

Programming OpenMP

Task Affinity

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Improving Tasking Performance: Task Affinity

Motivation



- Techniques for process binding & thread pinning available
 - →OpenMP thread level: OMP_PLACES & OMP_PROC_BIND
 - →OS functionality: taskset -c

OpenMP Tasking:

- In general: Tasks may be executed by any thread in the team
 - → Missing task-to-data affinity may have detrimental effect on performance

OpenMP 5.0:

affinity clause to express affinity to data

affinity clause



- New clause: #pragma omp task affinity (list)
 - → Hint to the runtime to execute task closely to physical data location
 - →Clear separation between dependencies and affinity

Expectations:

- → Improve data locality / reduce remote memory accesses
- → Decrease runtime variability
- Still expect task stealing
 - →In particular, if a thread is under-utilized

Code Example



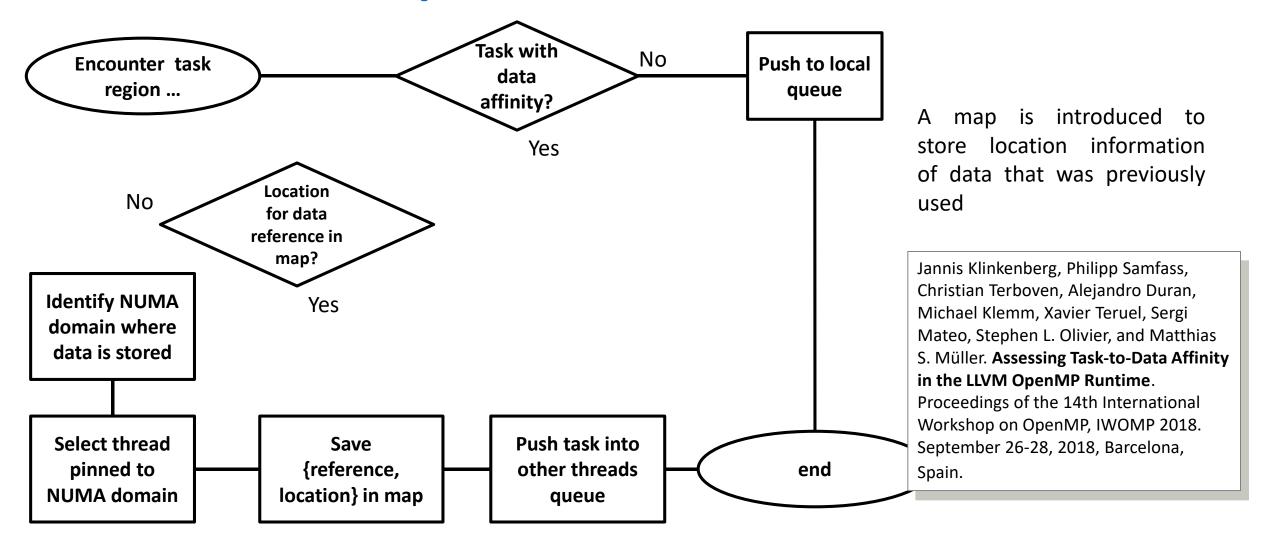
Excerpt from task-parallel STREAM

```
1  #pragma omp task \
2     shared(a, b, c, scalar) \
3     firstprivate(tmp_idx_start, tmp_idx_end) \
4     affinity( a[tmp_idx_start] )
5     {
6        int i;
7      for(i = tmp_idx_start; i <= tmp_idx_end; i++)
8        a[i] = b[i] + scalar * c[i];
9     }</pre>
```

- → Loops have been blocked manually (see tmp_idx_start/end)
- → Assumption: initialization and computation have same blocking and same affinity

Selected LLVM implementation details

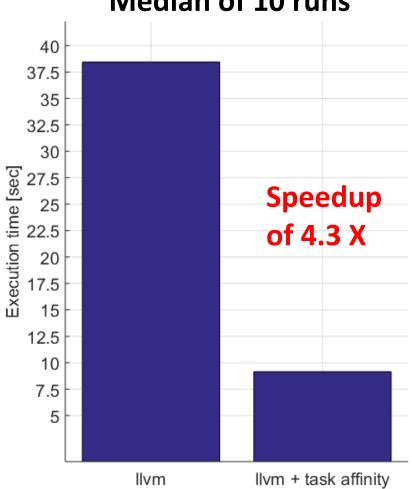




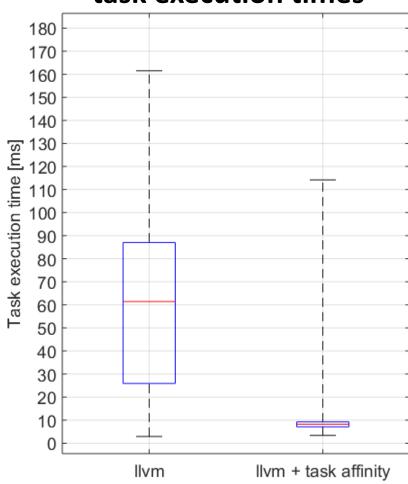
Evaluation

OpenMP

Program runtime Median of 10 runs



Distribution of single task execution times



LIKWID: reduction of remote data volume from 69% to 13%

Summary



- Requirement for this feature: thread affinity enabled
- The affinity clause helps, if
 - → tasks access data heavily
 - → single task creator scenario, or task not created with data affinity
 - →high load imbalance among the tasks

Different from thread binding: task stealing is absolutely allowed