



CS 423

Operating System Design: Scheduling

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Goals for Today



- Learning Objective:
 - Introduce goals, definitions, and policies related to uniprocessor and multiprocessor scheduling
 - Reason about advantages and disadvantages of different foundational scheduling algorithms
- Announcements, etc:
 - MP1 is out! Due **Feb 19**



Reminder: Please put away devices at the start of class

What Are Scheduling Goals?



- What are the goals of a scheduler?
- Scheduling Goals:
 - Generate illusion of concurrency
 - Maximize resource utilization (e.g., mix CPU and I/O bound processes appropriately)
 - Meet needs of both I/O-bound and CPU-bound processes
 - Give I/O-bound processes better interactive response
 - Do not starve CPU-bound processes
 - Support Real-Time (RT) applications



Definitions



- **Task/Job**
 - Something that needs CPU time: a thread associated with a process or with the kernel...
 - ... a user request, e.g., mouse click, web request, shell command, ...
- **Latency/response time**
 - How long does a task take to complete?
- **Throughput**
 - How many tasks can be done per unit of time?

Definitions



- **Overhead**
 - How much extra work is done by the scheduler?
- **Fairness**
 - How equal is the performance received by different users?
- **Predictability**
 - How consistent is the performance over time?
- **Starvation**
 - A task 'never' receives the resources it needs to complete
 - Not very fair :- (

Definitions



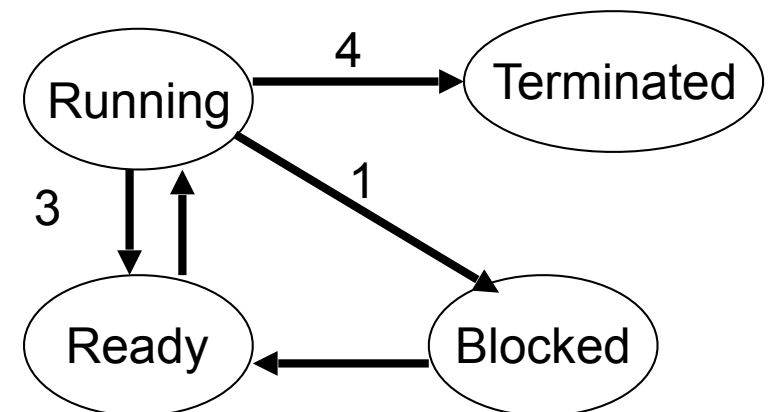
- **Workload**
 - Set of tasks for system to perform
- **Work-conserving**
 - Resource is used whenever there is a task to run
 - For non-preemptive schedulers, work-conserving is not always better

Definitions



- **Non-preemptive scheduling:**

- The running process keeps the CPU until it **voluntarily** gives up the CPU
 - process exits
 - switches to blocked state
 - 1 and 4 only (no 3)



- **Preemptive scheduling:**

- The running process can be interrupted and must release the CPU (can be **forced** to give up CPU)



- Scheduling algorithm
 - takes a workload as input
 - decides which tasks to do first
 - Performance metric (throughput, latency) as output
 - Only preemptive, work-conserving schedulers to be considered

First In First Out (FIFO)



- Schedule tasks in the order they arrive
 - Continue running them until they complete or give up the processor
- **Example: memcached**
 - Facebook cache of friend lists, ...
- On what workloads would FIFO be particularly bad?

First In First Out (FIFO)



