```
1 DP
1.1 SOS DP
/// 0(N * 2^N)
///memory optimized version
for (int i = 0; i < (1 << N); ++i) F[i] = A[i];
for (int i = 0; i < N; ++i) for (int mask = 0;
\rightarrow mask < (1 << N); ++mask)
 if (mask \& (1 << i)) F[mask] += F[mask ^ (1
 < < i)]:
/// How many pairs in ara[] such that (ara[i] &
\rightarrow ara[i]) = 0
/// N --> Max number of bits of any array
→ element
const int N = 20;
int inv = (1 << N) - 1;
int F[(1 << N) + 10];
int ara[MAX];
/// ara is 0 based
long long howManyZeroPairs(int n, int ara[]) {
CLR(F);
for (int i = 0; i < n; i++) F[ara[i]]++;
for (int i = 0; i < N; ++i) for (int mask = 0;
\rightarrow mask < (1 << N); ++mask) {
   if (mask \& (1 << i)) F[mask] += F[mask ^ (1
\rightarrow << i)];
 }
long long ans = 0;
for (int i = 0; i < n; i++) ans += F[ara[i] ^

   inv];

 return ans;
/// F[mask] = sum of A[i] given that
for (int i = 0; i < (1 << N); ++i) F[i] = A[i];
for (int i = 0; i < N; ++i)
for (int mask = (1 \ll N) - 1; mask >= 0;
→ --mask) {
 if (!(mask \& (1 << i))) F[mask] += F[mask]
\rightarrow (1 << i)];
/// Number of subsequences of ara[0:n-1] such

    that

/// sub[0] \& sub[2] \& ... \& sub[k-1] = 0
const int N = 20;
int inv = (1 << \hat{N}) - 1;
int F[(1 << N) + 10];
int ara[MAX];
int p2[MAX]; /// p2[i] = 2^i
///0 based array
int howManyZeroSubSequences(int n, int ara[]) {
CLR(F);
for (int i = 0; i < n; i++) F[ara[i]]++;</pre>
for (int i = 0; i < N; ++i)
 for (int mask = (1 \ll N) - 1; mask >= 0;
→ --mask) {
   if (!(mask & (1 << i)))
    F[mask] += F[mask | (1 << i)];
```

```
int ans = 0;
 for (int mask = 0; mask < (1 << N); mask++) {
  if ( builtin popcount(mask) \& 1) ans =

    sub(ans, p2[F[mask]]);

  else ans = add(ans, p2[F[mask]]);
 return ans;
/// Number of subsequences of ara[0:n-1] such
/// sub[0] | sub[2] | ... | sub[k-1] = Q
int F[(1 << 20) + 10], ara[MAX];
int p2[MAX]; /// p2[i] = 2^i
/// ara is 0 based
int howManySubsequences(int n, int ara[], int

→ m. int 0) {
 CLR(F);
 for (int i = 0; i < n; i++) F[ara[i]]++;
 if (Q == 0) return sub(p2[F[0]], 1);
 for (int i = 0; i < m; ++i)
  for (int mask = 0; mask < (1 << m); ++mask) {</pre>
   if (mask & (1 << i))
    F[mask] += F[mask ^ (1 << i)];
 int ans = 0;
 for (int mask = 0; mask < (1 << m); mask++) {
  if (mask & Q != mask) continue;
  if ( builtin popcount(mask ^{\circ} Q) \& 1) ans =

    sub(ans, p2[F[mask]]);

  else ans = add(ans, p2[F[mask]]);
 return ans;
```

## 2 Data Structures

#### 2.1 2D Fenwick Tree

```
2.2 Fenwick Tree
int bit[1000], arra[1000];
int n;
void update( int idx, int val ) {
for ( int i = idx; i <= n; i += i & (-i) )
→ bit[i] += val:
 return;
int query( int idx ) {
int sum = 0;
 for ( int i = idx; i > 0; i -= i & (-i) ) sum
→ += bit[i];
 return sum;
2.3 LIS
int lis(vector<int> a) {
int n = a.size();
 vector<int>d(n + 1, INF);
d[0] = -INF;
 for (int i = 0; i < n; i++) {
 int j = upper bound(d.begin(), d.end(), a[i])
- d.begin();
 if (d[j - 1] < a[i] and a[i] < d[j]) d[j] =
\rightarrow a[i];
 int ret = 0;
 for (int i = 1; i <= n; i++)
 if (d[i] < INF) ret = i;
 return ret;
```

### 2.4 MO's Algo

```
const int mx = 100005;
const int sz = 100005;
struct query {
int l, r, id;
bool operator<(const query &a) const {</pre>
 int x = l / sz; int y = a.l / sz;
 if (x != y) return x < y;
 if (x \% 2) return r < a.r;
 return r > a.r;
void add(int indx) { }
void baad(int indx) { }
void solve() {
int l = 0;
int r = -1;
sort(ques + 1, ques + q + 1);
for (int i = 1; i \le q; i++) {
 while (l > ques[i].l) add(--l);
 while (r < ques[i].r) add(++r);</pre>
 while (l < ques[i].l) baad(l++);
 while (r > ques[i].r) baad(r--);
```

```
ans[ques[i].id] = sum[now];
 for (int i = 1; i <= q; i++) cout << ans[i] <<
2.5 PBDS
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace gnu pbds;
template<class T>
using ordered set = tree<T, null type, less<T>,
      rb tree tag,
tree order statistics node update> ;
0 based indexing
1) insert(value)
2) erase(value)
3) order of key(value) // Number of items
→ strictly smaller than value
4) *find by order(k) : K-th element in a

    set(counting from zero)

2.6 Persistent Seg Tree (Pointer)
```

```
struct node {
ll sum;
node *left, *right;
node(ll sum = 0) {
 sum = \overline{s}um;
 left = right = NULL;
void build(int l, int r) {
 if (l == r) return;
 left = new node();
  right = new node();
  int mid = (l + r) / 2;
  left->build(l, mid);
  right->build(mid + 1, r);
node *update(int l, int r, int i, int x) {
 if (l > i \mid | r < i) return this;
  if (l == r) return new node(x);
  int mid = (l + r) / 2;
  node *ret = new node();
  ret->left = left->update(l, mid, i, x);
  ret->right = right->update(mid + 1, r, i, x);
  ret->sum = ret->left->sum + ret->right->sum;
  return ret;
 ll query(int tL, int tR, int rL, int rR) {
 if (tL > rR || tR < rL) return 0;
if (tL >= rL && tR <= rR) return sum;</pre>
  int mid = (tL + tR) / 2;
  ll a = left->query(tL, mid, rL, rR);
  ll b = right->query(mid + 1, tR, rL, rR);
  return a + b;
```

```
int size() { return sizeof(*this) /

¬ sizeof(node*); }

const int mx = 2e5 + 5;
node *root[mx];
void solve() {
root[0] = new node();
 root[0]->build(1, n);
 for (int i = 1; i <= n; i++) {
 int x;
  cin >> x;
  root[0] = root[0] -> update(1, n, i, x);
 int sz = 0;
 while (q--) {
  int op; cin >> op;
  if (op == 1) {
  int version, i, x;
   cin >> version >> i >> x;
   version--;
   root[version] = root[version]->update(1, n,
  i, x);
  else if (op == 2) {
   int version, l, r;
   cin >> version >> l >> r;
   version--:
   cout << root[version]->query(1, n, l, r) <<</pre>
  else {
  int version;
  cin >> version;
   version--;
   root[++sz] = root[version];
2.7 RMQ
template<typename T>
```

```
struct sparse_table {
  vector<T>ara;
vector<int>logs;
vector<vector<T>>table;
sparse table(int n) {
 ara.resize(n + 1);
  logs.resize(n + 1);
f func(T a, T b) { }
void build(int n) {
 logs[1] = 0;
  for (int i = 2; i \le n; i++) logs[i] = logs[i]
 table.resize(n + 1, vector<T>(logs[n] + 1));
 for (int i = 1; i \le n; i++) table[i][0] =
    ara[i]:
```

```
for (int j = 1; j <= logs[n]; j++) {</pre>
  int sz = 1 << j;
   for (int i = 1; i + sz - 1 \le n; i++)
    table[i][j] = func(table[i][j - 1], table[i
   + sz / 2][i - 1]);
 T query(int l, int r) {
  int d = logs[r - l + 1];
  return func(table[l][d], table[r - (1 << d) +</pre>
   1][d]);
}
};
```

```
2.8 Seg Tree(Inz)
class SegmentTree {
#define Lc(idx) idx * 2
#define Rc(idx) idx * 2 + 1
public:
 struct node {
 int value;
  int lazy;
  node() {
  this->value = ??;
   this->lazv = ??:
 vector <node> seqT;
 vector <int> A;
 SegmentTree(int sz) {
// need to clear!
  segT.resize(4 * sz + 10);
  A.resize(sz + 1);
node Merge(node L, node R) {
 node F;
 F = ??
      return F;
void Relax(int L, int R, int idx) {
//Do somethina
 segT[idx].lazy = ??; //after Relaxing
 void MakeSegmentTree(int L, int R, int idx) {
 if (L == R) {
   segT[idx].value = ??;
   return;
  int M = (L + R) / 2;
 MakeSegmentTree(L, M, Lc(idx));
MakeSegmentTree(M + 1, R, Rc(idx));
  segT[idx] = Merge(segT[Lc(idx)],

    segT[Rc(idx)]);

 node RangeQuery(int L, int R, int idx, int l,
```

```
Relax(L, R, idx);
 node F;
 if (L > r || R < l) return F;
  if (L >= l \& R <= r) return segT[idx];
 int M = (L + R) / 2;
 F = Merge(RangeQuery(L, M, Lc(idx), l, r),
\rightarrow RangeQuery(M + 1, R, Rc(idx), l, r));
  segT[idx] = Merge(segT[Lc(idx)],
→ segT[Rc(idx)]): //is it useful?
 return F;
void RangeUpdate(int L, int R, int idx, int l,

    int r, int lz) {

 Relax(L, R, idx);
 if (L > r || R < l) return;
  if (L >= l \& R <= r) {
// Do something
   seqT[idx].lazy = ??;
   Relax(L, R, idx);
   return;
 int M = (L + R) / 2;
RangeUpdate(L, M, Lc(idx), l, r, lz);
 RangeUpdate(M' + 1, R, Rc(idx'), l, r, lz);
  segT[idx] = Merge( segT[Lc(idx)],
   segT[Rc(idx)]);
```

