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1 DP

1.1 Convex Hull Trick

```
const ll is query=-(1LL<<62);
const ll inf=1e18L;</pre>
struct Line {
ll m, b;
mutable function<const Line*()> succ;
 bool operator<(const Line& rhs) const
  if (rhs.b != is query) return m < rhs.m;</pre>
  const Line* s = succ();
  if (!s) return 0;
  int x = rhs.m;
  return b - s - b < (s - m - m) * x:
};
/* will maintain upper hull for maximum
 * for minimum insert line as (-m,-b) !!
 * then the ans would be -ans
struct HullDynamic : public multiset<Line>
 bool bad(iterator y)
  auto z=next(y);
  if(y==begin())
   if(z==end()) return 0;
   return y->m == z->m \&\& y->b <=z->b;
  auto x=prev(y);
  if(z=end()) return y->m == x->m \&\& y->b <= x->b;
    int128 left= int128(x->b - y->b)*(z->m - y->m);
  =int128 right==int128(y->b - z->b)*(y->m -
  \rightarrow x->m):
  return left>=right;
 void insert line(ll m,ll b)
  auto y=insert({m,b});
  y -> succ = [=]  { return next(y) == end()?0: \&*next(y);
  if(bad(y)) { erase(y); return; }
  while (\text{next}(y)) = \text{end}() \& \text{bad}(\text{next}(y))
  - erase(next(y));
while(y!=begin() && bad(prev(y))) erase(prev(y));
 ll query(ll x)
  auto l=*lower bound((Line){x,is query});
  return l.m*x+T.b;
```

1.2 DnC

```
const ll inf=1e15L;
const int mx=1e5+5;
ll ara[mx], sum[mx];
ll dp[105][mx];
ll C(int i, int j,int k)
{
   return 1LL*k*(sum[j]-sum[i-1]);
```

```
void compute(int groupNo,int l,int r,int optL,int

→ optR)

if(l>r) return;
int mid=(l+r)/2:
dp[groupNo][mid]=-inf;
int optNow=optL;
 for(int endOfLast=optL;endOfLast<=optR &&</pre>

    endOfLast<mid;endOfLast++)</pre>
  Il ret=dp[groupNo-1][endOfLast]+C(endOfLast+1,mid_
      ,groupNo);
  if(ret>=dp[groupNo][mid])
   dp[groupNo][mid]=ret;
   optNow=endOfLast;
compute(groupNo,l,mid-1,optL,optNow);
compute(groupNo, mid+1, r, optNow, optR);
void solve()
 //your code goes here
for(int groupNo=2;groupNo<=k;groupNo++)</pre>
  compute(groupNo,1,n,1,n);
  for(int i=groupNo+1;i<=n;i++) dp[groupNo][i]=max(_</pre>
     dp[groupNo][i-1],dp[groupNo][i]);
cout<<dp[k][n]<<"\n";
```

1.3 Li Chao Tree

```
* Li Chao Tree for maximum query, represents line
   as mx+b
 * for minimum query initialize line with (0,inf)
* change the sign of update and query function
- too( > to <, max to min)
* for update call update(1,n,1,{m,b})</pre>
   for query at point x call query (1, n, 1, x)
 * This tree works for almost any functions
   just change the parameters of the Line
const ll inf=1LL<<62;</pre>
const int mx=1e6+5;
struct Line
Line(ll m=0,ll b=-inf) { m=_m; b=_b; }
ll f(Line line, int x) { return line.m*x+line.b; }
Line Tree[4*mx];
void update(int l,int r,int at,Line line)
int mid=(l+r)/2;
```

```
bool left=f(line,l)>f(Tree[at],l);
bool middle=f(line,mid)>f(Tree[at],mid);
if(middle) swap(Tree[at],line);
if(l==r) return;
if(left!=middle) update(l,mid,at*2,line);
else update(mid+1,r,at*2+1,line);
}

ll query(int l,int r,int at,int x)
{
    lt val=f(Tree[at],x);
    if(l==r) return val;
    int mid=(l+r)/2;
    if(x<=mid) return max(val,query(l,mid,at*2,x));
    return max(val,query(mid+1,r,at*2+1,x));
}</pre>
```

1.4 SOS DP

```
/// 0(N * 2^N)
///memory optimized version
for(int i = 0; i < (1 << N); ++i) F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask <
   (1 << N); ++ mask)
if(mask & (1 << i)) F[mask] += F[mask^(1 << i)];
/// How many pairs in ara[] such that (ara[i] &
\rightarrow ara[i]) = 0
/// N --> Max number of bits of any array element
const int N = 20;
int inv = (1 << N) - 1;
int F[(1 << N) + 10];
int ara[MAX]:
/// ara is 0 based
long long howManyZeroPairs(int n,int ara[]) {
    CLR(F);
    for(int i=0;i<n;i++) F[ara[i]]++;</pre>
    for(int i = 0; i < N; ++i) for(int mask = 0;
        mask < (1 << N): ++ mask)
        if(mask & (1 << i)) F[mask] += F[mask^(1 << i)];
    long long ans = 0;
    for(int i=0;i<n;i++) ans += F[ara[i] ^ inv];</pre>
    return ans;
/// F[mask] = sum of A[i] given that (i&mask)=mask
for(int i = 0; i < (1 << N); ++i) F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = (1 << N) - 1;
   mask >= 0: --mask)
if (!(mask & (1<<i))) F[mask] += F[mask | (1<<i)];
/// Number of subsequences of ara[0:n-1] such that
/// sub[0] \& sub[2] \& ... \& sub[k-1] = 0
const int N = 20;
int inv = (1 << N) - 1;
int F[(1<<N) + 10];</pre>
int ara[MAX];
int p2[MAX]; /// p2[i] = 2^i
///0 based array
```

```
int howManyZeroSubSequences(int n,int ara[]) {
    CLR(F);
    for(int i=0;i<n;i++) F[ara[i]]++;</pre>
    for(int i = 0; i < N; ++i)
    for(int mask = (1<<N)-1; mask >= 0; --mask){
             if (!(mask & (1<<i)))
                 F[mask] += F[mask \mid (1 << i)];
    int ans = 0;
    for(int mask=0; mask<(1<< N); mask++) {
        if( builtin popcount(mask) \& 1) ans =
        sub(ans, p2[F[mask]]);
else ans = add(ans, p2[F[mask]]);
    return ans;
/// Number of subsequences of ara[0:n-1] such that
/// sub[0] | sub[2] | ... | sub[k-1] = 0
int F[(1 << 20) + 10], ara[MAX];
int p2[MAX]; /// p2[i] = 2^i
/// ara is 0 based
int howManySubsequences(int n, int ara[], int m,
→ int Q) {
    CLR(F);
    for(int i=0;i<n;i++) F[ara[i]]++;</pre>
    if(0 == 0) return sub(p2[F[0]], 1);
    for(int i = 0; i < m; ++i)
         for(int mask = 0; mask < (1<<m); ++mask){</pre>
             if (mask & (1<<i))
                 F[mask] += F[mask ^ (1<<i)];
    int ans = 0:
    for(int mask=0; mask<(1<<m); mask++) {</pre>
         if(mask & Q != mask) continue;
         if( builtin popcount(mask ^{\circ} Q) & 1) ans =

    sub(ans, p2[F[mask]]);

         else ans = add(ans, p2[F[mask]]);
    return ans:
```

2 Data Structures

2.1 2D fenwick tree

```
OLIN | ITTEHAD | INZAMAM
                bit[i][i] += delta;
   }
2.2 Fenwick tree
int bit[1000], arra[1000];
int n;
void update( int idx,int val )
    for( int i = idx; i \le n; i + i\delta(-i) ) bit[i]
    return;
int query( int idx )
    int sum = 0:
    for( int i = idx; i > 0; i -= i\delta(-i) ) sum +=
    → bit[i]:
    return sum;
2.3 GP Hash Table
// policy based hash table. use random to reduce
 - collisions and pass anti-hash tests
#include <ext/pb ds/assoc container.hpp>
const int RANDOM = chrono::high resolution clock::n_
→ ow().time since epoch().count();
struct chash ₹
    int operator()(int x) const { return x ^
    → RANDOM; }
gp hash table<int, int, chash>table;
struct chash {
    int operator()(pii x) const { return x.first*
    \sim 31 + x.second; }
gp hash table<pii, int, chash> table; //for pairs
2.4 LIS
```

```
2.5 MO s Algo
const int mx=
const int sz=
struct query
    int l,r,id;
    bool operator<(const query &a) const
  int x=l/sz; int y=a.l/sz;
  if(x!=y) return x<y;</pre>
  if(x%2) return r<a.r;</pre>
  return r>a.r;
} ques[mx];
void add(int indx) { }
void baad(int indx) { }
void solve()
//write code here
int l=0:
int r=-1;
 sort(ques+1, ques+q+1);
 for(int i=1;i<=q;i++)
  while(l>ques[i].l) add(--l);
  while(r<ques[i].r) add(++r);</pre>
  while(l<ques[i].l) baad(l++);</pre>
  while(r>ques[i].r) baad(r--);
  ans[ques[i].id]=sum[now];
for(int i=1;i<=q;i++) cout<<ans[i]<<" ";
```

2.6 PBDS

2.7 Persistent Segment Tree(pointer) **struct** node ll sum; node *left,*right; node(ll sum=0) sum= sum; left=riaht=NULL: void build(int l,int r) if(l==r) return; left=new node() right=**new** node(); int mid=(l+r)/2;left->build(l,mid); right->build(mid+1,r); node *update(int l,int r,int i,int x) if(l>i || r<i) return this;</pre> if(l==r) return new node(x); int mid=(l+r)/2; node *ret=new node(); ret->left=left->update(l,mid,i,x); ret->right=right->update(mid+1,r,i,x); ret->sum=ret->left->sum+ret->right->sum; return ret; ll query(int tL,int tR,int rL,int rR) if(tL>rR || tR<rL) return 0;</pre> if(tL>=rL'&& tR<=rR) return sum;</pre> int mid=(tL+tR)/2; ll a=left->query(tL,mid,rL,rR); ll b=right->query(mid+1,tR,rL,rR); return a+b; int size() { return sizeof(*this)/sizeof(node*); } const int mx=2e5+5; node *root[mx]; void solve() int n,q; cin>>n>>q; root[0]=new node(); root[0]->build(1,n);for(int i=1;i<=n;i++)</pre> int x; cin>>x; root[0] = root[0] - supdate(1,n,i,x);

```
int sz=0:
while(g--)
 int op;
 cin>>op;
 if(op==1)
  int version,i,x;
  cin>>version>>i>>x;
  version--;
  root[version]=root[version]->update(1,n,i,x);
 else if(op==2)
  int version.l.r:
  cin>>version>>l>>r;
  version--
  cout<<root[version]->query(1,n,l,r)<<"\n";</pre>
 else
  int version:
  cin>>version;
  version--:
  root[++sz]=root[version];
```

2.8 RMQ

```
template<typename T>
struct sparse table
vector<T>ara;
vector<int>logs;
vector<vector<T>>table;
sparse table(int n)
 ara.resize(n+1);
  logs.resize(n+1);
T func(T a,T b) { }
void build(int n)
  for(int i=2;i<=n;i++) logs[i]=logs[i/2]+1;</pre>
 table.resize(n+1,vector<T>(logs[n]+1));
  for(int i=1;i<=n;i++) table[i][0]=ara[i];</pre>
 for(int j=1; j<=logs[n]; j++)</pre>
  int sz=1<<j;
  for(int i=1;i+sz-1<=n;i++)
    table[i][j]=func(table[i][j-1],table[i+sz/2][j-

→ 1]);
  query(int l,int r)
```

```
int d=logs[r-l+1];
  return func(table[l][d],table[r-(1<<d)+1][d]);</pre>
2.9 Segment Tree (Inz)
class SegmentTree {
                          idx * 2
    #define Lc(idx)
                          idx * 2 + 1
    #define Rc(idx)
    public:
        struct node {
            int value:
            int lazy;
            node() {
                 this->value = ??;
                 this->lazy = ??;
        };
        vector <node> segT;
        vector <int> A;
        SegmentTree(int sz) {
             // need to clear!
            segT.resize(4 * sz + 10);
            A.resize(sz + 1);
        node Merge(node L, node R) {
            node F;
            F = ??
            return F;
        void Relax(int L, int R, int idx) {
            //Do something
            segT[idx].lazy = ??; //after Relaxing
        void MakeSegmentTree(int L, int R, int idx)
            if (L == R)
                 segT[idx].value = ??;
                 return;
            int M = (L + R) / 2;
            MakeSegmentTree(L, M, Lc(idx));
MakeSegmentTree(M + 1, R, Rc(idx));
            seqT[idx] = Merge(seqT[Lc(idx)],

    segT[Rc(idx)]);

        node RangeQuery(int L, int R, int idx, int
         → l, int r) {
            Relax(L, R, idx);
            node F;
            if (L > r || R < l)
                                     return F;
            if (L >= l \& R <= r) return segT[idx];
            int M = (L + R) / 2;
            F = Merge(RangeQuery(L, M, Lc(idx), l,
                r), RangeQuery(M + 1, R, Rc(idx),
```

```
segT[idx] = Merge(segT[Lc(idx)],

→ segT[Rc(idx)]); //is it useful?

           return F;
       }
        void RangeUpdate(int L, int R, int idx, int
           l, int r, int lz) {
           Relax(L, R, idx);
           if (L > r || R < l)
                                   return:
           if (L >= l && R <= r) {
                // Do something
                segT[idx].lazy = ??;
               Relax(L, R, idx);
                return;
           int M = (L + R) / 2;
           RangeUpdate(L, M, Lc(idx), l, r, lz);
           RangeUpdate(M + 1, R, Rc(idx), l, r,
            segT[idx] = Merge( segT[Lc(idx)],

    segT[Rc(idx)]);

       }
};
```

2.10 Segment Tree (Point Update, Range Query)

2.11 Segment Tree (Range Update, Range Query)

```
if(lazy[rt]) push(rt, l, r);
 if(l > e or r < b or b > e) return;
 if(l >= b \text{ and } r <= e) 
  lazv[rt] += v:
  return push(rt, l, r);
 int m = l + r >> 1, lc = rt << 1, rc = lc | 1;
 update(lc, l, m, b, e, v); update(rc, m + 1, r, b, l)
 tree[rt] = merge(tree[lc], tree[rc]);
ll query(int rt, int l, int r, int b, int e){
if(lazy[rt]) push(rt, l, r);
 if(l > e or r < b or b > e) return 0;
 if(l >= b and r <= e) return tree[rt];</pre>
 int m = l + r >> 1, lc = rt << 1, rc = lc | 1;
 return merge(query(lc, l, m, b, e), query(rc, m +
 \rightarrow 1, r, b, e));
2.12 Sqrt Decomposition
int block size = ??;
int Block[block size + 5];
```

```
int getBlock(int idx)
    return (idx + block size - 1) / block size;
     → //for 1-base index
    return idx / block size;
     → //for 0-base index
int getQueryAns(int L, int R) //0-base index
    int ANS = 0;
int CL = L / block_size;
    int CR = R / block size;
    if (CL == CR) {
        for (int i = L; i <= R; ++i)
            ANS += ArrName[i];
    else {
        for (int i = L, LM = (CL + 1) * block size
         \rightarrow - 1; i <= LM; ++i)
            ANS += ArrName[i];
        for (int i = CL + 1; i \le CR - 1; ++i)
            ANS += Block[i];
        for (int i = CR * block size; i <= R; ++i)</pre>
            ANS += ArrName[i];
    return ANS;
//Update :
             Block[ idx / block size ] += ??
```