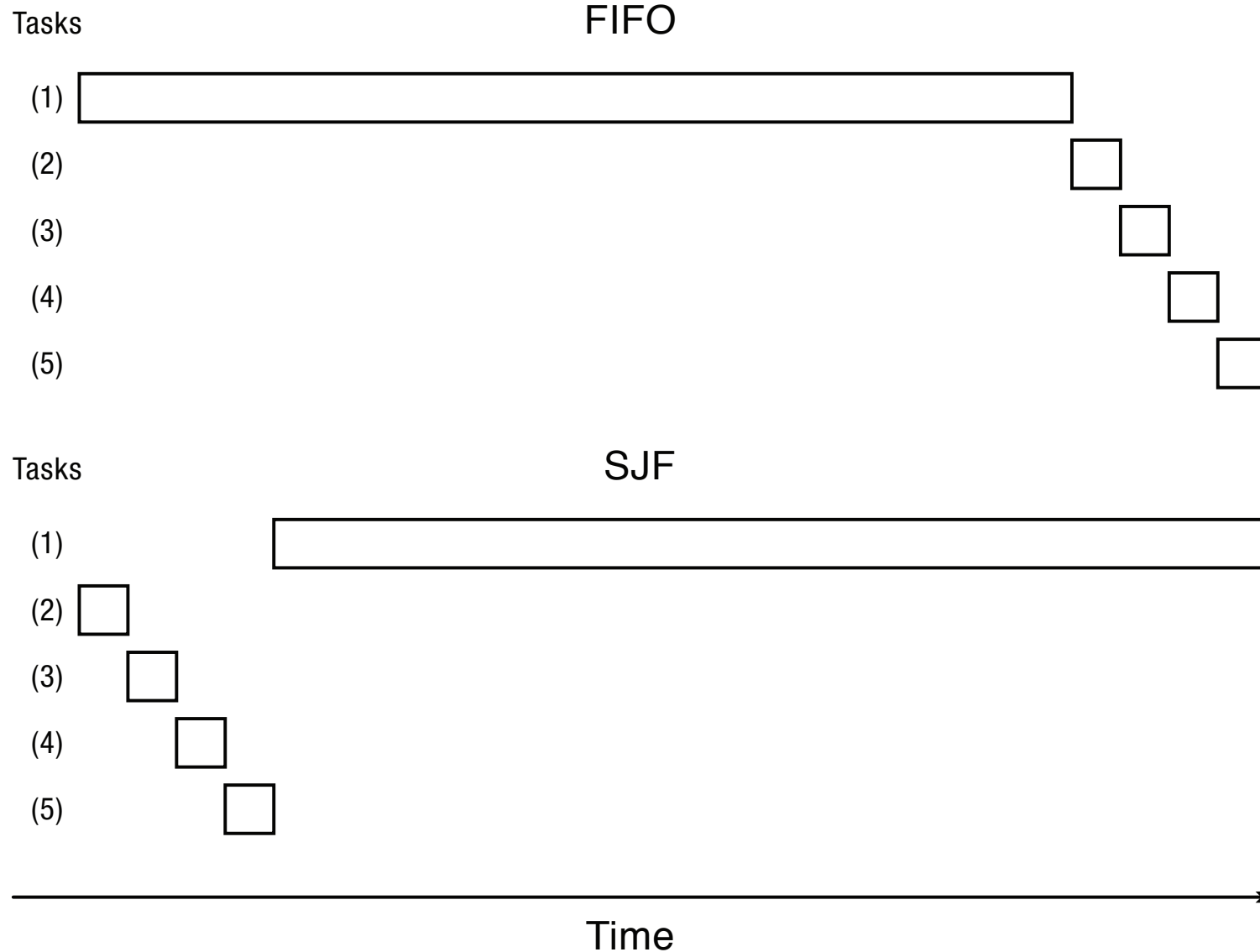
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- **Example: memcached**
 - Facebook cache of friend lists, ...
- On what workloads would FIFO be particularly bad?

Shortest Job First (SJF)



- Always do the task that has the shortest remaining amount of work to do
 - Often called Shortest Remaining Time First (SRTF)
- Suppose we have five tasks arrive one right after each other, but the first one is much longer than the others
 - Which completes first in FIFO? Next?
 - Which completes first in SJF? Next?

