

CSO-Recitation 07

CSCI-UA 0201-007

R07: Assessment 05 & Assembly & lab2

Today's Topics

- Assessment 05
- Assembly
- More about lab2
 - More debugging
 - Some valuable questions asked
- Some exercises

Assembly

C is for people

Why Assembly

- In the real world, computers don't "understand" code
- They only "understand" a set of instructions
- To run code
 - 1. The CPU fetches an instruction from the memory at the PC(program counter)
 - 2. The CPU decodes that instruction
 - 3. If needed, the CPU fetches data from memory
 - 4. The CPU performs computations
 - 5. If needed, the CPU writes data to memory
 - 6. The CPU increments the PC to the next instruction

Why Assembly

- Computers don't "understand" assembly either, but assembly maps much more closely to machine instructions than C code
- Assembly code involves instruction "mnemonics"
 - For x86_64, These are things like `addq`, `movq`, `imul`

X86 general purpose registers

- Accessing memory is very, very slow compared to the rest of what a CPU can do
- Registers are fast temporary storage
- X86-64 ISA: 16 8-byte general purpose registers
- Originally there were 8, all 16-bits large
 - `%ax, %bx, %cx, %dx, %si, %di, %bp, %sp`
 - These have 32-bit counterparts – `add an e`, eg `%eax, %esp`
 - These also have 64-bit counterparts – `add an r`, eg `%rax, %rsp`
- With 64 bits came 8 more registers, `%r8` to `%r15`
 - These have 32-bit counterparts - `add a d`, eg `%r8d`
 - These have 16-bit counterparts – `add a w`, eg `%r8w`
- All registers also allow you to access their lowest 8 bits
- `%ax, %bx, %cx`, and `%dx`, allow you to access their upper 8 bits

Important Instructions

Instruction	What it does
mov <code>src</code> , <code>dest</code>	<code>dest = src</code>
add <code>src</code> , <code>dest</code>	<code>dest = dest + src</code>
sub <code>src</code> , <code>dest</code>	<code>dest = dest - src</code>
imul <code>src</code> , <code>dest</code>	<code>dest = dest * src</code>
inc <code>dest</code>	<code>dest = dest + 1</code>

More about lab2

Debugging & Some valuable questions

More on debugging

1. Program received signal SIGSEGV, Segmentation fault.
 - GDB will tell you where your code segfaulted
 - GDB can tell you what values are what
 - why your code segfaulted

Debugging a crash

- *run* your program
- Use *bt* to see the call stack
 - You can also use *where* to see where you were last running
- Use *frame* to go to where your code was last running
- Use *list* to see the code that ran
- Check the locals (*info locals*) and args (*info args*) to see if they are bad

More on debugging

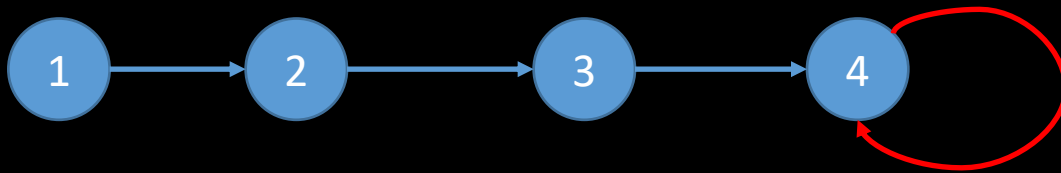
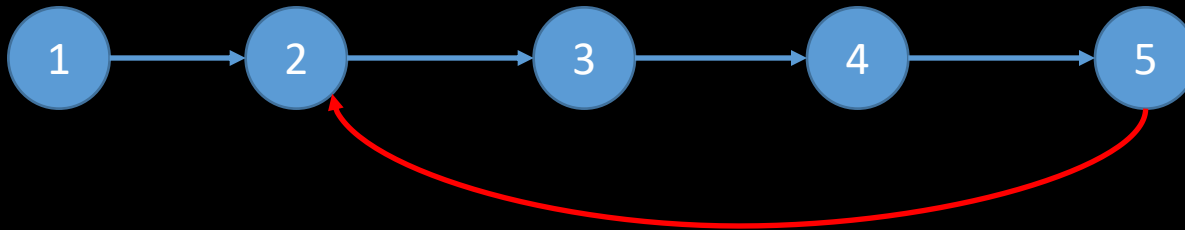
1. Program received signal SIGSEGV, Segmentation fault.
 - GDB will tell you where your code segfaulted
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 - why your code segfaulted
2. Program get stuck
 - infinite loop

Debugging an infinite loop

- Just *run* it inside gdb and hit *control-c* (signal)
- *list* the code
 - This is so you can see the loop condition
- *step* over the code
- Check (*print*) the values involved in the loop condition
 - Are they changing the right way? Are the variables changing at all?

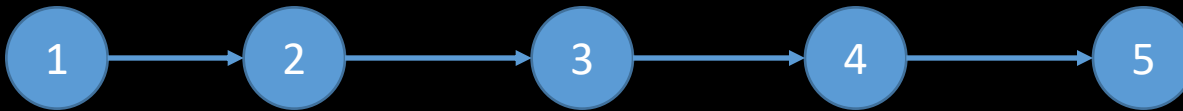
Loop in a linked list

- When I insert a node into the linked list, what will cause a loop?