2.13 segment tree(lazy propagation)

```
int n, q, arra[100005];
struct idk{
   int sum, prop;
} tree[300005];
```

```
void init( int node,int b,int e )
    if( b == e ){
        tree[node].sum = arra[b];
        return;
    int left = node*2:
    int right = node*2 + 1;
    int mid = (b+e)/2;
    init(left,b,mid);
    init(right, mid+1, e);
    tree[node].sum = tree[left].sum +

    tree[right].sum;

    return;
void update( int node,int b,int e,int i,int j,int
    val )
    if( b > j || e < i ) return;
    if( b >= i && e <= j ){
        tree[node].sum += (e-b+1)*val;
        tree[node].prop += val;
        return;
    int left = node*2;
    int right = node*2 + 1;
    int mid = (b+e)/2;
update( left,b,mid,i,j,val );
    update( right,mid+1,e,i,j,val );
    tree[node].sum = tree[left].sum +

    tree[right].sum + (e-b+1)*tree[node].prop;

    return;
int query( int node, int b, int e, int i, int j, int
    carry )
    if( b > j | | e < i ) return 0;
    if( b \ge i \&\& e \le j ) return tree[node].sum +
     \leftarrow (e-b+1)*carry;
    int left = node*2:
    int right = node*2 + 1;
    int mid = (b+e)/2;
    int p1 = query(
     left,b,mid,i,j,carry+tree[node].prop );
    int p2 = query(
     right,mid+1,e,i,j,carry+tree[node].prop );
    return p1+p2;
```

3 Graph

3.1 Articulation Point Detection

```
vector<int> adj[N];
bool vis[N], articulation[N];
int low[N], tin[N], taim;
void dfs(int node, int par = -1)
{
    vis[node] = 1;
    tin[node] = low[node] = taim++;
    int children = 0;
    for(int x : adj[node])
}
```

```
if(x==par) continue;
        if(vis[x]) low[node] =

→ min(low[node],tin[x]);
            dfs(x,node);
             low[node] = min(low[node],low[x]);
            if(low[x] > = tin[node] \&\& par! = -1)
                 articulation[node] = 1;
             children++;
        }
    if(children>1 and par==-1) articulation[node] =
int main()
    int n,m,i;
    while(1)
        scanf("%d %d", \&n, \&m);
        if(!n and !m) break;
        for(i=0;i<=n;i++) adj[i].clear();</pre>
        memset(vis,0,sizeof vis);
        memset(articulation, 0, sizeof articulation);
        memset(tin,-1,sizeof tin);
        memset(low, -1, sizeof low);
        taim = 0:
        while(m--)
            int a,b;
scanf("%d %d",&a,&b);
            adj[a].push back(b);
            adj[b].push_back(a);
        for(i=1;i<=n;i++)
            if(!vis[i]) dfs(i);
        printf("%d\n",count(articulation+1,articula;
         \rightarrow tion+1+n,1));
```

3.2 Bridge Detection

```
dfs(to, v);
             low[v] = min(low[v], low[to]);
             if (low[to] > tin[v]) IS BRIDGE(v, to);
    }
int main()
    int n,m,i,t,tc;
    scanf("%d",&t);
    for(tc=1;tc<=t;tc++)
        scanf("%d %d",&n,&m);
        for(i=0;i<=n;i++) adj[i].clear();</pre>
        timer = 0:
        memset(tin,-1,sizeof tin);
        memset(low, -1, sizeof low);
        memset(visited,0,sizeof visited);
        bridges.clear();
        while(m--)
            int a,b;
scanf("%d %d",&a,&b);
             adj[a].push back(b);
             adj[b].push_back(a);
        for(i=1;i<=n;i++)
            if(!visited[i]) dfs(i);
        sort(bridges.begin(),bridges.end());
```

3.3 Centroid Decomposition

```
// Builds a centroid tree of height O(logn) in
 \rightarrow O(nloan).
const int M = 2e5 + 3;
int sz[M], done[M], cpar[M], root;
|vector<<mark>int</mark>>ctree[M];
void go(int u, int p=-1){
 sz[u] = 1;
 for(int \lor : g[u]){
  if(v == p or done[v]) continue;
  qo(v, u);
  sz[u] += sz[v];
int find centroid(int v, int p, int n){
 for(int \times : a[v]){
  if (x != p \text{ and } !\text{done}[x] \text{ and } sz[x] > n/2) return

    find centroid(x, v, n);

 return v;
void decompose(int v=0, int p=-1){
 int c = find centroid(v, -1, sz[v]);
 if(p == -1) root = c;
 done[c] = 1;
 cpar[c] = p;
```

```
if(p != -1) ctree[p].push back(c);
for(int \times : g[c])
 if(!done[x]) decompose(x, c);
3.4 DSU On Tree
vector <int> G[mx]; /// adjacency list of the tree
int sub[mx]; /// subtree size of a node
int color[mx]; /// color of a node
int freq[mx];
int n;
void calcSubSize(int s,int p) {
    sub[s] = 1;
    for(int \times : G[s]) {
        if(x==p) continue;
        calcSubSize(x,s);
        sub[s] += sub[x];
void add(int s,int p,int v,int bigchild = -1) {
    freq[color[s]] += v;
    for(int \times : G[s]) {
        if(x==p | | x==bigchild) continue;
        add(x,s,v);
void dfs(int s,int p,bool keep) {
    int biaChild = -1:
    for(int \times : G[s]) {
        if(x==p) continue;
        if(bigChild==-1 | | sub[bigChild] < sub[x] )</pre>
         \rightarrow bigChild = x;
    }
    for(int \times : G[s]) {
        if(x==p | | x==bigChild) continue;
        dfs(x,s,0);
    if(bigChild!=-1) dfs(bigChild,s,1);
    add(s,p,1,bigChild);
    /// freq[c] now contains the number of nodes in
    /// the subtree of 'node' that have color c
    /// Save the answer for the gueries here
    if(keep==0) add(s,p,-1);
int main() {
    input color
    construct G
    calcSubSize(root,-1);
    dfs(root, -1,0);
    return 0;
```

3.5 Dijkstra

```
#define pii pair<long long,int>
vector<int>Edges[100005];
vector<long long>Cost[100005];
long long dis[100005];
int vis[100005]:
void dijkstra( int source )
     priority queue< pii, vector<pii>, greater<pii> >0;
     0.push(\overline{p}ii(0,source));
     dis[source] = 0;
    pii q;
while( !Q.empty() ){
          q = 0.top();
0.pop();
          int u = q.second;
          if( vis[u] != -1 ) continue; // *idk why*
vis[u] = 1; // *idk why*
for( int i = 0; i < Edges[u].size(); i++ ){</pre>
               int v = Edges[u][i];
               if( vis[v] != -1 ) continue: // *idk

→ whv*

               if( dis[u] + Cost[u][i] < dis[ v ] ){
    dis[ v ] = dis[u] + Cost[u][i];</pre>
                     Q.push( pii( dis[v],v ) );
     return;
```

3.6 Dinic

```
// O(V^2 E), solves SPOJ FASTFLOW
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
struct edge {
  int u, v;
ll cap, flow;
edge () {}
  edge (int u, int v, ll cap) : u(u), v(v),
  \rightarrow cap(cap), flow(0) {}
struct Dinic {
  int N;
  vector <edge> E;
  vector <vector <int>> q;
  vector <int> d, pt;
  Dinic (int N) : N(N), E(\theta), g(N), d(N), pt(N) {}
  void AddEdge (int u, int v, ll cap) {
    if (u ^ v) {
      E.emplace back(u, v, cap);
      g[u].emplace back(E.size() - 1);
      E emplace ba\overline{c}k(v, u, 0);
      g[v].emplace back(E.size() - 1);
```

```
bool BFS (int S, int T) {
    queue <int> q({S});
    fill(d.begin(), d.end(), N + 1);
    d[S] = 0;
    while (!q.empty()) {
      int u = q.front(); q.pop();
      if (u == T) break:
      for (int k : g[u]) {
        edge &e = E[k];
        if (e.flow < e.cap and d[e.v] > d[e.u] + 1)
          d[e.v] = d[e.u] + 1;
           q.emplace(e.v);
    } return d[T] != N + 1;
  ll DFS (int u, int T, ll flow = -1) {
    if (u == T or flow == 0) return flow;
    for (int \&i = pt[u]; i < g[u].size(); ++i) {
      edge &e = E[g[u][i]];
edge &oe = E[g[u][i] ^ 1];
      if(d[e.v] == d[e.u] + 1) {
        ll amt = e.cap - e.flow;
        if (flow != -1 and amt > flow) amt = flow;
        if (ll pushed = DFS(e.v, T, amt)) {
           e.flow += pushed;
           oe.flow -= pushed;
           return pushed;
    } return 0;
  ll_MaxFlow (int S, int T) {
    ll total = 0;
    while (BFS(S, T)) {
  fill(pt.begin(), pt.end(), 0);
      while (ll flow = DFS(S, T)) total += flow;
    return total;
int main() {
  int N, E;
  scanf("%d %d", &N, &E);
  Dinic dinic(N);
  for (int i = 0, u, v; i < E; ++i) {
   il cap;
scanf("%d %d %lld", &u, &v, &cap);
dinic.AddEdge(u - 1, v - 1, cap);
    dinic.AddEdge(v - 1, u - 1, cap);
  printf("%lld\n", dinic.MaxFlow(0, N - 1));
  return 0:
```

3.7 Heavy Light Decomposition

```
vector < int > List[ ?? ]; // Tree's Adj List ->
→ Need to Clear??
class HeavyLightDecomposition
```

```
#define L R ? ?
|public:
    vector<int> ValueOfNode;
    vector<int> Position;
    vector<int> Parent;
    vector<int> Depth;
    vector<int> Heavy;
    vector<int> Head;
    int CurrentPosition = 1; // 0/1 - index based
            segmentTree ST = segmentTree( ?? ) /

→ AnyQueryTree;

            HeavyLightDecomposition(int NN)
                ValueOfNode.resize(NN);
                Position.resize(NN);
                Parent.resize(NN, -1);
                Depth.resize(NN, 0);
                Heavy.resize(NN, -1);
                Head.resize(NN);
            int DFS(int Vertex)
                int TotalSize = 1;
                int MaxChildSize = 0:
                for (int i = 0; i <

    List[Vertex].size(); ++i)

                    int Child = List[Vertex][i];
                    if (Child != Parent[Vertex])
                        Parent[Child] = Vertex;
                        Depth[Child] =
                        → Depth[Vertex] + 1;
                        int ChildSize = DFS(Child);
                        TotalSize += ChildSize;
                        if (ChildSize >
                            MaxChildSize)
                            MaxChildSize =

→ ChildSize;

                            Heavy[Vertex] = Child;
                return TotalSize;
            void TreeDecompose(int Vertex, int Hd)
                Head[Vertex] = Hd;
                ST.A[CurrentPosition] =
                → ValueOfNode[Vertex];
                Position[Vertex] =
                if (Heavy[Vertex] != -1)
                    TreeDecompose(Heavy[Vertex],

→ Hd);

                for (int i = 0; i < 0

    List[Vertex].size(); ++i)
```

```
int Child = List[Vertex][i];
       if (Child != Parent[Vertex] &&
        TreeDecompose(Child, Child);
void MakeQueryTree()
{ // ?? = Number of Node in Tree;
    // Build Query Data Structure
int Query(int NodeA, int NodeB)
   int Res = 0:
   while (Head[NodeA] != Head[NodeB])
       if (Depth[Head[NodeA]] >
        → Depth[Head[NodeB]])
           swap(NodeA, NodeB);
       int CurrentPathResult =
           ST.rangeQuery(L R,
           Position[Head[NodeB]],
           Position[NodeB]).Value;
       Res = ? ? (Res,
        NodeB = Parent[Head[NodeB]];
   if (Depth[NodeA] > Depth[NodeB])
    swap(NodeA, NodeB);
   int LastHeavyPathResult =
       ST.rangeQuery(L R,
       Position[NodeA]
    Position[NodeB]).Value;
   Res = ? ? (Res,
    return Res;
int Update(int NodeA, int NodeB, int X)
   while (Head[NodeA] != Head[NodeB])
       if (Depth[Head[NodeA]] >
        → Depth[Head[NodeB]])
           swap(NodeA, NodeB);
       ST.rangeUpdate(L R,
           Position[Head[NodeB]],
           Position[NodeB], X);
       NodeB = Parent[Head[NodeB]];
   if (Depth[NodeA] > Depth[NodeB])
       swap(NodeA, NodeB):
   ST.rangeUpdate(L R,
       Position[NodeA],
       Position[NodeB], X);
```

3.8 Hopcroft Karp

};

```
vector<int>q[M];
int Lmatch[M], Rmatch[M], dist[M];
|bool bfs(int n){
 queue<int>q;
 for(int u=1; u<=n; u++){
  if(Lmatch[u] == NIL) dist[u] = 0, q.push(u);</pre>
  else dist[u] = INF;
 dist[NIL] = INF;
 while(!q.empty()){
  int u = q.front();
  q.pop();
  if(dist[u] < dist[NIL]){</pre>
   for(int \lor : g[u]){
    if(dist[Rmatch[v]] == INF){
  dist[Rmatch[v]] = dist[u] + 1;
      q.push(Rmatch[v]);
 return dist[NIL] != INF;
bool dfs(int u){
 if(u == NIL) return true;
 for(int \lor : g[u]){
  if(dist[Rmatch[v]] == dist[u] + 1 and

    dfs(Rmatch[v])){
   Rmatch[v] = u;
   Lmatch[u] = v;
   return true;
 dist[u] = INF;
 return false;
int HopcroftKarp(int n, int m){
  fill(Lmatch, Lmatch+n+1, 0);
 fill(Rmatch, Rmatch+m+1, 0);
 int res = 0;
 while(bfs(n)){
  for(int u=1; u<=n; u++){</pre>
   if(Lmatch[u]==NIL and dfs(u)) res++;
 return res:
3.9 Hungarian
```

```
/*returns maximum/minimum weighted bipartite
    matching. Complexity : O(N^2 * M)
flag = -1 minimizes, flag = 1 maximizes. */
#define CLR(a) memset(a, 0, sizeof a)

ll weight[N][M];
int used[M], P[M], way[M], match[M];
ll U[M], V[M], minv[M], ara[N][M];

ll hungarian(int n, int m, int flag){
    CLR(U), CLR(V), CLR(P), CLR(ara), CLR(way);
    for(int i=1; i<=n; i++){
        for(int j=1; j<=m; j++){</pre>
```

```
ara[i][j] = -flag * weight[i][j];
}
if(n > m) m = n;
int a, b, d;
ll r, w;
for(int i=1; i<=n; i++){</pre>
    P[0] = i, b = 0;
    for(int j=0; j<=m; j++) minv[j] = INF,</pre>

    used[i] = false;

    do{
         used[b] = true;
         a = P[b], d = 0, w = INF;
         for(int j=1; j<=m; j++){
   if(!used[j]){</pre>
                   r = ara[a][j] - U[a] - V[j];
                   if(r < minv[j]) minv[j] = r,</pre>
                    \rightarrow way[j] = b;
                   if(minv[j] < w) w = minv[j], d
                    \rightarrow = j;
         for(int j=0; j<=m; j++){
   if(used[j])_U[P[j]] += w, V[j] -= w;</pre>
              else minv[j] -= w;
         \dot{b} = d;
     }while(P[b] != 0);
     do{
          d = wav[b]:
         P[b] = P[d], b = d;
    }while(b != 0);
for(int j=1; j<=m; j++) match[P[j]] = j;</pre>
return flag*V[0];
```

