

Binary Arithmetic

- Subtrahend: what is being subtracted
- Minuend: What it is being subtracted from
- Minuend – Subtrahend = Result
- $0 + 0 = 0$
- $0 + 1 = 1$
- $1 + 0 = 1$
- $1 + 1 = 0$ carry over 1
- $1 + 1 + 1 = 1$ carry over 1

Example: Perform the binary subtraction $1010101 - 11100$

A'		0	1	10				
A'		1	0	0	10			
A		1	0	1	0	1	0	1
B			-	1	1	1	0	0
Diff		0	1	1	1	0	0	1
Step		7	6	5	4	3	2	1

A				1	1	1	0	1
B					x	1	1	1
				1	1	1	0	1
			1	1	1	0	1	
		1	1	1	0	1		
Answer	1	1	0	0	1	0	1	1
Carry	1	10	10	1	1			

- Division = Dividend / Divisor = Quotient + (remainder / divisor)
- Use long division.

Overflow

- When the result of an arithmetic operation is too large or too small to fit into the resultant bit-group (e.g. 9 cannot fit into 4-bits in 2's C)
- Normally left to programmer to deal with this

Two's Complement - Addition

- Add the values
- Discard all carry out bits

Two's Complement – Overflow through Addition

- Occurs if and only if 2 Two's Complement numbers are added and they both have the same sign and the result has the opposite sign
- E.g $(-7) + (-6)$ using 4 bits
- $+7 = 0111$
- $-7 = 1001$
- $+6 = 0110$
- $-6 = 1010$
- $(-7) + (-6) = 1\ 0011 = \text{Overflow} = +3$

Two's Complement – Subtraction

- Negate the subtrahend and add to the minuend
- Discard all carry out bits
- e.g. $8 - 5 = 8 + (-5)$ = using 8 bits
- $+8 = 00001000$
- $+5 = 00000101$
- $-5 = 11111011$
- $1\ 00000011 = 00000011$

Two's Complement – Overflow through Subtraction

- Occurs if and only if 2 Two's Complement numbers are subtracted, and their signs are different, and the result has the same sign as the subtrahend
- e.g. using 4 bits = $7 - (-6)$
- $7 = 0111$
- $6 = 0110$
- $-6 = 1010$

- $-6 = 0110$
- $7 + (-6) = 1101 = \text{Overflow} = -3$

Two's Complement – Multiplication and Division

Cannot be accomplished using the standard technique

Example: consider $X * (-Y)$

- Two's complement of $-Y$ is $2^n - Y \rightarrow X * (-Y) = X * (2^n - Y) = 2^n X - XY$
- Expected result should be $2^{2n} - XY$

Sign extension to 2^{2n} and omission of the leading n bits works but is extremely inefficient

+3 * -2 (4-bit)

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0000 0011 x 1111 1110
0000 0011
0000 0011
0000 0011
0000 0011
0000 0011
0000 0011
0000 0011
0000 0011
0000 0011
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0000 1011 1111 1010

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