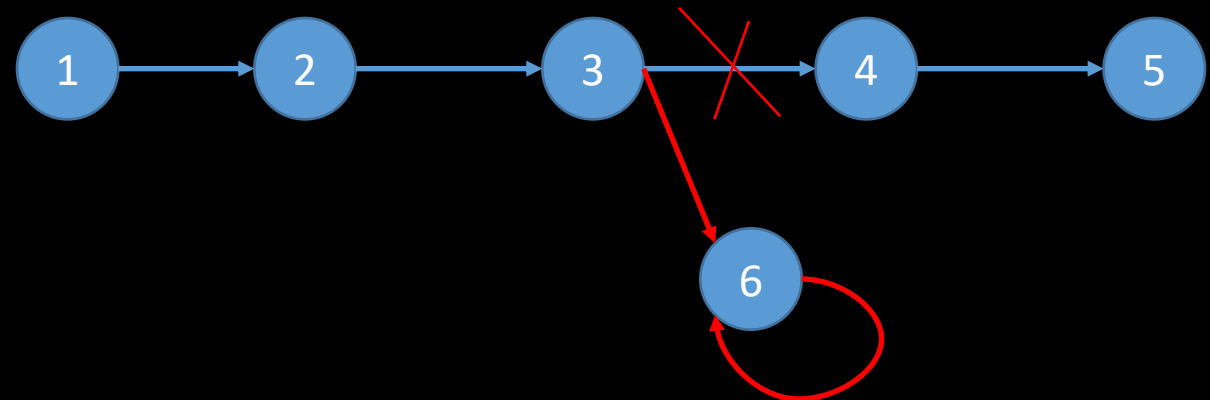
Loop in a linked list

- When I insert a node into the linked list, what will cause a loop?



- I need to insert n6 between n3 and n4:

- suppose our *head* now points to n3
- head -> next = n6
- n6 -> tuple.key = keys
- n6 -> tuple.value = value
- n6 -> next = head->next



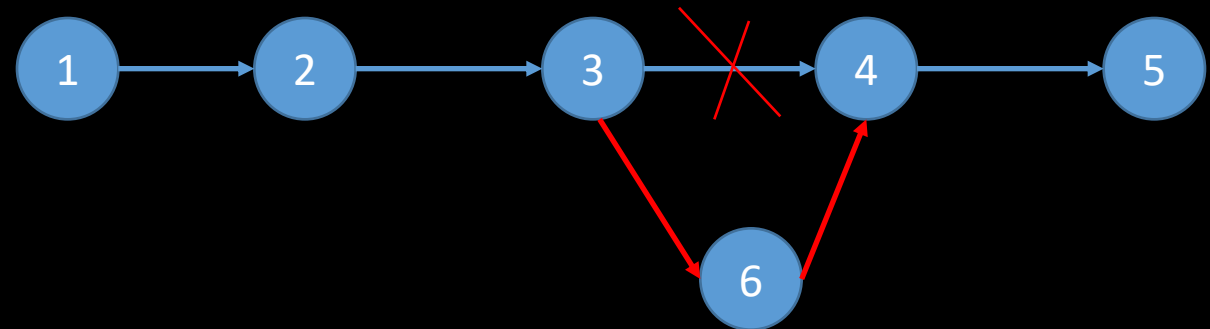
Loop in a linked list

- When I insert a node into the linked list, what will cause a loop?



- I need to insert n6 between n3 and n4:

- suppose our *head* now points to n3
- $n6 \rightarrow \text{next} = \text{head} \rightarrow \text{next}$
- $n6 \rightarrow \text{tuple.key} = \text{keys}$
- $n6 \rightarrow \text{tuple.value} = \text{value}$
- $\text{head} \rightarrow \text{next} = n6$



More on debugging

1. Program received signal SIGSEGV, Segmentation fault.
 - GDB will tell you where your code segfaulted
 - GDB can tell you what values are what
 - why your code segfaulted
2. Program get stuck
 - infinite loop
3. GDB can print structs!
 - `p *head` will print the fields of the struct pointed to by *head*
4. GDB can interpret numbers however you tell it to!
 - Use the `x` command to view the data at a memory address
 - `x buf` means print the value at *buf*
 - `x/10b` means print 10 (**10**) bytes (**b**) – can be used to “`x/8b $rdi`”
 - use `p` command with `/x` to print number in hex notation: `p /x val`

Function pointer

- In lab2, we invoke the function by function pointer *accum*
- Like normal data pointers (int *, char *, ..), we can have pointers to functions
 - A function pointer points to code, not data. Typically a function pointer stores the start of executable code
 - A function's name can also be used to get functions' address
 - In general, function pointer refer to functions of any signature. return type does not necessarily have to be void
 - Like normal data pointers, a function pointer can be passed as an argument and can also be returned from a function
 - why this useful?

Function pointer

- Like normal data pointers, a function pointer can be passed as an argument and can also be returned from a function
- We can invoke different functions into one function by using a function pointer
 - as long as the different functions using the same parameters and have the same return types
 - In C, we can use function pointers to avoid code redundancy

Testing

- When you fail in one test case, it does not mean you can only have bugs in this function implementation
 - Even if you have passed the previous test cases..
- No one test can help you test all possible bugs in your code
- Led to an interesting research topic:
 - Proof of Program Correctness