

4.3 Substitution Method

Solve: $T(n) = \begin{cases} 1 & \text{if } n=0 \\ 2T(\lfloor \frac{n}{3} \rfloor) + n^2 & \text{if } n>0 \end{cases}$

We guess that $T(n) = O(n^2)$
and try to prove it.

We want $c \in \mathbb{R}^+$ and $n_0 \in \mathbb{N}$ such that
 $T(n) \leq cn^2$ for all $n \geq n_0$

Assume that $T(m) \leq cm^2$ for all $m \leq n$

Show: $T(n) \leq cn^2$

$$\begin{aligned}
 T(n) &= 2T(\lfloor \frac{n}{3} \rfloor) + n^2 \\
 &\leq 2c(\lfloor \frac{n}{3} \rfloor)^2 + n^2 \quad (\text{because } \lfloor \frac{n}{3} \rfloor < n) \\
 &\leq 2c\left(\frac{n}{3}\right)^2 + n^2 \\
 &= \left(\frac{2}{9}c + 1\right)n^2
 \end{aligned}$$

We want $c \in \mathbb{R}^+$ such that

$$\left(\frac{2}{9}c + 1\right)n^2 \leq cn^2.$$

$$\Leftrightarrow \frac{2}{9}c + 1 \leq c$$

$$\Leftrightarrow 1 \leq \frac{7}{9}c$$

$$\Leftrightarrow \frac{9}{7} \leq c$$

For the base case, consider $n=0$:

$$T(0) = 1 \leq c(0)^2$$

?!

$$\therefore 1 \leq 0$$

Doesn't work for any $c > 0$.

We can choose a different base case

$$\begin{aligned} T(1) &= 2T(\lfloor \frac{1}{3} \rfloor) + 1^2 \\ &= 2T(0) + 1^2 \\ &= 2(1) + 1^2 \\ &= 3 \end{aligned}$$

we want $T(1) \leq c(1)^2$

$$\therefore 3 \leq c$$

Note: $T(2) = 2T\left(\lfloor \frac{2}{3} \rfloor\right) + 2^2$

$$\begin{aligned} &= 2T(0) + 2^2 \\ &= 2(1) + 2^2 \\ &= 6 \end{aligned}$$

We want $T(2) \leq c(2)^2$

$$6 \leq 4c$$

$c = 3$ is okay

so the base cases are $n=1$ and $n=2$

$$\text{and } T(n) \leq 3n^2 \text{ for all } n \geq 1$$

$$\therefore T(n) = O(n^2)$$