## 2.9 Seg Tree(Lazy Prop)

```
int n, q, arra[100005];
struct idk {
int sum, prop;
} tree[300005];
void init( int node, int b, int e ) {
if ( b == e ) {
 tree[node].sum = arra[b];
  return;
int left = node * 2;
int right = node * 2 + 1;
int mid = (b + e) / 2;
init(left, b, mid);
init(right, mid + 1, e);
tree[node].sum = tree[left].sum +

    tree[right].sum;

return;
void update( int node, int b, int e, int i, int
→ j, int val ) {
if ( b > j || e < i ) return;
if (b \ge i \& e \le j)
 trèe[node].sum += (e - b + 1) * val;
 tree[node].prop += val;
 return;
int left = node * 2;
int right = node * 2 + 1;
```

```
int mid = (b + e) / 2;
update( left, b, mid, i, j, val );
update( right, mid + 1, e, i, j, val );
tree[node].sum = tree[left].sum +
   tree[right].sum + (e - b + 1) *
  tree[node].prop;
return;
int query( int node, int b, int e, int i, int

    j, int carry ) {

if ( b > j || e < i ) return 0;
if ( b \ge i \& \& e \le j ) return tree[node].sum
\rightarrow + (e - b + 1) * carry;
int left = node * 2:
int right = node * 2 + 1;
int mid = (b + e) / 2;
int p1 = query(left, b, mid, i, j, carry +

    tree[node].prop );

int p2 = query(right, mid + 1, e, i, j, carry
+ tree[node].prop );
return p1 + p2;
```

## 2.10 Seg Tree(point update, range query)

## 2.11 Seg Tree(range update, range query)

```
ll tree[4 * N], lazy[4 * N];
inline ll merge(ll a, ll b) { return a + b; }
void push(int rt, int l, int r) {
   if (l ^ r) {
      lazy[rt << 1] += lazy[rt];
      lazy[rt << 1 | 1] += lazy[rt];
   }
   tree[rt] += (r - l + 1) * lazy[rt];
   lazy[rt] = 0;
}
void update(int rt, int l, int r, int b, int e,
      ll v) {
   if (lazy[rt]) push(rt, l, r);</pre>
```

## 2.12 Sqrt Decomposition

```
int block size = ??;
int Block[block size + 5];
int getBlock(int idx) {
return (idx + block size - 1) / block size;
→ //for 1-base index
return idx / block size; //for 0-base index
int getQueryAns(int L, int R) //0-base index
int ANS = 0;
int CL = L / block size;
int CR = R / block_size;
if (CL == CR) {
 for (int i = L; i <= R; ++i)
  ANS += ArrName[i];
else {
 for (int i = L, LM = (CL + 1) * block size -
\rightarrow 1; i <= LM; ++i)
  ANS += ArrName[i];
 for (int i = CL + 1; i \le CR - 1; ++i)
  ANS += Block[i];
 for (int i = CR * block size; i <= R; ++i)
  ANS += ArrName[i];
return ANS;
//Update : Block[ idx / block size ] += ??
```

## 3 Geometry

## 3.1 2D point line- segment

```
const double PI = acos(-1.0);
const double EPS = 1e-12;
/***
```

```
u \cdot v = |u| * |v| * cos(theta)
= u.x*v.\dot{x} + \dot{u}.\dot{y}*v.y
= How much parallel they are
= Dot product does not change if one vector

→ move perpendicular to the other

u \times v = |u| * |v| * sin(theta)
= u.x*v.y - v.x*u.y
= How much perpendicular they are
= Cross product does not change if one vector

    move parallel to the other

dot(a-b,a-b) returns squared distance between
\underset{***}{\smile} pt a and pt b
struct pt {
double x, y,
pt() {}
 pt(double x, double y) : x(x) , y(y) {}
 pt operator + (const pt &p) const {
 return pt(x + p.x, y + p.y);
 pt operator - (const pt \delta p) const {
 return pt( x - p.x , y - p.y );
 pt operator * (double c) const {
 return pt( x * c , y * c );
pt operator / (double c) const {
 return pt( x / c , y / c );
bool operator == (const pt &p) const {
 return ( fabs( x - p.x ) < EPS && fabs( y -
   p.v ) < EPS );
bool operator != (const pt \&p) const {
 return !(pt(x, y) == p);
ostream& operator << (ostream& os, pt p) {
return os << "(" << p.x << "," << p.y << ")";
// u.v = |u|*|v|*cos(theta)
inline double dot(pt u, pt v) {
return u.x * v.x + u.y * v.y;
// a x b = |a|*|b|*sin(theta)
inline double cross(pt u, pt v) {
return u.x * v.y - u.y * v.x;
// returns |u|
inline double norm(pt u) { return sqrt(dot(u,
// returns angle between two vectors
inline double angle(pt u, pt v) {
double cosTheta = dot(u, v) / norm(u) /
return acos(max(-1.0, min(1.0, cosTheta))); //

    keeping cosTheta in [-1, 1]

// returns ang radian rotated version of vector
```

```
// ccw rotation if angle is positive else cw

→ rotation

inline pt rotate(pt u, double ang) {
return pt( u.x * cos(ang) - u.y * sin(ang) ,
\rightarrow u.x * sin(ang) + u.y * cos(ang) );
// returns a vector perpendicular to v
inline pt perp(pt u) { return pt( -u.y , u.x );
// returns 2*area of triangle
inline double triArea2(pt a, pt b, pt c) {
return cross(b - a, c - a);
// compare function for angular sort around
   point P0
inline bool comp(pt P0, pt a, pt b) {
 double d = triArea2(P0, a, b);
 if (d < 0) return false;</pre>
if (d == 0 \&\& dot(P0 - a, P0 - a) > dot(P0 - a)
→ b, P0 - b) ) return false;
return true;
if line equation is, ax + by = c
v --> direction vector of the line (b,-a)
c --> v cross p
p --> Any point(vector) on the line
side(p) = (v cross p) - c
= triArea2(origin,v,p)
if side(p) is,
positive --> p is above the line
zero --> p is on the line
negative --> p is below the line
struct line {
pt v;
 double c;
 line(pt v, double c) : v(v), c(c) {}
 / From equation ax + by = c
line(double a, double b, double c) : v( {b,
\rightarrow -a}), c(c) {}
// From points p and q
line(pt p, pt q) : v(q - p), c(cross(v, p)) {}
// |v| * dist
// dist --> distance of p from the line
 double side(pt p) { return cross(v, p) - c; }
 / better to using sqDist than dist
 double dist(pt p)
  return abs(side(p)) / norm(v);
 double sqDist(pt p) {
  return side(p) * side(p) / dot(v, v);
   perpendicular line through point p
// 90deg ccw rotated line
 line perpThrough(pt p) {
  return {p, p + perp(v)};
// translates a line by vector t(dx,dy)
```

```
// every point (x,y) of previous line is
 \rightarrow translated to (x + dx, y + dy)
 line translate(pt t) {
     return {v, c + cross(v, t)};
// for every point
// distance between previous position and
 line shiftLeft(double dist) {
    return {v, c + dist * norm(v)};
// projection of point p on the line
 pt projection(pt p) -
    return p - perp(v) * side(p) / dot(v, v);
// reflection of point p wrt the line
 pt reflection(pt p) {
  return p - perp(v) * side(p) * 2.0 / dot(v,
       v);
inline bool lineLineIntersection(line l1, line

→ l2, pt &out) {
 double d = cross(l1.v, l2.v);
 if (d == 0) return false;
  out = (l2.v * l1.c - l1.v * l2.c) / d;
 return true;
// interior = true for interior bisector
// interior = false for exterior bisector
inline line bisector(line l1, line l2, bool
 → interior) {
 assert(cross(l1.v, l2.v) != 0); // l1 and l2
 double sign = interior ? 1 : -1;
 return \{l\tilde{2}.v / norm(l2.v) + (l1.v * sign) / (l2.v) + (l1.v * sign) / (l2.v) + (
 \rightarrow norm(l1.v),
                    l2.c / norm(l2.v) + (l1.c * sign) /
 \rightarrow norm(l1.v)};
/*** Segment ***/
/// C --> A circle which have diameter ab
/// returns true if point p is inside C or on

    the border of C

inline bool inDisk(pt a, pt b, pt p) {
 return
       , | dot(a - p, b - p) <= 0;
/// returns true if point p is on the segment
inline bool onSegment(pt a, pt b, pt p)
 return triArea2(a, b, p) == 0 \&\& inDisk(a, b,
 → p);
inline bool segSegIntersection(pt a, pt b, pt
 if (onSegment(a, b, c)) return out = c, true;
 if (onSegment(a, b, d)) return out = d, true;
 if (onSegment(c, d, a)) return out = a, true;
```

```
if (onSegment(c, d, b)) return out = b, true;
double oa = triArea2(c, d, a);
double ob = triArea2(c, d, b);
double oc = triArea2(a, b, c);
 double od = triArea2(a, b, d);
if (oa * ob < 0 && oc * od < 0) {
 out = (a * ob - b * oa) / (ob - oa);
 return true;
return false:
// returns distance between segment ab and

→ point p

inline double segPointDist(pt a, pt b, pt p) {
if ( norm(a - b) == 0 ) {
 line l(a, b);
  pt pr = l.projection(p);
 if (onSegment(a, b, p)) return l.dist(p);
return min(norm(a - p), norm(b - p));
// returns distance between segment ab and

→ segment cd

inline double segSegDist(pt a, pt b, pt c, pt
→ d) {
double oa = triArea2(c, d, a);
double ob = triArea2(c, d, b);
double oc = triArea2(a, b, c);
double od = triArea2(a, b, d);
if (oa * ob < 0 && oc * od < 0) return 0; //
→ proper intersection
// If the segments don't intersect, the result
→ will be minimum of these four
return min({segPointDist(a, b, c),

    segPointDist(a, b, d)

             segPointDist(c, d, a),
   seqPointDist(c, d, b)
            });
```

