

Chapter 17

17.1-1

<u>No.</u>	A sequence like		Cost
		Multipush(S, n)	n
		Multipop(S, n)	n
		Multipush(S, n)	n
		Multipop(S, n)	n
		⋮	⋮
		Multipop(S, n)	n

n operations

The cost of the sequence of n operations is n^2 , so the amortized cost is $\Theta(\frac{n^2}{n}) = \underline{\underline{\Theta(n)}}$

17.1-2

once the counter is at:

$\overbrace{\quad\quad\quad}^{k \text{ bits}}$

011---11

incr.

100---00

cost = k

cost = k

decr.

011---11

cost = k

incr.

100---00

⋮

n operations, each with cost k , gives $\underline{\underline{\Theta(nk)}}$.

17.2-1

Note that only Push, Pop are allowed
& then COPY which copies all items
on the stack to backup drive.

Use amortized cost

$$\begin{array}{rcl} \hat{c}_i = 2 & \text{if } c_i \text{ is Push} \\ & \& \text{if } c_i \text{ is Pop} \\ \hline \hat{c}_i = 0 & \text{if } c_i \text{ is COPY.} \end{array}$$

The amortized cost includes the
actual cost of 1 for Push or Pop
& the remaining 1 is kept aside
to use for the copy operation.

Then there is always enough credit
to copy stack items.

Thus any n operations cost $O(2n)$
 $= \underline{\underline{O(n)}}$.

17.3 - 1

Define the function Φ' by

$$\Phi'(D_i) = \Phi(D_i) - \Phi(D_0).$$

Then $\Phi'(D_0) = 0$

& $\Phi'(D_i) \geq 0$ for all i ,

so Φ' is a potential function.

The amortized costs associated with Φ' are:

$$\begin{aligned}\hat{c}'_i &= c_i + \Phi'(D_i) - \Phi'(D_{i-1}) \\ &= c_i + (\Phi(D_i) - \Phi(D_0)) - (\Phi(D_{i-1}) - \Phi(D_0)) \\ &= c_i + \Phi(D_i) - \Phi(D_{i-1}) \\ &= \hat{c}_i\end{aligned}$$

So the amortized costs for Φ' and Φ are the same.

17.3-4

Since there are s_0 elements on the stack at the start, and no credits to begin with, it costs 1 for each element popped off of the original s_0 elements.

In the worst case, all s_0 elements are popped off using a single `MULTIPOP`, which has a cost of s_0 .

Then the remaining $n-1$ operations are as before, ~~etc~~ with 2 for `PUSH`, 0 for `POP` and 0 for `MULTIPOP`.

Thus the cost of n operations is

$$O(s_0 + 2(n-1)).$$

which is $O(n)$.

If the initial elements are popped off one at a time, then s_0 pops are needed each with cost 1. Then n operations costs

$$O(s_0 + 2(n-s_0))$$

which is also just $O(n)$.

Note that the number of elements s_n on the stack at the end of n operations is irrelevant.