Machine Program: Data

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Last lecture

- How x86 supports procedure calls
 - call (pushes return address on stack; jump to function)
 - ret (pops return address from stack; jump to return address)
- C/UNIX calling convention (location of args/return val)
 - First 6 args are stored in regs: %rdi, %rsi, %rdx, %rcx, %r8, %r9
 - Rest of arguments are stored on the stack
 - Return value (if there's one) is stored in %rax

Today's lesson plan

- C/UNIX calling convention (caller vs. callee-save reg)
- Program data storage and manipuation

Calling convention: Caller vs. callee-save registers

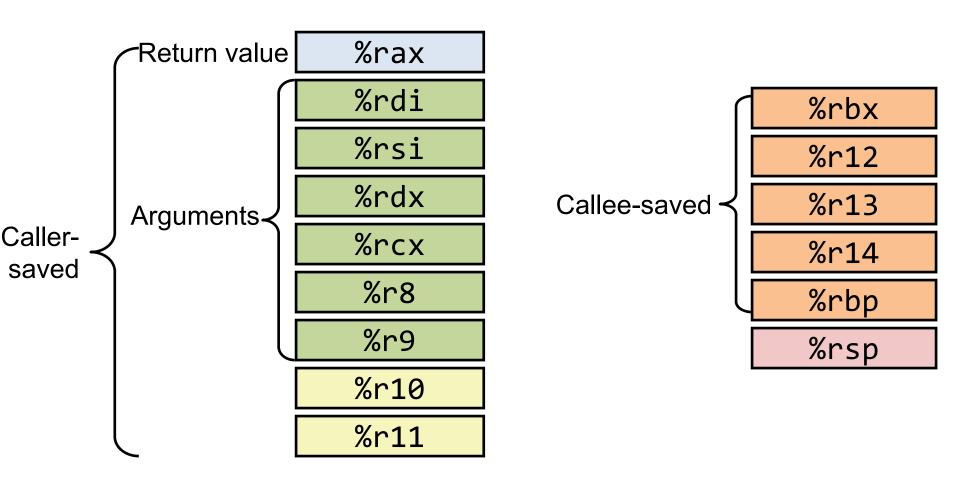
 What can the caller assume about the content of a register across function calls?

```
int foo() {
   int a;    // suppose a is stored in %r12
   a = .... // compute result of a
   int r = bar();
   int result = r + a; // does %r12 still store the value of a?
   return result;
}
```

Calling convention: register saving

- Caller-save registers
 - Register X's value may change across the function call.
 - It is caller's responsibility to save X on the stack and restore X after function returns (if caller needs X's value)
- Callee-save registers
 - Register Y's value must remain unchanged across the function call
 - It is callee's responsibility to save Y on the stack and restore Y before function return (if callee wants to change Y)

Calling convention: Register saving



Callee can use these regs without save/restore

Caller can assume these regs are unchanged across function calls

Example

```
int add2(int a, int b)
                               add2:
                                           (%rdi,%rsi), %eax
                                   leal
 return a + b;
                                   ret
                               add3:
int add3(int a, int b, int c)
                                          %rbx
                                   pushq
                                  movl
                                          %edx, %ebx
 int r = add2(a, b);
                                  movl $0, %eax
 r = r + c;
                                  call add2
 return r;
                                   addl
                                          %ebx, %eax
                                          %rbx
                                   popq
                                   ret
```

Registers

First 6 Arguments: %rdi, %rsi, %rdx, %rcx, %r8, %r9

Return value: %rax

Example

```
int add2(int a, int b)
                                     add2:
                                                   (%rdi,%rsi), %eax
                                         leal
  return a + b;
                                          ret
                                                         save %rbx (callee-save)
                                                         before overwriting it
                                     add3:
int add3(int a, int b, int c)
                                                  %rbx
                                         pushq
                                                  %edx, %ebx
                                         movl
  int r = add2(a, b);
                                                  $0, %eax
                                         movi
  r = r + c;
                                         call
                                                  add2
                                                                  c is copied to %ebx,
  return r;
                                          addl
                                                  %ebx, %eax
                                                                  which is callee save
            %rdx (contains c) is caller save,
                                                   %rbx
                                         popq
               i.e. may be changed by add2
                                         ret
```

Registers

First 6 Arguments: %rdi, %rsi, %rdx, %rcx, %r8, %9

Return value: %rax

restore %rbx before ret

Today's lesson plan

- C/UNIX calling convention (caller vs. callee-save reg)
- Program data storage and manipuation

