# C - Functions, Pointers

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#### **Lesson Plan**

- Control Flow
  - goto
- Function & variable storage
- Pointers

#### C's Control flow

- Same as Java
- Conditional:
  - if ... else if... else
  - switch
- Loops: while, do while, for
  - continue
  - break

# goto statements allow jump anywhere

```
while (cond) {
       Any control flow primitive can be
       expressed as a bunch of goto's.
A:
 if (cond = false) goto B;
 goto A
B:
```

```
There's no goto in Java or core Python
```

```
for(...) {
   for(...) {
       for(...) {
           goto error
error:
    code handling error
```

The only acceptable scenario for using goto

There's no try/catch or try/except in C

# Avoid goto's whenever possible

#### Edgar Dijkstra: Go To Statement Considered Harmful

#### Go To Statement Considered Harmful

Key Words and Phrases: go to statement, jump instruction, branch instruction, conditional clause, alternative clause, repetitive clause, program intelligibility, program sequencing CR Categories: 4.22, 5.23, 5.24

#### EDITOR:

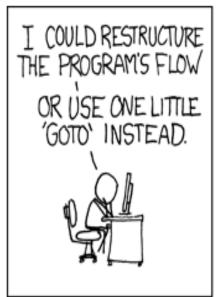
For a number of years I have been familiar with the observation that the quality of programmers is a decreasing function of the density of go to statements in the programs they produce. More recently I discovered why the use of the go to statement has such disastrous effects, and I became convinced that the go to statement should be abolished from all "higher level" programming languages (i.e. everything except, perhaps, plain machine code). At that time I did not attach too much importance to this discovery; I now submit my considerations for publication because in very recent discussions in which the subject turned up, I have been urged to do so.

My first remark is that, although the programmer's activity ends when he has constructed a correct program, the process taking place under control of his program is the true subject matter of his activity, for it is this process that has to accomplish the desired effect; it is this process that in its dynamic behavior has to satisfy the desired specifications. Yet, once the program has been made, the "making" of the corresponding process is deledynamic progress is only characterized when we also give to which call of the procedure we refer. With the inclusion of procedures we can characterize the progress of the process via a sequence of textual indices, the length of this sequence being equal to the dynamic depth of procedure calling.

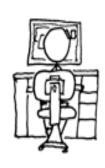
Let us now consider repetition clauses (like, while B repeat A or repeat A until B). Logically speaking, such clauses are now superfluous, because we can express repetition with the aid of recursive procedures. For reasons of realism I don't wish to exclude them: on the one hand, repetition clauses can be implemented quite comfortably with present day finite equipment; on the other hand, the reasoning pattern known as "induction" makes us well equipped to retain our intellectual grasp on the processes generated by repetition clauses. With the inclusion of the repetition clauses textual indices are no longer sufficient to describe the dynamic progress of the process. With each entry into a repetition clause, however, we can associate a so-called "dynamic index," inexorably counting the ordinal number of the corresponding current repetition. As repetition clauses (just as procedure calls) may be applied nestedly, we find that now the progress of the process can always be uniquely characterized by a (mixed) sequence of textual and/or dynamic indices.

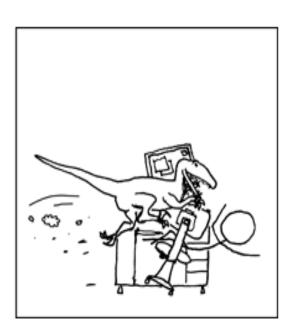
The main point is that the values of these indices are outside programmer's control; they are generated (either by the write-up of his program or by the dynamic evolution of the process) whether

# Avoid goto's whenever possible









# C is a procedural-language

- C program consists of functions
  - Also called procedures or subroutines
- Why breaking code into functions?
  - Readability and Reusability
- Keep functions short
  - General rules of thumb:
    - Small enough to keep the code in your head
    - Fits on screen without scrolling

# Variable and its storage: local variable

# Local variables are defined inside a function

You can view function arguments as local variables defined in function body

```
int add(int a, int b)
{
  int r = a + b;
  r's scope is in function add
  return r;
}
```

#### Local variable scope:

- Within the function/block the local variable is declared
- Local variables with the same name in different scopes are unrelated
- Shadowing: Nested variable "hides" the outer variable in areas where they are both in scope

```
int add(int a, int b)
{
    int r = a;
    {
        int r;
        r += b;
    }
    return r;
}
```



What does add compute?

