

6.033 Spring 2019

Lecture #6

- **Monolithic kernels vs. Microkernels**
- **Virtual Machines**

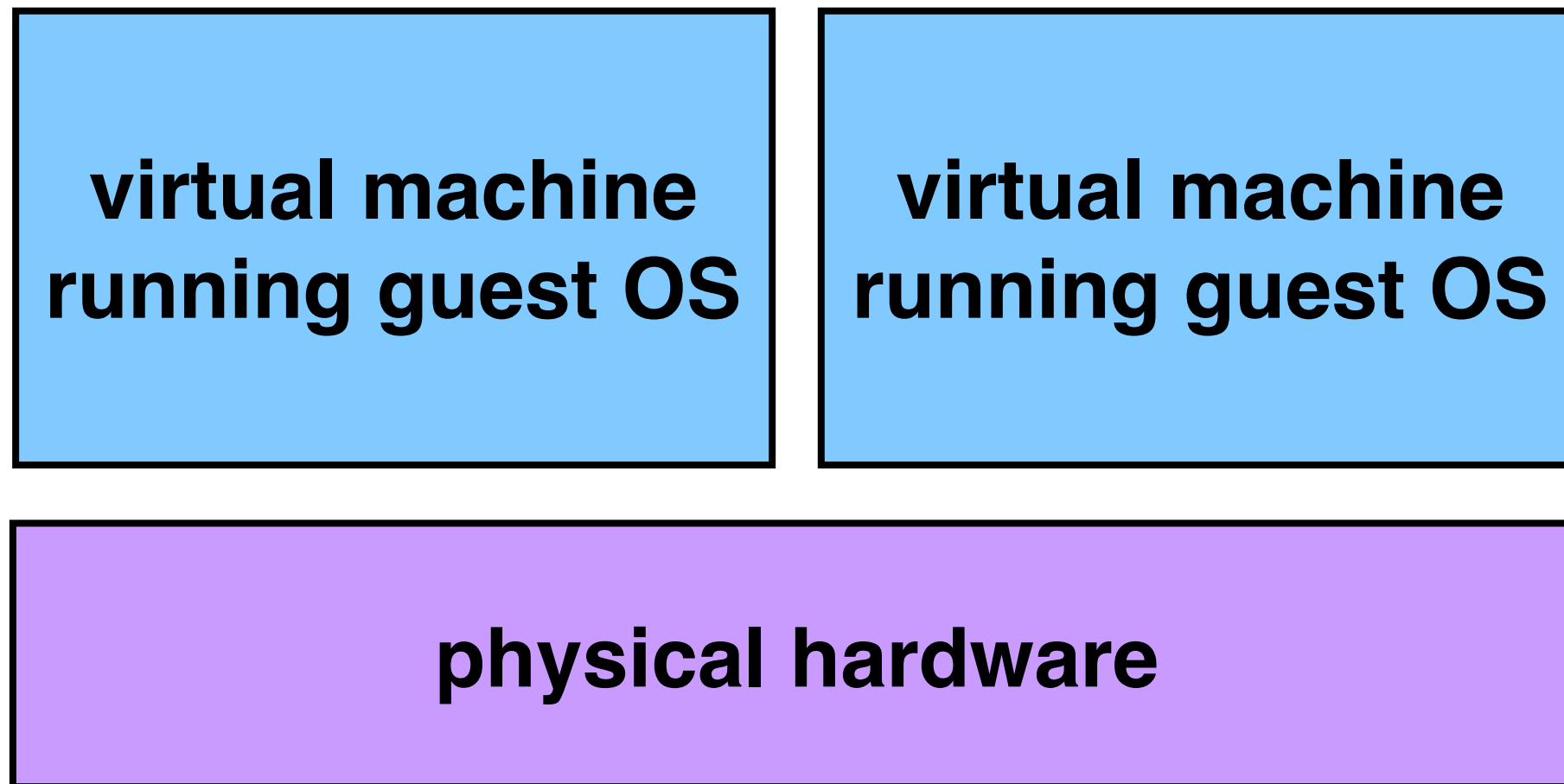
operating systems enforce modularity on a single machine using **virtualization**

in order to enforce modularity + build an effective operating system

- 1. programs shouldn't be able to refer to (and corrupt) each others' **memory** → **virtual memory**
- 2. programs should be able to **communicate** → **bounded buffers**
(virtualize communication links)
- 3. programs should be able to **share a CPU** without one program halting the progress of the others → **threads**
(virtualize processors)

today: running multiple OSes at once
(and dealing with kernel bugs)

