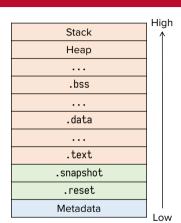
Nested attestation

Typical enclave attestation:

Boot time only

How to...

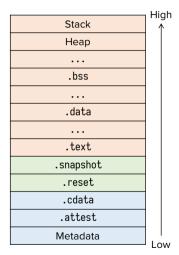
- Prove the reset indeed took place?
- Prove the reset is correct?
- User workload attestation?



Nested attestation

New attestation module!

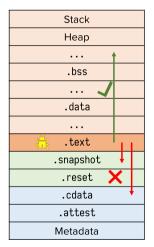
- Public-private key pair
- Reports with reset info
- User payload info



Multi-Layer Intra-Enclave Compartmentalisation

Observations

- Runtime (.text) can be buggy
- Must not touch snapshots and attestation data
- Layers of security



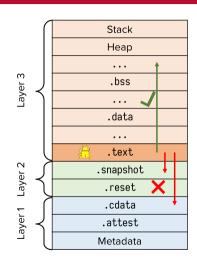
Multi-Layer Intra-Enclave Compartmentalisation

- A higher-security layer can access lower one's data
- Not vice versa!

Software-Fault Isolation (SFI)

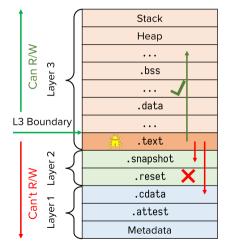
- Inspired by SGX-Shield [1]
- Compiler techniques

[1] J. Seo, et al, 'SGX-Shield: Enabling Address Space Layout Randomization for SGX Programs,' NDSS 2017.



Shepherded memory access

- R/W boundary for each layer
- Only allow R/W above the boundary



Shepherded memory access

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- Only allow R/W above the boundary

Steps

Step 1: Get offset from boundary

Step 2: Make offset positive

Step 3: Add back to boundary

Step 4: Access

Before
mov %rax, (%rdx)

After

	mov	%rdx, %r14
Step 1	sub	%r15, %r14
Step 2	shl	\$1 , %r14
	shr	\$1 , %r14
Step 3	add	%r15, %r14
Step 4	mov	%rax, (%r14)

Problem...

Branching to arbitrary address can bypass the SMA

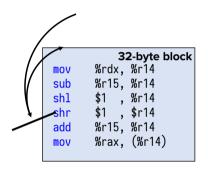
```
mov %rdx, %r14
sub %r15, %r14
sh1 $1 , %r14
shr $1 , $r14
add %r15, %r14
mov %rax, (%r14)
```

Problem...

Branching to arbitrary address can bypass the SMA

Aligned branching

- Branching to arbitrary address can bypass the SMA
- Emit code into fixed-size blocks (e.g., 32 bytes)
- Force all branching aligned to the block size



Problem...

Branching to arbitrary address can bypass the SMA

Aligned branching

- Branching to arbitrary address can bypass the SMA
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```
Before jmp *%rax
```

