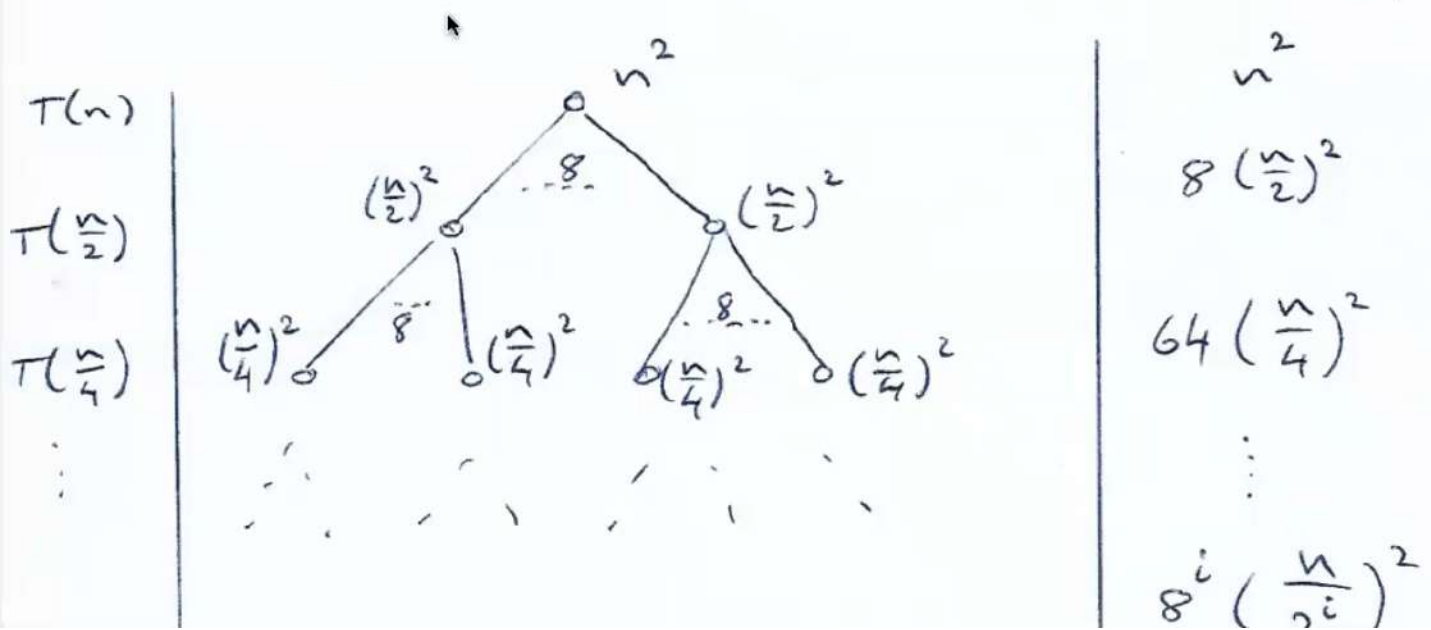


#### 4.4 Recursion-Tree Method

Example: Solve  $T(n) = \begin{cases} c & \text{if } n=1 \\ 8T(\frac{n}{2}) + n^2 & \text{if } n>1 \end{cases}$



