## **EE3980 Algorithms**

## **HW4 Trading Stock, II**

104061212 馮立俞

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## **Introduction**

In this assignment, we're required to improve the brute-force approach for maximum sum array problem in HW03 to reach  $O(n^2)$  time complexity. Then, we'll discuss whether it's possible to device an algorithm having lower complexity than  $O(n \log n)$ .

### **Approach**

#### **Recap: Brute-force Approach**

```
1. Algorithm MaxSubArrayBF(A, n, low, high) // Find low and high to
                             maximize \Sigma A[i], low \le i \le high.
2. {
3.
       max: = 0;low: = 1;high: = n;
4.
       for j: = 1 to n do { // Try all possible ranges: A[j : k].
           for k: = j to n do {
5.
                    sum = price[high] - price[low]; // n^2 complexity version
                    sum: = 0;
                    for i: = j to k do {
                            sum: = sum + A[i]; // n^3 complexity version,
                                                  // could choose either one
9.
                    if (sum > max) then {
                       // Record the maximum value and range.
10.
                        max = sum;
11.
                        low = j;
12.
                        high = k;
13.
14.
           }
15.
       }
   return max;
```

Since we've known the prices of stock in the given time, the price change over a certain period can be obtained by subtracting the high price with low one. Doing so would save us a loop; thus improve the overall complexity from  $O(n^3)$  to  $O(n^2)$ . Space complexity would remain the same nonetheless, i.e. O(n).

### **Linear-Complexity Approach**

```
    Algorithm LinearApproach (Prices[], N) {

2.
        max_sum: = 0,
3.
        sum: = 0;start: = 0,
4.
        end: = 0,
        temp_start: = 0;
5.
6.
        for i from 0 to N - 1: {
            sum +=(a[i];
7.
                             What is this?
8.
            if (sum < 0) {
                sum: = 0; temp_start: = i + 1;
9.
10.
            } //restart record
11.
            if (sum > max_sum) {
12.
                max_sum = sum;
13.
                start = temp_start;
14.
                end = i;
15.
            } //update record
16.
17.
        if (start >= end)
18.
            do something // array all negative, error handling
19.
        return max_sum, start, end;
20.}
```

In this implementation, only one single loop is used. Obviously the time complexity is O(n). Plus, no other large memory space is required, so the space complexity is O(n).

?

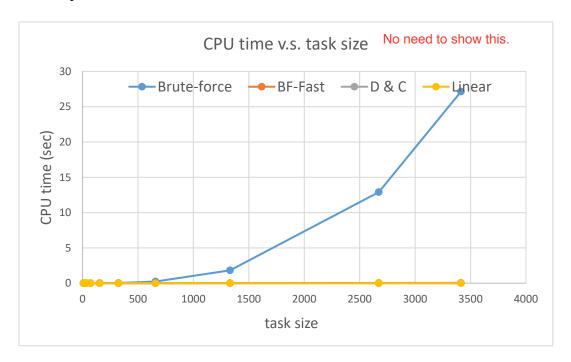
In the algorithm, we keep adding <a href="mailto:new term">new term</a> to current sum, restart our record if current sum is less than zero, and update <a href="mailto:max\_sum">max\_sum</a> and start & end if current sum is greater than <a href="mailto:max\_sum">max\_sum</a>. Simple as that, yet it works (at least intuitively and empirically for this homework)!

# **Results and analysis**

Table 1. Algorithm CPU time (in seconds) vs. task size

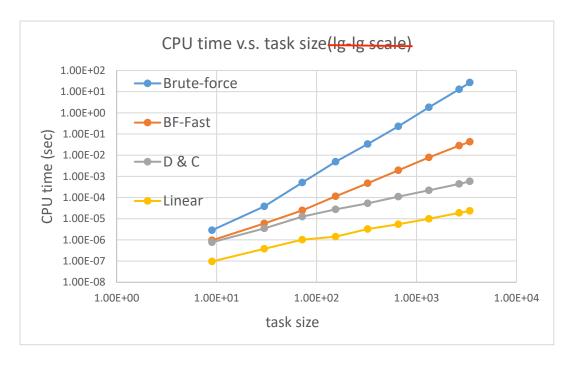
task size	9	30	72	155	325	658	1331	2672	3414
Brute-	2.86E-06	3.81E-05	0.000509	0.0049	0.0336	0.23	1.82	12.9	27.1378
force									
BF-Fast	9.54E-07	5.96E-06	2.48E-05	1.14E-04	4.73E-04	1.93E-03	7.84E-03	2.83E-02	4.27E-02
D&C	7.72E-07	3.48E-06	1.25E-05	2.76E-05	5.33E-05	1.10E-04	2.19E-04	4.41E-04	5.79E-04
Linear	9.61E-08	3.78E-07	1.02E-06	1.42E-06	3.26E-06	5.46E-06	9.93E-06	1.88E-05	2.35E-05