#### Variable and its storage: local variable scope

```
int find(int start, int end)
{
    for (int i = 0; i < a; i++) {
        if ((i % 3) == 0)
            break;
    }
    return i;
}</pre>
```

```
$gcc -std=c99 test.c
```



```
test3.c: In function 'find':
test3.c:8:10: error: 'i' undeclared (first use in this function)
return i;
^
```

#### Variable and its storage: local variable scope

```
int find(int start, int end)
{
    int r;
    for (int i = 0; i < a; i++) {
        if ((i % 3) == 0) {
           r = i;
           break;
    return r;
```

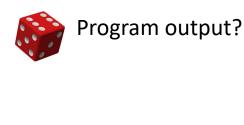


#### Variable and its storage: local variable storage

#### Local variable storage:

- allocated upon function invocation
- deallocated upon function return

```
void add(int a, int b, int result)
    int result = a + b;
    return;
int main()
    int result;
    add(1, 2, result);
    printf("r=%d\n", result);
```



# Variable and its storage: global variable

Global variables are defined outside any function

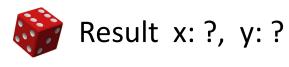
```
#include <stdio.h>
int r = 0;
int sum(int a, int b)
    r = a + b;
int main()
    sum(1,2);
    int r = 0;
    printf("r=%d\n", r);
```

- Global variable scope:
  - Can be accessed from within any function
  - May be shadowed
- Global variable storage
  - Allocated upon program start, deallocated when entire program exits

#### **Function invocation**

C (and Java) passes arguments by value

```
void swap(int a, int b)
{
    int tmp = a;
     a = b;
     b = tmp;
int main()
{
   int x = 1;
    int y = 2;
    swap(x, y);
    printf("x: %d, y: %d", x, y);
}
```



#### **Function invocation**

C (and Java) passes arguments by value

```
void swap(int x, int y)
{
    int tmp = x;
    x = y;
    y = tmp;
int main()
{
   int x = 1;
    int y = 2;
    swap(x, y);
    printf("x: %d, y: %d", x, y);
}
```

```
Result x: ? y: ?
```

#### **Function invocation**

```
void swap(int a, int b)
{
    int tmp = x;
    x = y;
    y = tmp;
int main()
{
   int x = 1;
   int y = 2;
   swap(x, y);
    printf("x: %d, y: %d", x, y);
}
```

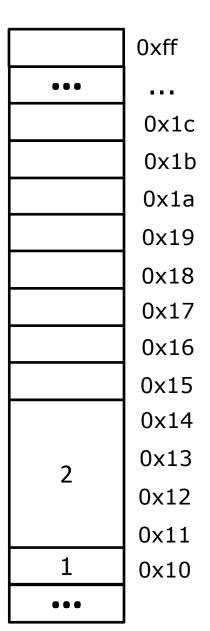
main.x:	
main.y:	
	•••
swap.x:	
swap.y:	
swap.tmp:	

Pointer is a memory address

	0xff
•••	• • •
	0x1c
	0x1b
	0x1a
	0x19
	0x18
	0x17
	0x16
	0x15
	0x14
	0x13
	0x12
	0x11
1	0x10

