

Game Theory with Computer Science Applications

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Lecture 12
May 7, 2024

Main content of this lecture

Revenue Optimal Mechanism

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Problem Setting

- 1 seller 1 item
- N buyer private valuation θ_i f_i , where $\theta_i \in [0, \theta_{i,\max}]$, each θ_i is independent.

Mechanism Design

- Allocation

$$\pi_i(b_i, b_{-i}) \quad s.t. \pi_i \geq 0 \quad \sum_{i=1}^N \pi_i \leq 1$$

- Payment $q_i(b_i, b_{-i})$

The expected allocation and payment:

$$\begin{aligned} \alpha_i(\theta_i) &= E_{\theta_{-i}}[\pi_i(\theta_i, \theta_{-i}) | \theta_i] \\ m_i(\theta_i) &= E_{\theta_{-i}}[q_i(\theta_i, \theta_{-i}) | \theta_i] \end{aligned}$$

- Payoff

$$\begin{aligned} &\theta_i \alpha_i(\hat{\theta}_i) - m_i(\hat{\theta}_i) \\ \max &\sum_{i=1}^N E[m_i(\theta_i)] \end{aligned}$$

We need to solve for $\{\pi_i\}$ and $\{q_i\}$, and then

s.t. (IC) and (IR)

$$\begin{aligned} (IC) : & \theta_i \alpha_i(\theta_i) - m_i(\theta_i) \geq \theta_i \alpha_i(\hat{\theta}_i) - m_i(\hat{\theta}_i) \quad \forall i, \theta_i, \hat{\theta}_i \\ (IR) : & \theta_i \alpha_i(\theta_i) - m_i(\theta_i) \geq 0 \quad \forall i, \theta_i \end{aligned}$$

Translate the conditions of IC and IR into three equivalent constraints:

$$(IC) (IR) \iff$$

$$(i) \quad m_i(\theta_i) = m_i(0) + \theta_i \alpha_i(\theta_i) - \int_0^{\theta_i} \alpha_i(\theta) d\theta$$

$$(ii) \quad \alpha_i \text{ is a non-decreasing function}$$

$$(iii) \quad m_i(0) \leq 0$$

Now we can remove the conditions of IC and IR and solve the new optimization problem.
Our objective for parameter optimization is:

$$\begin{aligned}
 E(m_i(\theta_i)) &= m_i(0) + \int_0^{\theta_{i,\max}} \alpha_i(\theta_i) \left(\theta_i - \frac{1 - F_i(\theta_i)}{f_i(\theta_i)} \right) d\theta_i \\
 &\quad \downarrow \text{red arrow} \\
 &\quad \int \pi_i(\theta_i, \theta_{-i}) \underbrace{f_{-i}(\theta_{-i})}_{f_1(\theta_1) f_2(\theta_2) \cdots f_{i-1}(\theta_{i-1}) f_{i+1}(\theta_{i+1}) \cdots f_N(\theta_N)} d\theta_{-i}
 \end{aligned}$$

Therefore, our optimization objective becomes:

$$\Rightarrow \sum_{i=1}^N E[m_i(\theta_i)] = \sum_{i=1}^N m_i(0) + \sum_{i=1}^N \int \pi_i(\theta_i, \theta_{-i}) \varphi_i(\theta_i) f(\vec{\theta}) d\vec{\theta}$$

Optimal Allocation rule is

$$t_i(\theta_i, \theta_{-i}) = 1 \quad \text{iff} \quad \varphi_i(\theta_i) = \max_j \varphi_j(\theta_j) \quad \text{and} \quad \varphi_i(\theta_i) \geq 0$$

Payment rule is :

$$\begin{cases} q_i(\theta_i, \theta_{-i}) = \partial_i \pi_i(\theta_i, \theta_{-i}) - \int^{\alpha_i} \pi_i(\theta_i; \theta_{-i}) d\theta_i. \\ q_v(\theta_v, \theta_{-i}) = 0 \Rightarrow m_v(0) = 0 \end{cases}$$

Assumption φ_i is a strictly increasing function \rightarrow regular

A sufficient condition for $\varphi_v(\theta_i)$ to be non-decreasing.