FIFO vs. SJF

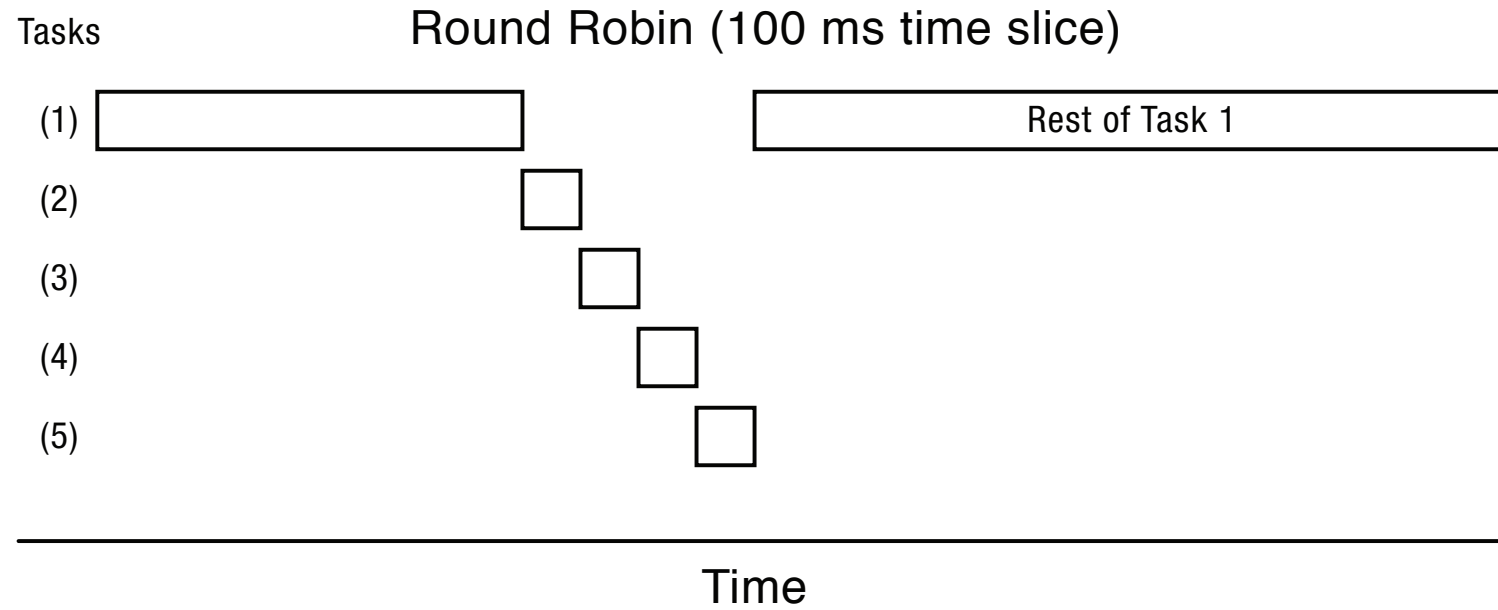


Round Robin (RR)

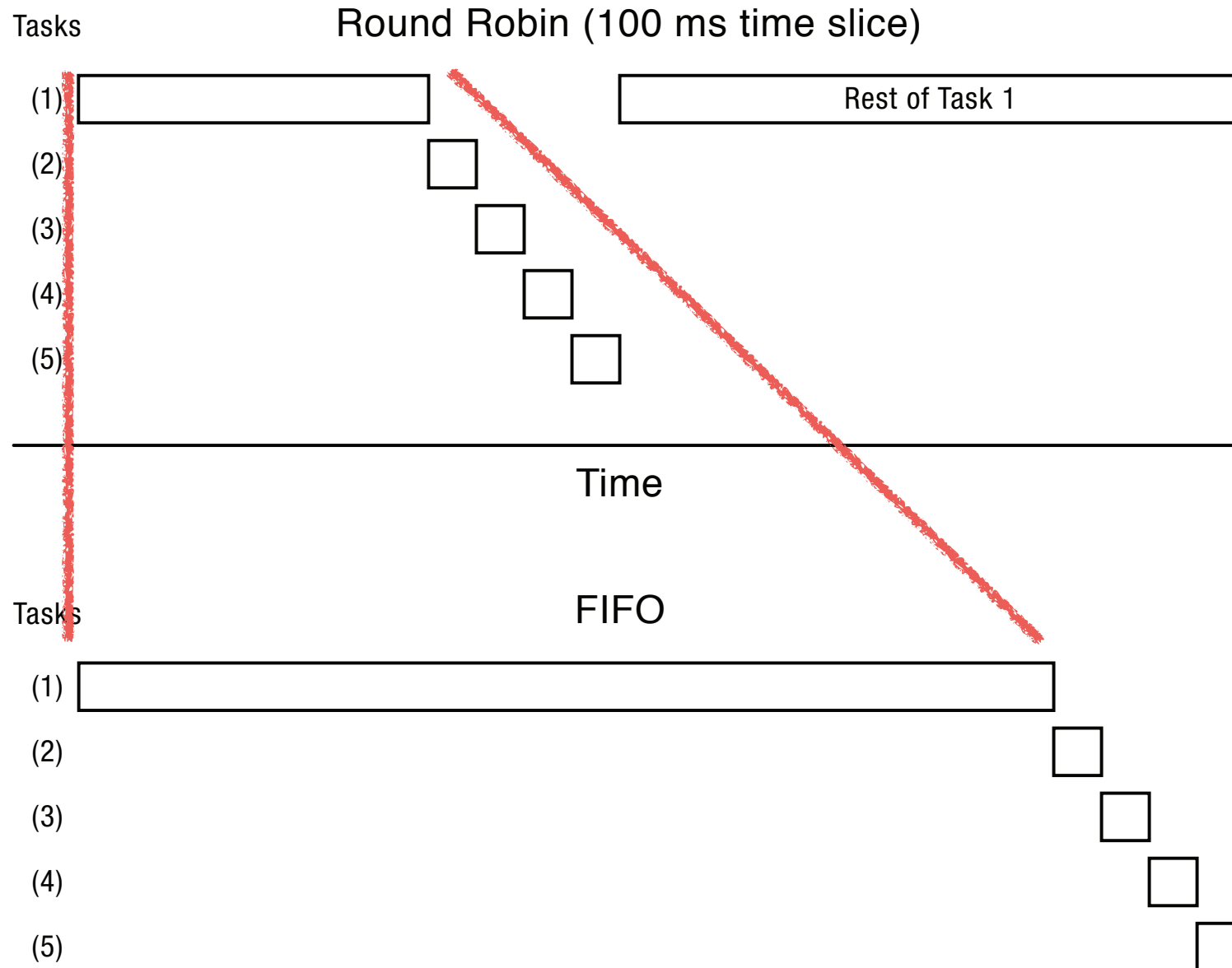


- Each task gets resource for a fixed period of time (time quantum)
 - If task doesn't complete, it goes back in line
- Performance changes based on time quantum size
 - What if time quantum is too short?
 - One instruction?
 - What if time quantum is too long?
 - Infinite?