Local variables

- For primitive data types, use registers whenever possible
- Allocate local array/struct variables on the stack

```
main:
                                                                   array
int main() {
                                                   $48, %rsp
                                        subq
                                                                   allocation
    int a[10];
                                         movl
                                                   $10, %esi
    clear_array(a, 10);
                                                   %rsp, %rdi
                                         movq
    return 0;
                                         call
                                                   clear array
                                                   $0, %eax
                                         movl
                                                                   array
                                         addq
                                                   $48, %rsp
                                                                   de-allocation
                                         ret
```

Global variables

- Allocated in a memory region called "data" segment
 - Statically allocated; compiler determines each global variable's location in data segment.

```
int count = 0;

void inc() {
   count++;
}

int main() {
   inc();
}
```

```
inc:
    addl $0x1, count(%rip)
    ret

main:
    ...
    call add
    movl $0, %eax
    ...
```

Dynamically allocated space

- Allocated in a memory region called "heap"
 - Allocated by malloc library using sophisticated algorithms (discussed in later lecture)

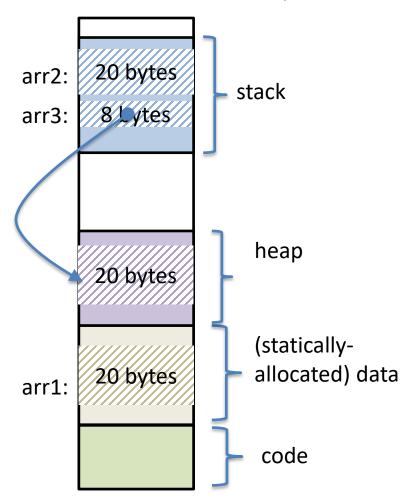
```
int main() {
   int *x;
   x=malloc(100*sizeof(int));
   ...
}
main:
   movl $400 %edi
   call malloc
   ...
```

Lots of code in this function

A process' memory regions

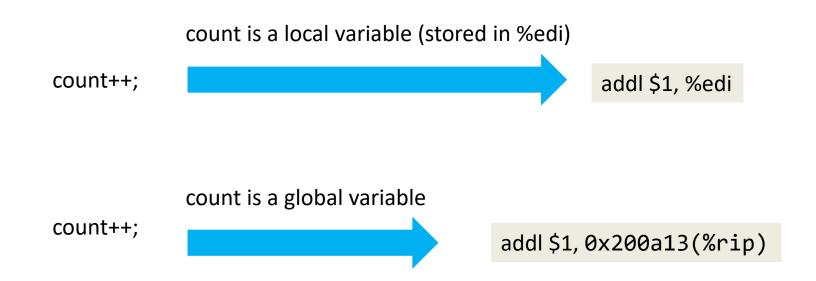
 A running program (process)'s memory consists of code, data, stack, heap (and code/data of its shared libraries)

```
int arr1[5];
int main() {
    int arr2[5];
    int *arr3;
    arr3 = malloc(sizeof(int)*5);
}
```