We can plot above table as follows

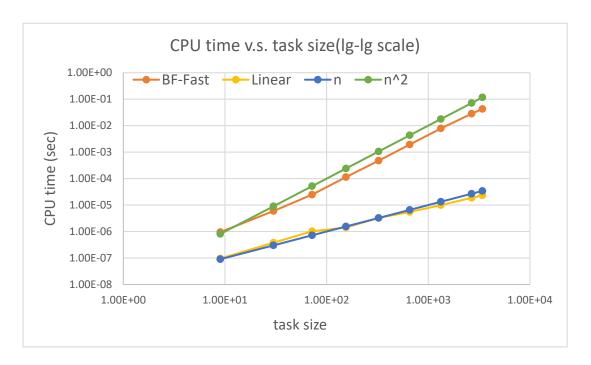


Well, the complexity of the four algorithms are basically not on the same league.

It's pretty hard to plot them on linear scale without some curves being suppressed.



We expect their complexity to be  $O(n^3)$ ,  $O(n^2)$ ,  $O(n \log n)$ , O(n) respectively. And the curves above are quite fit.



We can further examine the complexity by plotting  $O(n^2)$  and O(n) curves with them, and they fit well too.

# **Observations and Conclusion**

"Brevity is the soul of wit", and the same philosophy applies for algorithms too.