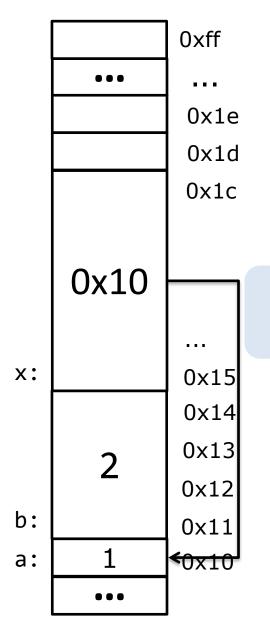
char a = 1;

Addresses should be 8-byte long, but for the same of simplicity, picture only shows one byte



b:

a:



```
char a = 1;
int b = 2;
char *x;
x = &a;
```

```
Same as: char* x;
```

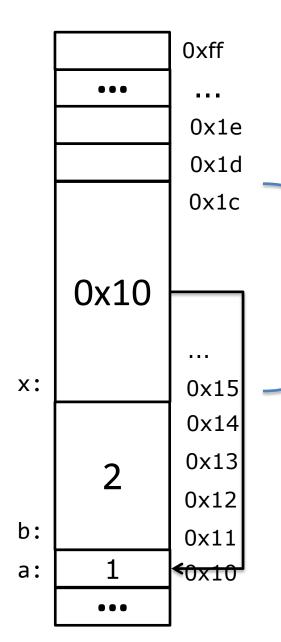
& gives address of variable

Can be combined as: char \*x = &a;

> what happens if I write char x = &a;or

type mismatch!

int \*x = &a;



```
char a = 1;
int b = 2;
char *x = &a;
printf("x=%p\n", x);
```



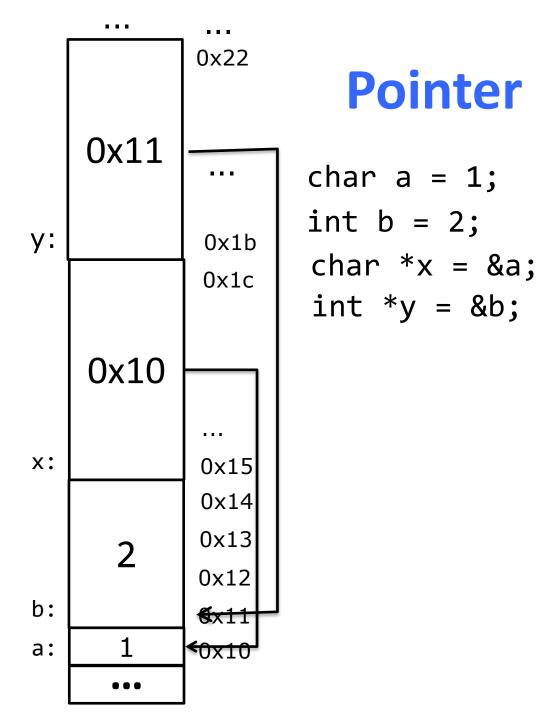
You can print the value of a pointer variable with %p (leading zeros are not printed)

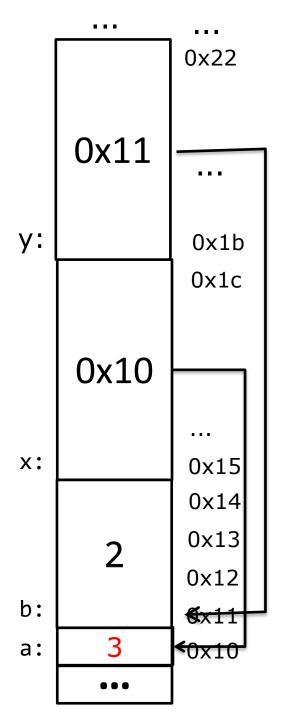
0x10



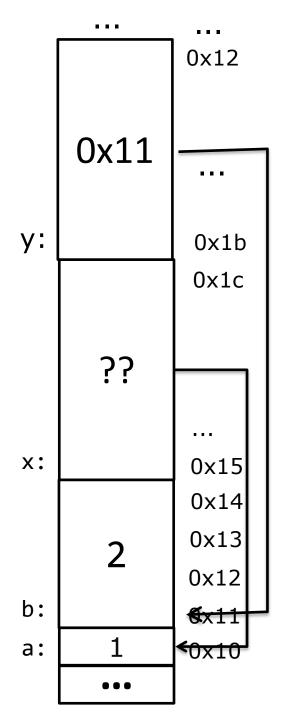
Size of pointer on a 64-bit machine?