## 3.2 Circle-line intersection

```
struct Point {
  double x, y;
  Point(double px, double py) {
    x = px;
    y = py;
  }
  Point sub(Point p2) {
    return Point(x - p2.x, y - p2.y);
  }
  Point add(Point p2) {
    return Point(x + p2.x, y + p2.y);
  }
  double distance(Point p2) {
    return sqrt((x - p2.x) * (x - p2.x) + (y - p2.y));
  }
  return normal() {
```

```
double length = sqrt(x * x + y * y)
  return Point(x / length, y / length);
 Point scale(double s) {
  return Point(x * s, y * s);
struct line // Creates a line with equation ax
\rightarrow + bv + c = 0
 double a, b, c;
 line() {}
 line( Point p1, Point p2 ) {
  a = p1.y - p2.y;
  b = p2.x - p1.x;
  c = p1.x * p2.y - p2.x * p1.y;
|inline bool eq(double a, double b) {
return fabs(a - b) < eps;
struct Circle {
double x, y, r, left, right;
Circle () {}
 Circle(double cx, double cy, double cr) {
 x = cx;
y = cy;
r = cr;
  left = x - r;
  right = x + r;
 pair<Point, Point> intersections(Circle c) {
  Point P0(x, y);
  Point P1(c.x, c.y);
  double d, a, h;
  d = P0.distance(P1);
  a = (r * r - c.r * c.r + d * d) / (2 * d);
  h = \operatorname{sqrt}(r * r - a * a);
  Point P2 = P1.sub(P0).scale(a / d).add(P0);
  double x3, y3, x4, y4;
  x3 = P2.x + h * (P1.y - P0.y) / d;
  y3 = P2.y - h * (P1.x - P0.x) / d;
  x4 = P2.x - h * (P1.y - P0.y) / d;
  y4 = P2.y + h * (P1.x - P0.x) / d;
  return pair<Point, Point>(Point(x3, y3),
→ Point(x4, y4));
inline double Distance( Point a, Point b ) {
return sqrt( ( a.x - b.x ) * ( a.x - b.x ) + (
\rightarrow a.y - b.y ) * (a.y - b.y ));
inline double Distance( Point P, line L ) {
return fabs( L.a * P.x + L.b * P.y + L.c ) /

    sqrt( L.a * L.a + L.b * L.b );
bool intersection(Circle C, line L, Point &p1,
→ Point &p2) {
if ( Distance( {C.x, C.y}, L ) > C.r + eps )

→ return false;
```

```
double a, b, c, d, x = C.x, y = C.y;
d = C.r * C.r - x * x - y * y;
if ( eq( L.a, 0) ) {
p1.y = p2.y = -L.c / L.b;
a = 1;
 b = 2 * x;
 c = p1.y * p1.y - 2 * p1.y * y - d;

d = b * b - 4 * a * c;
 d = sqrt(fabs(d));
p1.x = (b + d) / (2 * a);

p2.x = (b - d) / (2 * a);
else {
 a = L.a * L.a + L.b * L.b;
 b = 2 * (L.a * L.a * y - L.b * L.c - L.a *
c = L.c * L.c + 2 * L.a * L.c * x - L.a * L.a
 d = b * b - 4 * a * c:
 \ddot{d} = sqrt(fabs(d));
 p1.v = (b + d) / (2 * a);
 p2.y = (b - d) / (2 * a);
p1.\dot{x} = (-L.b * p1.y - L.c) / L.a;
 p2.x = (-L.b * p2.y - L.c) / L.a;
return true;
```

### 3.3 Circle

```
struct circle {
pt c;
double r;
circle() {}
circle(pt c, double r) : c(c) , r(r) {}
/* returns circumcircle of a triangle
the radius of circumcircle --> intersection
→ point of the perpendicular
bisectors of the three sides */
circle circumCircle(pt a, pt b, pt c) {
b = b - a, c = c - a; // consider coordinates

→ relative to point a

assert(cross(b, c) != 0); // no circumcircle
→ if A, B, C are co - linear
// detecting the intersection point using the
   same technique used in line line
   intersection
pt center = a + (perp(b * dot(c, c) - c *
   dot(b, b) ) / cross(b, c) / 2 );
return {center, norm(center - a)};
int sqn(double val) {
if (val > 0) return 1;
else if (val == 0) return 0;
else return -1;
/* returns number of intersection points
    hetween a line and a circle
```

```
0 --> Center
I,J --> Intersection points
P' -- > Projection of 0 onto line 1
IP = JP = h , OP = d */
int circleLineIntersection(circle c, line l,
   pair<pt, pt> &out) {
double h2 = c.r * c.r - l.sqDist(c.c); // h^2
if (h2 >= 0) { // the line touches the circle
  pt p = l.proj(c.c); // point P
 pt h = l.v * sqrt(h2) / norm(l.v); // vector

→ parallel to l, of length h

 out = \{p - h, p + h\}; // \{I, J\}
return 1 + sgn(h2); // number of intersection
→ points
/* returns number of intersection points between
 two circles
0 i --> Center of circle i
I, J --> Intersection points
P -- > Projection of O onto line IJ
IP = JP = h , 0 10 2 = d */
int circleCircleIntersection(circle c1, circle

    c2, pair<pt, pt> &out) {
pt d = c2.c - c1.c; double d2 = dot(d, d); //
→ d^2
if (d2 == 0) { // concentric circle
 assert(c1.r != c2.r); // same circle
 return 0;
double pd = (d2 + c1.r * c1.r - c2.r * c2.r) /
- 2; // = | 0 1P | * d
double h2 = c1.r * c1.r - pd * pd / d2; // =
→ h^2
if (h2 >= 0) {
 pt p = c1.c + d * pd / d2, h = perp(d) *
\rightarrow sqrt(h2 / d2);
 out = \{p - h, p + h\};
return 1 + sqn(h2);
/* inner --> if true returns inner tangents
* if the radius of c2 is 0, returns tangents

→ that go through the center

of circle c2 (value of inner is does not matter

    in this case)

* if there are 2 tangents, it fills out with

    → two pairs of points: the pairs

of tangency points on each circle (P1; P2), for

→ each of the tangents

* if there is 1 tangent, the circles are
→ tangent to each other at some point
P, out just contains P 4 times, and the tangent
→ line can be found as
line(c1.c,p).perpThrough(p)
* if there are 0 tangents, it does nothing
* if the circles are identical, it aborts. */
int tangents(circle c1, circle c2, bool inner,

    vector < pair <pt, pt> > &out) {
```

```
if (inner) c2.r = -c2.r;
 pt d = c2.c - c1.c;
 double dr = c1.r - c2.r, d2 = dot(d, d), h2 =
 \rightarrow d2 - dr * dr;
 if (d2 == 0 || h2 < 0) {//assert(h2 != 0);}
  return 0;
 for (double sign : { -1, 1}) {
  pt v = (d * dr + perp(d) * sqrt(h2) * sign) /
  out.push back(\{c1.c + v * c1.r, c2.c + v *
 \leftarrow c2.r\});
 return 1 + (h2 > 0):
3.4 Convex Hull
struct Point {
int x, y;
|Point p0;
Point nextToTop(stack<Point> &S) {
 Point p = S.top();
 S.pop();
 Point res = S.top();
 S.push(p);
 return res;
void swap(Point &p1, Point &p2) {
 Point temp = p1;
 p1 = p2;
 p2 = temp;
int distSq(Point p1, Point p2) {
return (p1.x - p2.x) * (p1.x - p2.x) +
         (p1.y - p2.y) * (p1.y - p2.y);
int orientation(Point p, Point q, Point r) {
 int val = (q.y - p.y) * (r.x - q.x) -
            (\dot{q}.\dot{x} - \dot{p}.\dot{x}) * (r.y - \dot{q}.y);
 if (val == 0) return 0;
 return (val > 0) ? 1 : 2;
int compare(const void *vp1, const void *vp2) {
 Point *p1 = (Point *)vp1;
 Point *p2 = (Point *)vp2;
 int o = orientation(p0, *p1, *p2);
 if (o == 0)
  return (distSq(p0, *p2) >= distSq(p0, *p1))?
 -1 : 1;
 return (o == 2) ? -1 : 1;
void convexHull(Point points[], int n) {
 int ymin = points[0].y, min = 0;
 for (int i = 1; i < n; i++) {
  int y = points[i].y;
  if ((y < ymin) || (ymin == y && points[i].x <</pre>
    points[min].x)
```

```
ymin = points[i].y, min = i;
swap(points[0], points[min]);
p0 = points[0];
qsort(&points[1], n - 1, sizeof(Point),

    compare);

int m = 1;
for (int i = 1; i < n; i++) {
 while (i < n - 1 && orientation(p0,
\rightarrow points[i], points[i + 1]) == 0)
  1++:
 points[m] = points[i];
 m++;
if (m < 3) return;</pre>
stack<Point> S:
S.push(points[0]);
S.push(points[1]);
S.push(points[2]);
for (int i = 3; i < m; i++) {
 while (S.size() > 1 \&\&
   orientation(nextToTop(S), S.top(),
→ points[i]) != 2)
  S.pop();
 S.push(points[i]);
while (!S.empty()) {
 Point p = S.top();
 cout << "(" << p.x << ", " << p.v << ")" <<
→ endl;
 S.pop();
int main() {
Point points[100005];
int n;
scanf("%d", &n);
for ( int i = 0; i < n; i++ )scanf("%d %d",
convexHull(points, n);
return 0;
```

## 3.5 Point inside Poly (Ray Shooting)

```
// if strict, returns false when a is on the
    boundary
inline bool insidePoly(pt *P, int np, pt a,
    bool strict = true) {
    int numCrossings = 0;
    for (int i = 0; i < np; i++) {
        if (onSegment(P[i], P[(i + 1) % np], a))
            return !strict;
        numCrossings += crossesRay(a, P[i], P[(i + 1) % np]);
        % np]);</pre>
```

```
return (numCrossings & 1); // inside if odd

→ number of crossings
```