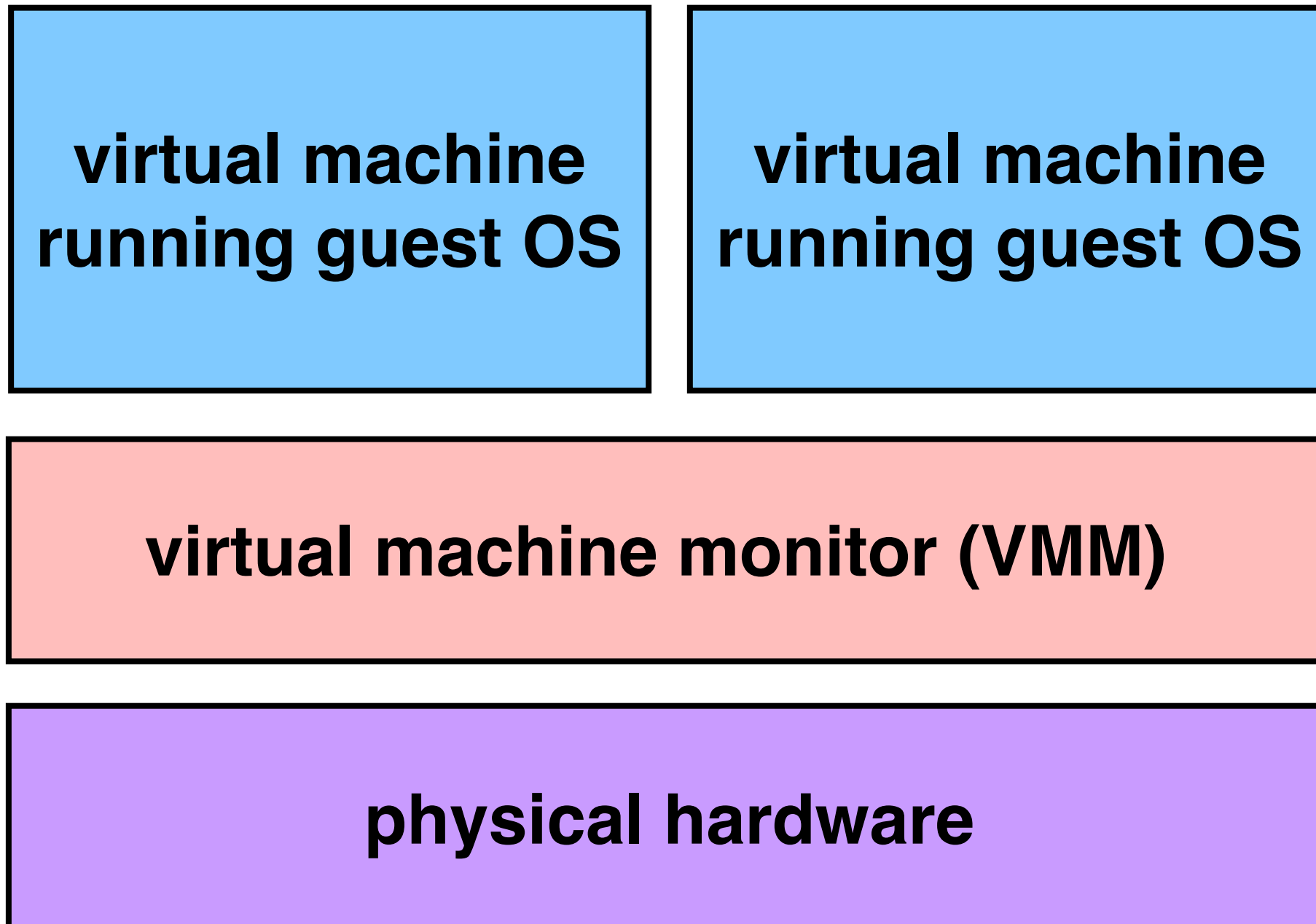
Virtual Machines



problem: how to (safely) share physical hardware?

Virtual Machines

VMM runs in kernel-mode on hardware



guest OS

guest OS

virtual hardware

U/K
PTR
page table
...