3.10 LCA(sparse table)

```
vector<int>Edges[10000];
int p[10005][17], level[10005], n, lg;
bool vis[10005];
void DFS( int par,int node )
    vis[node] = 1;
    if( par != -1 ) level[node] = level[par] + 1;
    p[node][0] = par;
    for( int i = 1; i <= lg; i++ ){
        if( p[node][i-1] != -1 ) p[node][i] = p[
         \rightarrow p[node][i-1] ][i-1]:
    for( int i = 0; i < Edges[node].size(); i++ ){</pre>
        if( vis[ Edges[node][i] ] == 0 ) DFS(

¬ node, Edges[node][i] );

    return:
int LCA( int u,int v )
    if( level[u] < level[v] ) swap(u,v);</pre>
    for( int i = lg; i >= 0; i-- ){
```

```
int par = p[u][i];
   if( level[par] >= level[v] ){
        u = par;
   }
}
if( u == v ) return u;
for( int i = lg; i >= 0; i-- ){
   int U = p[u][i];
   int V = p[v][i];
   if( U != V ){
        u = U;
        v = V;
   }
}
return p[u][0];
```

3.11 MCMF(SPFA)

```
* 1 BASED NODE INDEXING
    * call init at the start of every test case
* Complexity --> E*Flow (A lot less actually,
     * Maximizes the flow first, then minimizes the
     * The algorithm finds a path with minimum cost
 → to send one unit of flow
      and sends flow over the path as much as
 → possible. Then tries to find
       another path in the residual graph.
     * SPFA Technique :
         The basic idea of SPFA is the same as
 → Bellman Ford algorithm in that each
         vertex is used as a candidate to relax its
 → adjacent vertices. The improvement
         over the latter is that instead of trying
 → all vertices blindly, SPFA maintains
         a queue of candidate vertices and adds a
   vertex to the queue only if that vertex is relaxed. This process repeats until no
 - more vertex can be relaxed.
         This doesn't work if there is a negative
← cycle in the graph
***/
namespace mcmf {
    using T = int;
    const T INF = 0x3f3f3f3f; // 0x3f3f3f3f or
     → 0x3f3f3f3f3f3f3f3fLL
    const int MAX = 204; // maximum number of nodes
    int n, src, snk;
    T dis[MAX], mCap[MAX];
int par[MAX], pos[MAX];
bool vis[MAX];
    struct Edge{
         int to, rev pos;
         T cap, cost, flow;
    vector <Edge> ed[MAX];
    void init(int _n, int _src, int _snk) {
    n = _n, src = _src, snk = _snk;
         for(int i=1;i<=n;i++) ed[i].clear();</pre>
     void addEdge(int u, int v, T cap, T cost) {
```

```
Edge a = \{v, (int)ed[v].size(), cap, cost,
        Edge b = \{u, (int)ed[u].size(), 0, -cost,
         → 0};
        ed[u].pb(a);
        ed[v].pb(b);
    inline bool SPFA(){
        CLR(vis);
        for(int i=1; i<=n; i++) mCap[i] = dis[i] =</pre>

→ INF;

        aueue <int> a:
        dis[src] = 0;
        vis[src] = true; /// src is in the queue now
        q.push(src);
        while(!q.empty()){
            int u = q.front();
            q.pop();
            vis[u] = false; /// u is not in the

→ queue now

            for(int i=0; i<(int)ed[u].size(); i++) {</pre>
                 Edge &e = ed[u][i];
                 int v = e.to;
                 if(e.cap > e.flow && dis[v] >

    dis[u] + e.cost){
                     dis[v] = dis[u] + e.cost;
                     par[v] = u;
                     pos[v] = i;
                     mCap[v] = min(mCap[u], e.cap -
                     ← e.flow):
                     if(!vis[v]) {
                         vis[v] = true;
                         q.push(v);
            }
        return (dis[snk] != INF);
    inline pair <T, T> solve() {
        T F = 0, C = 0, f;
        int u, v;
        while(SPFA()){
            u = snk:
            f = mCap[u];
            F += f;
            while(u!=src){
                 v = par[u];
                 ed[v][pos[u]].flow += f; // edge of

→ v-->u increases

                 ed[u][ed[v][pos[u]].rev pos].flow
                → -= f;
u = v:
            \dot{C} += dis[snk] * f;
        return make pair(F,C);
ll arr[103] ;
int main() {
    ios::sync with stdio(0);
    ll i,j,k,\overline{l}, \overline{m}, n;
    cin > n > m;
    mcmf::init(n + 2, 1, n + 2);
```

```
for(i = 1; i <= n; i++){
    cin >> arr[i+1];
}
for(i = 0; i < m; i++){
    ll u,v,c;
    cin >> v >> c;
    u++;
    w+minum in the second in the seco
```

```
3.12 MCMF(normal)
#include<bits/extc++.h>
using T=int;
const T kInf=numeric limits<T>::max()/4;
struct mcf graph
struct Edge { int to,from,nxt; T flow,cap,cost; };
vector<Edae>edaes:
int n;
vector<T>dist,pi;
vector<int>par, graph;
 mcf graph(int n):
n(\overline{n}), dist(n), \overline{p}i(n,0), par(n), graph(n,-1) {}
void addEdge(int from,int to,T cap,T cost)
 edges.push back(Edge{to,from,graph[from],0,cap,co_
 graph[from]=edges.size()-1;
void add edge(int from,int to,T cap,T cost)
   addEdge(from, to, cap, cost);
   addEdge(to,from,0,-cost);
bool dijkstra(int s,int t)
 fill(dist.begin(),dist.end(),kInf);
fill(par.begin(),par.end(),-1);
    gnu pbds::priority queue<pair<T, int>>q;
  vector<decltype(q)::point iterator>its(n);
  dist[s]=0; q.push(\{0,s\});
  while(!q.empty())
   int node; T d;
   tie(d,node)=q.top(); q.pop();
   if(dist[node]!=-d) continue;
   for(int i=graph[node];i>=0;)
    const auto &e=edges[i];
    T now=dist[node]+pi[node]-pi[e.to]+e.cost;
    if(e.flow<e.cap && now<dist[e.to])</pre>
     dist[e.to]=now;
```

```
par[e.to]=i;
     if(its[e.to]==q.end())
      its[e.to]=q.push({-dist[e.to],e.to});
     else q.modify(its[e.to], {-dist[e.to], e.to});
    i=e.nxt;
  for(int i=0;i<n;i++)</pre>

→ pi[i]=min(pi[i]+dist[i],kInf);
  return par[t]!=-1;
 pair<T,T> flow(int s,int t)
  T flow=0,cost=0;
  while(dijkstra(s,t))
   T now=kInf;
   for(int node=t;node!=s;)
    int ei=par[node];
    now=min(now,edges[ei].cap-edges[ei].flow);
    node=edges[ei^1].to;
   for(int node=t;node!=s;)
    int ei=par[node];
    edges[ei].flow+=now;
    edges[ei^1].flow-=now;
    cost+=edges[ei].cost*now;
    node=edges[ei^1].to;
   flow+=now;
  return {flow,cost};
//use add edge(from, to, cap, cost) for adding edge
//use flow(s,t) for finding max flow and minimum

→ cost
```

3.13 Max Flow-1

```
return 0;
int max flow( int s,int d )
    int total flow = 0;
    for( int \bar{i} = 1; i \le n; i++){
        for( int j = 1; j <= n; j++ ) rgraph[i][j]</pre>
         \rightarrow = graph[i][i];
    int mn;
    while (bfs(s,d) == 1)
        mn = INT MAX;
        for( int child = d; child != s; child =
         → par[child] ){
            int P = par[child];
            mn = min(mn,rgraph[P][child] );
        for( int child = d; child != s; child =
         → par[child] ){
            int P = par[child];
            rgraph[P][child] -= mn;
            rgraph[child][P] += mn;
        total flow += mn;
    return total flow;
```

3.14 Maximum flow - Edmonds Karp

```
int n:
vector<vector<int>> capacity;
vector<vector<int>> adj;
int bfs(int s, int t, vector<int> &parent)
    fill(parent.begin(), parent.end(), -1);
    parent[s] = -2;
    queue<pair<int, int>> q;
q.push({s, INF});
    while (!q.empty())
        int cur = q.front().first;
        int flow = q.front().second;
        q.pop();
        for (int next : adj[cur])
             if (parent[next] == -1 \&\&
                capacity[cur][next])
                 parent[next] = cur;
                 int new flow = min(flow,

    capacity[cur][next]);
                 if (next == t)
                     return new flow:
                 q.push({next, new flow});
    }
    return 0;
```

```
int maxflow(int s, int t)
{
    int flow = 0;
    vector<int> parent(n);
    int new_flow;
    while (new_flow = bfs(s, t, parent))
    {
        flow += new_flow;
        int cur = t;
        while (cur != s)
        {
            int prev = parent[cur];
                capacity[prev][cur] -= new_flow;
                capacity[cur][prev] += new_flow;
                 cur = prev;
        }
    }
    return flow;
}
```

3.15 Online Bridge

```
vector<int> par, dsu 2ecc, dsu cc, dsu cc size;
int bridges;
int lca iteration;
vector<int> last visit;
void init(int n) {
    par.resize(n):
    dsu 2ecc.resize(n);
    dsu cc.resize(n);
    dsu cc size.resize(n);
    lca_{iteration} = 0:
    las\overline{t} visit.assign(n, 0);
    for \overline{(int i=0; i< n; ++i)} {
         dsu \ 2ecc[i] = i;
         dsu^-cc[i] = i;
         dsu^-cc size[i] = 1;
         parTil = -1;
    bridges = 0;
int find 2ecc(int v) {
    if (\overline{v} == -1)
         return -1:
    return dsu 2ecc[v] == v ? v : dsu 2ecc[v] =

    find 2ecc(dsu 2ecc[v]);

int find cc(int v) {
    v = find 2ecc(v);
    return d\overline{s}u cc[v] == v ? v : dsu cc[v] =
     \rightarrow find \overline{cc}(dsu cc[v]);
void make root(int v) {
    v = find 2ecc(v);
    int root = v;
    int child = -1:
    while (v != -1) {
```

```
int p = find 2ecc(par[v]);
        par[v] = chiTd;
        dsu cc[v] = root;
        chiTd = v;
        v = p;
    dsu cc size[root] = dsu cc size[child];
void merge path (int a, int b) {
    ++lca iteration;
    vector<int> path_a, path_b;
    int lca = -1;
    while (lca == -1) {
        if (a != -1) {
            a = find 2ecc(a);
            path a.push back(a);
            if (\overline{last visit[a]} == lca iteration){
                 lca ≡ a;
                 break;
            last visit[a] = lca iteration;
            a = \overline{par[a]};
        if (b != -1) {
            b = find Žecc(b):
            path b.push back(b);
            if (Tast visit[b] == lca iteration){
                 lca = b;
                 break;
            last_visit[b] = lca_iteration;
            b = \overline{par}[b];
        }
   for (int v : path_a) {
        dsu 2ecc[v] = lca;
        if (v == lca)
            break:
        --bridges;
    for (int v : path b) {
        dsu 2ecc[v] = lca;
        if (v == lca)
            break:
        --bridges;
void add edge(int a, int b) {
    a = find 2ecc(a);
    b = find^2 ecc(b);
    if (a == b)
        return;
    int ca = find cc(a);
    int cb = find cc(b);
   if (ca != cb) {
        ++bridges;
        if (dsu cc size[ca] > dsu cc size[cb]) {
            swap(a, b);
            swap(ca, cb);
        make root(a);
```

```
par[a] = dsu cc[a] = b;
        dsu cc size[cb] += dsu cc size[a];
    } else ₹
        merge path(a, b);
3.16 SCC
/*In a directed graph, an SCC is a connected
    component where all nodes are pairwise
condesation graph is the DAG built on a directed
— graph by compressing each SCC into a node.
define M */
|vector<<mark>int</mark>>g[M], gr[M];
|set<<mark>int</mark>>gc[M];
int vis[M], id[M], sz[M];
vector<int>order, comp, roots;
namespace SCC{
    void addEdge(int u, int v){
        g[u].push back(v), gr[v].push back(u);
    void dfs1(int u){
        vis[u] = 1;
        for(int \times : q[u]){
            if(!vis[x]) dfs1(x);
        order.push back(u);
    void dfs2(int u){
        vis[u] = 1;
        comp.push_back(u);
        for(int \times red : gr[u])
            if(!vis[x]) dfs2(x);
    }
    void condense(int n){
        fill(vis, vis+n+1, 0);
        for(int i=1; i<=n; i++){
            if(!vis[i]) dfs1(i);
        reverse(order.begin(), order.end());
        fill(vis, vis+n+1, 0);
        for(int u : order){
            if(!vis[u]){
                 dfs2(u); //this part of the code
                  _ processes components, returns
                  → them in comp
                 for(int v : comp) id[v] = u;
                 sz[u] = (int)comp.size();
                 roots.push back(u);
                 comp.clear();
        fill(vis, vis+n+1, 0);
        for(int u=1; u<=n; u++){
            for(int v : g[u]){
    if(id[u] != id[v]){
                     gc[id[u]].insert(id[v]);
```

```
}
         }
     void reset(int n){
         order.clear(), comp.clear(), roots.clear();
         for(int i=1; i<=n; i++){
              g[i].clear(), gr[i].clear(),

    gc[i].clear();

              id[i] = vis[i] = sz[i] = 0;
     }
4 Math
4.1 Geometry
4.1.1 2D Point Line _ Segment
const double PI = acos(-1.0);
const double EPS = 1e-12;
    u \cdot v = |u| * |v| * cos(theta)
            = \dot{u}.\dot{x}*\dot{v}.\dot{x} + u.\dot{y}*\dot{v}.\dot{y}
            = How much parallel they are
            = Dot product does not change if one

    vector move perpendicular to the other

    u x v = |u|*|v|*sin(theta)
= u.x*v.y - v.x*u.y
= How much perpendicular they are
= Cross product does not change if one
→ vector move parallel to the other
    dot(a-b,a-b) returns squared distance between
→ pt a and pt b
***/
struct pt {
    double x, y;
     pt() {}
     pt(double x, double y) : x(x) , y(y) {}
     pt operator + (const pt &p) const { return pt(
     \rightarrow x+p.x , y+p.y ); }
     pt operator - (const pt &p) const { return pt(
    - x-p.x , y-p.y ); }
pt operator * (double c) const { return pt( x*c
     \rightarrow , y*c ); }
     pt operator / (double c) const { return pt( x/c
     \rightarrow , y/c); }
```

bool operator == (const pt δp) const { return (

bool operator != (const pt &p) const { return

return os << "("<< p.x << "," << p.y << ")";

inline double dot(pt u, pt v) { return u.x*v.x +

EPS); }

 \rightarrow !(pt(x,y) == p); }

// u.v = |u|*|v|*cos(theta)