## 3.6 pointInPolygon

```
// Test if a point is inside a convex polygon
\rightarrow in O(\lg n) time
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
struct segment {
  point P1, P2;
  segment () {}
  segment (point P1, point P2) : P1(P1), P2(P2)
inline ll ccw (point A, point B, point C) {
 return (B.x - A.x) * (C.y - A.y) - (C.x - A.y)
\rightarrow A.x) * (B.y - A.y);
inline bool pointOnSegment (segment S, point P)
 ll^{x} = P.x, y = P.y, x1 = S.P1.x, y1 =
\rightarrow S.P1.y, x2 = S.P2.x, y2 = S.P2.y;
 ll a = x - x1, b = y - y1, c = x2 - x1, d =
y2 - y1, dot = a * c + b * d, len = c * c +
 - d * d:
 if (x1 == x2 \text{ and } y1 == y2) return x1 == x and
\rightarrow y1 == y;
  if (dot < 0 or dot > len) return 0;
  return x1 * len + dot * c == x * len and y1 *
\rightarrow len + dot * d == y * len;
const int M = 17;
const int N = 10010;
struct polygon {
  int n; // n > 1
  point p[N]; // clockwise order
  polygon () {}
  polygon (int n, point *T) {
    for (int i = 0; i < n; ++i) p[i] = T[i];
  bool contains (point P, bool strictlyInside)
    int lo = 1, hi = n - 1;
    while (lo < hi) {
      int mid = lo + hi >> 1:
      if (ccw(p[0], P, p[mid]) > 0) lo = mid +
      else hi = mid;
    if (ccw(p[0], P, p[lo]) > 0) lo = 1;
    if (!strictlyInside and
    pointOnSegment(segment(p[0], p[n - 1]), P))
    return 1
```

```
if (!strictlyInside and
    pointOnSegment(segment(p[lo], p[lo - 1]),
    P)) return 1;
    if (lo == 1 or ccw(p[0], P, p[n - 1]) == 0)
   return 0;
    return ccw(p[lo], P, p[lo - 1]) < 0;
|point P;
polygon p;
3.7 tmp
const double EPS = 1e-9, pi = acos(-1.0);
//try to use point i whenever possible
|struct point i {
int x, y;
point_i() {x = y = 0;}
 point i(int x, int y): x(x), y(y) {}
struct point {
 double x, y;
 point() \{x = y = 0.0;\}
 point(double x, double y) : x(x), y(y) {}
// operator overloading to sort the points
 bool operator < (point other) const {</pre>
  if (fabs(x - other.x) > EPS)
   return x < other.x;</pre>
  return y < other.y;</pre>
 // to check if the points are equal
 bool operator == (point other) const {
  return (fabs(x - other.x) < EPS && (fabs(y -
→ other.v) < EPS));</pre>
|<mark>double dist</mark>(point p1, point p2) {
return (p1.x - p2.x) * (p1.x - p2.x) + (p1.y - p2.x)
\rightarrow p2.y) * (p1.y - p2.y);
double DEG to RAD(double theta) { return theta
→ * pi / 180.0;}
// rotate a point by theta degree
point rotate (point p, double theta) { //theta
→ is in degree
 double rad = DEG to RAD(theta);
 return point(p.x<sup>-</sup>* cos(rad) - p.y * sin(rad)
               p.x * sin(rad) + p.y * cos(rad));
//don't know how it works
struct line {
 double a, b, c;
 // ax + by + c = 0, but b = 1.0, so y = -ax - by
line pointsToLine (point p1, point p2) {
line l;
if (fabs(p1.x - p2.x) < EPS) { // vertical line</pre>
  l.a = 1.0; l.b = 0.0 ; l.c = -p1.x;
```

```
} else {
 l.a = -(double)(p1.y = p2.y) / (p1.x - p2.x);
 l.b = 1.0;
 l.c = -(double)(l.a * p1.x) - p1.y;
bool areParallel(line l1, line l2) {
return (fabs(l1.a - l2.a) < EPS) && (fabs(l1.b)
- 12.b) < EPS):</p>
<mark>bool</mark> areSame(line l1, line l2) {
return areParallel(11, 12) && (fabs(11.c -
\rightarrow l2.c) < EPS);
// check areParallel before calling this
point areIntersect(line l1, line l2) {
point p;
p.x = (l2.b * l1.c - l1.b * l2.c) / (l2.a * l2.c) / (l2.a * l2.c) / (l2.a * l2.c)
☐ l1.b - l1.a * l2.b);
// test for vertical line
if (fabs(l1.b) > EPS) p.y = -(l1.a * p.x +

    l1.c);

else p.y = -(12.a * p.x + 12.c);
return p;
struct vec {
double x, y;
vec(double x, double y) : x(x), y(y) {}
vec toVec(point a, point b) { // convert 2
→ points to vector a->b
return vec(b.x - a.x, b.y - a.y);
vec scale(vec v, double s) { // nonnegative s =
→ [<1 .. 1 .. >1]
// shorter.same.longer
return vec(v.x * s, v.y * s);
point translate(point p, vec v) { // translate

→ p according to v

return point(p.x + v.x , p.y + v.y);
double dot(vec a, vec b) { return (a.x * b.x +
\rightarrow a.v * b.v); }
double norm sq(vec v) { return v.x * v.x + v.y

→ * v.v; }

// returns the distance from p to the line

→ defined by

// two points a and b (a and b must be
→ different)
// the closest point is stored in the 4th
   parameter (byref)
double distToLine(point p, point a, point b,
→ point &c) {
// formula: c = a + u * ab
vec ap = toVec(a, p), ab = toVec(a, b);
double u = dot(ap, ab) / norm sq(ab);
c = translate(a, scale(ab, u)); // translate a
→ to c
```

```
return dist(p, c);
} // Euclidean distance between p and c
double angle(point a, point o, point b) { //

→ returns angle aob in rad

vec oa = toVec(o, a), ob = toVec(o, b);
return acos(dot(oa, ob) / sqrt(norm sq(oa) *
\rightarrow norm sq(ob)));
double cross(vec a, vec b) { return a.x * b.y -
\rightarrow a.y * b.x; }
// note: to accept collinear points, we have to
// returns true if point r is on the left side

→ of line pa

//counter clock wise test
bool ccw(point p, point q, point r) {
return cross(toVec(p, q), toVec(p, r)) > 0;
// returns true if point r is on the same line
— as the line pa
bool collinear(point p, point q, point r) {
return fabs(cross(toVec(p, q), toVec(p, r))) <</pre>
// circles
int insideCircle(point_i p, point_i c, int r) {
→ // all integer version
int dx = p.x - c.x, dy = p.y - c.y;
int Euc = dx * dx + dy * dy, rSq = r * r; //
→ all integer
return Euc < rSq ? 0 : Euc == rSq ? 1 : 2;
      //inside/border/outside
//inscribed circle or incircle radius
double area(double ab, double bc, double ca) {
double s = (ab + bc + ca) / 2.0;
return sqrt(s * (s - ab) * (s - bc) * (s -

    ca));
//returns radius of incircle
double rInCircle(double ab, double bc, double
return area(ab, bc, ca) / (0.5 * (ab + bc +

    ca));
double rInCircle(point a, point b, point c) {
return rInCircle(dist(a, b), dist(b, c),
\rightarrow dist(c, a));
// returns 1 if there is an inCircle center,

→ returns 0 otherwise

// if this function returns 1, ctr will be the

→ inCircle center

// and r is the same as rInCircle
int inCircle(point p1, point p2, point p3,

→ point &ctr, double &r) {
r = rInCircle(p1, p2, p3);
if (fabs(r) < EPS) return 0; // no inCircle</pre>

→ center
```

```
line l1, l2; // compute these two angle
 → bisectors
 double ratio = dist(p1, p2) / dist(p1, p3);
 point p = translate(p2, scale(toVec(p2, p3),
 \rightarrow ratio / (1 + ratio)));
 l1 = pointsToLine(p1, p);
 ratio = dist(p2, \dot{p}1) / dist(p2, p3);
 p = translate(p1, scale(toVec(p1, p3), ratio /
 (1 + ratio));
 12 = pointsToLine(p2, p);
 ctr = areIntersect(l1, l2); // get their

    intersection point

 return 1:
//radius of circumcircle
double rCircumCircle(double ab, double bc,
 → double ca) {
 return ab * bc * ca / (4.0 * area(ab, bc, ca)); // line segment p-g intersect with line A-B.
double rCircumCircle(point a, point b, point c)
 return rCircumCircle(dist(a, b), dist(b, c),
 \rightarrow dist(c, a));
//polygon
1//\text{vector } P = \text{set of all points of a polygon}
// P[0] = P[n - 1]
|double perimeter(const vector<point> &P) {
 double result = 0.0;
 for (int i = 0; i < (int)P.size() - 1; i++) //
 \rightarrow remember that P[0] = P[n-1]
  result += dist(P[i], P[i + 1]);
 return result;
// returns the area, which is half the
 determinant
double area(const vector<point> &P) {
 double result = 0.0, x1, y1, x2, y2;
 for (int i = 0; i < (int)P.size() - 1; i++) {
  x1 = P[i].x; x2 = P[i + 1].x;
  y1 = P[i].y; y2 = P[i + 1].y;
  result += (x1 * y2 - x2 * y1);
 return fabs(result) / 2.0;
// returns true if all three consecutive
 → vertices of P form the same turns
bool isConvex(const vector<point> &P) {
 int sz = (int)P.size();
 if (sz <= 3) return false; // a point/sz=2 or</pre>
 → a line/sz=3 is not convex
 bool isLeft = ccw(P[0], P[1], P[2]); //

→ remember one result

 for (int i = 1; i < sz - 1; i++) // then
 if (ccw(P[i], P[i+1], P[(i+2) == sz?1:
 \rightarrow i + 2]) != isLeft)
   return false; // different sign -> this
 polygon is concave
```

```
return true;
// returns true if point p is in either

→ convex/concave polygon P

bool inPolygon(point pt, const vector<point>
→ &P) {
if ((int)P.size() == 0) return false;
double sum = 0; // assume the first vertex is
→ equal to the last vertex
for (int i = 0; i < (int)P.size() - 1; i++) {</pre>
 if (ccw(pt, P[i], P[i + 1]))
  sum += angle(P[i], pt, P[i + 1]); // left
  turn/ccw
 else sum -= angle(P[i], pt, P[i + 1]);
} // right turn/cw
return fabs(fabs(sum) - 2 * pi) < EPS;</pre>
point lineIntersectSeg(point p, point q, point
→ A, point B) {
double a = B.y - A.y;
double b = A.x - B.x;
double c = B.x * A.y - A.x * B.y;
double u = fabs(a * p.x + b * p.y + c);
double v = fabs(a * q.x + b * q.y + c);
return point((p.x * v + q.x * u) / (u + v),
(p.y * v + q.y * u) / (u + v));
// cuts polygon Q along the line formed by

→ point a -> point b

// (note: the last point must be the same as

→ the first point)

vector<point> cutPolygon(point a, point b,

    const vector<point> &Q) {

vector<point> P;
for (int i = 0; i < (int)0.size(); i++) +</pre>
 double left1 = cross(toVec(a, b), toVec(a,
\rightarrow Q[i])), left2 = 0;
 if (i != (int)Q.size() - 1) left2 =
\rightarrow cross(toVec(a, b), toVec(a, Q[i + 1]));
 if (left1 > -EPS) P.push back(Q[i]); // Q[i]

    is on the left of ab

 if (left1 * left2 < -EPS) // edge (0[i],
   O[i+1]) crosses line ab
   P.push back(lineIntersectSeg(Q[i], Q[i + 1],
   a, b));
if (!P.empty() && !(P.back() == P.front()))
 P.push back(P.front()); // make Ps first
\rightarrow point = Ps last point
return P;
// convex hull
point pivot(0, 0);
bool angleCmp(point a, point b) { //
   angle-sorting function
if (collinear(pivot, a, b)) // special case
```

```
return dist(pivot, a) < dist(pivot, b); //</pre>
double d1x = a.x - pivot.x, d1y = a.y -
   pivot.y;
double d2x = b.x - pivot.x, d2y = b.y -

→ pivot.y;

return (atan2(d1y, d1x) - atan2(d2y, d2x)) < 0;
} // compare two angles
vector<point> CH(vector<point> P) { // the
int i, j, n = (int)P.size();
if (n <= 3) {
 if (!(P[0] == P[n - 1])) P.push back(P[0]);
→ // safeguard from corner case
 return P;
} // special case, the CH is P itself
// first, find P0 = point with lowest Y and if

    → tie: rightmost X

int P0 = 0;
for (i = 1; i < n; i++)
 if (P[i].y < P[P0].y || (P[i].y == P[P0].y &&</pre>
\rightarrow P[i].x > P[P0].x))
  P0 = i;
point temp = P[0]; P[0] = P[P0]; P[P0] = temp;

→ // swap P[P0] with P[0]

// second, sort points by angle w.r.t. pivot P0
pivot = P[0]; // use this global variable as

→ reference

sort(++P.begin(), P.end(), angleCmp); // we do
→ not sort P[0]
// third, the ccw tests
vector<point> S;
S.push back(P[n - 1]); S.push back(P[0]);
S.push back(P[1]); // initial S
i = 2; // then, we check the rest
while (i < n) { // note: N must be >= 3 for

→ this method to work

 j = (int)S.size() - 1;
 if (ccw(S[j - 1], S[j], P[i]))
S.push back(P[i++]); // left turn, accept
 else S.pop back();
} // or pop the top of S until we have a left
return S; // return the result
```