virtual hardware

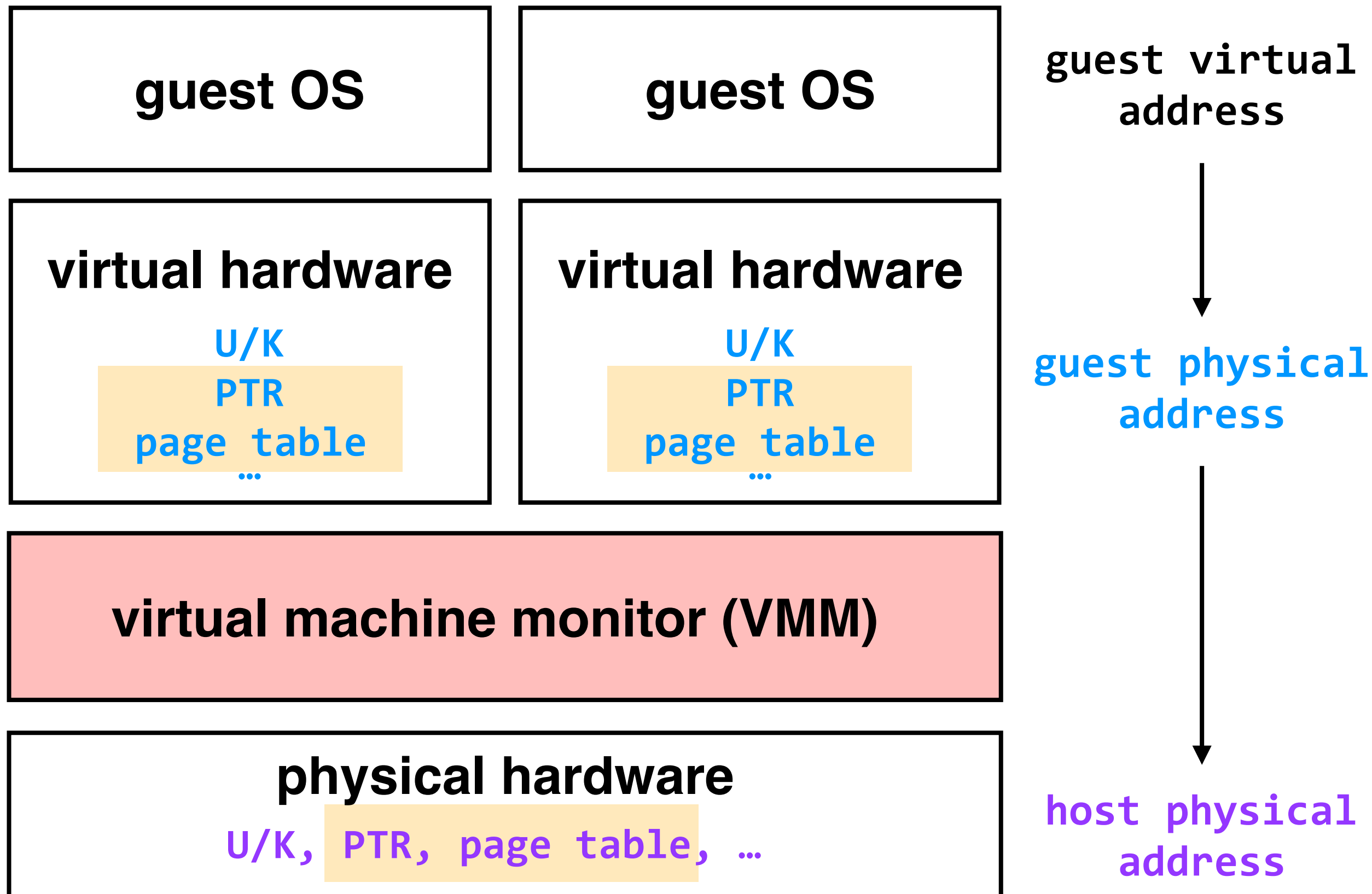
U/K
PTR
page table
...