// a x b = |a|*|b|*sin(theta)

ostream& operator << (ostream& os, pt p) {

fabs(x - p.x) < EPS && fabs(y - p.y) <

```
inline double cross(pt u, pt v) {return u.x*v.y -

    u.y*v.x;}
// returns |u|
inline double norm(pt u) { return sqrt(dot(u,u)); }
// returns angle between two vectors
inline double angle(pt u,pt v) {
    double cosTheta = dot(u,v)/norm(u)/norm(v);
    return acos(max(-1.0, min(1.0, cosTheta))); //
    → keeping cosTheta in [-1.1]
// returns ang radian rotated version of vector u
// ccw rotation if angle is positive else cw
inline pt rotate(pt u,double ang) {
    return pt( u.x*cos(ang) - u.y*sin(ang) ,
    \rightarrow u.x*sin(ang) + u.y*cos(ang) );
// returns a vector perpendicular to v
inline pt perp(pt u) { return pt( -u.y , u.x ); }
// returns 2*area of triangle
inline double triArea2(pt a,pt b,pt c) { return
cross(b-a,c-a); }
// compare function for angular sort around point P0 }
inline bool comp(pt P0,pt a, pt b) {
    double d = triArea2(P0, a, b);
    if(d < 0) return false;</pre>
   if(d == 0 \&\& dot(P0-a, P0-a) > dot(P0-b, P0-b)
    → ) return false;
    return true;
    if line equation is, ax + by = c
        v --> direction vector of the line (b,-a)
        c --> v cross p
        p --> Any point(vector) on the line
   side'(p) = (v cross p) - c)
            = triArea2(origin, v, p)
   if side(p) is,
        positive --> p is above the line
               --> p is on the line
        negative --> p is below the line
***/
struct line {
    pt v;
    double c:
   line(pt v, double c) : v(v), c(c) {}
    // From equation ax + by = c
    line(double a, double b, double c) : v({b,-a}),
     - c(c) {}
    // From points p and q
    line(pt p, pt q) : v(q-p), c(cross(v,p)) {}
    // |v| * dist
    // dist --> distance of p from the line
    double side(pt p) { return cross(v,p)-c; }
    // better to using sqDist than dist
    double dist(pt p) { return abs(side(p)) /

    norm(v);
]
    double sqDist(pt p) {return side(p)*side(p) /

    dot(v,v);}
    // perpendicular line through point p
    // 90deg ccw rotated line
   line perpThrough(pt p) { return {p, p +
       perp(v)}; }
    // translates a line by vector t(dx,dy)
```

```
// every point (x,y) of previous line is
     \rightarrow translated to (x+dx,y+dy)
    line translate(pt t) {return {v, c +

    cross(v,t)};

    // for every point
    // distance between previous position and
     line shiftLeft(double dist) { return {v, c +
        dist*norm(v)}; }
    // projection of point p on the line
    pt projection(pt p) { return p -
        perp(v)*side(p)/dot(v,v); }
    // reflection of point p wrt the line
    pt reflection(pt p) { return p -
     \rightarrow perp(v)*side(p)*2.0/dot(v,v); }
inline bool lineLineIntersection(line l1, line l2,
    pt &out) {
    double d = cross(l1.v, l2.v);
    if (d == 0) return false;
    out = (l2.v*l1.c - l1.v*l2.c) / d;
    return true;
// interior = true for interior bisector
// interior = false for exterior bisector
linline line bisector(line l1, line l2, bool
   interior) {
    assert(cross(l1.v, l2.v) != 0); // l1 and l2
     double sign = interior ? 1 : -1;
    return {l2.v/norm(l2.v) + (l1.v *

    sign)/norm(l1.v),

            l2.c/norm(l2.v) + (l1.c *

    sign)/norm(l1.v)};
/*** Segment ***/
/// C --> A circle which have diameter ab
/// returns true if point p is inside C or on the
→ border of C
inline bool inDisk(pt a, pt b, pt p) { return
\rightarrow dot(a-p, b-p) <= 0; }
/// returns true if point p is on the segment
inline bool onSegment(pt a, pt b, pt p)
    return triArea2(a,b,p) == 0 \&\& inDisk(a,b,p);
inline bool segSegIntersection(pt a,pt b,pt c,pt
   d,pt &out) {
    if(onSegment(a,b,c)) return out = c, true;
    if(onSegment(a,b,d)) return out = d, true;
    if(onSegment(c,d,a)) return out = a, true;
    if(onSegment(c,d,b)) return out = b, true;
    double oa = triArea2(c,d,a);
    double ob = triArea2(c,d,b);
    double oc = triArea2(a,b,c);
    double od = triArea2(a,b,d);
    if (oa*ob < 0 \&\& oc*od < 0) {
        out = (a*ob - b*oa) / (ob-oa);
        return true:
    return false;
// returns distance between segment ab and point p
inline double segPointDist(pt a,pt b,pt p) {
    if( norm(a-b) == 0 ) {
```

```
line l(a,b);
       pt pr = l.projection(p);
       if(onSegment(a,b,p)) return l.dist(p);
    return min(norm(a-p),norm(b-p));
// returns distance between segment ab and segment
inline double segSegDist(pt a, pt b, pt c, pt d) {
    double oa = triArea2(c,d,a);
    double ob = triArea2(c,d,b);
    double oc = triArea2(a,b,c);
    double od = triArea2(a,b,d);
    if (oa*ob < 0 && oc*od < 0) return 0; // proper
      intersection
    // If the segments don't intersect, the result
    → will be minimum of these four
    return min({segPointDist(a,b,c),

→ segPointDist(a,b,d),

               segPointDist(c,d,a),

    segPointDist(c,d,b)});
```

4.1.2 Circle Line Intersection

```
struct Point {
    double x, y;
Point(double px, double py) {
        x = px;

y = py;
    Point sub(Point p2) {
        return Point(x - p2.x, y - p2.y);
    Point add(Point p2) {
        return Point(x + p2.x, y + p2.y);
    double distance(Point p2) {
        return sqrt((x - p2.x)*(x - p2.x) + (y -
         \rightarrow p2.y)*(y - p2.y));
    Point normal() {
        double length = sqrt(x*x + y*y)
        return Point(x/length, y/length);
    Point scale(double s) {
        return Point(x*s, y*s);
struct line // Creates a line with equation ax +
   by + c = 0
    double a, b, c;
    line() {}
    line( Point p1,Point p2 ) {
        a = p1.y - p2.y;
        b = p2.x - p1.x;
        c = p1.x * p2.y - p2.x * p1.y;
inline bool eq(double a, double b) {
    return fabs(a - b) < eps;
struct Circle {
    double x, y, r, left,right;
```

```
Circle () {}
    Circle(double cx, double cy, double cr) {
        x = cx;
y = cy;
r = cr;
         left = x - r;
         right = x + r;
    pair<Point, Point> intersections(Circle c) {
         Point PO(x, y);
         Point P1(c.x, c.y);
         double d, a, h;
         d = P0.distance(P1);
         a = (r*r - c.r*c.r + d*d)/(2*d);
         h = sqrt(r*r - a*a);
         Point P2 = P1.sub(P0).scale(a/d).add(P0);
         double x3, y3, x4, y4;
         x3 = P2.x + h*(P1.y - P0.y)/d;
         y3 = P2.y - h*(P1.x - P0.x)/d;
         x4 = P2.x - h*(P1.y - P0.y)/d;
         y4 = P2.y + h*(P1.x - P0.x)/d;
         return pair<Point, Point>(Point(x3, y3),
            Point(x4, y4));
inline double Distance( Point a, Point b ) {
    return sqrt( ( a.x - b.x ) * ( a.x - \dot{b}.\dot{x} ) + (
     \rightarrow a.y - b.y ) * (a.y - b.y ));
inline double Distance( Point P, line L ) {
    return fabs( L.a * P.x + L.b * P.y + L.c ) /

    sqrt( L.a * L.a + L.b * L.b );
bool intersection(Circle C, line L, Point &p1, Point
    if( Distance(\{C.x,C.y\}, L) > C.r + eps)

→ return false:

    double a, b, c, d, x = C.x, y = C.y;
    d = C.r*C.r - x*x - y*y;
    if( eq( L.a, 0) ) {
   p1.y = p2.y = -L.c / L.b;
        a = 1;
b = 2 * x;
        c = p1.y * p1.y - 2 * p1.y * y - d;
d = b * b - 4 * a * c;
         d = sqrt(fabs(d));
         p1.x = (b + d) / (2 * a);

p2.x = (b - d) / (2 * a);
    else {
         a = L.a *L.a + L.b * L.b;
         b = 2 * (L.a * L.a * y - L.b * L.c - L.a *
         \rightarrow L.b * x);
         c = L.c * L.c + 2 * L.a * L.c * x - L.a *
          d = b * b - 4 * a * c;
        d = sqrt( fabs(d) );
p1.y = ( b + d ) / ( 2 * a );
         p2.y = (b - d) / (2 * a);
        p1.x = ( -L.b * p1.y -L.c ) / L.a;
p2.x = ( -L.b * p2.y -L.c ) / L.a;
    return true:
```

```
4.1.3 Circle
struct circle {
    pt c;
    double r;
    circle() {}
    circle(pt c, double r) : c(c) , r(r) {}
   returns circumcircle of a triangle
   the radius of circumcircle --> intersection
 → point of the perpendicular
                                     bisectors of
   the three sides */
circle circumCircle(pt a, pt b, pt c) {
    b = b-a, c = c-a; // consider coordinates

→ relative to point a

    assert(cross(b,c) != 0); // no circumcircle if
       A,B,C are co-linear
    // detecting the intersection point using the
        same technique used in line line
       intersection
    pt center = a + (perp(b*dot(c,c) - c*dot(b,b))
     \rightarrow )/cross(b,c)/2);
    return {center, norm(center-a)};
int sqn(double val) {
    if(val>0) return 1;
    else if(val == 0) return 0;
    else return -1:
    returns number of intersection points between a
    line and a circle
    0 --> Center
    I,J --> Intersection points
    P \rightarrow Projection of 0 onto line l
    IP = JP = h , OP = d */
int circleLineIntersection(circle c, line l,
    pair<pt,pt> &out) {
    double h2 = c.r*c.r - l.sqDist(c.c); // h^2
    if (h2 >= 0) { // the line touches the circle
        pt p = l.proj(c.c); // point P
        pt h = l.v*sqrt(h2)/norm(l.v); // vector
         → parallel to l, of length h
        out = \{p-h, p+h\}; //\{I,J\}
    return 1 + sqn(h2); // number of intersection
     → points
    returns number of intersection points between
    two circles
    0 i --> Center of circle i
    I, J --> Intersection points
    P'-- > Projection of O onto line IJ
    IP = JP = h , 0 10 2 = d */
int circleCircleIntersection(circle c1, circle c2,
    pair<pt.pt> &out) {
    pt d = c2.c - c1.c; double d2 = dot(d,d); // d^2
    if (d2 == 0) { // concentric circle
        assert(c1.r != c2.r); // same circle
        return 0;
    double pd = (d2 + c1.r*c1.r - c2.r*c2.r)/2; //
     - = |0 \ 1P| * d
    double h\overline{2} = c1.r*c1.r - pd*pd/d2: // = h^2
    if (h2 >= 0) {
```

```
pt p = c1.c + d*pd/d2, h =

→ perp(d)*sqrt(h2/d2);

                      out = \{p-h, p+h\};
           return 1 + sqn(h2);
        inner --> if true returns inner tangents
           * if the radius of c2 is 0, returns tangents
         that go through the center
                 of circle c2 (value of inner is does not
          matter in this case)
           * if there are 2 tangents, it fills out with
          two pairs of points: the pairs
                 of tangency points on each circle (P1; P2),
          for each of the tangents
           * if there is 1 tangent, the circles are
          tangent to each other at some point
                 P, out just contains P 4 times, and the
        tangent line can be found as
                 line(c1.c,p).perpThrough(p)
           * if there are 0 tangents, it does nothing
           * if the circles are identical, it aborts. */
int tangents(circle c1, circle c2, bool inner,
 → vector < pair <pt,pt> > &out) {
           if (inner) c2.r = -c2.r:
           pt d = c2.c-c1.c;
           double dr = c1.r-c2.r, d2 = dot(d,d), h2 =
                   d2-dr*dr;
           if (d2 == 0 | | h2 < 0) {
                      //assert(h2 != 0):
                       return 0;
           for (double sign : {-1,1}) {
                      pt v = (d*dr + perp(d)*sqrt(h2)*sign)/d2;
                       out push back(\{c1.c + v*c1.r, c2.c + v*c1.r, c3.c + v*c1.r, c3.c
                        \rightarrow v*c2.r}):
           return 1 + (h2 > 0);
```

DU HereForTheFood

4.1.4 Closest Point Pair

```
void solve()
{
   int n;
    cin>>n;
   vector<pair<pii,int>>vec(n);
   for(int i=0;i<n;i++)
   {
      cin>>vec[i].first.first>>vec[i].first.second;
      vec[i].second=i;
   }
   sort(vec.begin(),vec.end());
   ll ans=le18L;
   int a=-1,b=-1;
   set<pair<pii,int>>st;
   for(int i=0,j=0;i<n;i++)
   {
   int d=ceil(sqrt(ans));
   while(j<i &&
      vec[j].first.first+d<vec[j].first.first)
   {
}</pre>
```

```
st.erase({{vec[j].first.second,vec[j].first.firs|

    t},vec[j].second});
 auto it1=st.lower bound({{vec[i].first.second-d,v}
    ec[i].first.first},-1});
 auto it2=st.lower_bound({{vec[i].first.second+d,v}
     ec[i].first.first},-1});
 for(auto it=it1;it!=it2;it++)
  int dx=vec[i].first.first-it->first.second;
  int dy=vec[i].first.second-it->first.first;
  ll curr=1LL*dx*dx+1LL*dy*dy;
  if(curr<ans)</pre>
  cans=curr;
   a=vec[i].second;
   b=it->second;
 st.insert({{vec[i].first.second,vec[i].first.firs
 → t},vec[i].second});
if(a>b) swap(a,b);
cout<<a<<" "<<b<<"
cout<<fixed<<setprecision(6)<<sqrt(ans)<<"\n":
```

