

# Chapter 17

17.1-1

No.	A sequence like	Cost
=	Multipush( $S, n$ )	$n$
	Multipop( $S, n$ )	$n$
	Multipush( $S, -$ )	$n$
	Multipop( $S, -$ )	$n$
	:	:
	Multipop( $S, n$ )	$n$

*n operations*

The cost of the sequence of  $n$  operations is  $n^2$ , so the amortized cost is  $\underline{\underline{\Theta(n)}}$

17.1-2

Once the counter is at:

$\overbrace{011\cdots11}^{k \text{ bits}}$

incr.       $100\cdots00$

cost =  $k$

decr.       $011\cdots11$

cost =  $k$

incr.       $100\cdots00$

cost =  $k$

:

:

$n$  operations, each with cost  $k$ , gives  $\underline{\underline{\Theta(nk)}}$

17.2 - 1

Note that only Push, Pop are allowed  
& the copy which copies all items  
on the stack to backup drive.

Use amortized cost

$$\hat{c}_i = \begin{cases} 2 & \text{if } c_i \text{ is Push} \\ & \& \text{if } c_i \text{ is Pop} \\ 0 & \text{if } c_i \text{ is COPY.} \end{cases}$$

The amortized cost includes the  
actual cost of 1 for Push or Pop  
& the remaining 1 is kept aside  
to use for the copy operation.

Then there is always enough credit  
to copy stack items.

Thus any  $n$  operations cost  $\mathcal{O}(2n)$   
 $= \underline{\mathcal{O}(n)}$ .

17.3 - 1

Define the function  $\bar{\Phi}'$  by

$$\bar{\Phi}'(D_i) = \bar{\Phi}(D_i) - \bar{\Phi}(D_0)$$

Then  $\bar{\Phi}'(D_0) = 0$

&  $\bar{\Phi}'(D_i) \geq 0$  for all  $i$ ,

so  $\bar{\Phi}'$  is a potential function.

The amortized costs associated with  $\bar{\Phi}'$  are:

$$\begin{aligned}\hat{c}'_i &= c_i + \bar{\Phi}'(D_i) - \bar{\Phi}'(D_{i-1}) \\ &= c_i + (\bar{\Phi}(D_i) - \bar{\Phi}(D_0)) - (\bar{\Phi}(D_{i-1}) - \bar{\Phi}(D_0)) \\ &= c_i + \bar{\Phi}(D_i) - \bar{\Phi}(D_{i-1}) \\ &= \hat{c}_i\end{aligned}$$

So the amortized costs for  $\bar{\Phi}'$  and  $\bar{\Phi}$  are the same.

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17.3-4

Since there are  $s_0$  elements on the stack at the start, and no credits to begin with, it costs 1 for each element popped ~~at~~ the original  $s_0$  element.

In the worst case, all  $s_0$  elements are popped off using a single multipop, which has a cost of  $s_0$ .

Then the remaining  $n-1$  operations are as before, ~~at~~ with 2 for Push, 0 for Pop and 0 for multipop.

Thus the cost of  $n$  operations is

$$O(s_0 + 2(n-1)).$$

which is  $\underline{\underline{O(n)}}$

If the initial elements are popped off one at a time, the  $s_0$  pops are needed each with cost 1. Then  $n$  operations cost

$$O(s_0 + 2(n-s_0))$$

which is also just  $\underline{\underline{O(n)}}$

Note that the number of elements in a the stack at the end of  $n$  operations is irrelevant.