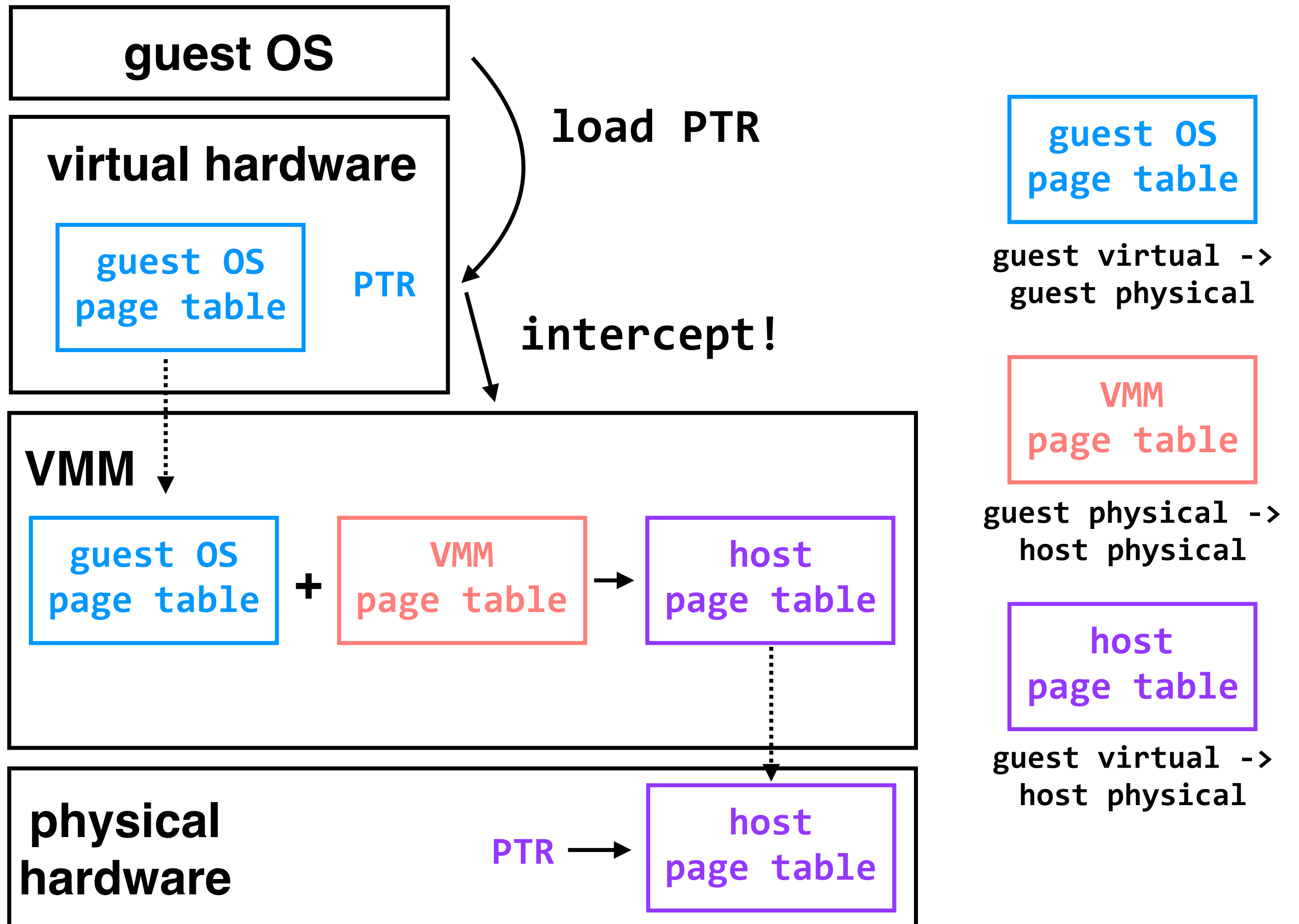
virtual machine monitor (VMM)

