# **6.033 Spring 2019**Lecture #6

- Monolithic kernels vs. Microkernels
- Virtual Machines

## **operating systems** enforce modularity on a single machine using **virtualization**

in order to enforce modularity + build an effective operating system

- programs shouldn't be able to refer to (and corrupt) each others' memory
- virtual memory

2. programs should be able to **communicate** 

- **bounded buffers**(virtualize communication links)
- 3. programs should be able to **share a CPU** without one program halting the progress of the others
- threads
  (virtualize processors)

today: running multiple OSes at once (and dealing with kernel bugs)

### Virtual Machines

virtual machine running guest OS

virtual machine running guest OS

physical hardware

problem: how to (safely) share physical hardware?

### Virtual Machines

### VMM runs in kernel-mode on hardware

virtual machine running guest OS

virtual machine running guest OS

virtual machine monitor (VMM)

physical hardware

### guest OS

### guest OS

### virtual hardware

U/K PTR page table

### virtual hardware

U/K
PTR
page table

### virtual machine monitor (VMM)

### physical hardware

U/K, PTR, page table, ...

VMM's goal: virtualize hardware

### guest OS

### guest OS

### guest virtual address

### virtual hardware

```
U/K
PTR
page table
```

### virtual hardware

```
U/K
PTR
page table
```

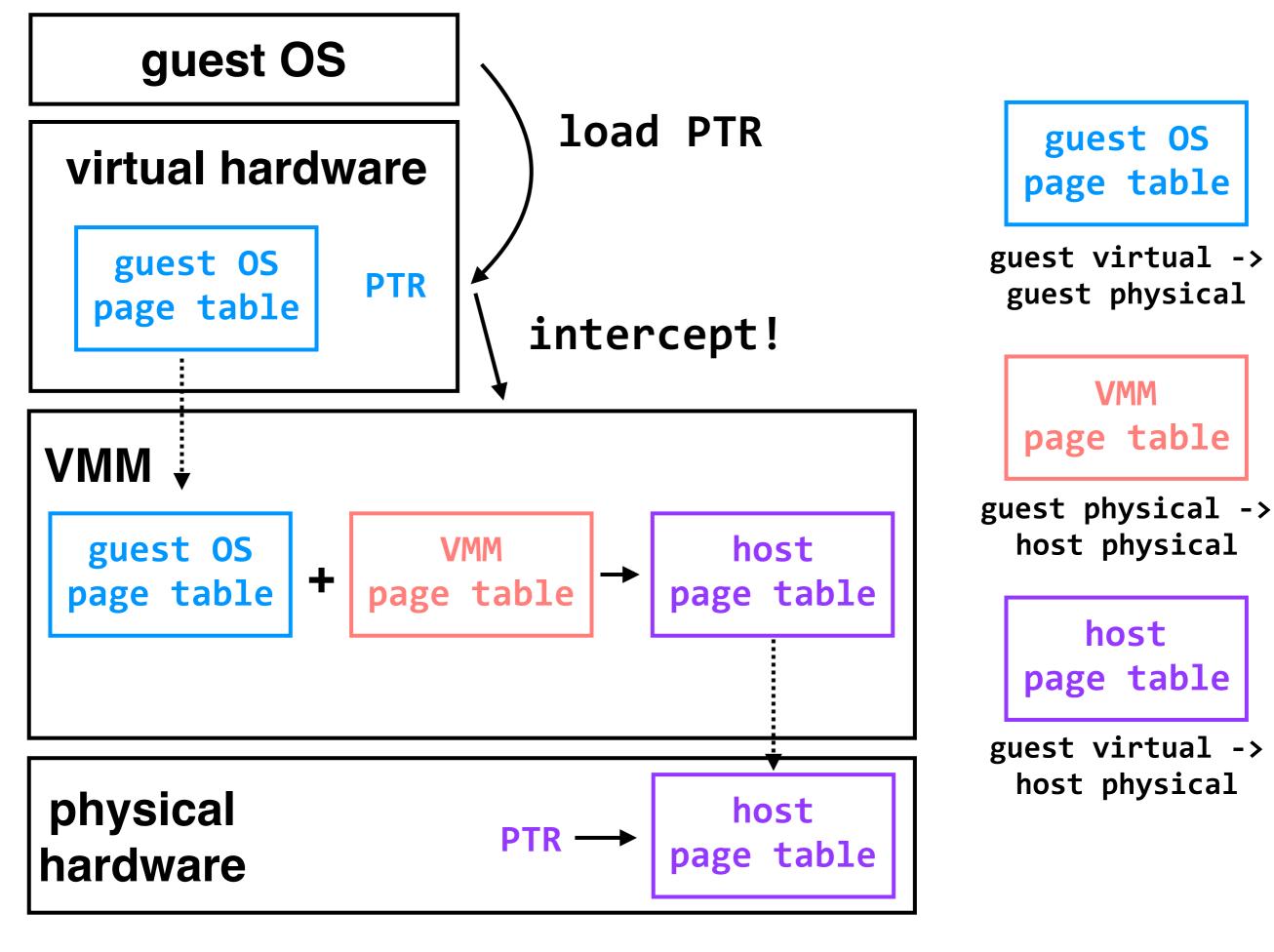
guest physical
 address

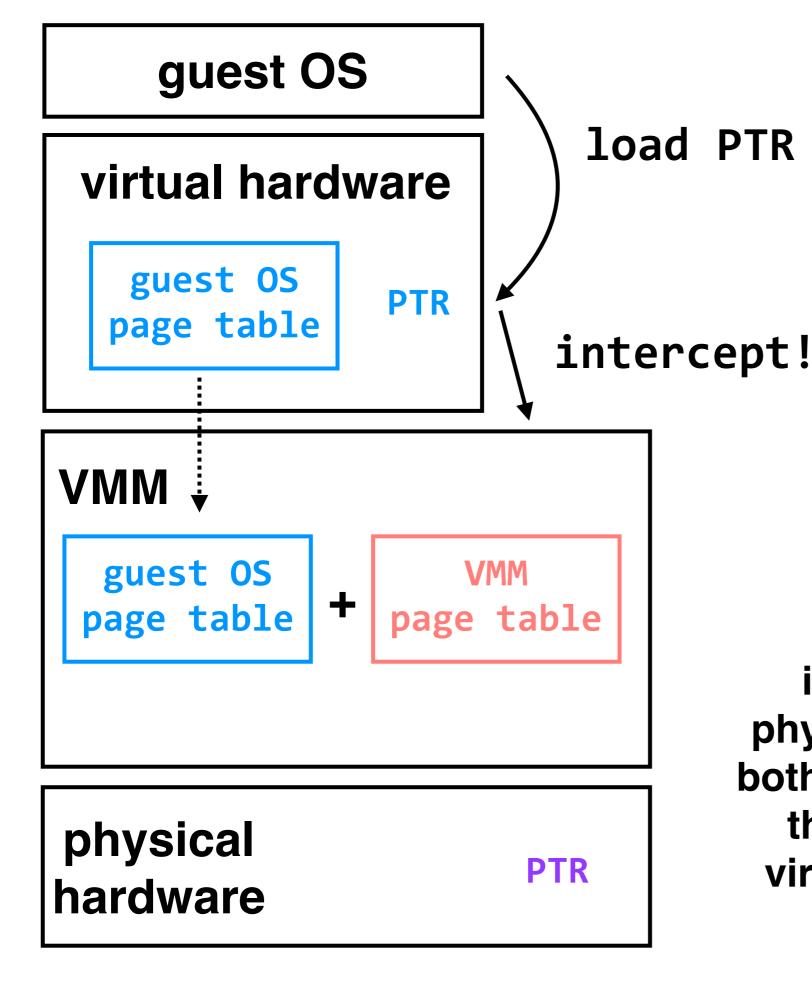
### virtual machine monitor (VMM)

### physical hardware

U/K, PTR, page table, ...

host physical address





guest OS
page table

guest virtual ->
 guest physical

VMM page table

guest physical ->
 host physical

in modern hardware, the physical hardware is aware of both page tables, and performs the translation from guest virtual to host physical itself