4.1.5 Convex_Hull

```
#include<bits/stdc++.h>
using namespace std;
struct Point
int x, y;
Point p0;
Point nextToTop(stack<Point> &S)
 Point p = S.top();
S.pop();
 Point res = S.top();
S.push(p);
return res;
void swap(Point &p1, Point &p2)
Point temp = p1;
p1 = p2;
p2 = temp;
int distSq(Point p1, Point p2)
return (p1.x - p2.x)*(p1.x - p2.x) +
  (p1.y - p2.y)*(p1.y - p2.y);
int orientation(Point p, Point q, Point r)
int val = (q.y - p.y) * (r.x - q.x) -
   (q.x - p.x) * (r.y - q.y);
if (val == 0) return 0;
return (val > 0)? 1: 2;
```

```
int compare(const void *vp1, const void *vp2)
Point *p1 = (Point *)vp1;
Point *p2 = (Point *)vp2;
int o = orientation(p0, *p1, *p2);
|if (o == 0)
 return (distSq(p0, *p2) >= distSq(p0, *p1))? -1:
 return (o == 2)? -1: 1;
void convexHull(Point points[], int n)
int ymin = points[0].y, min = 0;
for (int i = 1; i < n; i++)
 int y = points[i].y;
 if ((y < ymin) \mid | (ymin == y \&\&
  points[i].x < points[min].x))</pre>
  ymin = points[i].y, min = i;
swap(points[0], points[min]);
p0 = points[0];
qsort(&points[1], n-1, sizeof(Point), compare);
int m = 1:
for (int i=1; i<n; i++)</pre>
 while (i < n-1 && orientation(p0, points[i],
         points[i+1]) == 0
  i++;
 points[m] = points[i];
|if (m < 3) return;
|stack<Point> S;
S.push(points[0]);
S.push(points[1]);
S.push(points[2]):
for (int i = 3; i < m; i++)
 while (S.size()>1 && orientation(nextToTop(S),
 S.pop();
 S.push(points[i]);
while (!S.empty())
 Point p = S.top();
 cout << "(" << p.x << ", " << p.y <<")" << endl;
 S.pop();
|int main()
 Point points[100005];
```

4.1.6 Half Plain Intersection

```
const double eps=1e-9;
template<class T>
struct Point
typedef Point P;
T x, y;
 Point(T x=0,T _y=0) { x=_x; y=_y; }
 bool operator<(P p) const { return</pre>

    tie(x,y)<tie(p.x,p.y);
</pre>
 bool operator==(P p) const { return
 \rightarrow tie(x,y)==tie(p.x,p.y); }
 P operator+(P p) const { return P(x+p.x,y+p.y);
 P operator-(P p) const { return P(x-p.x,y-p.y); }
P operator*(T d) const { return P(x*d,y*d); } P operator/(T d) const { return P(x/d,y/d); }
T dot(P p) const { return x*p.x+y*p.y; }
T cross(P p) const { return x*p.y-y*p.x; }
T cross(P a,P b) const { return
     (a-*this).cross(b-*this); }
T dist2() const { return x*x+y*y; }
 double dist() const { return

    sqrt(double(dist2())); }

 double angle() const { return atan2(y,x); }
 P unit() const { return *this/dist(); }
 P perp() const { return P(-v,x); }
 P normal() const { return perp().unit(); }
 P rotate(double a) const
  return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a));
template<class P>
int lineIntersection(const P &s1,const P &e1,const
   P &s2, const P &e2, P &r)
if((e1-s1).cross(e2-s2)) //if not parallel
  r=s2-(e2-s2)*(e1-s1).cross(s2-s1)/(e1-s1).cross(e<sub>1</sub>
   \rightarrow 2-s2);
  return 1;
return -((e1-s1).cross(s2-s1)==0 | | s2==e2);
typedef Point<ld> P;
struct Line
P P1, P2;
 //right hand side of the ray P1 --> P2
Line(P = P(), P = P())
 P1=a;
  P2=b;
```

```
P intpo(Line y)
  assert(lineIntersection(P1,P2,y.P1,y.P2,r)==1);
  return r;
P dir() { return P2-P1; }
bool contains(P x) { return
 \leftarrow (P2-P1).cross(x-P1)<eps; }
bool out(P x) { return !contains(x); }
template<class T>
bool mycmp(Point<T>a,Point<T>b)
if(a.x*b.x<0) return a.x<0;
if(abs(a.x)<eps)</pre>
  if(abs(b.x)<eps) return a.y>0 && b.y<0;
  if(b.x<0) return a.v>0:
 if(b.x>0) return true;
if(abs(b.x)<eps)
 if(a.x<0) return b.y<0;</pre>
 if(a.x>0) return false;
return a.cross(b)>0;
bool cmp(Line a,Line b) { return
mycmp(a.dir(),b.dir()); }
ld Intersection Area(vector<Line>b)
sort(b.begin(),b.end(),cmp);
int n=b.size():
int q=1, h=0;
vector<Line>c(b.size()+10);
for(int i=0:i<n:i++)</pre>
 while(q<h && b[i].out(c[h].intpo(c[h-1]))) h--;
while(q<h && b[i].out(c[q].intpo(c[q+1]))) q++;
  c[++h]=b[i];
  if (q < h \&\& abs(c[h].dir().cross(c[h-1].dir())) < eps)
   if(b[i].out(c[h].P1)) c[h]=b[i];
while (q < h-1 \& \& c[q].out(c[h].intpo(c[h-1]))) h--;
while (q < h-1 \& \& c[h] . out(c[q] . intpo(c[q+1]))) q++;
if(h-q<=1) return 0;</pre>
c[h+1]=c[q];
vector<P>s;
for(int i=q;i<=h;i++)</pre>
 s.push back(c[i].intpo(c[i+1]));
s.push bac\overline{k}(s[0]);
ld ans=0:
for(int i=0;i<int(s.size())-1;i++)</pre>

¬ ans+=s[i].cross(s[i+1]):

ans/=2.0;
return ans;
void solve()
int n;
cin>>n:
```

4.1.7 Point Inside Poly (Ray Shooting)

4.2 Matrices

4.2.1 Gauss-Jordan Elimination in GF(2)

```
const int SZ = 105;
const int MOD = 1e9 + 7;
bitset <SZ> mat[SZ];
int where[SZ];
bitset <SZ> ans;
ll bigMod(ll a,ll b,ll m){
    ll ret = 1LL:
    a \% = m;
    while (b){
        if (b & 1LL) ret = (ret * a) % m;
        a = (a * a) % m;
        b >>= 1LL;
    return ret:
/// n for row, m for column, modulo 2
int GaussJordan(int n,int m) {
    SET(where); /// sets to -1
    for(int r=0, c=0; c<m && r<n; c++) {
        for(int i=r; i<n; i++)</pre>
             if( mat[i][c] ){
                 swap(mat[i],mat[r]); break;
        if( !mat[r][c] ) continue;
        where [c] = r;
        for (int i=0; i<n; ++i) if (i != r &&
            mat[i][c]) mat[i] ^= mat[r]:
```

4.2.2 Gauss-Jordan Elimination in GF(P)

```
const int SZ = 105:
const int MOD = 1e9 + 7:
int mat[SZ][SZ], where[SZ], ans[SZ];
ll bigMod(ll a,ll b,ll m){
    ll ret = 1LL;
    a \%= m;
    while (b){
         if (b & 1LL) ret = (ret * a) % m;
         a = (a * a) % m:
         b >>= 1LL:
    return ret;
int GaussJordan(int n,int m,int P) {
    SET(where); /// sets to -1
    for(int r=0, c=0; c<m && r<n; c++) {
         int mx = r:
         for(int i=r; i<n; i++) if( mat[i][c] >
          \rightarrow mat[mx][c] ) mx = i;
         if( mat[mx][c] == 0 ) continue;
         if(r != mx) for(int j=c; j<=m; j++)

    swap(mat[r][i],mat[mx][i]);

         where [c] = r;
         int mul, minv = bigMod(mat[r][c],P-2,P);
         int temp;
         for(int i=0; i<n; i++){</pre>
              if( i!=r && mat[i][c]!=0){
    mul = ( mat[i][c] * (long long)
                    → minv ) % P:
                  for(int j=c; j<=m; j++) {
   temp = mat[i][j];
   temp -= ( ( mul * (long long))</pre>
                        → mat[r][j] ) % P );
                        temp += P;
                        if( temp >= P ) temp -= P;
                        mat[i][j] = temp;
             }
         }
r++;
```

```
for(int j=0; j<m; j++) {
    if(where[j]!=-1) ans[j] = (
        mat[where[j]][m] * 1LL *
        bigMod(mat[where[j]][j],P-2,P) ) % P;
    else ans[j] = 0;
}
for(int i=0; i<n; i++) {
    int sum = 0;
    for(int j=0; j<m; j++) {
        sum += (ans[j] * 1LL * mat[i][j] ) % P;
        if(sum >= P) sum -= P;
}
if( sum != mat[i][m] ) return 0; /// no
        - solution
}
int cnt = 0;
for(int j=0; j<m; j++) if (where[j]==-1) cnt++;
return bigMod(P,cnt,MOD);
}</pre>
```

4.2.3 Gauss-Jordan Elimination

```
* mat is 0 based
    * In every test case, clear mat first and then

→ do the changes

    * For solving problems on graphs with
 → probability/expectation, make sure the graph
      is connected and a single component. If not,
 → then re-number the vertex and solve
      for each connected component separately.
    * Complexity --> O( min(n,m) * nm )
const int SZ = 105;
const double EPS = 1e-9;
double mat[SZ][SZ], ans[SZ];
int where[SZ];
int GaussJordan(int n,int m) {
    SET(where); \/// sets to -1
    for(int r=0, c=0; c< m && r< n; c++) {
        int mx = r;
        for(int i=r; i<n; i++) if( abs(mat[i][c]) >
         \rightarrow abs(mat[mx][c]) ) mx = i;
        if( abs(mat[mx][c]) < EPS ) continue;</pre>
        if(r != mx) for(int j=c; j<=m; j++)
         → swap(mat[r][j],mat[mx][j]);
        where [c] = r;
        for(int i=0; i<n; i++) if( i!=r )
    double mul = mat[i][c]/mat[r][c];
    for(int j=c; j<=m; j++) mat[i][j] -=</pre>

    mul*mat[r][j];

        f++;
    for(int j=0; j<m; j++) {
    if(where[j]!=-1) ans[j] =</pre>

    mat[where[j]][m]/mat[where[j]][j];

        else ans[j] = 0;
    for(int i=0; i<n; i++){
        double sum = 0;
        for(int j=0; j<m; j++) sum += ans[j] *</pre>
            mat[i][i]:
```

```
if( abs(sum - mat[i][m]) > EPS ) return 0;
         → /// no solution
    for(int j=0; j<m; j++) if (where[j]==-1) return</pre>

→ INF;

    return 1;
4.2.4 Space of Binary Vectors
// A vector can be added to the space at any moment
// Following queries can be made on the current
 → basis at anv moment
const int B = ?;
struct space {
    int base[B]:
    int sz;
    void init() {
        CLR(base); sz = 0;
    // if the vector val is not currently in the
       vector space,
    // then adds val as a basis vector
    void add(int val) {
        for(int i = B-1; i >= 0; i--) {
            if( val & (1<<i) ) {
                if(!base[i]) {
   base[i] = val; ++sz; return;
                else val ^= base[i]:
    int getSize() { return sz; }
    // returns maximum possible ret such that, ret
     \rightarrow = (val ^ x)
    // and x is also in the vector space of the
     int getMax(int val) {
        int ret = val;
        for(int i = B - 1; i >= 0; i--) {
            if(ret & (1 << i)) continue;
            ret ^= base[i];
        return ret:
    // returns minimum possible ret such that, ret
     \rightarrow = (val ^ x)
    // and x is also in the vector space of the
     int getMin(int val) {
        int ret = val;
        for(int i = B - 1; i >= 0; i--) {
            if( !(ret \& (!<<i)) ) continue;
            ret ^= base[i];
        return ret;
    }
// returns true if val is in the vector space
    bool possible(int val) {
        for(int i = B - 1; i >= 0; i - -) {
            if(val & (1<<i)) val ^= base[i];
        return (val == 0);
```

```
// returns the k'th element of the current

→ vector space

    // if we sor't all the elements according to

→ their values

    int query(int k) {
         int ret = 0;
         int tot = 1 << getSize();</pre>
        for(int i = B - 1; i >= 0; i--) {
   if(!base[i]) continue;
             int low = tot >> 1;
             if ((low < k \&\& (ret \& 1 << i) == 0) | |
              \rightarrow (low >= k && (ret & 1 << i) > 0))
                  ret ^= base[i]:
             if (low < k) k -= low;
             tot /= 2:
        return ret;
};
```

4.2.5 matrix_exponentiation

```
#include<bits/stdc++.h>
using namespace std;
long long a, b, n, m, F[2][2], f[2][2];
long long p = 1e9 + 7;
void multiply( long long a[2][2],long long b[2][2])
    long long g[2][2];
    for( int i = 0; i < 2; i++){
       for( int j = 0; j < 2; j++ ){
           g[i][j] = 0; for( int k = 0; k < 2; k++
            ) g[i][j] = ((g[i][j]%p) +
            for( int i = 0; i < 2; i++){ for( int j = 0; j
    \prec < 2; j++ ) F[i][j] = q[i][j];}
void power( long long N )
    if(N == 1) return:
   if( N%2 == 0 ){ power( N/2 ); multiply(F,F); }
    else{power(N-1); multiply(F,f);}
    return:
```

4.3 Modular Arithmatic

4.3.1 Chinese Remainder Theorem

```
/***

X = a 1 % m 1

X = a 2 % m 2

X = a 3 % m 3

m_1,m_2,m_3 are pair wise co-prime

M = m_1*m_2*m_3

u_i = Modular inverse of (M/m_i) with respect

m i
```

```
X = (a1 * (M/m 1) * u 1 + a 2 * (M/m 2) * u 2
   + a 3 * (M/m 3) * u 3 ) % M
ll inv(ll a, ll m)
    ll m0 = m, t, q;
    11 \times 0 = 0, \times 1 = 1;
    if (m == 1) return 0;
    while (a > 1)
        q = a / m; t = m; m = a % m, a = t; t = x0;
         \rightarrow x0 = x1 - q * x0; x1 = t;
    if (x1 < 0) x1 += m0;
    return x1;
Îl findMinX(ll num[], ll rem[], ll k)
    ll prod = 1;
    for (ll i = 0; i < k; i++) prod *= num[i];
    ll result = 0;
    for (ll i = 0; i < k; i++)
        ll pp = prod / num[i];
        result += rem[i] * inv(pp, num[i]) * pp;
    return result % prod;
int main()
    ll num[15],rem[15],n,t,i,j;
scanf("%lld",&t);
    for(i=1;i<=t;i++)
        scanf("%lld",&n);
        for(j=0; j<n; j++)
            scanf("%lld %lld",&num[j],&rem[j]);
        printf("Case %lld:
         → %lld\n",i,findMinX(num,rem,n));
```