Round Robin



Round Robin




Round Robin



Tasks

Round Robin (1 ms time slice)

(1) 

 Rest of Task 1

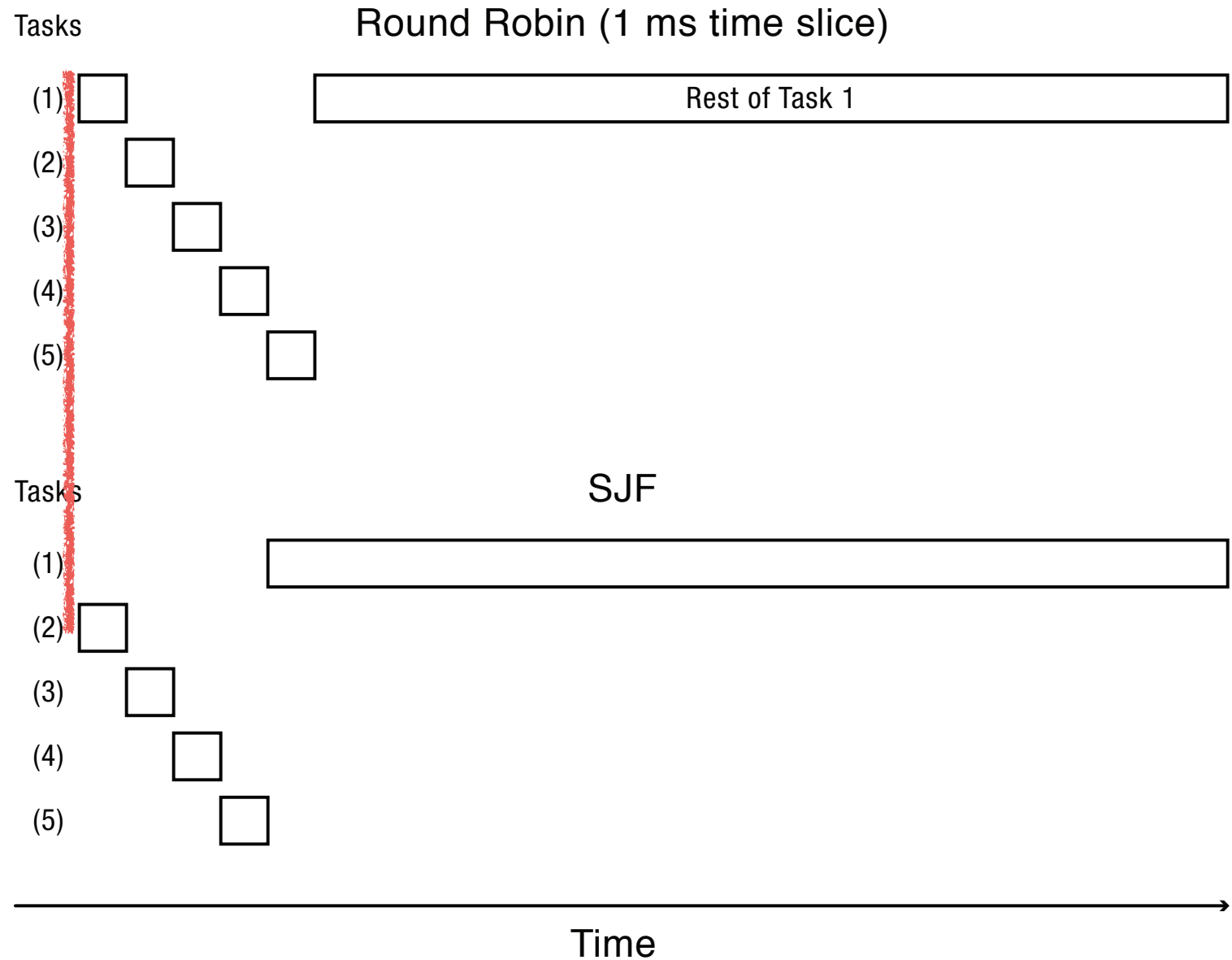
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Round Robin



