

4.5

Master Method

- used for solving recurrences like:

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

where $a > 1$ and $b > 1$.

and $f(n)$ is an asymptotically positive function -

The Master theorem says :

1. If $f(n) = O(n^{\log_b a - \epsilon})$ for some $\epsilon > 0$
then $T(n) = \Theta(n^{\log_b a})$
2. If $f(n) = \Theta(n^{\log_b a})$
then $T(n) = \Theta(n^{\log_b a} \log n)$
3. If $f(n) = \Omega(n^{\log_b a + \epsilon})$ for some $\epsilon > 0$
~~then~~ and $a f(\frac{n}{b}) \leq c f(n)$ break

3. If $f(n) = \mathcal{R}(n^{\log_b a + \varepsilon})$ for some $\varepsilon > 0$
~~then~~ and $a f(\frac{n}{b}) \leq c f(n)$ for ~~c < 1~~

then $T(n) = \Theta(f(n))$

$$\text{Solve : } T(n) = 4T\left(\frac{n}{3}\right) + n$$

$$\begin{aligned}a &= 4 \\b &= 3 \\f(n) &= n\end{aligned}$$

compare $f(n)$ with $n^{\log_b a}$

$$n < n^{\log_3 4}$$

$$\therefore f(n) = O(n^{\log_b a - \varepsilon})$$

\therefore case 1.

$$T(n) = \underline{\underline{\Theta(n^{\log_3 4})}}$$

$$\text{Solve : } T(n) = 25T\left(\frac{n}{5}\right) + n^2$$

$$\begin{array}{l} a = 25 \\ b = 5 \\ f(n) = n^2 \end{array} \quad \left| \begin{array}{l} \text{Compare } n^2 \text{ with } n^{\log_5 25} \\ \qquad \qquad \qquad \approx n^2 \end{array} \right. \quad \therefore f(n) = \Theta(n^{\log_5 25})$$

∴ Case 2.

$$\begin{aligned} \therefore T(n) &= \Theta(n^{\log_5 25} \cdot \log n) \\ &= \Theta(n^2 \log n) \end{aligned}$$

Solve : $T(n) = 15T\left(\frac{n}{4}\right) + n^2$

$$a = 15$$

$$b = 4$$

$$f(n) = n^2$$

Compare n^2 with $n^{\log_4 15}$
 $n^2 > n^{\log_4 15}$.

$$\therefore f(n) = \Omega(n^{\log_4 15})$$

∴ Case 3.

check : $a f\left(\frac{n}{b}\right) \leq c f(n)$ for $c < 1$.
some

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$$15\left(\frac{n}{4}\right)^2 \leq cn^2$$

$$\Leftrightarrow \frac{15}{16}n^2 \leq cn^2$$

$$\Leftrightarrow \frac{15}{16} \leq c. \quad \text{choose } c = \frac{15}{16} < 1.$$

$$\therefore T(n) = \Theta(n^2)$$