# 4 Graph

### 4.1 Articulation Point Detection

```
vector<int> adj[N];
bool vis[N], articulation[N];
int low[N], tin[N], taim;
void dfs(int node, int par = -1) {
vis[node] = 1;
tin[node] = low[node] = taim++;
int children = 0;
for (int x : adj[node]) {
if (x == par) continue;
```

```
Ittehad | Nafi | Inzamam
  if (vis[x]) low[node] = min(low[node],
 \rightarrow tin[x]);
  else {
   dfs(x, node);
   low[node] = min(low[node], low[x]);
   if (low[x] >= tin[node] \&\& par != -1) {
    articulation[node] = 1;
   children++;
 if (children > 1 and par == -1)
   articulation[node] = 1;
4.2 BlockCutTree
const int N = 300010:
|bitset <N> art, good;
vector <int> g[N], tree[N], st, comp[N];
int n, m, ptr, cur, in[N], low[N], id[N];
void dfs (int u, int from = -1) {
  in[u] = low[u] = ++ptr;
  st.emplace back(u);
  for (int v : g[u]) if (v ^ from) {
      if (!in[v]) {
         dfs(v, u);
         low[u] = min(low[u], low[v]);
        if (low[v] >= in[u]) {
           art[u] = in[u] > 1 \text{ or } in[v] > 2;
          comp[++cur].emplace back(u);
          while (comp[cur].back() ^ v) {
             comp[cur].emplace back(st.back());
             st.pop back();
      } else {
         low[u] = min(low[u], in[v]);
void buildTree() {
  for (int i = 1; i \le n; ++i) {
    if (art[i]) id[i] = ++ptr;
  for (int i = 1; i <= cur; ++i) {
    int x = ++ptr;
    for (int u : comp[i]) {
      if (art[u]) {
        tree[x].emplace back(id[u]);
        tree[id[u]].empTace back(x);
      } else {
        id[u] = x;
```

int main() {

```
cin >> n >> m;
while (m--) {
  int u, v;
scanf("%d %d", &u, &v);
  g[u].emplace back(v);
  g[v].emplace back(u);
for (int i = 1; i \le n; ++i)
  if (!in[i]) dfs(i);
buildTree();
```

## 4.3 Bridge Detection

```
vector<int> adj[N];
bool visited[N];
int low[N], tin[N], timer;
vector<pair<int, int>> bridges;
void IS BRIDGE(int a, int b) {
bridges.push back({min(a, b), max(a, b)});
void dfs(int v, int p = -1) {
visited[v] = true;
tin[v] = low[v] = timer++;
for (int to : adj[v]) {
 if (to == p) continue;
 if (visited[to]) low[v] = min(low[v],

    tin[to]);

 else {
  dfs(to, v);
  low[v] = min(low[v], low[to]);
  if (low[to] > tin[v]) IS BRIDGE(v, to);
```

#### 4.4 BridgeTree

```
const int N = 300005;
vector<pair<int, int> > edge[N]; // {adjacent

→ edge, index}
vector<int> dfsTime(N), low(N), comp(N),
→ bridgeTree[N];
vector<bool> vis(N, 0), isBridge(N, 0);
int timer = 0:
//to find bridges
void dfs(int u, int par) {
vis[u] = 1:
dfsTime[u] = low[u] = ++timer;
 for (int i = 0; i < edge[u].size(); i++) {</pre>
 int v = edge[u][i].first, ind =

    edge[u][i].second;

 if (v == par)
   continue; // don't visit the parent node
  if (vis[v]) { // cross edge
  low[u] = min(low[u], dfsTime[v]);
```

```
} else {
  dfs(v, u);
  low[u] = min(low[u], low[v]);
  if (low[v] > dfsTime[u]) { // checking

→ among the back edges

    // u->v is a bridge
    isBridge[ind] = 1;
// to assign unique component number to each

→ component and its children

void dfs2(int u, int comp number) {
vis[u] = 1;
comp[u] = comp_number;
for (int i = 0; i < edge[u].size(); i++) {</pre>
 int v = edge[u][i].first, ind =

→ edge[u][i].second;

 if (!vis[v] && !isBridge[ind])
  dfs2(v, comp_number);
void make bridge tree(int n) {
// assign unique component number to each

→ component

vis.assign(n + 1, 0);
comp.assign(n + 1, -1);
int comp number = 1;
for (int^{-}i = 1; i \le n; i++) {
 if (!vis[i]) {
  dfs2(i, i);
  //i will be the root of its component
for (int i = 0; i <= n;

    i++)bridgeTree[i].clear();

//creating bridge tree
for (int i = 1; i <= n; i++) {
 for (int j = 0; j < edge[i].size(); j++) {</pre>
  int v = edge[i][j].first;
  if (comp[i] != comp[v]) {
   bridgeTree[comp[i]].push_back(comp[v]);
    bridgeTree[comp[v]].push_back(comp[i]);
void find bridges(int n) {
timer = \overline{0};
vis.assign(n + 1, 0);
low.assign(n + 1, -1);
dfsTime.assign(n + 1, -1);
for (int i = 1; i <= n; i++)
 if (!vis[i])dfs(i, -1);
```

```
4.5 Centroid Decomposition
// Builds a centroid tree of height O(logn) in
\rightarrow O(nlogn).
const int M = 2e5 + 3;
int sz[M], done[M], cpar[M], root;
vector<<mark>int</mark>>ctree[M];
void qo(int u, int p = -1) {
sz[u] = 1;
for (int v : g[u]) {
  if (v == p or done[v]) continue;
  go(v, u);
  sz[u] += sz[v];
int find centroid(int v, int p, int n) {
for (int x : g[v]) {
 if (x != p \text{ and } !done[x] \text{ and } sz[x] > n / 2)

→ return find centroid(x, v, n);

return v;
void decompose(int v = 0, int p = -1) {
qo(v);
 int c = find centroid(v, -1, sz[v]);
 if (p == -1) root = c;
 done[c] = 1;
 cpar[c] = p;
if (p != -1) ctree[p].push_back(c);
 for (int x : q[c]) {
 if (!done[x]) decompose(x, c);
4.6 DSU on Tree
vector <int> G[mx]; /// adjacency list of the
int sub[mx]; /// subtree size of a node
int color[mx]; /// color of a node
int freq[mx];
void calcSubSize(int s, int p) {
sub[s] = 1;
for (int x : G[s]) {
  if (x == p) continue;
  calcSubSize(x, s);
  sub[s] += sub[x];
void add(int s, int p, int v, int bigchild =
freq[color[s]] += v;
for (int x : G[s]) {
 if (x == p || x == bigchild) continue;
  add(x, s, v);
void dfs(int s, int p, bool keep) {
```

```
int bigChild = -1;
for (int x : G[s]) {
 if (x == p) continue;
 if (bigChild == -1 || sub[bigChild] < sub[x]</pre>
→ ) biaChild = x:
for (int x : G[s]) {
 if (x == p || x == bigChild) continue;
 dfs(x, s, 0);
if (bigChild != -1) dfs(bigChild, s, 1);
add(s, p, 1, bigChild);
/// freq[c] now contains the number of nodes in
/// the subtree of 'node' that have color c
/// Save the answer for the queries here
if (keep == 0) add(s, p, -1);
int main() {
input color
construct G
calcSubSize(root, -1);
dfs(root, -1, 0);
return 0;
```