```
After and ~$32, %rax imp *%rax
```

Problem...

Not all functions can be instrumented...

- Security:
 - Boundary setup
 - RWX-granting
- Performance
 - memcpy





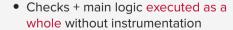
Security

Performance

Problem...

Not all functions can be instrumented...

- Security:
 - Boundary setup
 - RWX-granting
- Performance
 - memcpy



⇒ Control Flow Integrity (CFI)





Security

Performance

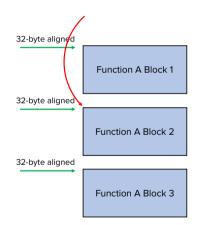
Traditional CFI

Trap & check

• Slow!



- Branches are aligned
- ⇒ Can't branch to unaligned target



Traditional CFI

Trap & check

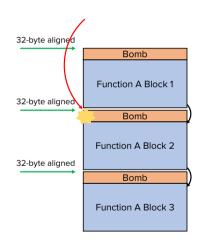
• Slow!

\triangle

- Branches are aligned
- ⇒ Can't branch to unaligned target

Solution

- Emit a bomb before each block
- Chain blocks with jumps



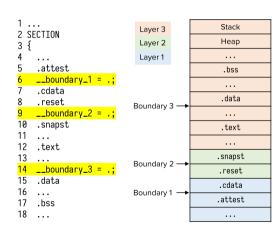
MLIEC: Multi-layer compartmentalisation

Why

- Least privilege principle
- Attestation > reset > runtime

How

- Fixeible boundary
- Linker script

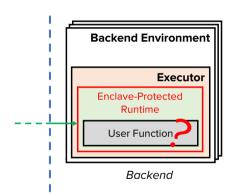


MLIEC: Dynamically-loaded code

Read-only code: Not enough for serverless

Ahead-of-Time (AoT)

- Bytecode to native binaries
- Good performance
- May contain any code/instructions...



MLIEC: Dynamically-loaded code

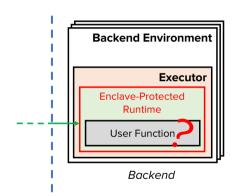
Read-only code: Not enough for serverless

Ahead-of-Time (AoT)

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Solution

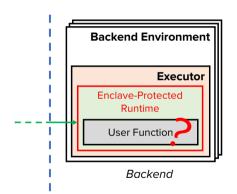
Use MLIEC techniques when compiling AoT binaries



MLIEC: Dynamically-loaded code

RWX granting function

- AoT requires RWX area
- Protect RWX granting with unaligned critical functions
- Disable it before user code execution



12/17

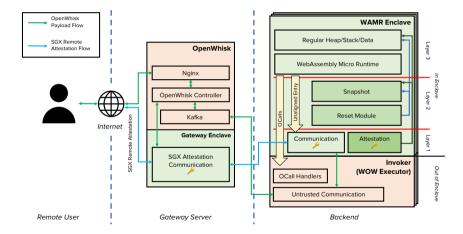
- MLIFC: LLVM-based toolchain
- Enclave: Intel SGX
- Frontend: OpenWhisk
 - Open source platform
 - Widely adopted
- Backend: WAMR
 - Open source
 - AoT mode

Code based on WebAssembly on OpenWhisk (WOW) [1]





Implementation



Implementation

• LLVM: 1070 LoC

OpenWhisk:

Action: 107 LoC

Gateway Untrusted: 1478 LoC

• Gateway Enclave: 1978 LoC

Executor:

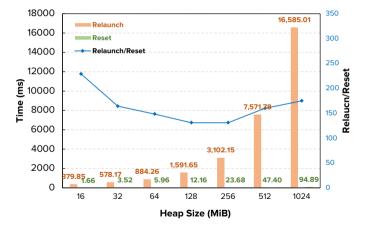
WOW: 1457 LoCEnclave: 4098 LoC

• Total: 10188 LoC

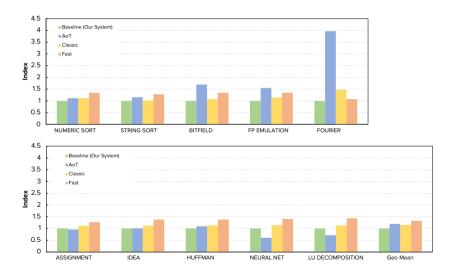




Relaunch vs. reuse

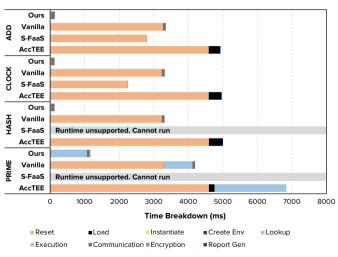


Instrumentation overheads



15/17 Motivations OOO Design Implementation Evaluation OO● Conclusion OO

Real-world end-to-end



[1] F. Alder, et al, 'S-FaaS: Trustworthy and Accountable Function-as-a-Service Using Intel SGX,' CCS 2019.

[2] D. Goltzsche, et al, 'AccTEE: A WebAssembly-Based Two-Way Sandbox for Trusted Resource Accounting,' Middleware 2019.

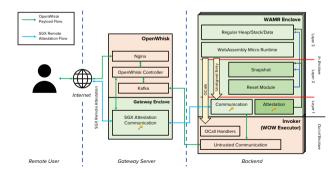
Conclusion



Reset benefits are significant

Solved cold start problem in confidential serverless with reusable encalves

- Enclave snapshot & rewinding
- Nested attestation
- MLIEC





Source Code

https://github.com/OSUSecLab/Reusable-Enclaves

SecLab @ OSU

https://go.osu.edu/seclab

Teecert Labs @ SUSTech

https://teecertlabs.com

NSEC @ SJTU

https://nsec.sjtu.edu.cn