4.3.2 Discrete log

```
int solve(int a, int b, int m) {
    a \%= m, b \%= m;
    int k = 1, add = 0, g;
    while ((g = gcd(a, m)) > 1) {
         if (b == k)
             return add:
        if (b % g)
             return -1;
        b /= g, m /= g, ++add;

k = (k * 111 * a / g) % m;
    int n = sqrt(m) + 1;
    int an = 1;
    for (int i = 0; i < n; ++i)
         an = (an * 111 * a) % m:
    unordered map<int, int> vals;
    for (int \overline{q} = 0, cur = b; q \le n; ++q) {
         vals[cur] = q;
```

```
OLIN | ITTEHAD | INZAMAM
        cur = (cur * 1ll * a) % m;
    for (int p = 1, cur = k; p \le n; ++p) {
        cur = (cur * 1ll * an) % m;
        if (vals.count(cur)) {
            int ans = n * p - vals[cur] + add;
            return ans;
    return -1;
4.3.3 Modular inverse EGCD
<u>#include<bits/std</u>c++.h>
using namespace std;
int gcdExtended(int a, int b, int* x, int* y)
    if (a == 0)
        *x = 0, *y = 1;
        return b;
    int x1, y1;
    int gcd = gcdExtended(b % a, a, &x1, &y1);
    *x = y1 - (b / a) * x1;
    *y = x1;
    return gcd;
void modInverse(int a, int m)
    int x, y;
    int g = gcdExtended(a, m, &x, &y);
    if (q != 1)
        printf("Inverse doesn't exist");
        int res = (x % m + m) % m;
        printf("Modular multiplicative inverse is
        int main()
    int a, m;
    cin >> a >> m:
    modInverse(a, m);
    return 0;
4.3.4 nCr Lucas
/*use this to calculate nCr modulo mod, when mod is
→ smaller than n and m. define MOD
Complexity: O(mod + log mod n) */
```

```
ll fact[MOD];
ll bigmod(int x, int p){
   ll res = 1;
   while(p){
        if(p \& 1) res = res * x % MOD;
```

```
DU HereForTheFood
        x = x * x % MOD;
        p >>= 1;
    return res;
ll modinv(ll x){
return bigmod(x, MOD-2);
void precalc(){ //run this
fact[0] = 1;
for(int i=1; i<MOD; i++){</pre>
 fact[i] = fact[i-1] * i % MOD;
int C(int n, int m){
if(m > n) return 0;
if(m == 0 \text{ or } m == n) \text{ return } 1;
ll ret = fact[n] * modinv(fact[m]) % MOD;
return ret * modinv(fact[n-m]) % MOD;
int nCr(int n, int m){
if(m > n) return 0;
if(m == 0) return 1;
return nCr(n/MOD, m/MOD)*C(n%MOD, m%MOD) % MOD;
4.4 Polynomial Multiplication
4.4.1 FFT
typedef cplx cd;
//define N as a power of two greater than the size

→ of any possible polynomial

using cd = complex<double>;
const double PI = acosl(-1);
int rev[N]; cd w[N];
static cd f[N];
void prepare(int &n){
    int sz = builtin ctz(n);
    for(int i=1; i<n; i++) rev[i] = (rev[i>>1] >>
    \rightarrow 1) | ((i & 1) << (sz - 1));
    w[0] = 0, w[1] = 1, sz = 1;
    while (1 \ll sz < n)
        cd w n = cd(cos(2*PI / (1<<(sz+1))),
         \rightarrow \sin(2*PI / (1 << (sz+1)));
```

for(int $i=1 << (sz-1); i < (1 << sz); i++){$ [w[i << 1] = w[i], w[i << 1 | 1] = w[i] *

if(i < rev[i]) swap(a[i], a[rev[i]]);

 \hookrightarrow w_n;

for(int i=1; i<n-1; i++){

for(int h=1; h<n; h<<=1){

for(int s=0; s<n; s+=h<<1){

SZ++;

void fft(cd *a, int n){

}

```
for(int i=0; i<h; i++){</pre>
                 cd \& u = a[s+i], \& v = a[s+i+h], t =

  v*w[h+i]:

                 v = u-t, u = u+t;
        }
vector<ll>multiply(vector<ll>a, vector<ll>b){
    int n = a.size(), m = b.size(), sz = 1;
    if(!n or !m) return {};
    while (sz < n+m-1) sz <<= 1;
    prepare(sz);
    for(int i=0; i<sz; i++) f[i] = cd(i < n ? a[i]</pre>
    \rightarrow : 0, i < m ? b[i] : 0);
    fft(f, sz);
    for(int i=0; i<=(sz>>1); i++){
        int j = (sz - i) \& (sz - 1)
        cd x = (f[i]*f[i] - conj(f[j]*f[j])) *
         \rightarrow cd(0, -0.25);
        f[j] = x, f[i] = conj(x);
   fft(f, sz);
    vector<ll>c(n+m-1);
    for(int i=0; i<n+m-1; i++) c[i] =

¬ round(f[i].real()/sz);

    return c;
```

4.4.2 NTT

```
const int G = 3;
const int MOD = 998244353;
const int N = ?; // (1 << 20) + 5; greater than</pre>
→ maximum possible degree of any polynomial
int rev[N], w[N], inv n;
int bigMod(int a, int e, int mod) {
    if(e == -1) assert(false);
    if(e == -1) e = mod - 2;
    int ret = 1;
    while (e) {
        if (e \& 1) ret = (ll) ret * a % mod;
        a = (ll) a * a % mod; e >>= 1;
    return ret;
void prepare(int &n) {
    int sz = abs(31 -
                        builtin clz(n));
    int r = bigMod(G, \overline{(MOD - 1)} / n, MOD);
    inv n = bigMod(n, \dot{M}OD - 2, \dot{M}OD), w[0] = w[n] =
    for (int i = 1; i < n; ++i) w[i] = (ll) w[i -

→ 1] * r % MOD;

    for (int i = 1; i < n; ++i) rev[i] = (rev[i >>
    \rightarrow 1] >> 1) | ((i & 1) << (sz - 1));
void ntt (int *a, int n, int dir) {
    for (int i = 1; i < n - 1; ++i) {
        if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
    for (int m = 2; m <= n; m <<= 1) {
        for (int i = 0; i < n; i += m) {
            for (int j = 0; j < (m >> 1); ++j) {
```

```
int \&u = a[i + j], \&v = a[i + j + j]
                  \rightarrow (m \gg 1)];
                 int t = (ll) v * w[dir ? n - n / m]
                 → * i : n / m * j] % MOD;
                 v = u - t < 0 ? u - t + MOD : u - t;
                 u = u + t >= MOD ? u + t - MOD : u
    if (dir) for (int i = 0; i < n; ++i) a[i] =
     |int f a[N], f b[N];
vector <int> multiply (vector <int> a, vector <int>
    int sz = 1, n = a.size(), m = b.size();
    while (sz < n + m - 1) sz <<= 1; prepare(sz);
    for (int i = 0; i < sz; ++i) f a[i] = i < n ?
    for (int i = 0; i < sz; ++i) f b[i] = i < m?
     \rightarrow b[i] : 0;
    ntt(f a, sz, 0); ntt(f b, sz, 0);
    for (int i = 0; i < sz; ++i) f a[i] = (ll)

    f a[i] * f b[i] % MOD;

    ntt(fa, sz, 1); return vector \langle int \rangle (fa, fa
     \rightarrow + n + m - 1);
I/I G = primitive root(MOD)
int primitive root (int p) {
    vector < int> factor;
    int tmp = p - 1;
    for (int i = 2; i * i <= tmp; ++i) {</pre>
        if (tmp \% i == 0) {
             factor.emplace back(i);
             while (tmp \% i^-== 0) tmp /= i:
    if (tmp != 1) factor.emplace back(tmp);
    for (int root = 1; ; ++root) [
        bool flag = true;
        for (int i = 0; i < (int) factor.size();</pre>

→ ++i) {
             if (bigMod(root, (p - 1) / factor[i],
             \rightarrow p) == 1) {
                 flag = false; break;
        if (flag) return root;
int main() { //(x+2)(x+3) = x^2 + 5x + 6
    vector \langle int \rangle a = \{2, 1\};
    vector < int > b = {3, 1};
    vector <int> c = multiply(a, b);
    for (int x : c) cout << x << " "; cout << endl;</pre>
    return 0:
4.5 2-SAT
    * 1 based index for variables
```

```
* F = (a \text{ op } b) \text{ and } (c \text{ op } d) \text{ and } \ldots (v \text{ op } z)
      a, b, c ... are the variables
      sat::satisfy() returns true if there is some
   assignment(True/False)
     for all the variables that make F = True
    * init() at the start of every case
namespace sat{
    #define CLR(ara,n) fill(ara+1,ara+n+1,0)
const int MAX = 200010; /// number of variables * 2
    bool vis[MAX]:
    vector <int> ed[MAX], rev[MAX];
    int n, m, ptr, dfs t[MAX], ord[MAX], par[MAX];
    inline int inv(int x){
        return ((x) \le n') (x + n) : (x - n);
    }
/// Call init once
    void init(int vars){
        n = vars, m = vars << 1;
        for (int i = 1; i \le m; i++){}
            ed[i].clear();
            rev[i].clear();
    /// Adding implication, if a then b ( a --> b )
    inline void add(int a, int b){
        ed[a].push back(b);
        rev[b].push back(a);
    /// (a or b) is true --> OR(a,b)
    /// (a or b) is true --> OR(inv(a),b)
    /// (a or b) is true --> OR(a,inv(b))
    /// (a or b) is true --> OR(inv(a),inv(b))
    inline void OR(int a, int b){
        add(inv(a), b);
        add(inv(b), a);
    /// same rule as or
    inline void AND(int a, int b){
        add(a, b);
        add(b, a);
    /// same rule as or
    void XOR(int a,int b){
        add(inv(b), a);
        add(a, inv(b));
        add(inv(a), b);
        add(b, inv(a));
    /// same rule as or
    inline void XNOR(int a, int b){
        add(a,b);
        add(b,a);
        add(inv(a), inv(b));
        add(inv(b), inv(a));
    /// (x <= n) means forcing variable x to be true
    /// (x = n + y) means forcing variable y to be

→ false

    inline void force true(int x){
        add(inv(x), x):
    inline void topsort(int s){
        vis[s] = true;
```

}

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```
for(int x : rev[s]) if(!vis[x]) topsort(x);
    dfs t[s] = ++ptr;
inline void dfs(int s, int p){
    par[s] = p;
    vis[s] = true;
    for(int x : ed[s]) if (!vis[x]) dfs(x, p);
void build(){
    CLR(vis,m);
    ptr = 0;
    for(int i=m;i>=1;i--) {
        if (!vis[i]) topsort(i);
        ord[dfs t[i]] = i;
    ĆLR(vis,m);
    for (int i = m; i >= 1; i - -){
        int x = ord[i];
        if (!vis[x]) dfs(x, x);
    }
/// Returns true if the system is 2-satisfiable
and returns the solution (vars set to true)
→ in vector res
bool satisfy(vector < int>& res){
    build();
    CLR(vis,m);
    for (int i = 1; i \le m; i++){
        int x = ord[i];
        if (par[x] == par[inv(x)]) return false;
        if (!vis[par[x]]){
            vis[par[x]] = true;
            vis[par[inv(x)]] = false;
    res.clear();
    for (int i = 1; i \le n; i++){
        if (vis[par[i]]) res.push back(i);
    return true;
#undef CLR
```

4.6 Catalan number

```
4.7 Diophantine Equation
int qcd extend(int a, int b,int& x, int& y)
     if (b == 0) {
         x = 1;
         v = 0;
         return a;
     else {
         int g = gcd extend(b,a % b, x, y);
         int x1 = x, y1 = y;
         x = y1;
         y = x1 - (a / b) * y1;
         return q;
void print solution(int a, int b, int c)
     int x, y;
     if (a == 0 && b == 0) {
         if (c == 0) {
   cout << "Infinite Solutions Exist" <<</pre>
               ← endl;
         else {
              cout << "No Solution exists" << endl;</pre>
     int gcd = gcd extend(a, b, x, y);
     if (\tilde{c} \% \text{ gcd} != 0) {
         cout<< "No Solution exists"<< endl;</pre>
         cout << "x = " << x * (c / gcd) << ", y = "
          \rightarrow << v * (c / acd)<< endl:
```

4.8 Euler_Totient

```
int phi(int n) {
    int result = n;
    for (int i = 2; i * i <= n; i++) {
        if (n % i == 0) {
            while (n \% i == 0)
                n /= i;
            result -= result / i;
    if (n > 1)
        result -= result / n;
    return result;
void phi 1 to n(int n) {
    vector<int> phi(n + 1);
    phi[0] = 0;
    phi[1] = 1;
    for (int i = 2; i \le n; i++)
        phi[i] = i;
    for (int i = 2; i <= n; i++) {
        if (phi[i] == i) {
            for (int j = i; j \le n; j += i)
                phi[j] -= phi[j] / i;
```

```
4.9 Extended Euclid
    c = qcd(a,b);
    ax + by = c;
    (bq + r)x + by = c;
    bq\dot{x} + rx + by = c;
    b(qx + y) + rx = c;
    bx' + ry' = c;
                         [r = a \% b]
    We get.
    x' = qx + y;
    V' = X
    So,
    y = x' - qx;
    y = x' - qy';
                         [v' = x]
    and
    x = v'
    If c is not the gcd then,
    actual x = x * (c/gcd)
actual y = y * (c/gcd)
    But if gcd doesn't divide c, there is no
    solution.
    returns (x,y) for ax + by = gcd(a,b)
/// keep in mind that if a or b or both are

→ negative, gcd(a,b) will be negative

PLL extEuclid(ll a,ll b)
    if(b==0LL) return mp(1LL,0LL);
    PLL ret,got;
got = extEuclid(b,a%b);
    ret = mp(got.yy,got.xx-(a/b)*got.yy);
    return ret:
    From one solution (x0,y0), we can obtain all
   the solutions of the given equation.
    Let g = gcd(a,b) and let x0, y0 be integers
   which satisfy the following:
    a*x0+b*v0 = c
    Now, we should see that adding b/g to x0 and at
   the same time subtracting a/g
    from y0 will not break the equality:
      a^*(x0 + b/g) + b^*(y0 - a/g)
    = a*x0 + b*y0 + (a*b)/q - (b*a)/q
    Obviously, this process can be repeated again,
    so all the numbers of the form:
    x = x0 + k * (b/g)
    y = y0 - k * (a/g)
    are solutions of the given Diophantine equation.
    In the solution returned by extEuclid:
        and |y| is minimized
        \leq b/2q
```

```
|v| <= a/2g
    Because we get a new x after every b/g amount
    and we get a new y after every a/g amount of
    Solution with minimum (x+y):
   x + y = x0 + y0 + k*(b/g - a/g)
   x + y = x0 + y0 + k*((b-a)/q)
    If b>a, we need to find the k with the minimum
    else we need to find the k with the maximum

    value

***/
/// Iterative Implementation
PLL extEuclid(ll a, ll b){
    ll s = 1, t = 0, st = 0, tt = 1;
    while(b) {
        s = s - (a/b)*st;
        swap(s,st);
        t = t - (a/b)*tt;
        swap(t,tt);
        a = a \% b:
        swap(a,b);
    return mp(s,t);
/// returns number of solutions for the equation ax
/// where minx <= x <= maxx and miny <= y <= maxy
ll numberOfSolutions(ll a, ll b, ll c, ll minx, ll
   maxx,ll miny,ll maxy)
    if(a==0 && b==0){
        if(c!=0) return 0;
        else return (maxx-minx+1)*(maxy-miny+1);
          ///  all possible (x,y) within the
         - ranges can be a solution
    il gcd = \underline{gcd(a,b)};
    if(c%gcd!=0) return 0;/// no solution ,

    gcd(a,b) doesn't divide c

    /// If b==0, x will be fixed, any y in the
    \rightarrow range can form a pair with that x
    if(b==0){
        c /= a;
        if(c>=minx && c<=maxx) return maxy-miny+1;</pre>
        else return 0;
    /// If a==0, x will be fixed, any x in the
    → range can form a pair with that y
    if(a==0){
        c /= b;
        if(c>=miny && c<=maxy) return maxx-minx+1;</pre>
        else return 0;
    /// gives a particular solution to the equation
       ax + by = gcd(a,b) \{gcd(a,b) \ can \ be

→ negative also?