## 4.7 Dijkstra

```
#define pii pair<long long,int>
vector<int>Edges[100005];
vector<long long>Cost[100005];
long long dis[100005];
int vis[100005];
void dijkstra( int source ) {
priority queue< pii, vector<pii>, greater<pii>
Q.push(pii(0, source));
dis[source] = 0;
pii q;
while ( !Q.empty() ) {
 q = Q.top();
 Q.pop();
 int u = q.second;
 if ( vis[u] != -1 ) continue; // *idk why*
 vis[u] = 1; // *idk why*
 for ( int i = 0; i < Edges[u].size(); i++ ) {</pre>
  int v = Edges[u][i];
  if ( vis[v] != -1 ) continue; // *idk why*
  if ( dis[u] + Cost[u][i] < dis[ v ] ) {</pre>
   dis[ v ] = dis[u] + Cost[u][i];
    Q.push( pii( dis[v], v ) );
return;
```

```
4.8 Dinic
// O(V^2 E), solves SPOJ FASTFLOW
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
struct edge {
 int u, v;
ll cap, flow;
 edge () {}
 edge (int u, int v, ll cap) : u(u), v(v),
 \rightarrow cap(cap), flow(0) {}
struct Dinic {
int N;
vector <edge> E;
vector <vector <int>> q;
vector <int> d, pt;
Dinic (int N) : \dot{N}(N), E(\theta), g(N), d(N), pt(N)
void AddEdge (int u, int v, ll cap) {
 if (u ^ v) {
  E.emplace_back(u, v, cap);
g[u].emplace_back(E.size() - 1);
   E.emplace back(v, u, 0);
   g[v].emplace back(E.size() - 1);
 bool BFS (int S, int T) {
  queue <int> q({S});
  fill(d.begin(), d.end(), N + 1);
  d[S] = 0;
  while (!q.empty()) {
   int u = q.front(); q.pop();
   if (u == T) break;
   for (int k : q[u]) {
    edge \&e = E[K]:
    if (e.flow < e.cap and d[e.v] > d[e.u] + 1)
     d[e.v] = d[e.u] + 1;
     q.emplace(e.v);
 } return d[T] != N + 1;
 ll DFS (int u, int T, ll flow = -1) {
  if (u == T or flow == 0) return flow;
  for (int &i = pt[u]; i < g[u].size(); ++i) {
  edge &e = E[g[u][i]];</pre>
  edğe &oe = E[g[u][i]'^1]
   if (d[e.v] == d[e.u] + 1) {
    ll amt = e.cap - e.flow;
    if (flow != -1 and amt > flow) amt = flow;
    if (ll pushed = DFS(e.v, T, amt)) {
     e.flow += pushed:
     oe.flow -= pushed;
     return pushed;
```

```
Ittehad | Nafi | Inzamam
 } return 0:
ll MaxFlow (int S, int T) {
 ll total = 0;
 while (BFS(S, T)) {
  fill(pt.begin(), pt.end(), 0);
  while (ll flow = DFS(S, T)) total += flow;
  return total;
int main() {
int N, E;
scanf("%d %d", &N, &E);
Dinic dinic(N);
 for (int i = 0, u, v; i < E; ++i) {
 ll cap;
 scanf("%d %d %lld", &u, &v, &cap);
  dinic.AddEdge(u - 1, v - 1, cap);
  dinic.AddEdge(v - 1, u - 1, cap);
printf("%lld\n", dinic.MaxFlow(0, N - 1));
 return 0;
4.9 HeavyLight Decomposition
→ -> Need to Clear ??
class HeavyLightDecomposition
```

```
vector < int > List[ ?? ]; // Tree's Adj List
#define L R ? ?
public :
vector<int> ValueOfNode;
vector<int> Position;
vector<int> Parent;
vector<int> Depth;
vector<int> Heavy;
 vector<int> Head;
int CurrentPosition = 1; // 0/1 - index based
 segmentTree ST = segmentTree( ?? ) /

¬ AnyQueryTree;

HeavyLightDecomposition(int NN) {
 ValueOfNode.resize(NN);
  Position.resize(NN);
  Parent.resize(NN, -1);
  Depth.resize(NN, 0);
 Heavy.resize(NN, -1);
 Head.resize(NN);
int DFS(int Vertex) {
 int TotalSize = 1;
  int MaxChildSize = 0;
  for (int i = 0; i < List[Vertex].size(); ++i)</pre>
  int Child = List[Vertex][i];
  if (Child != Parent[Vertex]) {
    Parent[Child] = Vertex;
    Depth[Child] = Depth[Vertex] + 1;
```

```
int ChildSize = DFS(Child);
    TotalSize += ChildSize;
    if (ChildSize > MaxChildSize) {
    MaxChildSize = ChildSize;
    Heavy[Vertex] = Child;
 return TotalSize:
void TreeDecompose(int Vertex, int Hd) {
 Head[Vertex] = Hd;
 ST.A[CurrentPosition] = ValueOfNode[Vertex];
 Position[Vertex] = CurrentPosition++;
  if (Heavy[Vertex] != -1)
  TreeDecompose(Heavy[Vertex], Hd);
  for (int i = 0; i < List[Vertex].size(); ++i)</pre>
   int Child = List[Vertex][i];
   if (Child != Parent[Vertex] && Child !=
   Heavy[Vertex])
    TreeDecompose(Child, Child);
void MakeQueryTree() { // ?? = Number of Node
→ in Tree;
// Build Query Data Structure ? ?
int Query(int NodeA, int NodeB) {
 int Res = 0;
 while (Head[NodeA] != Head[NodeB]) {
  if (Depth[Head[NodeA]] > Depth[Head[NodeB]])
    swap(NodeA, NodeB);
   int CurrentPathResult =
    ST.rangeQuery(L R, Position[Head[NodeB]],
  Position[NodeB]).Value;
   Res = ? ? (Res, CurrentPathResult);
  NodeB = Parent[Head[NodeB]];
  if (Depth[NodeA] > Depth[NodeB])
  swap(NodeA, NodeB);
  int LastHeavyPathResult =
    ST.rangeQuery(L R, Position[NodeA],
→ Position[NodeB]).Value;
 Res = ? ? (Res, LastHeavyPathResult);
  return Res;
int Update(int NodeA, int NodeB, int X) {
 while (Head[NodeA] != Head[NodeB]) {
  if (Depth[Head[NodeA]] > Depth[Head[NodeB]])
  swap(NodeA, NodeB);
ST.rangeUpdate(L_R, Position[Head[NodeB]],
→ Position[NodeB], X);
  NodeB = Parent[Head[NodeB]];
 if (Depth[NodeA] > Depth[NodeB])
  swap(NodeA, NodeB);
```

```
ST.rangeUpdate(L R, Position[NodeA],
 → Position[NodeB], X;
};
4.10 Hopcroft Karp
// Maximum biparite matching. Complexity :

→ 0(E*sqrt(V))

//define NIL (dummy vertex), M and INF
vector<int>q[M];
int Lmatch[M], Rmatch[M], dist[M];
bool bfs(int n) {
queue<int>q;
 for (int u = 1; u <= n; u++) {
  if (Lmatch[u] == NIL) dist[u] = 0, q.push(u);
  else dist[u] = INF;
 dist[NIL] = INF;
 while (!q.empty()) {
  int u = q.front();
  q.pop();
  if (dist[u] < dist[NIL]) {</pre>
   for (int v : g[u]) {
    if (dist[Rmatch[v]] == INF) {
     dist[Rmatch[v]] = dist[u] + 1;
     q.push(Rmatch[v]);
 return dist[NIL] != INF;
bool dfs(int u) {
if (u == NIL) return true;
 for (int v : q[u]) {
  if (dist[Rmatch[v]] == dist[u] + 1 and

    dfs(Rmatch[v])) {

   Rmatch[v] = u;
   Lmatch[u] = v;
   return true;
 dist[u] = INF;
 return false;
int HopcroftKarp(int n, int m) {
fill(Lmatch, Lmatch + n + 1, 0);
 fill(Rmatch, Rmatch + m + 1, 0);
 int res = 0;
 while (bfs(n)) {
  for (int u = 1; u <= n; u++) {
   if (Lmatch[u] == NIL and dfs(u)) res++;
 return res;
```

```
4.11 Hungarian
/*returns maximum/minimum weighted bipartite
, matching. Complexity : O(N^2 * M)
flag = -1 minimizes, flag = 1 maximizes. */
#define CLR(a) memset(a, 0, sizeof a)
ll weight[N][M];
int used[M], P[M], way[M], match[M];
ll U[M], V[M], minv[M], ara[N][M];
ll hungarian(int n, int m, int flag) {
CLR(U), CLR(V), CLR(P), CLR(ara), CLR(way);
 for (int i = 1; i <= n; i++) {
 for (int j = 1; j \le m; j++)
   ara[i][j] = -flag * weight[i][j];
 if (n > m) m = n;
 int a, b, d;
 ll r, w;
 for (int i = 1; i <= n; i++) {
  P[0] = i, b = 0;
  for (int j = 0; j <= m; j++) minv[j] = INF,

    used[j] = false;

   used[b] = true;
   a = P[b], d = 0, w = INF;
   for (int j = 1; j \le m; j++) {
    if (!used[j]) {
     r = ara[a][j] - U[a] - V[j];
     if (r < minv[j]) minv[j] = r, way[j] = b;
     if (minv[j] < w) w = minv[j], d = j;
   for (int j = 0; j \le m; j++) {
    if (used[j]) U[P[j]] += w, V[j] -= w;
    else minv[j] -= w;
   b = d;
  } while (P[b] != 0);
  do {
   d = way[b];
   P[b] = P[d], b = d;
 } while (b != 0);
 for (int j = 1; j \le m; j++) match[P[j]] = j;
 return flag * V[0];
4.12 LCA(sparse table)
vector<<mark>int</mark>>Edges[10000];
int p[10005][17], level[10005], n, lg;
bool vis[10005];
void DFS( int par, int node ) {
vis[node] = 1;
 if ( par != -1 ) level[node] = level[par] + 1;
 p[node][0] = par;
 for ( int i = 1; i <= lq; i++ )
```

```
if ( p[node][i - 1] != -1 ) p[node][i] =
\rightarrow p[p[node][i - 1]][i - 1];
for ( int i = 0; i < Edges[node].size(); i++ )</pre>
 if ( vis[ Edges[node][i] ] == 0 ) DFS( node,
  Edges[node][i] );
return;
int LCA( int u, int v ) {
if ( level[u] < level[v] ) swap(u, v);</pre>
for ( int i = lq; i >= 0; i-- ) {
 int par = p[u][i];
 if ( level[par] >= level[v] ) {
  u = par;
if ( u == v ) return u;
for ( int i = lq; i >= 0; i-- ) {
 int U = p[u][i];
 int V = p[v][i];
 if ( U != V ) {
  u = U: v = V:
return p[u][0];
```

## 4.13 Max Flow Edmond Karp

```
int n;
vector<vector<int>> capacity;
vector<vector<int>> adj;
int bfs(int s, int t, vector<int> &parent) {
fill(parent.begin(), parent.end(), -1);
parent[s] = -2;
queue<pair<int, int>> q;
q.push({s, INF});
while (!q.empty()) {
 int cur = q.front().first;
 int flow = q.front().second;
  q.pop();
 for (int next : adj[cur]) {
   if (parent[next] = -1 &&
   capacity[cur][next]) {
    parent[next] = cur;
    int new flow = min(flow,
   capacity[cur][next]);
    if (next == t) return new flow;
    q.push({next, new flow});
return 0;
int maxflow(int s, int t) {
int flow = 0;
```

```
University of Dhaka
 vector<int> parent(n);
 int new flow;
while (\overline{new} \text{ flow} = bfs(s, t, parent)) {
  flow += \overline{new} flow;
  int cur = t;
  while (cur != s) {
   int prev = parent[cur];
   capacity[prev][cur] -= new_flow;
   capacity[cur][prev] += new flow;
   cur = prev;
 return flow;
4.14 Max Flow-1
int graph[105][105];
int rgraph[105][105];
int par[105];
int n;
int bfs( int s, int d ) {
bool vis[105];
 memset( vis, 0, sizeof(vis) );
 queue<int>Q;
 Q.push(s);
 while ( !Q.empty() ) {
  int q = Q.front();
  Q.pop();
  for ( int i = 1; i <= n; i++ ) {
  if (vis[i] == 0 \&\& rgraph[q][i] > 0) {
    vis[i] = 1;
    par[i] = q;
    if ( i == d ) return 1;
    Q.push(i);
return 0;
int max flow( int s, int d ) {
int total flow = 0;
 for ( int i = 1; i <= n; i++ ) {
 for ( int j = 1; j <= n; j++ ) rgraph[i][j] =

¬ graph[i][j];
```

int mn;

mn = INT MAX;

par[child] ) {
int P = par[child];

par[child] ) ·

int P = par[child];

rgraph[P][child] -= mn;