8 bytes





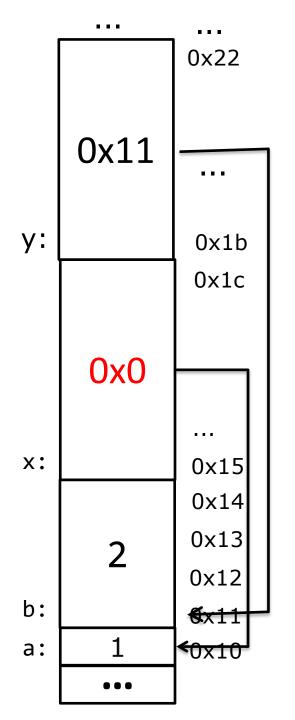
Value of variable a after this statement?



what if x is uninitialized?

$$*x = 3;$$

Dereferencing an arbitrary address value may result in "Segmentation fault" or a random memory write

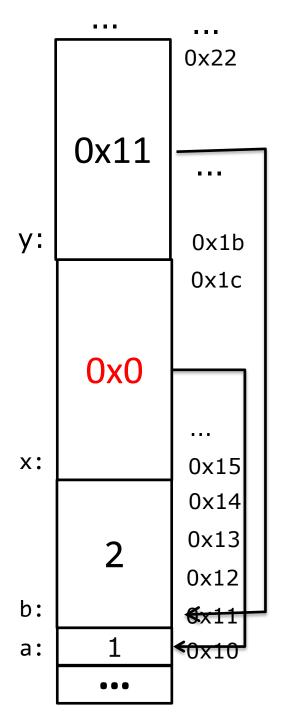


```
char a = 1;
int b = 2;
char *x = NULL;
int *y = &b;
```

Always initialize pointers!

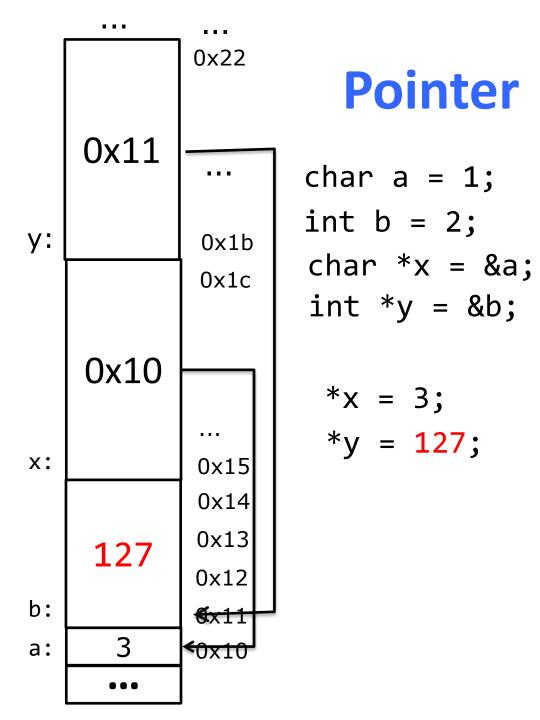
$$*x = 3;$$

Dereferencing NULL pointer definitely results in "Segmentation fault"



```
char a = 1;
int b = 2;
char *x = NULL;
int *y = &b;

*x = 3;
```