    PLL sol = extEuclid(a,b);
    a /= qcd;
    b /= qcd;
    c /= gcd;
```

```
ll x,y;
    x = sol.xx*c;
    y = sol.yy*c;
    il lx,ly,rx,ry;
/// lx -> minimum value of k such that sol.xx +
     \rightarrow k * (b/q) is in range[minx,maxx]
    /// rx -> maximum value of k such that sol.xx +
     \rightarrow k * (b/g) is in range[minx,maxx]
    if(x<minx) lx = ceil( (minx-x) / (double)abs(b)</pre>
     → );
    else lx = -floor( (x-minx) / (double)abs(b) );
    if(x<maxx) rx = floor((maxx-x) / (double)abs(b)</pre>
    else rx = -ceil((x-maxx) / (double)abs(b) );
/// Doing this I because I ignored sign of b
        before passing to getCeil/getFloor
    if(b<0){
         lx *= -1:
         rx *= -1;
         swap(lx,rx);
    if(lx>rx) return 0;
    /// ly -> minimum value of k such that sol.yy -
     \rightarrow k * (a/g) is in range[miny,maxy]
    /// ry -> maximum value of k such that sol.yy -
       k * (a/q) is in range[miny, maxy]
    if(y<miny) ly = ceil( (miny-y) / (double)abs(a)</pre>
    else ly = -floor( (y-miny) / (double)abs(a) );
    if(y<maxy) ry = floor( (maxy-y) /</pre>
     else ry = -ceil( (y-maxy) / (double)abs(a) );
/// Doing this because I ignored sign of a
         before passing to getCeil/getFloor
    if(a<0){
         ly *= -1;
         rý *= -1:
         swap(ly,ry);
    if(ly>ry) return 0;
    lv *= -1;
    ry *= -\bar{1};
    swap(ly,ry);
    /// getting the intersection between (x range)
         and (y range) of k
    ll li = max(lx,ly);
    ll ri = min(rx,ry);
    return max(`ri´-`li + 1 , 0LL );
4.10 Mobius Function
```

```
mu[1] = 1.
    mu[n] = 0 if n has a squared prime factor,
    mu[n] = 1 if n is square-free with even number
   of prime factors
    mu[n] = -1 if n is square-free with odd number
   of prime factors
    *** sum of mu[d] where d \mid n is 0 ( For n=1,
   sum is 1 )***
|int mu[MAX] = {0};
void Mobius(int N){
```

```
int i, j;
mu[1] = 1;
for (i = 1; i \le N; i++){
    if (mu[i]){
        for (j = i + i; j \le N; j += i){
            mu[i] -= mu[i];
}
```

```
4.11 Primes_stuff
const int N = 10000000+6;
vector<long long>primes;
bitset<N>flag;
vector<long long>v;
void siv()
    flag[1] = 1;
    for( int i = 2; i*i <= N; i++){
        if( flag[i] == 0 ){
            for (int j = i*i; j < N; j+= i)
             \rightarrow flag[i] = 1:
    for( int i = 2; i < N; i++ ){
        if( flag[i] == 0 ) primes.push back(i);
long long mul(long long a,long long b,long long mod)
    long long res = 0;
    a %= mod;
    while (b)
        if (b & 1) res = (res + a) % mod;
        a = (2 * a) % mod;
        b >>= 1; // b = b / 2
    return res;
long long gcd( long long x,long long y )
    if( x < 0 ) x = -x;
    if(y < 0) y = -y;
    if( !x || !y ) return x+y;
    long long temp;
    while( x%y ){
        temp = x;
        x = y;
y = temp%y;
    return y;
long long mod inverse( long long n,long long p )
    long long x, y, g;
    g = gcd_extended( n,p,x,y );
    if( q < 0 ) x = -x;
    return (x%p +p)%p;
long long mpow( long long x,long long y,long long
```

```
long long ret = 1;
    while( y ){
        if( y&1 ) ret = mul(ret,x,mod);
        y \gg 1, x = mul(x,x,mod);
    return ret%mod;
int isPrime( long long p )
   if( p < 2 || !(p&1) ) return 0;
if( p == 2 ) return 1;</pre>
    long long q = p-1, a, t;
    int k = 0, b = 0;
    while (!(q\&1)) q >>=1, k++;
    for( int it = 0; it < 2; it++ ){</pre>
        a = rand()\% (p-4) + 2;
        t = mpow(a,q,p);
b = (t==1) || (t == p-1);
        for (int i = 1; i < k \&\& !b; i++)
            t = mul(t,t,p);
            if( t == p-1 ) b = 1;
        if( b == 0 ) return 0;
    return 1;
long long pollard rho( long long n,long long c )
    long long x = 2, y = 2, i = 1, k = 2, d;
    while( 1 ){
        x = (mul(x,x,n) + c);
        if( x >= n ) x -= n;
        d = gcd(x-y,n);
        if(d > 1) return d;
        if( ++i == k ) y = x, k <<= 1;
    return n;
map<long long,int>mp;
void factorize( long long n )
    int l = primes.size();
    for( int i = 0; primes[i]*primes[i] <= n && i <</pre>
    → l; i++ ){
        if( n%primes[i] == 0 ){
            mp[primes[i]] = 1;
            while( n%primes[i] == 0 ) n/= primes[i];
    if( n != 1 ) mp[n] = 1;
void lfactorize( long long n )
    if( n == 1 ) return;
    if( n < 1e9 ){
        factorize(n);
        return:
    if( isPrime(n) ){
        mp[n] = 1;
        return;
    long long d = n;
    for( int i = 2; d == n; i++ ) d =
       pollard rho(n.i):
```

```
lfactorize(d);
    lfactorize(n/d);
long long f(long long r,vector<long long> v1){
    int sz=v1.size();
    long long res=0;
    for(long long i=1; i<(1<<sz); i++){
        int ct=0:
        long long mul=1;
for(int j = 0; j < sz; j++ ) {</pre>
            if(i&(1<<j)){
                ct++;
                mul *= v1[j];
        long long sign=-1;
        if(ct&1)sign=1;
        res += sign*(r/mul);
    return r-res;
4.12 Xor basis
int basis[d]; // basis[i] keeps the mask of the

→ vector whose f value is i

int sz; // Current size of the basis
void insertVector(int mask) {
for (int i = 0; i < d; i++) {
  if ((mask \& 1 << i) == 0) continue; // continue
  \rightarrow if i != f(mask)
  if (!basis[i]) { // If there is no basis vector
  with the i'th bit set, then insert this

    vector into the basis

   basis[i] = mask;
   ++SZ;
   return;
  mask ^= basis[i]; // Otherwise subtract the basis

→ vector from this vector

4.13 stirling_number_of_2nd_kind
long long p = 1e9 + 7;
long long fact[1000005];
int n, m, k;
long long s( long long N,long long R )
    if( N == 0 && R == 0 ) return 1;
    if( N == 0 \mid \mid R == 0 ) return 0;
    long long ans = 0:
    for( int i = 1; i <= R; i++ ){
        long long par;
        if((R-i)\%2 == 0) par = 1;
        else par = -1;
        par = (par+p)%p;
```

```
long long temp = (ncr(R,i) * bm(i,N))%p;
temp = (temp%p * par%p)%p;
ans = (ans%p + temp%p)%p;
}
return (ans*bm( fact[R],p-2 ))%p;
}
```

5 Misc

5.1 Compilation

```
//pragma
#pragma GCC optimize("03")
#pragma GCC optimize("unroll-loops")
compile: g++ -std=c++17 -I . -Dakifpathan -o "%e"
build: g++ -std=c++17 -DHFTF -Wshadow -o "%e" "%f"
   -g -fsanitize=address -fsanitize=undefined
  -D GLIBCXX DEBUG
run: "./%e'
//for sublime
"cmd" : ["g++ -std=c++14 $file name -o
   $file base name && timeout 6s
./$file base name<in>out"],
"selector" : "source.c, source.cpp, source.Cc",
"shell": true,
"working dir" : "$file path"
//windows
cmd": ["g++.exe", "-std=c++14", "${file}", "-o",
   "${file base name}.exe", "&&"
"${file base name}.exe<in>out"],
"shell":true,
"working dir":"$file path",
"selector": "source.cpp, source.c, source.c++,

→ source.cc¹
```

5.2 Ternary Search

```
while(hi>=lo)
int mid1=lo+(hi-lo)/3; int mid2=hi-(hi-lo)/3;
if(f(mid1)>f(mid2)) { }//change
else //change
}//ittehad
double x1, why1, z1, x2, y2, z2, x, y , z;
double f( double t )
    double xt = x1 + (x2-x1)t;
    double yt = why1 + (y2-why1)t;
    double zt = z1 + (z2-z1)t;
    return ((xt-x)(xt-x) + (yt-y)(yt-y) +
     \rightarrow (zt-z)(zt-z));
double Tsearch()
    double low = 0, high = 1, mid;
    int step = 64:
    while ( step-- ) {
```

```
double t1 = (2low + high)/3;
        double t2 = (low + 2high)/3;
        double d1 = f(t1);
        double d2 = f(t2);
        if(d1 < d2) high = t2;
        else low = t1;
    return low;
6 String
6.1 Aho-Corasick
struct vartex
    int next[30];
    int endmark;
    int link;
    vector<int>dlink;
    vartex()
        memset(next, -1, sizeof(next));
        endmark=-1;
        link=0:
};
void addstring(string& s,vector<vartex>&trie)
    int v=0:
    for(auto \times : s)
        if(trie[v].next[x-'a']==-1)
            trie[v].next[x-'a']=trie.size();
            trie.emplace back();
        v=trie[v].next[x-'a'];
    trie[v].endmark=0;
void fail(vector<vartex>&trie)
    int v=0;
    trie[v].link=0;
    queue<int>q;
    q.push(0);
    while(!q.empty())
        v=q.front();
        q.pop();
        for(int i=0;i<26;i++)
            if(trie[v].next[i]!=-1)
                if(v==0)
                    trie[trie[v].next[i]].link=0;
                else
                    int x=trie[v].link;
                    while (x!=0 &&
                         trie[x].next[i]==-1)
```

```
x=trie[x].link;
                     if(trie[x].next[i]==-1)
                         trie[trie[v].next[i]].link=
                     else
                         trie[trie[v].next[i]].link=_

    trie[x].next[i];

                 q.push(trie[v].next[i]);
    }
void dictionary link(vector<vartex>&trie)
    queue<int>q;
    q.push(0);
    while (!q.empty())
        int u=q.front();
        q.pop();
        for(int i=0;i<26;i++)
            if(trie[u].next[i]!=-1)
                 q.push(trie[u].next[i]);
        int k=u;
        while (k!=0)
            if(trie[k].endmark!=-1 && k!=u)
                trie[u].dlink.push back(k);
            k=trie[k].link;
        debug(u,trie[u].dlink);
int search(string& s,vector<vartex>&trie)
    int v=0:
    for(auto \times : s)
        v=trie[v].next[x-'a'];
    return trie[v].endmark;
6.2 Hashing_without_inv
#include<bits/stdc++.h>
using namespace std;
long long h[400005];
long long MOD[400005];
int L;
```

```
void pre hash( string s )
    long long p = 31;
    long long m = 1e9 + 9;
    long long power = 1;
    long long hash = 0;
    int^z = 0;
    for( int i = s.size()-1; i >= 0; i-- ){
        hash = (hash*p+(s[i]-'A'+1))%m;
        h[i] = hash;
        MOD[z] = power;
        power = (power*p)%m;
   }
long long f( int l,int r )
    long long val = h[r], m = 1e9 + 9;
    if( l != L-1 )
        long long val2 = (h[l+1]%m *MOD[l-r+1]%m
        → )%m;
        val -= val2;
        val += m;
        val %= m;
    if( val < 0 ) val = (val+m)%m;
    return val;
int main()
    string s;
    cin >> s;
   L = s.size();
    pre hash(s);
    int q;
    cin >> q;
    while( q-- ){
        int l, r;
        cin >> l >> r;
        cout \ll f(l,r) \ll endl;
    return 0;
```

6.3 KMP

```
#include<bits/stdc++.h>
using namespace std;
#define pii pair<int,int>
vector<int> prefix_function (string Z) {
   int n = (int) Z.length();
   vector<int> F (n);
   F[0]=0;
   for (int i=1; i<n; ++i) {
      int j = F[i-1];
      while (j > 0 && Z[i] != Z[j])
```

v[i] << " ";
</pre>

cout << endl;</pre>

return 0;

6.4 Manacher

```
/// When i is even, pal[i] = largest palindromic
   → substring centered from str[i / 2]
/// When i is odd, pal[i] = largest palindromic
               substring centered between str[i / 2] and str[i
vector <int> manacher(char *str){
                 int i, j, k, l = strlen(str), n = l << 1;</pre>
                 vector <int> pal(n);
                 for (i = 0, j = 0, k = 0; i < n; j = max(0, j - 1)
                    \rightarrow k), i += k
                                  while (j \le i \&\& (i + j + 1) < n \&\& str[(i +
                                                 -i) >> 1] == str[(i + j + 1) >> 1])
                                     _ j++;
                                  for (k = 1, pal[i] = j; k <= i \&\& k <=
                                                 pal[i] \&\& (pal[i] - k) != pal[i - k];
                                                 pal[i + k] = min(pal[i - k], pal[i] -