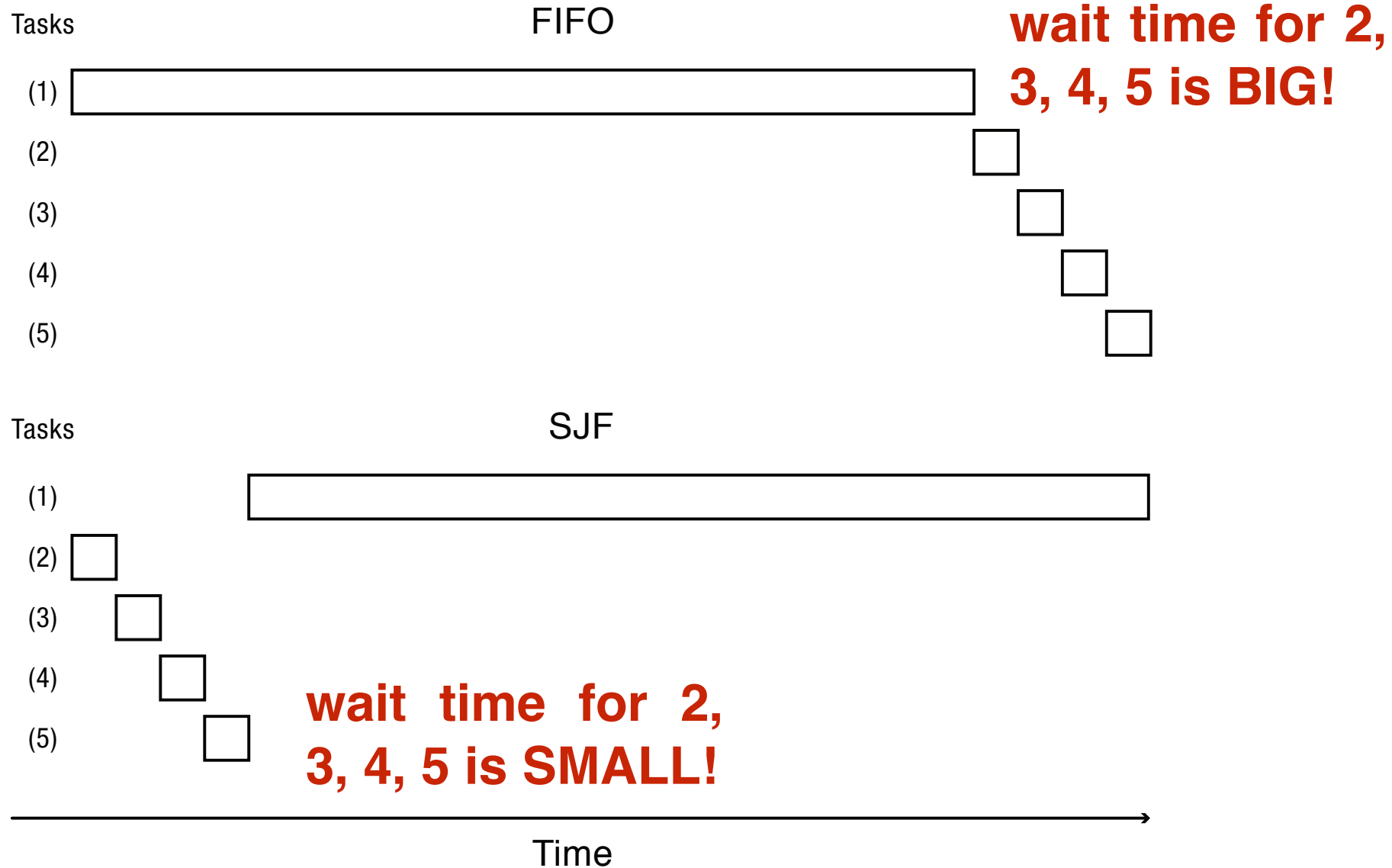


- Basic scheduling algorithms
 - FIFO (FCFS)
 - Shortest job first
 - Round Robin



- Basic scheduling algorithms
 - FIFO (FCFS)
 - Shortest job first
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- What is an optimal algorithm in the sense of maximizing the number of jobs finished (i.e., minimizing average response time)?