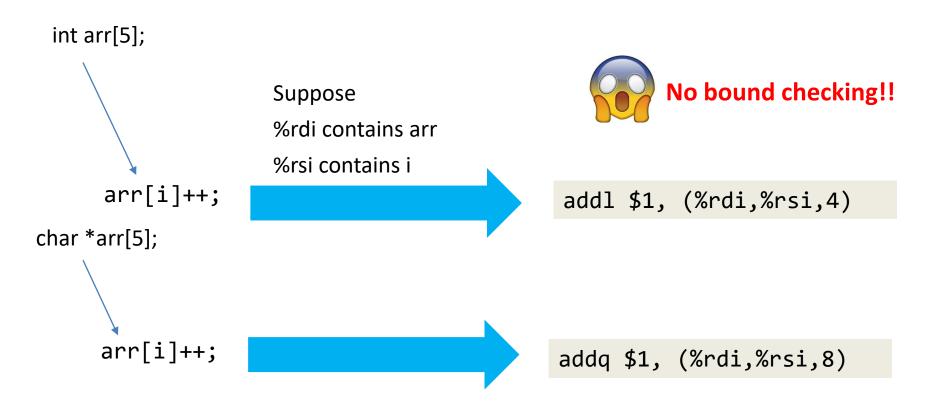
Accessing program data: primitive types

 Local variables of primitive data types are commonly stored in regs



Accessing program data: arrays

 Arrays are always stored in the memory (stack, heap or data)



```
void mystery(int *arr, int n) {
    ???
}
```

```
movl $0, %eax
  jmp .L3
.L4:
  movslq %eax, %rdx
  addl $1, (%rdi,%rdx,4)
  addl $1, %eax
.L3:
  cmpl %esi, %eax
  jl .L4
  ret
```

```
void mystery(int *arr, int n) {
    ???
}
```

```
movl $0, %eax
  jmp .L3
.L4:
  movslq %eax, %rdx
  addl $1, (%rdi, %rdx, 4)
  addl $1, %eax
.L3:
  cmpl %esi, %eax
  jl .L4
  ret
```

```
a = 0;
goto .L3
```

```
void mystery(int *arr, int n) {
    ???
}
```

```
movl $0, %eax
  jmp .L3
.L4:
  movslq %eax, %rdx
  addl $1, (%rdi, %rdx, 4)
  addl $1, %eax
.L3:
  cmpl %esi, %eax
  jl .L4
  ret
```

```
a = 0;
goto .L3

.L3:
    if a < n
        goto .L4

    return</pre>
```

```
void mystery(int *arr, int n) {
    ???
}
```

```
movl $0, %eax
  jmp .L3
.L4:
  movslq %eax, %rdx
  addl $1, (%rdi,%rdx,4)
  addl $1, %eax
.L3:
  cmpl %esi, %eax
  jl .L4
  ret
```

```
a = 0;
goto .L3
.L4
arr[a] = arr[a] + 1
a++

.L3:
if a < n
goto .L4
return</pre>
```

```
void mystery(int *arr, int n) {
  for( int i = 0; i < n; i++)
  {
    arr[i] = arr[i] + 1;
  }
}</pre>
```

```
movl $0, %eax
  jmp .L3
.L4:
  movslq %eax, %rdx
  addl $1, (%rdi,%rdx,4)
  addl $1, %eax
.L3:
  cmpl %esi, %eax
  jl .L4
  ret
```

```
a = 0;
goto .L3
.L4
arr[a] = arr[a] + 1
a++
.L3:
if a < n
    goto .L4
return</pre>
```

```
?? mystery(char *s) {
     ???
}
```

%rdi contains s

```
movl $0x0,%eax
jmp L1.
L2.
  addl $0x1,%eax
L1.
  movslq %eax,%rdx
  cmpb $0x0,(%rdi,%rdx,1)
  jne L2.
  ret
```

```
?? mystery(char *s) {
     ???
}
```

%rdi contains s

```
movl $0x0,%eax
jmp L1.
L2.
   addl $0x1,%eax
L1.
   movslq %eax,%rdx
   cmpb $0x0,(%rdi,%rdx,1)
   jne L2.
   ret
```

int a = 0;
goto L1;

```
?? mystery(char *s) {
     ???
}
```

%rdi contains s

```
movl $0x0,%eax
jmp L1.
L2.
   addl $0x1,%eax
L1.
   movslq %eax,%rdx
   cmpb $0x0,(%rdi,%rdx,1)
   jne L2.
   ret
```

```
int a = 0;
goto L1;
```

L1. long
$$d = a$$
;

```
?? mystery(char *s) {
     ???
}
```

%rdi contains s

```
goto L1;

L1.
   long d = a;
   if(0 != s[d])
      goto L2;
```

int a = 0;

```
?? mystery(char *s) {
     ???
}
```

%rdi contains s

```
movl $0x0,%eax
jmp L1.
L2.
   addl $0x1,%eax
L1.
   movslq %eax,%rdx
   cmpb $0x0,(%rdi,%rdx,1)
   jne L2.
   ret
```

```
int a = 0;
  goto L1;
L2.
  a = a + 1;
L1.
  long d = a;
  if(0 != s[d])
    goto L2;
```

```
int mystery(char *s) {
    int a = 0;
    while(s[a]) {
        a = a + 1;
    }
    return a;
}
```

%rdi contains s

```
int a = 0;
  goto L1;
L2.
  a = a + 1;
L1.
  long d = a;
  if(0 != s[d]) {
     goto L2;
  }
  ret;
```

Accessing Program Data: struct

- Struct is stored in the memory
 - Fields are contiguous in the order they are declared in struct
 - There may be padding (gaps) between fields

```
typedef struct node {
    long id;
                                     id
                                                        next
                                              name
    char *name;
                                                    16
                                                             2.4
    struct node *next;
}node;
                                                   $10, (%rdi)
                     %rdi contains n
                                         movq
n->id = 10;
                                                   $0, 8(%rdi)
                                         movq
n->name = NULL;
n-next = n;
                                                   %rdi, 16(%rdi)
                                         movq
```

```
?? mystery(node *n, long id) {
    ???
}
```

```
.L1
 jmp
.L3:
       %rsi, (%rdi)
 cmpq
 jne
       .L2
       8(%rdi), %rax
 mova
 ret
.L2:
         16(%rdi), %rdi
 movq
.L1:
 testq %rdi, %rdi
 jne
       .L3
 movq $0, %rax
 ret
```

```
%rdi has the value of n%rsi has the value of id%rax is to contain return value
```

```
?? mystery(node *n, long id) {
    ???
}
```

```
.L1
 jmp
.L3:
        %rsi, (%rdi)
 cmpq
 jne
        .L2
       8(%rdi), %rax
 movq
 ret
.L2:
         16(%rdi), %rdi
 movq
.L1:
 testq %rdi, %rdi
 jne
        .L3
      $0, %rax
 movq
 ret
```