rgraph[child][P] += mn;

while ( bfs( s, d ) == 1 ) {

for ( int child = d; child != s; child =

for ( int child = d; child != s; child =

mn = min(mn, rgraph[P][child] );

```
total flow += mn:
return total flow;
4.15 Online Bridge
vector<int> par, dsu 2ecc, dsu cc, dsu cc size;
int bridges;
int lca iteration;
vector<int> last visit;
void init(int n) {
par.resize(n);
 dsu 2ecc.resize(n);
 dsu cc.resize(n);
 dsu cc size.resize(n);
 lcaiteration = 0;
 las\overline{t} visit.assign(n, 0);
 for (int i = 0; i < n; ++i) {
 dsu_2ecc[i] = i;
  dsu^{-}cc[i] = i;
  dsu^{-}cc size[i] = 1;
  par[i] = -1
 bridges = 0;
int find 2ecc(int v) {
if (v = -1)
 return -1;
 return dsu 2ecc[v] == v ? v : dsu 2ecc[v] =

    find 2ecc(dsu 2ecc[v]);

int find cc(int v) {
v = find 2ecc(v);
return d\overline{s}u cc[v] == v ? v : dsu cc[v] =

    find cc(dsu cc[v]);

void make root(int v) {
v = find_2ecc(v);
int root = v:
int child = -1;
 while (v != -1) {
  int p = find 2ecc(par[v]);
  par[v] = chiTd;
  dsu cc[v] = root;
  chiTd = v;
  v = p;
 dsu cc size[root] = dsu cc size[child];
void merge path (int a, int b) {
++lca iteration;
vector<int> path a, path b;
int lca = -1;
 while (lca == -1) {
 if (a`!= -1) {
   a = find 2ecc(a);
   path a.p\overline{u}sh back(a);
```

```
if (last visit[a] == lca iteration) {
    lca = a;
    break;
   last visit[a] = lca iteration;
   a = \overline{par}[a];
 if (b != -1) {
   b = find 2ecc(b);
   path b.p\overline{u}sh back(b);
   if (\lambda ast visit[b] == lca iteration) {
    break;
   last_visit[b] = lca iteration;
   b = \overline{par[b]};
for (int v : path a) {
 dsu \ 2ecc[v] = lca;
 if (v == lca)
  break;
 --bridges;
for (int v : path_b) {
 dsu 2ecc[v] = lc\overline{a};
 if (v == lca)
  break:
 --bridges;
void add edge(int a, int b) {
a = find 2ecc(a);
b = find^2 ecc(b);
if (a == b) return;
int ca = find cc(a);
int cb = find cc(b);
if (ca != cb) -{
 ++bridges;
 if (dsu cc size[ca] > dsu cc size[cb]) {
  swap(a, b);
   swap(ca, cb);
 make root(a);
 par[\overline{a}] = dsu cc[a] = b;
 dsu cc size[cb] += dsu cc size[a];
} else ₹
 merge path(a, b);
```

4.16 SCC

/\*In a directed graph, an SCC is a connected component where all nodes are pairwise reachable. ,

```
condesation graph is the DAG built on a directed 4.17 centroid_root_decomposition
, graph by compressing each SCC into a node. define M */
vector<int>g[M], gr[M];
set<int>gc[M];
int vis[M], id[M], sz[M];
vector<int>order, comp, roots;
namespace SCC {
void addEdge(int u, int v) {
g[u].push back(v), gr[v].push back(u);
void dfs1(int u) {
vis[u] = 1;
for (int x : g[u])
 if (!vis[x]) dfs1(x);
 order.push back(u);
void dfs2(int u) {
vis[u] = 1;
 comp.push back(u);
for (int \overline{x} : gr[u])
 if (!vis[x]) dfs2(x);
void condense(int n) {
fill(vis, vis + n + 1, 0);
 for (int i = 1; i <= n; i++)
 if (!vis[i]) dfs1(i);
 reverse(order.begin(), order.end());
 fill(vis, vis + \bar{n} + 1, 0);
 for (int u : order) {
 if (!vis[u]) {
  dfs2(u); //this part of the code processes
- components, returns them in comp
  for (int v : comp) id[v] = u;
   sz[u] = (int)comp.size();
   roots push back(u);
   comp.clear();
 fill(vis, vis + n + 1, 0);
 for (int u = 1; u <= n; u++) {
  for (int v : g[u]) {
  if (id[u] != id[v])
    gc[id[u]].insert(id[v]);
void reset(int n) {
order.clear(), comp.clear(), roots.clear();
 for (int i = 1; i \le n; i++) {
 g[i].clear(), gr[i].clear(), gc[i].clear(); id[i] = vis[i] = sz[i] = 0;
```

```
const int MAXN = 100050;
const int LOGN = 17;
int par[LOGN][MAXN];
                     // par[i][v]: (2^i)th

→ ancestor of v

int level[MAXN], sub[MAXN]; // sub[v]: size of
int ctPar[MAXN], n; // ctPar[v]: parent of v in

→ centroid tree

vector<int> adj[MAXN];
bool vis[MAXN];
int ans[MAXN]; // ans[v]: shortest distance
_ between v and red nodes in subtree
long long INF = 1e18;
// calculate level by dfs
void dfsLevel(int node, int pnode) {
for(auto cnode : adj[node]) {
 if(cnode != pnode) {
  par[0][cnode] = node;
   level[cnode] = level[node] + 1;
   dfsLevel(cnode, node);
void preprocess() -
level[0] = 0; par[0][0] = 0;
dfsLevel(0, -1);
 for(int i = 1; i < LOGN; i++)
 for(int node = 0; node < n; node++)</pre>
  par[i][node] = par[i-1][par[i-1][node]];
int lca(int u, int v) {
if(level[u] > level[v]) swap(u, v);
int d = level[v] - level[u];
 // make u, v same level
 for(int i = 0; i < LOGN; i++) {
 if(d & (1 << i)) {
  v = par[i][v];
if(u == v) return u;
 // find LCA
 for(int i = LOGN - 1; i >= 0; i--) {
 if(par[i][u] != par[i][v]) {
  u = par[i][u];
  v = par[i][v];
return par[0][u];
int dist(int u, int v) {
return level[u] + level[v] - 2 * level[lca(u,

    ∨)];

/* Centroid decomposition */
// Calculate size of subtrees by dfs
void dfsSubtree(int node, int pnode)
```

```
sub[node] = 1;
for(auto cnode : adj[node]) {
 if(cnode != pnode && vis[cnode] == 0) {
  dfsSubtree(cnode, node);
   sub[node] += sub[cnode];
// find Centroid
int dfsCentroid(int node, int pnode, int size) {
for(auto cnode : adj[node]) {
 if(cnode != pnode && sub[cnode] > size / 2 &&

    vis[cnode] == 0 )
   return dfsCentroid(cnode, node, size);
return node;
// Centroid decomposition
void decompose(int node, int pCtr) {
dfsSubtree(node, -1);
int ctr = dfsCentroid(node, node, sub[node]);
vis[ctr] = 1;
if(pCtr == -1)
 pCtr = ctr; // root of centroid tree
 ctPar[ctr] = pCtr;
 for(auto cnode : adj[ctr]) {
 if( vis[cnode] == 0 ) decompose(cnode, ctr);
adj[ctr].clear();
// color node v red
void update(int v) {
int rNode = v;
while(1) {
 ans[v] = min(ans[v], dist(rNode, v));
 if(v == ctPar[v]) break;
 v = ctPar[v];
// reply query
int query(int v) {
int start = v;
int minD = INF;
while(1) {
 minD = min(minD, dist(start, v) + ans[v]);
 if(v == ctPar[v]) break;
 v = ctPar[v];
return minD;
int main() {
preprocess();
decompose(0, -1);
fill(ans, ans + n, INF);
update(0);
```

```
5 Grundy
5.1 grundy
int grundyValue[ Nn ];
int mexValue[ Nn ], MEX;
void calculateGrundyValue() {
      grundyValue[ 1 ] = grundyValue[ 2 ] = 0;
      for ( int i = 3; i <= 10000; ++i ) {
             for ( int j = 1; j + j < i; ++j )
                   mexValue[ grundyValue[ j ] ^
    grundyValue[ i - j ] ] = i;
            MEX = 0;
            while ( mexValue[ MEX ] == i )
                   MEX++;
            grundyValue[ i ] = MEX;
void solve( int t ) {
      int N, a;
      cin >> N;
      int XOR = 0;
      while ( N-- ) {
            cin >> a;
            XOR \( \text{ } = \text{ grundyValue[ a ];} \)
      cout << "Case " << t << ": ";
      if ( XOR ) cout << "Alice\n";</pre>
```

## 6 Math

## 6.1 Matrices

#### 6.1.1 Gauss-Jordan Elimination in GF(2)

```
const int SZ = 105;
const int MOD = 1e9 + 7;
bitset <SZ> mat[SZ];
int where[SZ];
bitset <SZ> ans;
ll bigMod(ll a, ll b, ll m) {
ll ret = 1LL;
a %= m;
while (b) {
  if (b & 1LL) ret = (ret * a) % m;
 a = (a * a) % m;
  b >>= 1LL;
return ret;
/// n for row, m for column, modulo 2
int GaussJordan(int n, int m) {
SET(where); /// sets to -1
for (int r = 0, c = 0; c < m \&\& r < n; c++) {
 for (int i = r; i < n; i++)
  if ( mat[i][c] ) {
    swap(mat[i], mat[r]); break;
 if ( !mat[r][c] ) continue;
  where [c] = r;
```

```
for (int i = 0; i < n; ++i) if (i != r &&
→ mat[i][c]) mat[i] ^= mat[r];
for (int j = 0; j < m; j++) {
 if (where[j] != -1) ans[j] = mat[where[j]][m]
  / mat[where[j]][j];
 else ans[i] = 0;
for (int i = 0; i < n; i++) {
 int sum = 0;
 for (int j = 0; j < m; j++) sum ^{=} (ans[j] &

    mat[i][i]);

 if ( sum != mat[i][m] ) return 0; /// no

→ solution

int cnt = 0;
for (int j = 0; j < m; j++) if (where[j] ==
\rightarrow -1) cnt++;
return bigMod(2, cnt, MOD); /// how many
  solutions modulo some other MOD
```