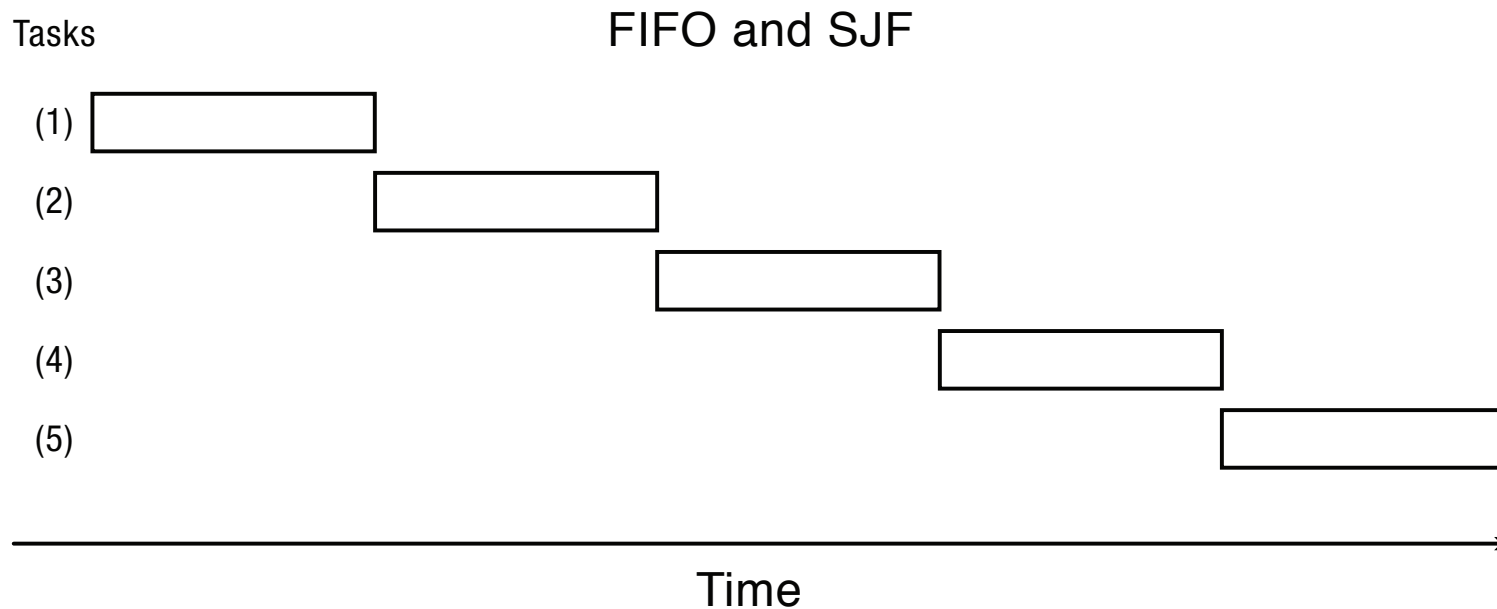
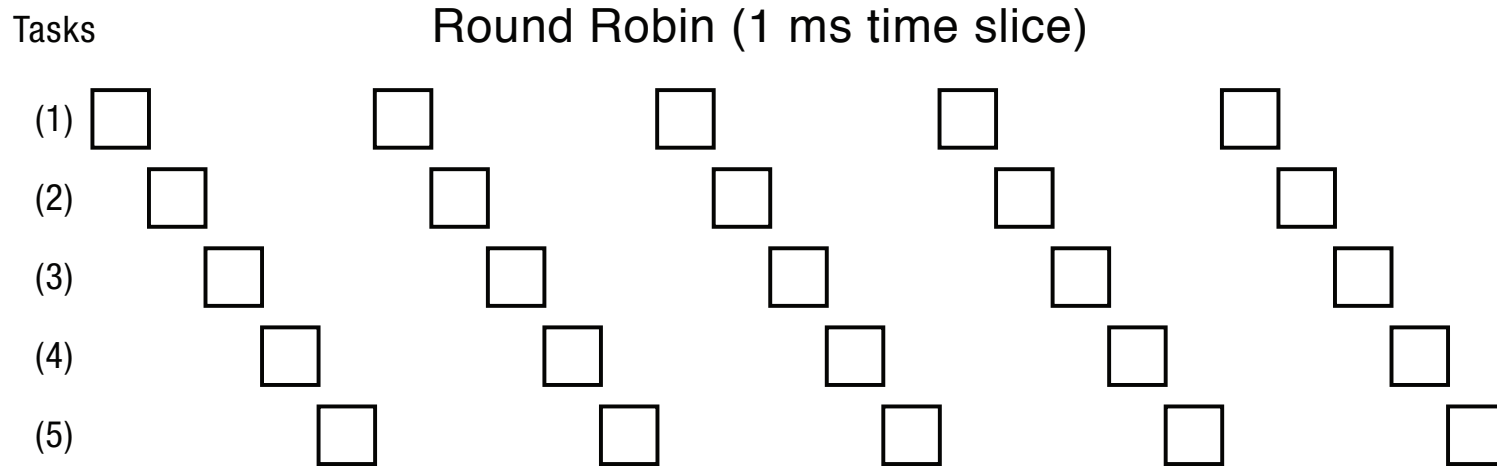
FIFO vs. SJF





- Basic scheduling algorithms
 - FIFO (FCFS)
 - Shortest job first
 - Round Robin
- Assuming zero-cost to time slicing, is Round Robin always better than FIFO?

RR v. FIFO (fixed size tasks)



Starvation, Sample Bias



- Suppose you want to compare two scheduling algorithms
 - Create some infinite sequence of arriving tasks
 - Start measuring
 - Stop at some point
 - Compute average response time as the average for completed tasks between start and stop
- Is this valid or invalid?

Sample Bias Solutions



- Measure for long enough that # of completed tasks \gg # of uncompleted tasks
 - For both systems!
- Start and stop system in idle periods
 - Idle period: no work to do
 - If algorithms are work-conserving, both will complete the same tasks

Round Robin = Fairness?