→ k);

                 pal.pop back();
                 return pal;
int main(){
                 char str[100];
                 while (scanf("%s", str)){
                                  auto v = manacher(str);
                                  for (auto it: v) printf("%d ", it);
                                 puts("");
                  return 0;
```

vector<int> v = prefix function(s);

for(int i = 0; i < v.size(); i++) cout <<

6.5 Palindromic Tree

```
#define CLR(a) memset(a,0,sizeof(a))
/***
    * str is 1 based
    Each node in the palindromic tree denotes a
    STRING
    Node 1 denotes an imaginary string of size -1
```

```
Node 2 denotes a string of size 0
    They are the two roots
    There can be maximum of (string length + 2)
   nodes in total
    It's a directed tree. If we reverse the
    direction of the suffix links, we get a dag. In
   this DAG, if node v is reachable from node u
   iff, u is a substring of v.
    * if ( tree[A].next[x] == B )
    then, B = xAx
* if (tree[A].suffixLink == B)
      Then B is the longest possible palindrome
→ which is a proper suffix of A
      (node 1 is an exception)
    * occ[i] contains the number of occurrences of
   the corresponding palindrome
    * st[i] denotes starting index of the first
   occurrence of the corresponding palindrome
    * st[] or occ[] or both can be ignored if not

→ needed
    * If memory limit is compact, a map has to be
   used instead of
      ed[MAXN][MAXC]. Swapping row and column of
   the matrix will save more memory.
      Example:
      map <int,int> ed[MAXC];
      ed[c][u] = v means, there is an edge from
      node v that is labeled character c.
namespace pt {
    const int MAXN = 100010; /// maximum possible

→ string size

    const int MAXC = 26; /// Size of the character
    int n; /// length of str
    char str[MAXN];
    int len[MAXN], link[MAXN], ed[MAXN][MAXC],
    → occ[MAXN]. st[MAXN]:
    int nc, suff, pos;
    /// nc -> node count
    /// suff -> Index of the node denoting the
        longest palindromic proper suffix of the

→ current prefix

    void init() {
        str[0] = -1;
        nc = 2; suff = 2;
        len[1] = -1, link[1] = 1;
        len[2] = 0, link[2] = 1;
        CLR(ed[1]), CLR(ed[2]);
        occ[1] = occ[2] = 0;
    inline int scale(char c) { return c-'a': }
    inline int nextLink(int cur) {
        while (str[pos - 1 - len[cur]] != str[pos])
        cur = link[cur];
        return cur;
```

```
inline bool addLetter(int p) {
        pos = p;
        int let = scale(str[pos]);
        int cur = nextLink(suff);
        if (ed[cur][let]) {
             suff = ed[cur][let];
             occ[suff]++;
             return false;
        suff = ++nc;
        CLR(ed[nc]);
        len[nc] = len[cur] + 2;
        ed[cur][let] = nc;
        occ[nc] = 1;
        if (len[nc] == 1) {
             st[nc] = pos;
             link[nc] = 2;
             return true;
        link[nc] = ed[nextLink(link[cur])][let];
        st[nc] = pos-len[nc] + 1;
        return true:
    void build(int n) {
        init();
        for(int i=1;i<=n;i++) addLetter(i);</pre>
        for(int i=nc;i>=3;i--) occ[link[i]] +=
         → occ[i]:
        occ[2] = occ[1] = 0;
    void printTree() {
        puts(str);
cout << "Node\tStart\tLength\tOcc\n";</pre>
        for(int i=3;i<=nc;i++) {
   cout << i << "\t" << st[i] << "\t" <<</pre>
              → len[i] << "\t" << occ[i] << "\n";</pre>
    }
int main() {
    scanf("%s",pt::str+1);
    pt::build(strlen(pt::str+1));
    return 0:
6.6 String Hashing
ll bigmod(ll x, ll p, ll md){
    ll res = 1;
    while(p){
        if(p \& 1) res = res * x % md;
        x = x * x % md;
        p >>= 1:
    return res;
ll modinv(ll x, ll md){
    return bigmod(x, md-2, md);
```

```
namespace Hash{
    ll pw[M][2];
    ll invpw[M][2];
    const int pr[] = {37, 53};
    const int md[] = {1000000007, 1000000009};
    void precalc(){
         pw[0][0] = pw[0][1] = 1;
         for(int i=1; i<M; i++){
             pw[i][0] = pw[i-1][0] * pr[0] % md[0];
pw[i][1] = pw[i-1][1] * pr[1] % md[1];
         invpw[M-1][0] = modinv(pw[M-1][0], md[0]);
invpw[M-1][1] = modinv(pw[M-1][1], md[1]);
         for(int i=M-2; i>=0; i--){
             invpw[i][0] = invpw[i+1][0] * pr[0] %
             invpw[i][1] = invpw[i+1][1] * pr[1] %
              \rightarrow md[1];
    pii get hash(const string &s){
         pii ret = \{0, 0\};
         for(int i=0; i<s.size(); i++){</pre>
              ret.first += (s[i]-'a'+1)*pw[i][0] %

→ md[0];

              ret.second += (s[i]-'a'+1)*pw[i][1] %
               \rightarrow md[1]:
             if(ret.first >= md[0]) ret.first -=
              \rightarrow md[0];
             if(ret.second >= md[1]) ret.second -=
              \rightarrow md[1];
         return ret;
    void prefix(const string &s, pii *H){
         H[0] = \{0, 0\};
         for(int i=1; i<=s.size(); i++){</pre>
             H[i].first = H[i-1].first +
                 (s[i-1]-'a'+1)*pw[i-1][0] % md[0];
             H[i].second = H[i-1].second +
              \rightarrow (s[i-1]-'a'+1)*pw[i-1][1] % md[1];
             if(H[i].first >= md[0]) H[i].first -=

    md[0]:

             if(H[i].second >= md[1]) H[i].second -=
              \rightarrow md[1]:
    void reverse prefix(const string &s, pii *H){
         int n = \overline{s.size()};
         for(int i=1; i<=s.size(); i++){</pre>
             H[i].first = H[i-1].first'+
              \rightarrow (s[i-1]-'a'+1)*pw[n-i][0] % md[0];
             H[i].second = H[i-1].second +
              \rightarrow (s[i-1]-'a'+1)*pw[n-i][1] % md[1];
             if(H[i].first >= md[0]) H[i].first -=
              \rightarrow md[0];
             if(H[i].second >= md[1]) H[i].second -=
              \rightarrow md[1];
    pii range hash(int L, int R, pii H[]){
         pii ret:
```

```
OLIN | ITTEHAD | INZAMAM
         ret.first = (H[R].first - H[L-1].first +
         \rightarrow md[0]) % md[0];
         ret.second = (H[R].second - H[L-1].second +
         \rightarrow md[1]) % md[1];
        ret.first = ret.first * invpw[L-1][0] %
        ret.second = ret.second * invpw[L-1][1] %
         \rightarrow md[1];
        return ret;
    pii reverse hash(int L, int R, pii H[], int n){
         ret.first = (H[R].first - H[L-1].first +
         \rightarrow md[0]) % md[0];
         ret.second = (H[R].second - H[L-1].second +
         \rightarrow md[1]) % md[1];
         ret.first = ret.first * invpw[n-R][0] %
        ret.second = ret.second * invpw[n-R][1] %
         \rightarrow md[1];
        return ret;
6.7 Suffix Array (n\log^2 n)
///if finding suffix array of vector, use const
→ vector<int> &s
vector<int> sa doubling(const string& s) {
    int n = int(s.size());
    vector<int> sa(n),rnk(n),tmp(n);
    for(int i=0;i<n;i++) rnk[i]=s[i];
    iota(sa.begin(), sa.end(), 0);
    for (int k = 1; k < n; k *= 2)
        auto cmp = [&](int x, int y) {
             if (rnk[x] != rnk[y]) return rnk[x] <</pre>
              \rightarrow rnk[y];
             int rx = x + k < n ? rnk[x + k] : -1;
             int ry = y + k < n ? rnk[y + k] : -1;
             return rx'< ry;</pre>
```

```
sort(sa.begin(), sa.end(), cmp);
        tmp[sa[0]] = 0;
        for (int i = 1; i < n; i++) {
            tmp[sa[i]] = tmp[sa[i - 1]] + (cmp(sa[i
             \rightarrow - 1], sa[i]) ? 1 : 0);
        swap(tmp, rnk);
    return sa;
template <class T>
std::vector<int> lcp array(const std::vector<T>& s,
                            const std::vector<int>&
                             → sa) {
    int n = int(s.size());
    assert(n >= 1);
    std::vector<int> rnk(n):
    for (int i = 0; i < n; i++) {
        rnk[sa[i]] = i;
    std::vector<int> lcp(n - 1):
    int h = 0;
```

```
for (int i = 0; i < n; i++) {
    if (h > 0) h--;
    if (rnk[i] == 0) continue;
    int j = sa[rnk[i] - 1];
    for (; j + h < n && i + h < n; h++) {
        if (s[j + h] != s[i + h]) break;
    }
    lcp[rnk[i] - 1] = h;
}
return lcp;
}</pre>
```

6.8 Suffix Array (nlogn)

```
struct suffix array
    vector<int> sa, lcp;
    suffix array(string &s, int \lim = 256)
        int n = s.size() + 1, k = 0, a, b;
        vector<int> x(s.begin(), s.end() + 1),
        \rightarrow y(n), ws(max(n, lim)), rank(n);
        sa = lcp = y, iota(sa.begin(), sa.end(), 0);
        for (int j = 0, p = 0; p < n; j = max(1, j)
        \rightarrow * 2), lim = p)
            p = j, iota(y.begin(), y.end(), n - j);
            for (int i = 0; i < n; i++)
                if (sa[i] >= j)
                    y[p++] = sa[i] - j;
            fill(ws.begin(), ws.end(), 0);
            for (int i = 0; i < n; i++)
                WS[X[i]]++;
            for (int i = 1; i < lim; i++)
                ws[i] += ws[i - 1];
            for (int i = n; i--;)
                sa[--ws[x[y[i]]]] = y[i];
            swap(x, y), p = 1, x[sa[0]] = 0;
            for (int i = 1; i < n; i++)
                a = sa[i - 1];
                b = sa[i];
                x[b] = (y[a] == y[b] \&\& y[a + j] ==
                 \vee v[b + i]) ? p - 1 : p++:
        for (int i = 1; i < n; i++)
            rank[sa[i]] = i;
        for (int i = 0, j; i < n - 1;
        \rightarrow lcp[rank[i++]] = k)
```

```
for (k \&\&k--, j = sa[rank[i] - 1]; s[i]
                + k] == s[j + k]; k++)
6.9 Suffix Automaton
// collected from cp algorithm
struct state {
    int len, link, cnt, firstpos; // cnt -> endpos

→ set size, link -> suffix link

    map <char, int> next;
const int MAXLEN = 100002;
state st[MAXLEN * 2];
struct SuffixAutomata { // O-based
    int sz, last;
    SuffixAutomata() { // init
    st[0].cnt = st[0].len = 0;
    st[0].link = -1;
        sz = 1, last = 0;
    void add(char c) { // add new char in automata
        int cur = sz++;
        st[cur].len = st[last].len + 1;
        st[cur].firstpos = st[cur].len - 1;
        st[cur].cnt = 1;
        int p = last;
        while (p != -1 \&\& !st[p].next.count(c)) {
             st[p].next[c] = cur;
            p = st[p].link;
        if (p == -1) {
    st[cur].link = 0;
        else {
            int q = st[p].next[c];
            if (st[p].len + 1 == st[q].len) {
                 st[cur].link = q;
            else { // clone state
                 int clone = sz++;
                 st[clone].len = st[p].len + 1;
                 st[clone].next = st[q].next;
                 st[clone].link = st[q].link;
                 st[clone] firstpos = st[q] firstpos;
                 st[clone].cnt = 0;
                 while (p != -1 && st[p].next[c] ==
                     st[p].next[c] = clone;
                     p = st[p].link;
                 st[q].link = st[cur].link = clone;
        }
        last = cur;
```

```
void occurrence() { // calculate number of

    occurrences of all possible substring

        vector <int> rank(sz);
        iota(all(rank), 0);
        sort(all(rank), [&](int i, int j) {
            return st[i].len > st[j].len;
        });
        for (int ii : rank) if (st[ii].link != -1)
            st[st[ii].link].cnt += st[ii].cnt;
    int count(string s) { // number of occurrences
        of string s. #prerequisite -> call
     → occurrence()
        int node = 0:
        for (char ch : s) {
            if (!st[node].next.count(ch)) return 0;
            node = st[node].next[ch];
        return st[node].cnt;
    int firstOcc(string s) { // first
        position(occurence) of string s
        int node = 0;
        for (char ch : s) {
            if (!st[node].next.count(ch)) return -1;
            node = st[node].next[ch];
        return st[node].firstpos + 2 -

    (int)s.size();

    void build(string S) { // build suffix
     → automata
        for (char ch : S) add(ch);
    bool find(string s) { // find string s in

→ automata
        int node = 0;
        for (char ch : s) {
            if (!st[node].next.count(ch)) return

→ false:

            node = st[node].next[ch];
        return true:
};
6.10 Trie
//define\ M,\ K = alphabet\ size
int trie[M][K], word[M*K+3], cnt[M*K+3], sz;
void Insert(string s){
    int node = 0;
    for(int i=0; i<s.size(); i++){</pre>
        int c = s[i] - 'a'
        if(!trie[node][c]){
            trie[node][c] = ++sz;
        node = trie[node][c];
```

```
cnt[node]++;
    word[node]++;
bool Search(string s){
    int node = 0, ret = 0;
    for(int i=0; i<s.size(); i++){</pre>
        int c = s[i] - a'
        if(!trie[node][c]) return false;
        node = trie[node][c];
    return (word[node] > 0);
void Delete(string s){
    int node = 0;
    vector<int>v(1, 0):
    for(int i=0; i<s.size(); i++){</pre>
        int c = s[i] - 'a'
        node = trie[node][c];
        cnt[node]--
        v.push back(node);
    word[node]--;
    for(int i=1; i<v.size(); i++){</pre>
        int c = s[i-1]-'a';
        if(!cnt[v[i]]){
            trie[v[i-1]][c] = 0;
```

6.11 Z algo

```
* z[i] denotes the maximum length of substring
* starting from position(i) which is also a prefix
  of the string
  call with Z zf(x) where x is the desired string
struct Z
int n;
string s;
vector<int>z;
Z(const string &a)
 n=a.size(); s=a; z.assign(n,0);
void z function()
 for (int i = 1, l = 0, r = 0; i < n; ++i)
  if (i \le r) z[i] = min (r - i + 1, z[i - l]);
  while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
   if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
```