

1. (6 marks) Java program is given:

```
class Person{
    int x;
    public Object work(){return new Person();}
}

class Human{
    int x;
    public Object work(){ return new Person();}
}

class Worker extends Person {
    int x;
    public Object work() { return new Worker();}
    public void overTime(int h) { x = x+h; }

    public static void main(String[] args){
        Person a = new Human();          // line 1
        Worker b = new Person();          // line 2
        Person c = new Worker();          // line 3
        c.overTime(5);                     // line 4
        c.x =5;                            // line 5
        Worker m = c.work();               // line 6
    }
}
```

For each line (line 1 to line 6), does it compile? If it does not compile, give the reason.

2. For the code below (a language with nested subroutine), the language uses a value model of variables.

```
program A(){  
  x, y, z: integer;  
  procedure B(){  
    y: integer;  
    y=0;  
    x=z+1;  
    z=y+2;  
  }
```

```
  procedure C(){  
    z: integer;  
    procedure D(){  
      x: integer;  
      x = z + 1;  
      y = x + 1;  
      call B();  
    }  
    z = 5;  
    call D();  
  }
```

```
  x = 10;  
  y = 11;  
  z = 12;  
  call C();  
  print x, y, z;
```

```
}
```

2.1 (3 marks) If the language uses static scoping, the printed result of x, y, and z is	2.2 (3 marks) If the language uses dynamic scoping, the printed result of x, y, and z is
x =	x =
y =	y =
z =	z =

3. Given the C++ code below.

```
class First {  
public:  
    First() { b = 10; }  
    virtual void display(int &x, int y) { x = x + y; cout << "b, x " << b << " " << x << endl; }  
private:  
    int b;  
};
```

```
class Second: public First {  
public:  
    Second() { d = 20; }  
    virtual void display(int &x, int y) { x = x * y; cout << "d, x " << d << " " << x << endl; }  
private:  
    int d;  
};
```

```
int main() {  
    First f, *p;  
    Second s;  
    int m = 1;  
    int *n = new int(2);  
    float o = 5.7;  
    p = &s;  
    p->display(m, o);           //line1  
    f = s;  
    f.display(m,o);           //line2  
    return 0;  
}
```

(1 mark) At line1, the method binding is ☐ static ☐ dynamic

(1 mark) At line2, the method binding is ☐ static ☐ dynamic

(1 mark) In the checking of the types of the method arguments at line2, the following rule(s) of the type system are used (you may choose 1 or more).

☐ type equivalence ☐ type compatibility ☐ type inference

4. A Java-like language uses left-to-right evaluation order. Its precedence and associativity rules are given below. (Precedence is from the highest down to the lowest.)

Operator	Description	Associativity
...	...	...
* / %	multiplicative	left to right
...	...	...
== !=	equality	left to right
...	...	...
&&	logical and	left to right
	logical or	left to right
...	...	...

- 4.1 (3.5 marks) Add parentheses to the expression below to show the effect of precedence and associativity to the grouping of operands to operators.

`c % 400 == 0 || c % 4 == 0 && c % 100 != 0`

- 4.2 (1.5 marks) If `c` is 1666, the result of the expression is .....

- 4.3 (3 marks) If this language has short circuiting, which of these subexpressions get evaluated in the question 4.2?

`c % 400 == 0`      ☐ yes      ☐ no

`c % 4 == 0`      ☐ yes      ☐ no

`c % 100 != 0`      ☐ yes      ☐ no