

# Arguments and validity

Consider the following argument

$p$ : Chariots of Fire is an Oscar Winning film.

$q$ : All Oscar winning films are great films.

$r$ : Chariots of Fire is a great film.

The argument is not *propositionally* valid... Why?

$$p, q \not\models r$$

Our experience with propositional logic tells us how to define ‘valid argument’ etc.

## Definition 4.11 (valid argument)

Let  $L$  be a signature and  $A_1, \dots, A_n, B$  be  $L$ -formulas.

An argument ‘ $A_1, \dots, A_n$ , therefore  $B$ ’ is *valid* if for any  $L$ -structure  $M$  and assignment  $h$  over  $M$ ,

if  $M, h \models A_1, M, h \models A_2, \dots$ , and  $M, h \models A_n$ , then  $M, h \models B$ .

We write  $A_1, \dots, A_n \models B$  in this case.

This says: in any situation (structure + assignment) in which  $A_1, \dots, A_n$  are all true,  $B$  must be true too.

Special case:  $n = 0$ . Then we write just  $\models B$ . It means that  $B$  is true in every  $L$ -structure under every assignment over it.

## Validity, satisfiability, equivalence

These are defined as in propositional logic. Let  $L$  be a signature.

#### Definition 4.12 (valid formula)

An  $L$ -formula  $A$  is *(logically) valid* if for every  $L$ -structure  $M$  and assignment  $h$  over  $M$ , we have  $M, h \models A$ .

We write ' $\models A$ ' (as above) if  $A$  is valid.

#### Definition 4.13 (satisfiable formula)

An  $L$ -formula  $A$  is *satisfiable* if for some  $L$ -structure  $M$  and assignment  $h$  over  $M$ , we have  $M, h \models A$ .

#### Definition 4.14 (equivalent formulas)

$L$ -formulas  $A, B$  are *logically equivalent* if for every  $L$ -structure  $M$  and assignment  $h$  over  $M$ , we have  $M, h \models A$  if and only if  $M, h \models B$ .

The links between these definitions that we have seen in propositional logic, also hold for predicate logic.

So (e.g.) the notions of valid/satisfiable formula, and equivalence, can all be expressed in terms of valid arguments.

## Which arguments are valid?

Some examples of valid arguments:

- $\forall x(\text{horse}(x) \rightarrow \text{animal}(x)) \models \forall x[\exists y(\text{headof}(x, y) \wedge \text{horse}(y)) \rightarrow \exists y(\text{headof}(x, y) \wedge \text{animal}(y))]$ .

A horse is an animal  $\models$  the head of a horse is the head of an animal.

Deciding if an argument  $A_1, \dots, A_n \models B$  is valid is extremely hard in general. We can't just check that all  $L$ -structures + assignments that make  $A_1, \dots, A_n$  true also make  $B$  true (like truth tables), because there are infinitely many  $L$ -structures (some are infinite!)

### Theorem 4.15 (Church, 1936)

No computer program can be written to identify precisely the valid arguments of predicate logic.

## Useful ways of validating arguments

In spite of Theorem 4.15, we can often verify in practice that a particular argument in predicate logic is valid. Ways to do it include:

- direct argument (the easiest)
- equivalences
- proof systems (like natural deduction)

The same methods work for showing a formula is valid. ( $A$  is valid if and only if  $\models A$ ) Truth tables no longer work (you can't tabulate all structures — there are infinitely many).

## Direct argument Examples

Let's show

$$\begin{aligned} &\forall x(\text{Human}(x) \rightarrow \text{Lecturer}(x)), \forall x(\text{PC}(x) \rightarrow \text{Lecturer}(x)), \\ &\quad \forall x(\text{Human}(x) \vee \text{PC}(x)) \models \forall x \text{Lecturer}(x) \end{aligned}$$

Let's show

$$\forall x(horse(x) \rightarrow animal(x)) \models \forall x[\exists y(headof(x, y) \wedge horse(y)) \rightarrow \exists y(headof(x, y) \wedge animal(y))].$$

Let's show  $\forall x \forall y (x = y \wedge \exists z R(x, z) \rightarrow \exists v R(y, v))$  is valid.

Take any structure  $M$ , and objects  $a, b$  in  $dom(M)$ .

We need to show

$$M \models a = b \wedge \exists z R(a, z) \rightarrow \exists v R(b, v).$$

So we need to show that

$$\text{IF } M \models a = b \wedge \exists z R(a, z) \text{ THEN } M \models \exists v R(b, v).$$

But IF  $M \models a = b \wedge \exists z R(a, z)$ , then  $a, b$  are the same object.

So  $M \models \exists z R(b, z)$ .

So there is an object  $c$  in  $dom(M)$  such that  $M \models R(b, c)$ .

Therefore,  $M \models \exists v R(b, v)$ .

We're done!

*Using direct argument, show that  $f(x, f(y, z)) = f(f(x, y), z)$  is satisfiable.*

All we need to do is find a structure  $M$  and assignment  $h$  over  $M$  in which the above formula is true.

Take an  $M$  where  $\text{dom}(M)$  is the set of natural numbers  $\mathbb{N}$  and the value of  $f$  in  $M$  is addition  $+$ .

Let  $h$  be an assignment in which the values of  $x, y, z$  are 2, 4, 9 respectively.

Then  $M, h \models f(x, f(y, z)) = f(f(x, y), z)$

**Exercise:**

*Using direct argument, show that  $\neg\forall x(\exists y R(x, y) \rightarrow R(x, f(x, x)))$  is satisfiable.*