## 0x15 XX: 0x22 0x1c 0x10 **x**: **40**×15 0x14 0x13 127 0x12 b: 0x11 3 a: **<**0x10

## **Double Pointer**

```
char a = 1;
int b = 2;
char *x = &a;
int *y = &b;
```

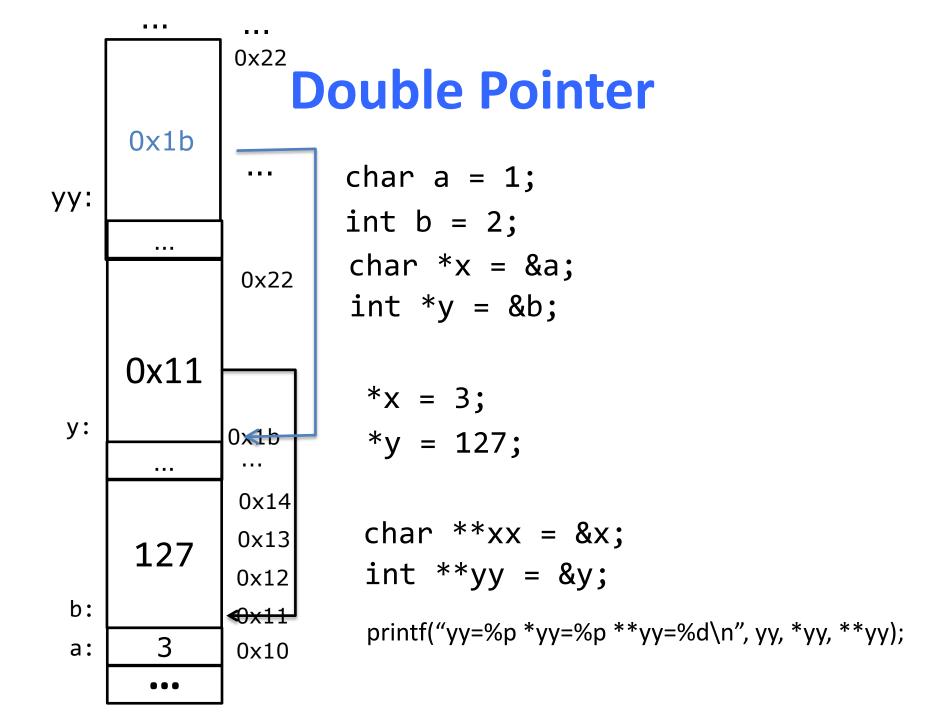
```
char **xx = &x;
```

Same as: char \*\*xx; xx = &x;

char \*\*xx is the same as char\*\* xx;

what if I write char\* xx; xx = &x;

```
printf("xx=%p *xx=%p **xx=%d\n", xx, *xx, **xx);
```



#### Common confusions on \*

- \* has two meanings!!
  - 1. part of a pointer type name, e.g. char \*, char \*\*, int \*
  - 2. the deference operator.

```
char a = 1;
char *p = &a;
*p = 2;

char *b, *c;
char **d,**e;

char *f=p, *g=p;
char **m=&p, **n=&p;
```

C's syntax for declaring multiple pointer variables on one line char\* b, c; does not work

C's syntax for declaring and initializing multiple pointer variables on one line

# Pass pointers to function

```
void swap(int* x, int* y)
                                                         0xf7
                                      main.x:
    int tmp = *x;
                                                         0xf4
    *x = *y;
                                                         0xf3
    *y = tmp;
                                      main.y:
                                                         0xf0
int main()
                                                  ??
    int x = 1;
                                      swap.x:
    int y = 2;
    swap(&x, &y);
    printf("x:%d, y:%d",x,y);
                                                  33
                                      swap.y:
 Size and value of x, y, tmp
                                                  55
 in swap upon function entrance?
                                    swap.tmp:
```

# Summary

- Control Flow
  - goto
- Local vs. global variable
  - Local variables allocated/deallocated upon function entrance/return
  - Global variable: always there
- Pointers
  - Pointer values are memory addresses
  - p =&x; (makes p points to variable x)
  - \*p ... (refers to the variable pointed to by p)