goto .L1

%rdi has the value of n%rsi has the value of id%rax is to contain return value

```
?? mystery(node *n, long id) {
    ???
}
```

```
.L1
 jmp
.L3:
       %rsi, (%rdi)
 cmpq
 jne
        .L2
       8(%rdi), %rax
 movq
 ret
.L2:
         16(%rdi), %rdi
 movq
.L1:
 testq %rdi, %rdi
 jne
       .L3
      $0, %rax
 movq
 ret
```

goto .L1

```
.L1:
if (n != 0)
goto .L3
```

%rdi has the value of n%rsi has the value of id%rax is to contain return value

```
?? mystery(node *n, long id) {
    ???
}
```

```
.L1
 jmp
.L3:
       %rsi, (%rdi)
 cmpq
 jne
       .L2
      8(%rdi), %rax
 movq
 ret
.L2:
         16(%rdi), %rdi
 mova
.L1:
 testq %rdi, %rdi
 jne
       .L3
 movq $0, %rax
 ret
```

```
.L1:
    if (n != 0)
        goto .L3
    return 0;
```

goto .L1

```
%rdi has the value of n%rsi has the value of id%rax is to contain return value
```

```
?? mystery(node *n, long id) {
   ???
}
```

```
.L1
 jmp
.L3:
       %rsi, (%rdi)
 cmpq
 jne
       .L2
       8(%rdi), %rax
 movq
 ret
.L2:
       16(%rdi), %rdi
 mova
.L1:
 testq %rdi, %rdi
 jne
       .L3
 movq $0, %rax
 ret
```

```
goto .L1
.L3:
    if (*((long *)n) != id)
        goto .L2

.L1:
    if (n != 0)
        goto .L3
    return 0;
```

```
%rdi has the value of n%rsi has the value of id%rax is to contain return value
```

```
?? mystery(node *n, long id) {
    ???
}
```

```
.L1
 jmp
.L3:
       %rsi, (%rdi)
 cmpq
 jne
       .L2
      8(%rdi), %rax
 movq
 ret
.L2:
       16(%rdi), %rdi
 mova
.L1:
 testq %rdi, %rdi
 jne
       .L3
 movq $0, %rax
 ret
```

```
goto .L1;
.L3:
 if (n->id != id)
     goto .L2;
  return n->name;
.L1:
 if (n != 0)
     goto .L3;
 return 0;
```

```
%rdi has the value of n%rsi has the value of id%rax is to contain return value
```

```
?? mystery(node *n, long id) {
    ???
}
```

```
.L1
 jmp
.L3:
       %rsi, (%rdi)
 cmpq
 jne
       .L2
      8(%rdi), %rax
 movq
 ret
.L2:
      16(%rdi), %rdi
 mova
.L1:
 testq %rdi, %rdi
 jne
       .L3
 movq $0, %rax
 ret
```

```
goto .L1;
.L3:
  if (n->id != id)
     goto .L2;
   return n->name;
.L2
    n = n-next;
.L1:
  if (n != 0)
     goto .L3;
  return 0;
```

```
%rdi has the value of n%rsi has the value of id%rax is to contain return value
```

```
char *mystery(node *n, long id) {
   while (n) {
     if (n->id == id)
        return n->name;
     n= n->next;
   }
   return NULL;
}
```

```
.L1
 jmp
.L3:
 cmpq %rsi, (%rdi)
 jne
       .L2
 movq 8(%rdi), %rax
 ret
.L2:
         16(%rdi), %rdi
 mova
.L1:
 testq %rdi, %rdi
 jne
      .L3
 movq $0, %rax
 ret
```

```
goto .L1;
.L3:
   if (n->id != id)
      goto .L2;

   return n->name;
.L2
   n = n->next;
.L1:
   if (n != 0)
      goto .L3;
   return 0;
```

Summary

- How program data is stored and accessed
 - Primitive data types
 - Arrays
 - Structs
- Separate memory regions for stack, heap, data