

Euler's Totient Function, denoted as $\phi(n)$

$\phi(n)$ = Num of positive integers less than n , that are relatively prime to n

$\phi(7) = ?$

We will check co-prime relationship of 7 with 1,2,3,4,5,6

$\text{GCD}(1,7) = 1 \Rightarrow 1 \text{ \& 7 co-prime}$

$\text{GCD}(2,7) = 1 \Rightarrow 2 \text{ \& 7 co-prime}$

$\text{GCD}(3,7) = 1 \Rightarrow 3 \text{ \& 7 co-prime}$

$\text{GCD}(4,7) = 1 \Rightarrow 4 \text{ \& 7 co-prime}$

$\text{GCD}(5,7) = 1 \Rightarrow 5 \text{ \& 7 co-prime}$

$\text{GCD}(6,7) = 1 \Rightarrow 6 \text{ \& 7 co-prime}$

So, $\phi(7) = 6$.

ut $\phi(367) = ?$

We will check co-prime relation of 1,2,3,4, ... 365,366 with 367 ????

No.....

Formula:

For $\phi(n)$:

1.If n is prime, then $\phi(n) = n-1$

2.If n can be divided as

$n = p \times q$, and p & q are primes,

then $\phi(n) = n-1$

3. If n can e divided as $n = a \times b$, where a or b or both are composite num, then

$\phi(n) = n \times (1 - 1/P_1) \times (1 - 1/P_2) \dots$, here P_1, P_2, \dots are distinct primes

2. Fermat's Little Theorem

Fermat's Little Theorem states that if p is a prime and a is an integer not divisible by p , then:

$$a^{(p-1)} \equiv 1 \pmod{p}$$

Example:

For $a = 2$ and $p = 7$:

$$2^{(7-1)} = 2^6 = 64$$

$$64 \bmod 7 = 1$$

Thus, $2^6 \equiv 1 \pmod{7}$.

3. Euler's Theorem

It states that if a and n are coprime, then:

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

Example:

For $a = 2$ and $n = 10$,

$$\phi(10) = 4: 2^4 = 16$$

$$16 \bmod 10 = 6$$

Thus, $2^4 \equiv 6 \pmod{10}$

4. RSA Key Generation for $p = 311$ and $q = 317$.

1. Calculate n :

$$n = p \times q = 311 \times 317 = 98587$$

2. Compute Euler's Totient Function $\phi(n)$:
 $\phi(n) = (p - 1) \times (q - 1) = (311 - 1) \times (317 - 1) = (310) \times (316) = 97960$
3. Choose Public Key Exponent e : We need to choose e such that $1 < e < \phi(n)$ and $\gcd(e, \phi(n)) = 1$.
 first 10 elements of Z_{97960}^* are: 1, 3, 9, 11, 17, 19, 23, 29, 31, 33.
 Let's select the 10th element of Z_{97960}^* . So choose $e = 33$
4. Find the Private Key Exponent d : Compute d using the Extended Euclidean Algorithm, which satisfies: $d \times e \equiv 1 \pmod{\phi(n)}$

Q	A	B	R	t1	t2	t
2968	97960	33	16	0	1	-2968
2	33	16	1	1	-2968	5937
16	16	1	0	-2968	5937	-97960
	1	0		5937	97960	

Public key: $(e,n)=(33,98587)$

Private key: $((d,n)=(5937,98587))$

Step 1: Encryption

The encryption formula is:

$$C = m^e \pmod{n}$$

Where:

- $m=10$ (the message)
- $e=33$ (public exponent)
- $n=98587$

$$C = 10^{33} \pmod{98587}$$

C = 18490 (How ?? Think ..)

Step 2: Decryption

The decryption formula is:

$$M = (c^d) \bmod n$$

Where:

- c is the ciphertext (calculated in step 1)
- d=5937 (private exponent)
- n=98587

$$m = 18490^{5937} \bmod 98587$$