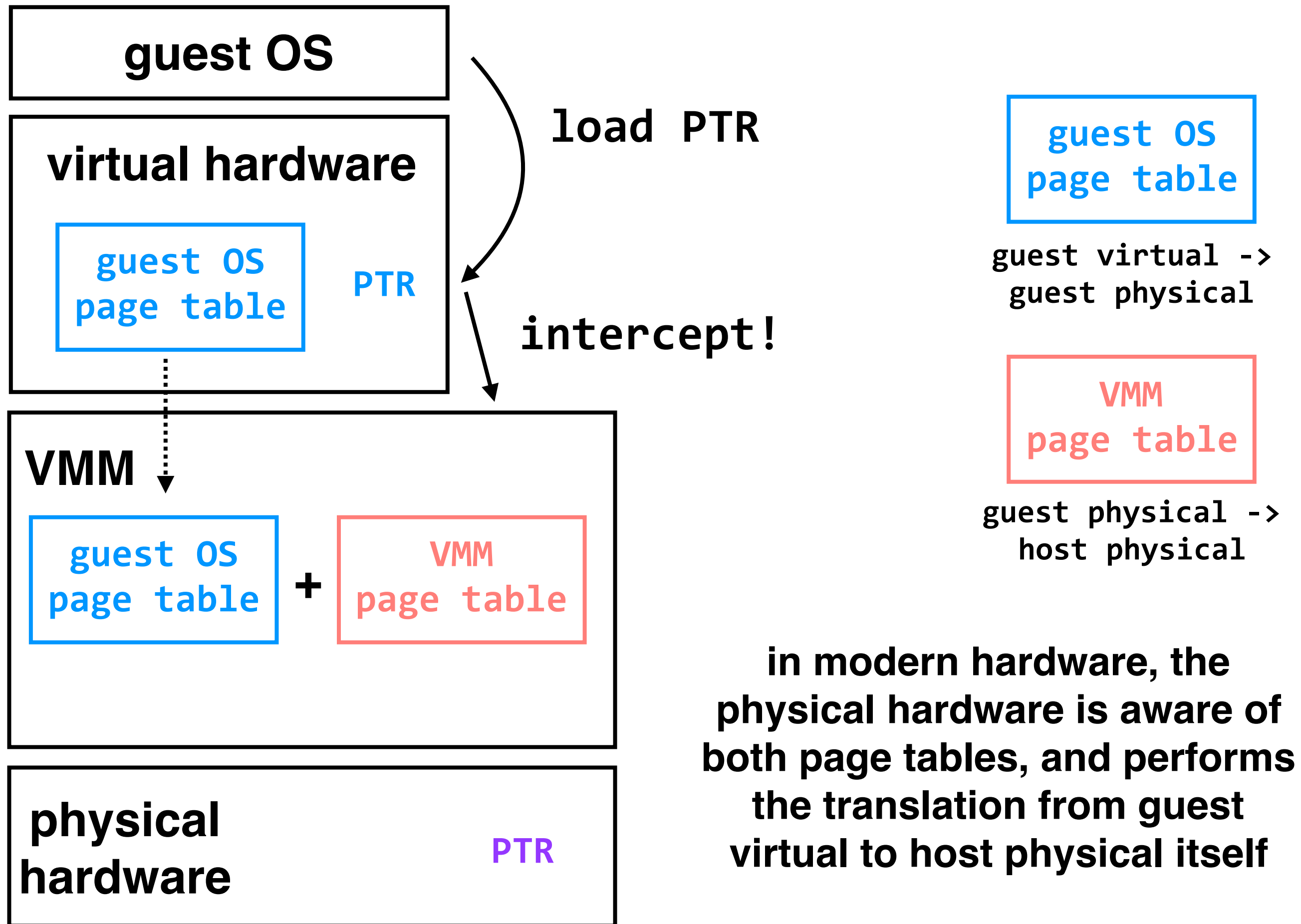
physical hardware

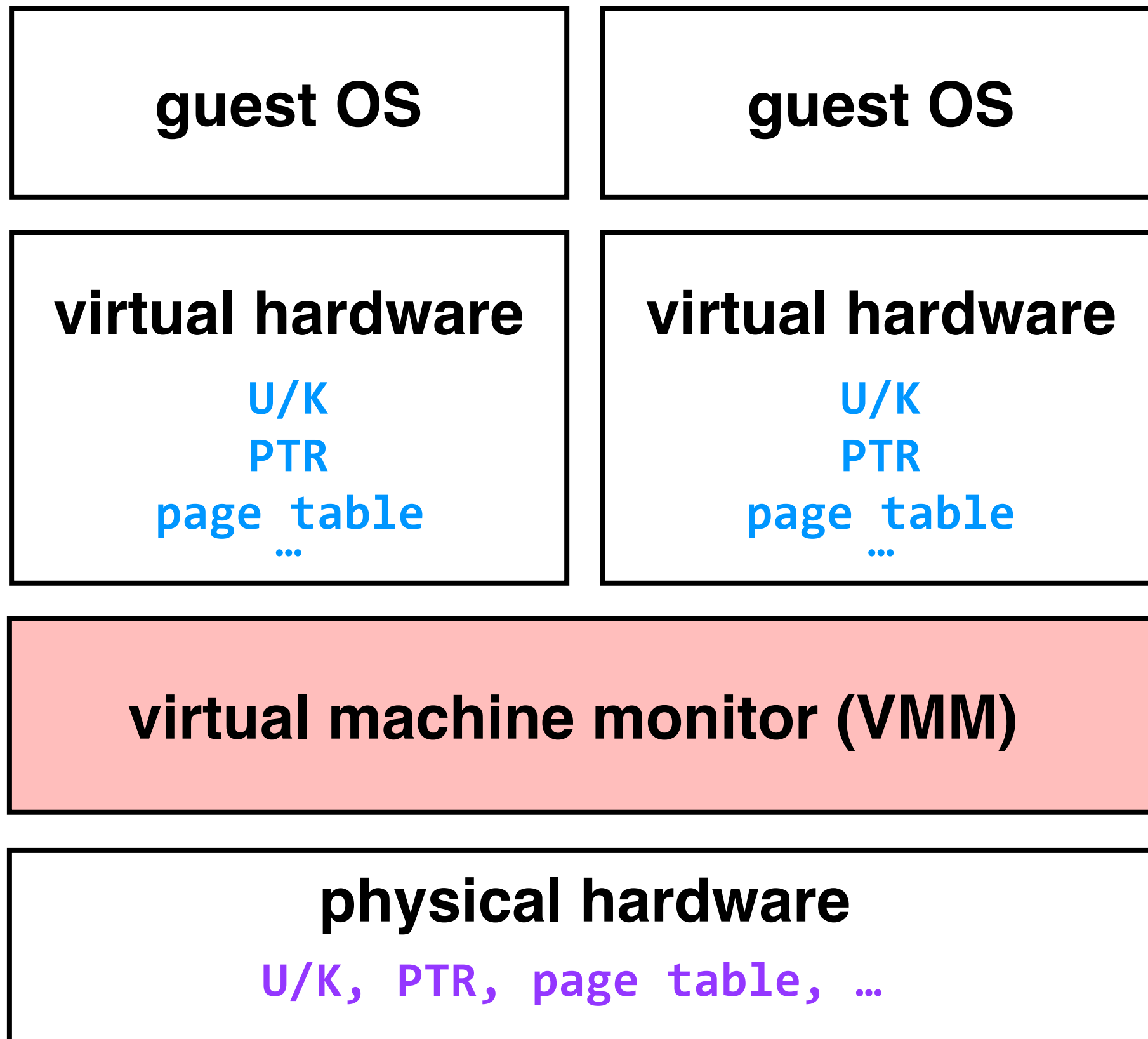
U/K, PTR, page table, ...