With some ingenuity, a great amount of time could be saved.

```
1 /***********
 2 EE3980 Algorithms HW04 Trading Stock, II
    Li-Yu Feng 104061212
    Date:2018/4/1
 7 #include <stdio.h>
 8 #include <stdlib.h>
9 #include <sys/time.h>
10
11 typedef struct sSTKprice {
                                 //store date, price & price change
      int year, month, day;
      double price, change;
13
14 } STKprice;
15
16 typedef struct retval{ //store low, high index and the maximum sum between
                         //for convenient return
      int low, high;
18
      double sum;
19 } RETval;
20
21 double GetTime(void);
22 void MaxSubArrayBF(STKprice *A, int n);
                                                        //Brute-force method
23 RETval MaxSubArray(STKprice *A, int begin, int end);
                                                        //D & C method
24 RETval MaxSubArrayXB(STKprice *A, int begin, int mid, int end);
                                                   //cross boundary
25
26 RETval MaxSubArrayFast(STKprice *A, int n);
                                                    //linear complexity
27 void PrintResult(STKprice *A, RETval result, double time,int flag); //as its n
  ame
28
29 double GetTime(void)
30 {
31
      struct timeval tv;
      gettimeofday(&tv,NULL);
33
      return tv.tv_sec+1e-6*tv.tv_usec;
34 }
35
36 void MaxSubArrayBF(STKprice *A, int n){
37
      double max = 0, sum;
      double t;
38
                             //time
      int low = 0, high = n-1;
40
      int i,j,k;
41
42
      t = GetTime();
43
      for (j = 0; j < n; j++){
                                     //try all n(n-1)/2 possible situations
44
          for (k = j; k < n; k++){
              sum = A[k].price - A[j].price; // n^2 complexity version
45
              //sum = 0;
46
47
              //for(i = j; i <= k; i++){
              // sum += A[i].change;
48
              //}
                                             //n^3 version
49
```

```
50
               if(sum > max){
51
                   max = sum;
                   low = j + 1;
52
                                                //Note: for alignment in printing
53
                   high = k;
                                                //process
54
               }
           }
55
56
       }
57
       t = GetTime() - t;
58
59
       //print result
       printf("Brute-force approach:\n");
60
       printf(" CPU time %.3g s\n", t);
61
       if(low != 0)printf(" Buy: %d/%d/%d at %g\n'', A[low-1].year,
62
                       A[low-1].month,A[low-1].day,A[low-1].price);
63
       else printf(" Buy: %d/%d/%d at $%g\n",A[0].year,
64
                       A[0].month,A[0].day,A[0].price );
65
                 Sell: %d/%d/%d at $%g\n",A[high].year,A[high].month,
66
       printf("
                                            A[high].day,A[high].price);
67
68
       printf(" Earning: $%g per share.\n",max);
69 }
70
71
72 RETval MaxSubArray(STKprice *A, int begin, int end){
74
       int mid;
       RETval Ans,lret,rret,xret;
75
                                        //store return values from divided aprts
76
       double lsum,rsum,xsum;
77
78
       if(begin == end){
79
           Ans.low = begin;
80
           Ans.high = end;
           Ans.sum = A[begin].change;
81
82
           return Ans;
83
       }
       mid = (begin + end) / 2;
84
85
       lret = MaxSubArray(A,begin,mid);
86
       rret = MaxSubArray(A,mid+1,end);
87
       xret = MaxSubArrayXB(A,begin,mid,end);
88
89
       lsum = lret.sum;
90
       rsum = rret.sum;
91
       xsum = xret.sum;
92
93
       if(lsum >= rsum && lsum >= xsum){ //left side returns maximum
94
           Ans.low = lret.low;
95
           Ans.high = lret.high;
96
           Ans.sum = lsum;
97
       }
       else if (rsum >= xsum){
                                        //right side
98
99
           Ans.low = rret.low;
```

```
100
            Ans.high = rret.high;
101
            Ans.sum = rsum;
102
        }
        else{
                                         //cross boundary
103
104
            Ans.low = xret.low;
            Ans.high = xret.high;
105
106
            Ans.sum = xsum;
107
        }
108
        return Ans;
109 }
111 RETval MaxSubArrayXB(STKprice *A, int begin, int mid, int end){
        double lsum = 0,rsum = 0;
112
113
        int low = mid;
        int high = mid + 1;
114
        double sum = 0;
115
116
        int i;
        RETval Ans;
117
118
        for(i = mid; i >= begin; i--){
                                             //find left side max sum
119
            sum = sum + A[i].change;
120
121
            if(sum > lsum){
122
                lsum = sum;
123
                low = i;
            }
124
125
        }
126
        sum = 0;
        for(i = mid + 1; i <= end; i++){</pre>
                                             //find at right side
127
128
            sum = sum + A[i].change;
129
            if(sum > rsum){
130
                rsum = sum;
131
                high = i;
132
            }
133
        }
134
        Ans.low = low;
135
136
        Ans.high = high;
137
        Ans.sum = lsum + rsum;
138
        return Ans;
139 }
140
141 void PrintResult(STKprice *A, RETval result, double time, int flag){
        int low = result.low;
143
        int high = result.high;
144
        double sum = result.sum;
145
146
                        printf("Divide and Conquer approach:\n");
        if(flag == 1)
147
        else printf("Fast approach:\n");
148
        printf(" CPU time %.3g s\n", time);
149
        if(low != 0)printf(" Buy: %d/%d/%d at %g\n",A[low-1].year,
                                                                          //to avoid
```

```
150
                        A[low-1].month, A[low-1].day, A[low-1].price);
                                                                          //segfault
151
        else printf(" Buy: %d/%d/%d at %g\n",A[0].year,
152
                        A[0].month,A[0].day,A[0].price );
                  Sell: d/d/d at g^n, A[high].year, A[high].month,
153
154
                                             A[high].day,A[high].price);
155
        printf("
                  Earning: $%g per share.\n",sum);
156
157 }
158
159 RETval MaxSubArrayFast(STKprice *A, int n){
        double sum = 0;
160
        int temp_start = 0;
161
        int i;
162
        RETval Ans;
163
164
        Ans.sum = 0;
165
        Ans.low = 0;
166
        Ans.high = 0;
167
168
        for (i = 0; i < n; i++){
169
            sum +=A[i].change;
170
171
            if (sum < 0){
172
173
                sum = 0;
174
                temp_start = i+1;
175
            }
                                     //restart record
176
177
            if (sum > Ans.sum){
178
                Ans.sum = sum;
179
                Ans.low = temp_start;
                Ans.high = i;
                                     //update record
180
181
            }
182
        }
183
        return Ans;
184 }
185
186
187
188 int main()
189 {
190
        int Ndays;
191
        int i;
192
        double t;
193
        STKprice *Prices;
194
        RETval result;
195
196
        scanf("%d",&Ndays);
        Prices = malloc(Ndays * sizeof(STKprice));
197
198
```

```
199
        for ( i = 0; i < Ndays; i++){
200
            scanf(" %d %d %d %lf", &Prices[i].year, &Prices[i].month,
201
                                     &Prices[i].day, &Prices[i].price );
202
        }
203
       Prices[0].change = 0;
                                    //calculate the price changes
204
205
        for ( i = 1; i < Ndays; i++)
            Prices[i].change = Prices[i].price - Prices[i-1].price;
206
207
208
        printf("N = %d\n",Ndays);
209
210
        MaxSubArrayBF(Prices,Ndays);
211
212
       t = GetTime();
       for ( i = 0; i < 1000; i++)
213
214
        {
215
            result = MaxSubArray(Prices,0,Ndays-1);
        }
216
217
        t = GetTime() - t;
        PrintResult(Prices, result, t/1000, 1);
218
219
220
       t = GetTime();
       for ( i = 0; i < 1000; i++)
221
222
            result = MaxSubArrayFast(Prices,Ndays);
223
224
        }
       t = GetTime() - t;
225
        PrintResult(Prices, result, t/1000, 2);
226
227
228
229
        return 0;
230 }
231
232
233
234
235
```

9				_	0
S	C	$\gamma$	e:	U	I()
$\sim$	v	$\mathcal{I}$	. •	U	U

[Pseudo code] on page 3 should be corrected.

- Your also should discuss arrays needed by those two algorithms are different.

[Correctness] of your algorithms should be provided.

- Either rigorous proof or compare the outputs to the old algorithms.