

4.3 Substitution Method

Solve:
$$T(n) = \begin{cases} 1 & \text{if } n = 0 \\ 2T(\lfloor \frac{n}{3} \rfloor) + n^2 & \text{if } n > 0 \end{cases}$$

We guess that $T(n) = O(n^2)$
and try to prove it.

We want $c \in \mathbb{R}^+$ and $n_0 \in \mathbb{N}$ such that
 $T(n) \leq cn^2$ for all $n \geq n_0$

Assume that $T(m) \leq cm^2$ for all $m \leq n$
show: $T(n) \leq cn^2$

$$\begin{aligned}
 T(n) &= 2 T(\lfloor \frac{n}{3} \rfloor) + n^2 \\
 &\leq 2c (\lfloor \frac{n}{3} \rfloor)^2 + n^2 \\
 &\leq 2c (\frac{n}{3})^2 + n^2 \\
 &= (\frac{2}{9}c + 1) n^2
 \end{aligned}$$

(because $\lfloor \frac{n}{3} \rfloor < n$)

We want $c \in \mathbb{R}^+$ such that

$$\left(\frac{2}{9}c + 1\right)n^2 \leq cn^2$$

$$\Leftrightarrow \frac{2}{9}c + 1 \leq c$$

$$\Leftrightarrow 1 \leq \frac{7}{9}c$$

$$\Leftrightarrow \frac{9}{7} \leq c$$

For the base case, consider $n=0$:

$$T(0) = 1 \leq c(0)^2$$

$$\therefore 1 \leq 0 \quad ?!$$

Doesn't work for any $c > 0$.

We can choose a different base case

$$\begin{aligned} T(1) &= 2T(\lfloor \frac{1}{3} \rfloor) + 1^2 \\ &= 2T(0) + 1^2 \\ &= 2(1) + 1^2 \\ &= 3 \end{aligned}$$

$$\begin{aligned} \text{We want } T(1) &\leq c(1)^2 \\ \therefore 3 &\leq c \end{aligned}$$

Note:
$$\begin{aligned} T(2) &= 2T\left(\left\lfloor \frac{2}{3} \right\rfloor\right) + 2^2 \\ &= 2T(0) + 2^2 \\ &= 2(1) + 2^2 \\ &= 6 \end{aligned}$$

We want $T(2) \leq c(2)^2$
 $6 \leq 4c$
 $\therefore \underline{c = 3}$ is okay

so the base cases are $n=1$ and $n=2$

$$\& \quad T(n) \leq 3n^2 \quad \text{for all } n \geq 1$$

$$\therefore T(n) = O(n^2)$$