VMM's goal: virtualize hardware

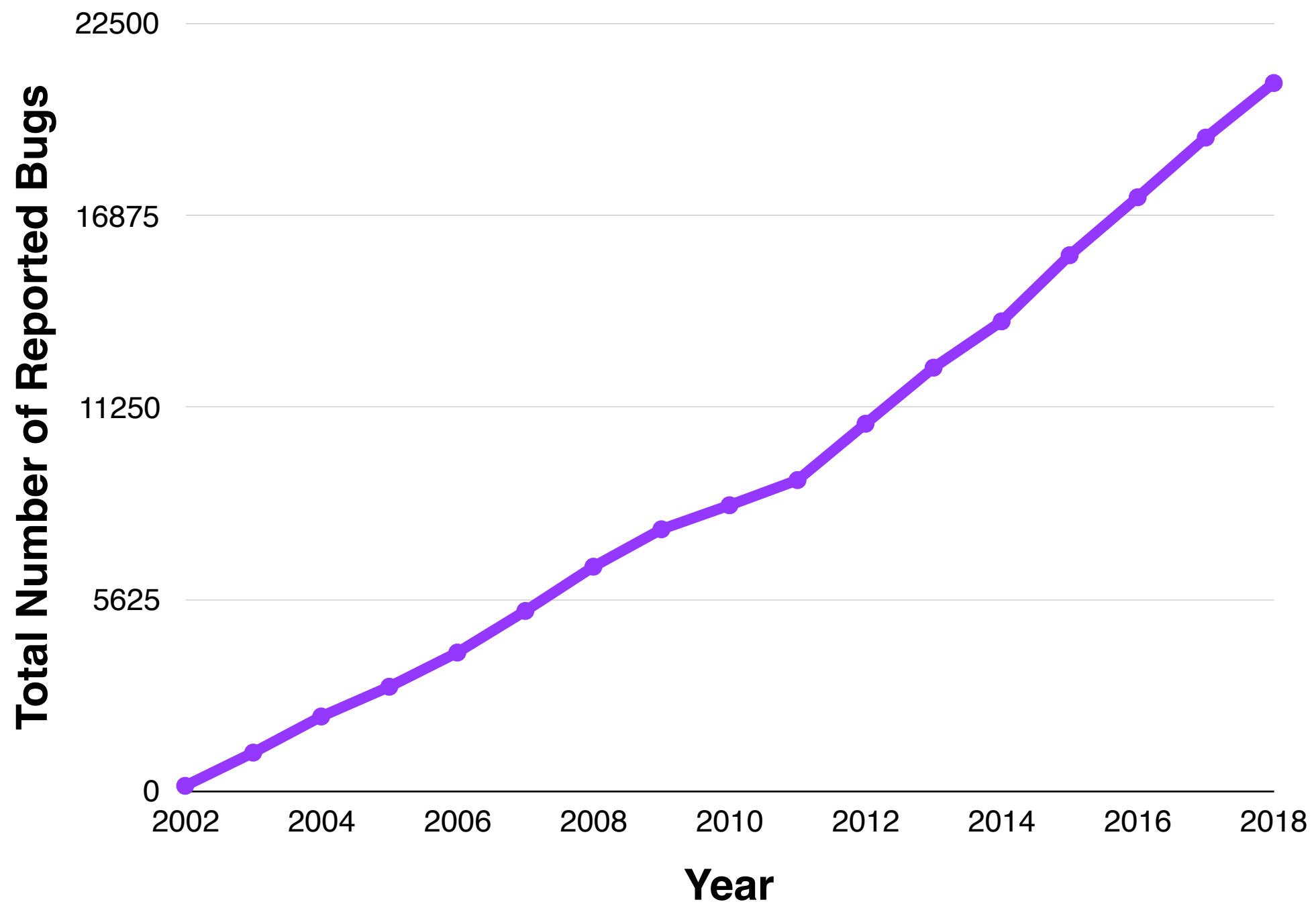






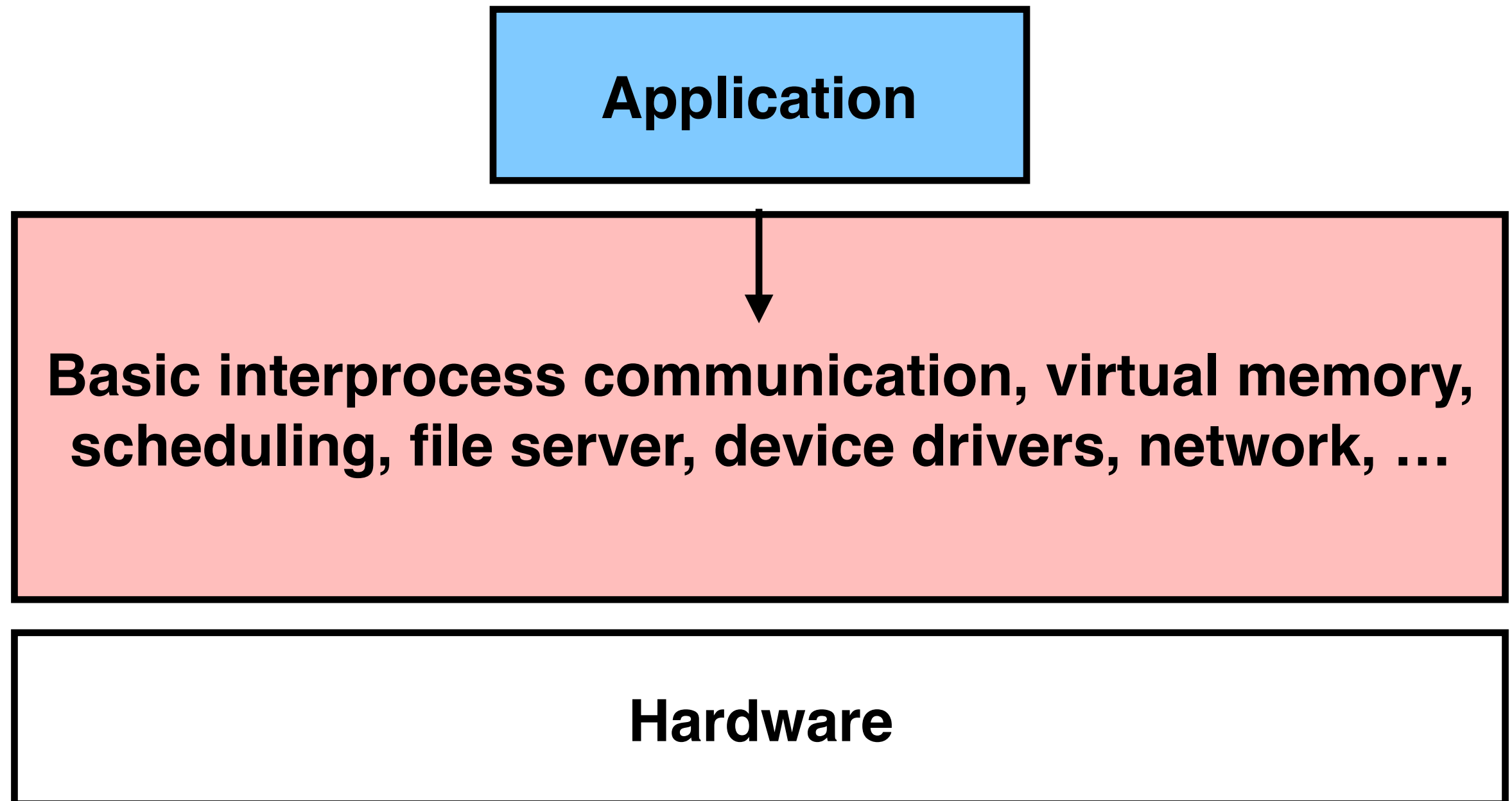


VMM's goal: virtualize hardware

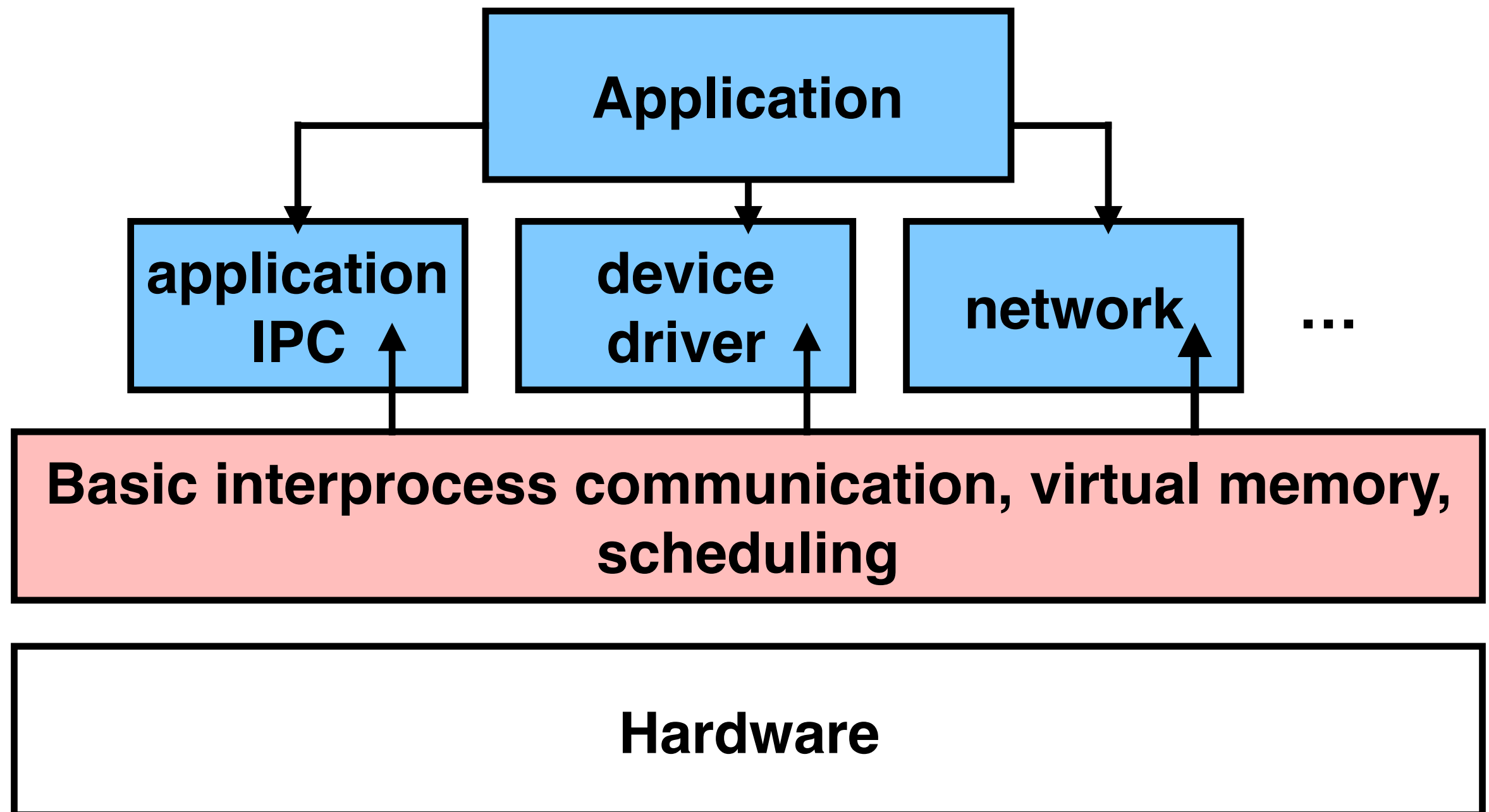


source: bugzilla.kernel.org, count of all bugs currently marked NEW, ASSIGNED, REOPENED, RESOLVED, VERIFIED, or CLOSED, by creation date
(rough estimate!)

monolithic kernels: no enforced modularity within the kernel itself



microkernels: enforce modularity by putting subsystems in user programs



- **Virtual Machines** allow us to run multiple **isolated** OSes on a single physical machine, similar to how we used an OS to run multiple programs on a single CPU. VMs must handle the challenges of virtualizing the hardware (examples: virtualizing memory, the U/K bit).
- **Monolithic kernels** provide no enforced modularity within the kernel. **Microkernels** do, but redesigning monolithic kernels as microkernels is challenging.