$$\frac{1 - F_i(\theta_i)}{f_r(\theta_i)} \text{ is non-decreasing}$$

$$\frac{f_r(\theta_i)}{1 - F_i(\theta_i)} \nearrow f(\theta) = \lambda e^{-\lambda \theta} \quad \theta \geq 0 \text{ for some } \lambda > 0.$$

Claim If φ_i is strictly increasing then (ii) is satisfied by our allocation rule

$$\text{if } \pi_i(\theta_i, \theta_{-i}) > 0 \text{ then } \bar{a}_i(\hat{\theta}_i, \theta_{-i}) = 1 \quad \text{if } \hat{\theta}_i > \theta_i$$

$$\Rightarrow \pi_v(\hat{\theta}_i, \theta_{-i}) \geq \pi_i(\theta_i, \theta_{-i})$$

$$\text{if } \pi_r(\theta_i, \theta_{-i}) = 0, \text{ then}$$

$$\Rightarrow \bar{a}_i(\hat{\theta}_i, \theta_{-i}) \geq \pi_i(\theta_i, \theta_{-i}) \quad \forall \hat{\theta}_i \geq \theta_{-i}$$

$$\alpha_i(\theta_i) = E_{\theta_{-i}}[\pi_i(\theta_i, \theta_{-i})] \quad \text{is non-decreasing}$$

Revenue

We take the maximum value among each φ_i and 0.

$$E[\max(\varphi_1(\theta_1), \varphi_2(\theta_2), \dots, \varphi_N(\theta_N), 0)]$$

$$r_i(\theta_{-i}) = \left\{ z_i : \varphi_i(z_i) \geq 0 \text{ and } \varphi_i(z_i) \geq \max_{j \neq i} \varphi_j(\theta_j) \right\}$$

r_i is the smallest valuation of the bidder i such that she is guaranteed to win the item. Under this situation, we can write the expressions for π_i and q_i as follows:

$$\begin{aligned} \pi_{ij}(\theta_i, \theta_{-i}) &= \begin{cases} 1 & \text{if } \theta_i > r_i(\theta_{-i}) \\ 0 & \text{if } \theta_i < r_i(\theta_{-i}) \end{cases} \\ q_i(\theta_i, \theta_{-i}) &= \theta_i \pi_r(\theta_i, \theta_{-i}) - \int_0^{\theta_i} \pi_i(\theta_v, \theta_{-i}) d\theta \\ &= \theta_i I_{\theta_i > r_i} - \int_0^{\theta_i} I_{\theta > r_i} d\theta \end{aligned}$$

Next, we will discuss the situation case by case.

If $\theta_i > r_i(\theta_{-i})$ then.

$$q_i(\theta_i, \theta_{-i}) = \theta_i - (\theta_i - r_i) = r_i(\theta_i)$$

If $\theta_i < r_i(\theta_{-i})$

$$q_i(\theta_i, \theta_{-i}) = 0$$

Revenue - optimal Mechanism

Next, we will provide a re-explanation for Revenue-Optimal Mechanism.

(1) The bidder with the largest virtual valuation $\varphi_i(\theta_i)$ wins

If $\max_j \varphi_j(\theta_j) \geq 0$

If $\max_j \varphi_j(\theta_j) < 0$ the seller does not sell the item.

(2) The winner (If there is one) pays a amount equal to the smallest value that would still win

Example If f_i are all identical

$$\varphi(\theta_i) = \theta_i - \frac{1 - F(\theta_i)}{f(\theta_i)}$$

$$\begin{aligned} r_i(\theta_i) &= \min \left\{ z_i : \varphi(z_i) \geq 0, \varphi(z_i) \geq \max_{j \neq i} \varphi(\theta_j) \right\} \quad \text{since non-decreasing} \\ &= \min \left\{ z_i : z_i \geq \varphi^{-1}(0), \quad z_i \geq \max_{j \neq i} \theta_j \right\} \end{aligned}$$

This is like a second price auction with a reserve price of $\varphi^{-1}(0)$.

A: $\varphi(\theta_i) < 0$

B: $\varphi(\theta_1) > 0, \varphi(\theta_i) < 0, i \neq 1$