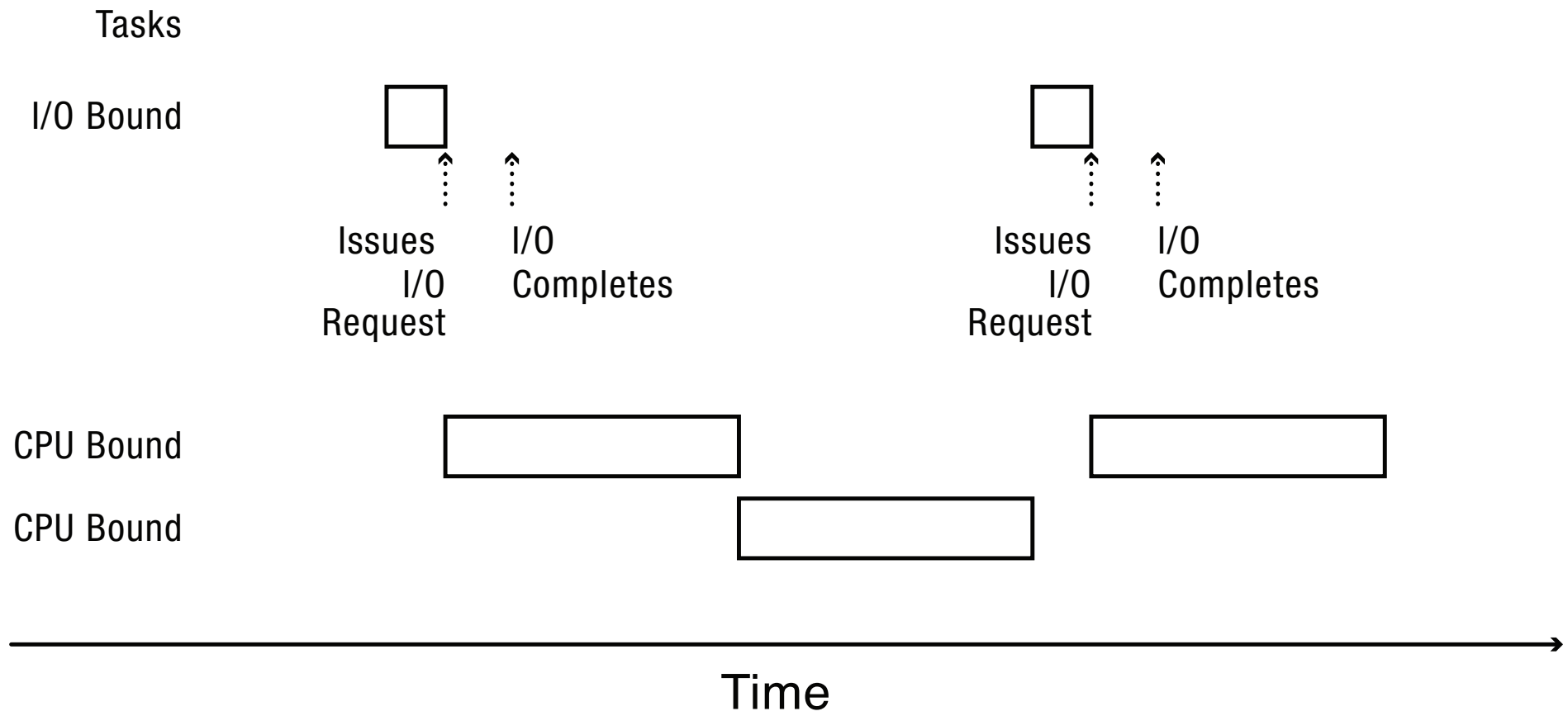


Is Round Robin the fairest possible algorithm?

What is fair?

- FIFO?
- Equal share of the CPU?
- What if some tasks don't need their full share?
- Minimize worst case divergence?
- Time task would take if no one else was running
- Time task takes under scheduling algorithm

Mixed Workloads??



Max-Min Fairness



- How do we balance a mixture of repeating tasks?
 - Some I/O bound, need only a little CPU
 - Some compute bound, can use as much CPU as they are assigned
- One approach: maximize the minimum allocation given to a task
 - If any task needs less than an equal share, schedule the smallest of these first
 - Split the remaining time using max-min
 - If all remaining tasks need at least equal share, split evenly