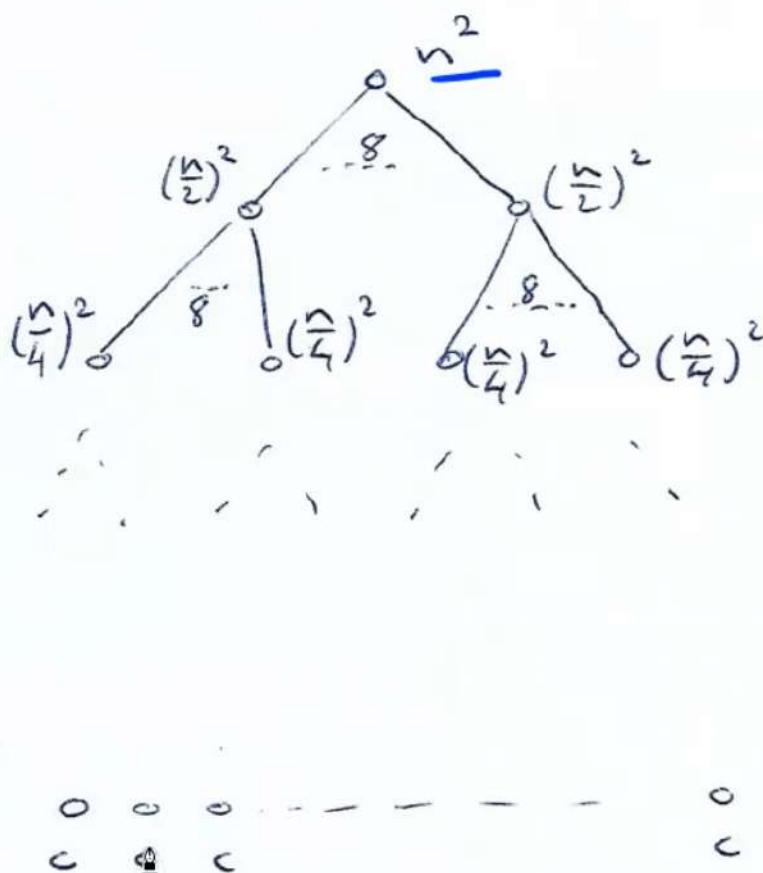
$T(n)$

$T(\frac{n}{2})$

$T(\frac{n}{4})$

$\vdots$

$T(1)$



$n^2$

$8(\frac{n}{2})^2$

$64(\frac{n}{4})^2$

$\vdots$

$8^i (\frac{n}{2^i})^2$

$\vdots$

$c \cdot (\# \text{leaves})$

For height: Solve  $\frac{n}{2^k} \leq 1$

$$\Leftrightarrow n \leq 2^k$$

$$\Leftrightarrow \log n \leq k \quad \therefore \text{height} = \log n$$

$$\begin{aligned} \# \text{leaves} &= (\text{branching factor})^{\text{height}} \\ &= 8^{\log n} \end{aligned}$$

$$T(n) = \sum_{i=0}^{\log n - 1} 8^i \left( \frac{n}{2^i} \right)^2 + c \cdot n^3$$

$$= n^2 \sum_{i=0}^{\log n - 1} \frac{8^i}{2^{2i}} + cn^3$$

$$= n^2 \sum_{i=0}^{\log n - 1} 2^i + cn^3$$

$$= n^2 \left( \frac{2^{\log n} - 1}{2 - 1} \right) + cn^3$$

$$= n^2 2^{\log n} - n^2 + cn^3$$

$$= n^3 - n^2 + cn^3$$

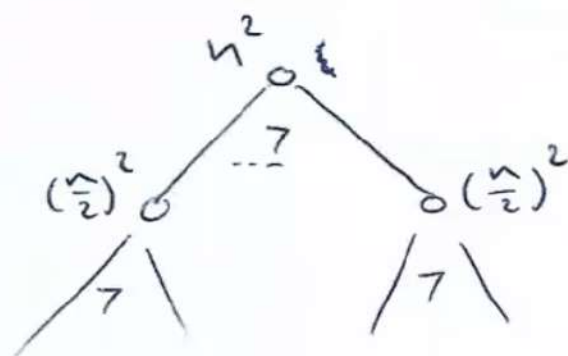
$$= \Theta(n^3)$$

Example: Solve  $T(n) = \begin{cases} c & \text{if } n=1 \\ 7T(\frac{n}{2}) + n^2 & \text{if } n>1 \end{cases}$

$$T(n)$$

$$T\left(\frac{n}{2}\right)$$

⋮



$$n^2$$

$$7\left(\frac{n}{2}\right)^2$$

⋮

$$7^i \left(\frac{n}{2^i}\right)^2$$

⋮

$$T(1)$$

o o o ... o  
c c c ... c

$c(\# \text{ leaves})$

$$\text{height} = \log n$$

$$\begin{aligned} \# \text{leaves} &= 7^{\log n} \\ &= n^{\log 7} \end{aligned}$$

$$\begin{aligned} T(n) &= \sum_{i=0}^{\log n - 1} 7^i \left( \frac{n}{2^i} \right)^2 + c \cdot n^{\log 7} \\ &= n^2 \sum_{i=0}^{\log n - 1} \left( \frac{7}{4} \right)^i + c \cdot n^{\log 7} \\ &= n^2 \left( \left( \frac{7}{4} \right)^{\log n} - 1 \right) + c n^{\log 7} \\ &= n^{\log 7} - n^2 + c n^{\log 7} \\ &= \Theta(n^{\log 7}) \approx \Theta(n^{2.81}) \end{aligned}$$