### guest OS

### guest OS

### virtual hardware

U/K PTR page table

### virtual hardware

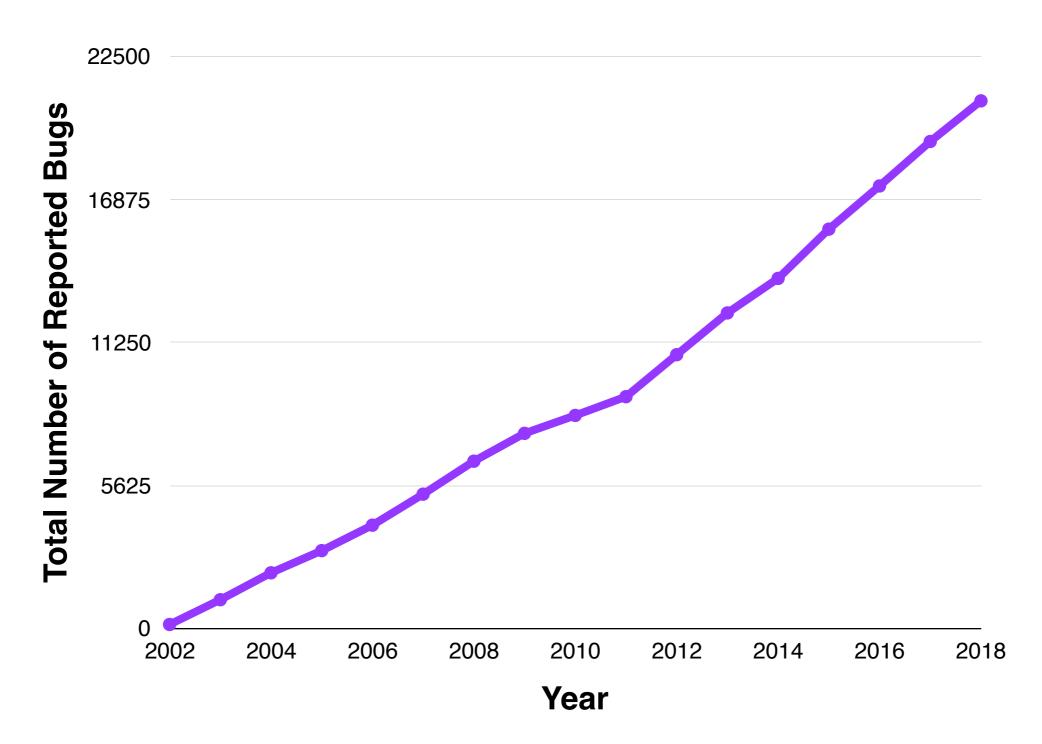
U/K
PTR
page table

### virtual machine monitor (VMM)

### physical hardware

U/K, PTR, page table, ...

VMM's goal: virtualize hardware



source: bugzilla.kernel.org, count of all bugs currently marked NEW, ASSIGNED, REOPENED, RESOLVED, VERIFIED, or CLOSED, by creation date (rough estimate!)

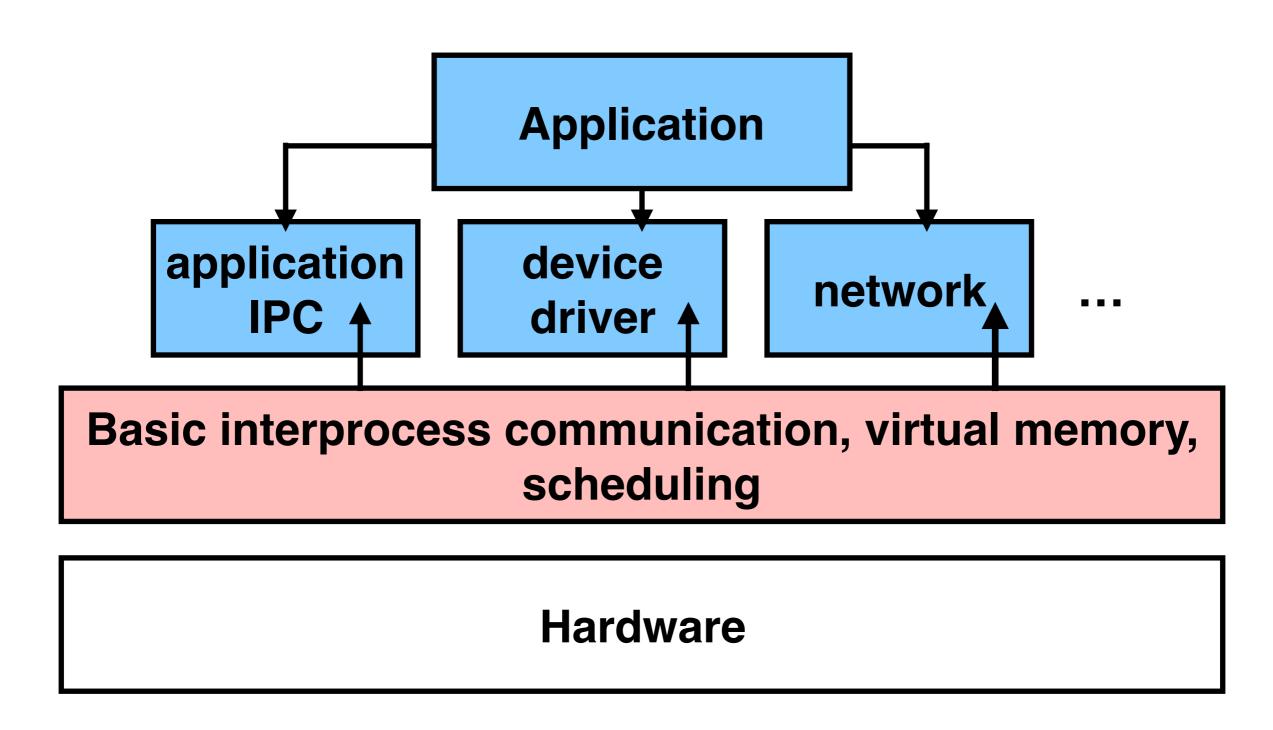
## monolithic kernels: no enforced modularity within the kernel itself

**Application** 

Basic interprocess communication, virtual memory, scheduling, file server, device drivers, network, ...

**Hardware** 

microkernels: enforce modularity by putting subsystems in user programs



- **Virtual Machines** allow us to run multiple **isolated** OSes on a single physical machine, similar to how we used an OS to run multiple programs on a single CPU. VMs must handle the challenges of virtualizing the hardware (examples: virtualizing memory, the U/K bit).
- Monolithic kernels provide no enforced modularity within the kernel. Microkernels do, but redesigning monolithic kernels as microkernels is challenging.