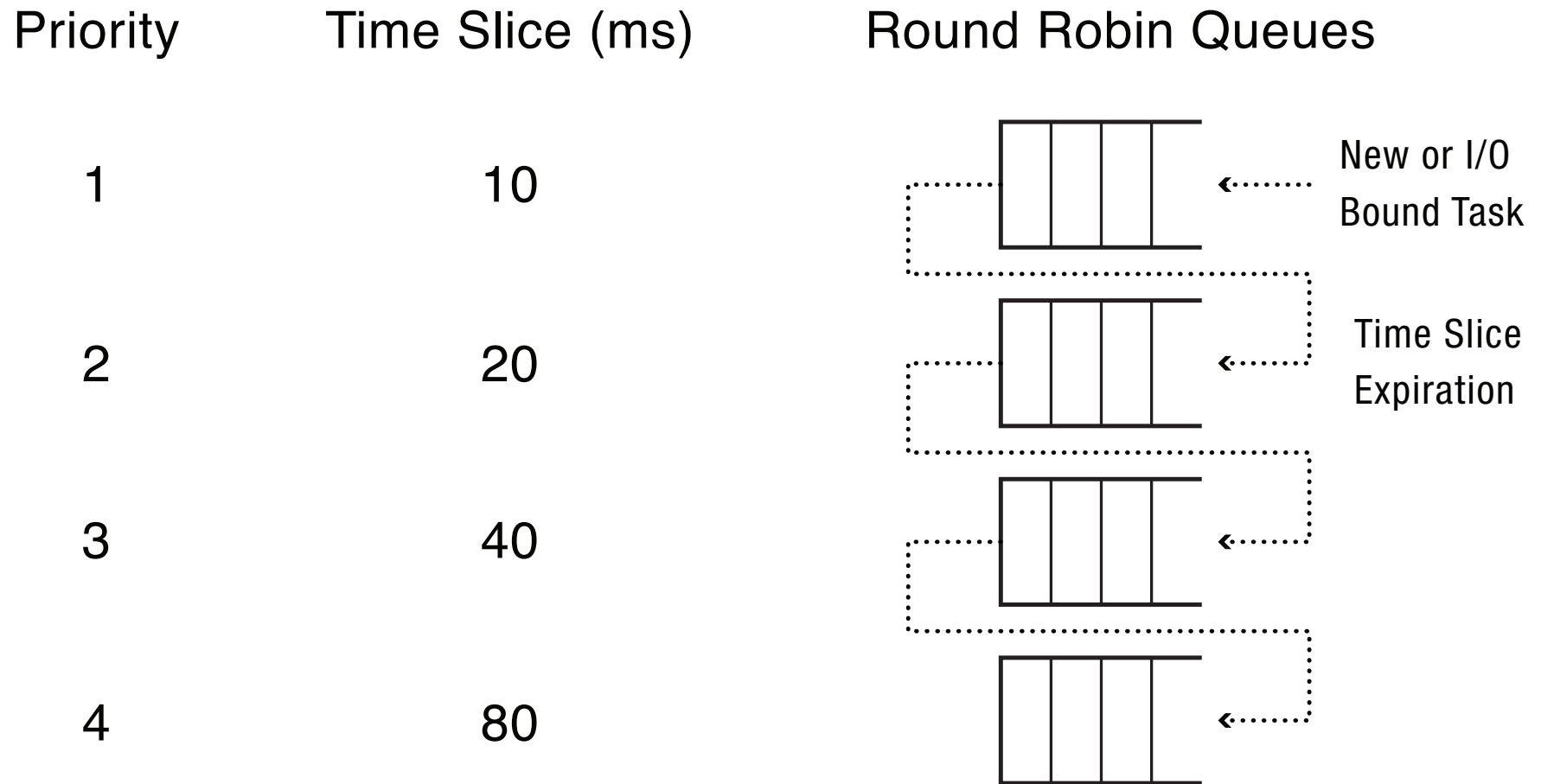
- **Goals:**
 - Responsiveness
 - Low overhead
 - Starvation freedom
 - Some tasks are high/low priority
 - Fairness (among equal priority tasks)
- **Not perfect at any of them!**
 - Used in Linux (and probably Windows, MacOS)

Multi-Level Feedback Queue



- Set of Round Robin queues
 - Each queue has a separate priority
- High priority queues have short time slices
 - Low priority queues have long time slices
- Scheduler picks first thread in highest priority queue
- Tasks start in highest priority queue
 - If time slice expires, task drops one level

Multi-Level Feedback Queue



Summary



- FIFO is simple and minimizes overhead.
- If tasks are variable in size, then FIFO can have very poor average response time.
- If tasks are equal in size, FIFO is optimal in terms of average response time.
- Considering only the processor, SJF is optimal in terms of average response time.
- SJF is pessimal in terms of variance in response time.

Summary



- If tasks are variable in size, Round Robin approximates SJF.
- If tasks are equal in size, Round Robin will have very poor average response time.
- Tasks that intermix processor and I/O benefit from SJF and can do poorly under Round Robin.



- Max-Min fairness can improve response time for I/O-bound tasks.
- Round Robin and Max-Min fairness both avoid starvation.
- By manipulating the assignment of tasks to priority queues, an MFQ scheduler can achieve a balance between responsiveness, low overhead, and fairness.