### 6.1.2 Gauss-Jordan Elimination in GF(P)

```
const int SZ = 105:
const int MOD = 1e9 + 7;
int mat[SZ][SZ], where[SZ], ans[SZ];
ll bigMod(ll a, ll b, ll m) {
ll ret = 1LL;
a %= m;
while (b) {
 if (b & 1LL) ret = (ret * a) % m;
 a = (a * a) % m;
 b >>= 1LL;
return ret;
int GaussJordan(int n, int m, int P) {
SET(where): /// sets to -1
for (int r = 0, c = 0; c < m \&\& r < n; c++) {
 for (int i = r; i < n; i++) if ( mat[i][c] >

→ mat[mx][c] ) mx = i;

 if ( mat[mx][c] == 0 ) continue;
 if (r != mx) for (int j = c; j <= m; j++)
→ swap(mat[r][j], mat[mx][j]);
 where [c] = r;
  int mul, minv = bigMod(mat[r][c], P - 2, P);
  int temp;
  for (int i = 0; i < n; i++) {
  if ( i != r && mat[i][c] != 0) {
   mul = ( mat[i][c] * (long long) minv ) % P;
   for (int j = c; j \le m; j++)
    temp = mat[i][j];
     temp -= ( ( mul * (long long) mat[r][j] )
→ % P ):
     temp += P;
```

```
if ( temp >= P ) temp -= P;
    mat[i][j] = temp:
for (int j = 0; j < m; j++) {
if (where [j] != -1) ans [j] =
   (mat[where[j]][m] * 1LL *
  bigMod(mat[where[j]][j], P - 2, P) ) % P;
 else ans[i] = 0;
for (int i = 0; i < n; i++) {
 int sum = 0:
 for (int j = 0; j < m; j++) {
  sum += ( ans[j] * 1LL * mat[i][j] ) % P;</pre>
  if (sum >= P) sum -= P;
 if ( sum != mat[i][m] ) return 0; /// no
  solution
for (int j = 0; j < m; j++) if (where[j] ==

→ -1) cnt++:

return bigMod(P, cnt, MOD);
```

#### 6.1.3 Gauss-Jordan Elimination

```
/*** mat is 0 based
* In every test case, clear mat first and then

→ do the changes

* For solving problems on graphs with
   probability/expectation, make sure the graph
is connected and a single component. If not,

    → then re-number the vertex and solve

for each connected component separately.
* Complexity --> O( min(n,m) * nm ) **/
const int SZ = 105;
const double EPS = 1e-9;
double mat[SZ][SZ], ans[SZ];
int where[SZ];
int GaussJordan(int n, int m) {
SET(where); /// sets to -1
for (int r = 0, c = 0; c < m \&\& r < n; c++) {
 int mx = r;
 for (int i = r; i < n; i++) if (
→ abs(mat[i][c]) > abs(mat[mx][c]) ) mx = i;
 if ( abs(mat[mx][c]) < EPS ) continue;</pre>
 if (r != mx) for (int j = c; j <= m; j++)
→ swap(mat[r][j], mat[mx][j]);
 where [c] = r;
 for (int i = 0; i < n; i++) if ( i != r ) {
    double mul = mat[i][c] / mat[r][c];
   for (int j = c; j <= m; j++) mat[i][j] -=
    mul * mat[r][i]:
```

#### 6.1.4 Matrix expo

```
long long a, b, n, m, F[2][2], f[2][2];
long long p = 1e9 + 7;
void multiply( long long a[2][2], long long
\rightarrow b[2][2]) {
long long q[2][2];
for ( int i = 0; i < 2; i++ ) {
 for ( int j = 0; j < 2; j++ ) {
  g[i][j] = 0; for ( int k = 0; k < 2; k++)
   g[i][j] = ((g[i][j] % p) + ((a[i][k] % p)
     (b[k][i] % p)) % p ) % p;
for ( int i = 0; i < 2; i++ ) {
 for ( int j = 0; j < 2; j++ ) F[i][j] =
  g[i][j];
void power( long long N ) {
if ( N == 1 ) return;
if (N \% 2 == 0) { power(N / 2);
→ multiply(F, F); }
else {power(N - 1); multiply(F, f);}
return;
```

#### 6.2 Modular Arithmetic

#### 6.2.1 Chinese Remainder Theorem

```
ll <mark>inv</mark>(ll a, ll m) {
ll m0 = m, t, q;
11 \times 0 = 0, \times 1 = 1;
if (m == 1) return 0;
while (a > 1) {
  q = a / m; t = m; m = a % m, a = t; t = x0;
\rightarrow x0 = x1 - q * x0; x1 = t;
if (x1 < 0) x1 += m0;
 return x1;
Il findMinX(ll num[], ll rem[], ll k) {
ll prod = 1;
 for (ll i = 0; i < k; i++) prod *= num[i];
ll result = 0;
 for (ll i = 0; i < k; i++) {
  ll pp = prod / num[i];
  result += rem[i] * inv(pp, num[i]) * pp;
return result % prod;
int main() {
ll num[15], rem[15], n, t, i, j; scanf("%lld", &t);
 for (i = 1; i <= t; i++) {
  scanf("%lid", \&n);
for (j = 0; j < n; j++)
  scanf("%lld %lld", &num[j], &rem[j]);
  printf("Case %lld: %lld\n", i, findMinX(num,
\rightarrow rem, n));
```

### 6.2.2 Discrete Log

```
int solve(int a, int b, int m) {
a \% = m, b \% = m;
int k = 1, add = 0, g;
while ((g = gcd(a, m)) > 1)  {
 if (b = k)
  return add;
 if (b % q)
  return -1;
  b /= g, m /= g, ++add;
 k = (k * 111 * a / g) % m;
int n = sqrt(m) + 1;
int an = 1;
for (int i = 0; i < n; ++i)
 an = (an * 111 * a) % m;
unordered map<int, int> vals;
for (int \overline{q} = 0, cur = b; q \le n; ++q) {
 vals[cur] = q;
 cur = (cur * 111 * a) % m;
for (int p = 1, cur = k; p \le n; ++p) {
 cur = (cur * 1ll * an) % m;
 if (vals.count(cur)) {
  int ans = n * p - vals[cur] + add;
```

```
return ans;
return -1:
6.2.3 Modular Inverse (EGCD)
int gcdExtended(int a, int b, int* x, int* y) {
if (a == 0) {
 *x = 0, *y = 1;
 return b;
int x1, y1;
 int gcd = gcdExtended(b % a, a, \&x1, \&y1);
 *x = y1 - (b / a) * x1;
*y = x1;
return gcd;
void modInverse(int a, int m) {
int x, y;
int q = qcdExtended(a, m, &x, &y);
 printf("Inverse doesn't exist");
else {
 int res = (x % m + m) % m;
 printf("Modular multiplicative inverse is
```

#### 6.2.4 Modular Inverse

```
int gcdExtended(int a, int b, int* x, int* y) {
if (a == 0) {
 *x = 0, *y = 1;
 return b;
int x1, y1;
int gcd = gcdExtended(b % a, a, &x1, &y1);
*x = y1 - (b / a) * x1;
*v = x1;
return qcd;
void modInverse(int a, int m) {
int x, y;
int q = gcdExtended(a, m, &x, &y);
if (q != 1)
 printf("Inverse doesn't exist");
 int res = (x % m + m) % m;
 printf("Modular multiplicative inverse is
```

# 6.2.5 nCr Lucas /\*use this to calculate nCr modulo mod, when → mod is smaller than n and m. define MOD Complexity: O(mod + log mod n) \*/ll fact[MOD]; ll bigmod(int x, int p) { ll res = 1;while (p) { if (p & 1) res = res \* x % MOD; $x = x \times x \times x \times MOD;$ p >>= 1;return res; 11 modinv(ll x) { return bigmod(x, MOD - 2); void precalc() { //run this fact[0] = 1;for (int i = 1; i < MOD; i++) fact[i] = fact[i - 1] \* i % MOD; int C(int n, int m) { if (m > n) return 0; if (m == 0 or m == n) return 1; ll ret = fact[n] \* modinv(fact[m]) % MOD; return ret \* modinv(fact[n - m]) % MOD; int nCr(int n, int m) { if (m > n) return 0; if (m == 0) return 1; return nCr(n / MOD, m / MOD) \* C(n % MOD, m % → MOD) % MOD;

# 6.3 Polynomial Multiplication

## 6.3.1 FFT

```
typedef cplx cd;
//define N as a power of two greater than the

    size

, of any possible polynomial
using cd = complex<double>;
const double PI = acosl(-1);
int rev[N]; cd w[N];
static cd f[N];
void prepare(int &n) {
int sz = builtin ctz(n);
for (int \overline{i} = 1; i < n; i++) rev[i] = (rev[i >>
-1] >> 1) | ((i & 1) << (sz - 1));
w[0] = 0, w[1] = 1, sz = 1;
while (1 << sz < n) {
 cd w n = cd(cos(2 * PI / (1 << (sz + 1))),
\rightarrow \sin(2 * PI / (1 << (sz + 1)));
  for (int i = 1 \ll (sz - 1); i < (1 \ll sz);
```

```
SZ++;
|void fft(cd *a, int n) {
for (int i = 1; i < n - 1; i++) {
 if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
 for (int h = 1; h < n; h <<= 1) {
 for (int s = 0; s < n; s += h << 1) {
  for (int i = 0; i < h; i++) {
    cd \& u = a[s + i], \& v = a[s + i + h], t = v
   * w[h + i];
    v = u - t, u = u + t;
vector<ll>multiply(vector<ll>a, vector<ll>b) {
int n = a.size(), m = b.size(), sz = 1;
if (!n or !m) return {};
while (sz < n + m - 1) sz <<= 1;
prepare(sz);
 for (int i = 0; i < sz; i++) f[i] = cd(i < n ?
\rightarrow a[i] : 0, i < m ? b[i] : 0);
 fft(f, sz);
 for (int i = 0; i \le (sz >> 1); i++) {
  int j = (sz - i) \& (sz - 1);
  cd x = (f[i] * f[i] - conj(f[j] * f[j])) *
\rightarrow cd(0, -0.25);
 f[j] = x, f[i] = conj(x);
 fft(f, sz);
 vector<ll>c(n + m - 1);
for (int i = 0; i < n + m - 1; i++) c[i] =
→ round(f[i].real() / sz);
return c;
```

# 6.3.2 NTT

```
const int G = 3;
                                               const int MOD = 998244353;
                                               const int N = ?; // (1 << 20) + 5; greater than</pre>
                                               , maximum possible degree of any polynomial
                                               int rev[N], w[N], inv n;
                                               int bigMod(int a, int e, int mod) {
                                                if (e == -1) assert(false);
                                                if (e == -1) e = mod - 2;
                                                int ret = 1;
                                                while (e) {
                                                 if (e & 1) ret = (ll) ret * a % mod;
                                                 a = (ll) a * a % mod; e >>= 1;
                                                return ret:
                                               void prepare(int &n) {
                                                                    builtin clz(n));
                                                int sz = abs(31 -
w[i << 1] = w[i], w[i << 1 | 1] = w[i] * w n; int r = bigMod(G, (MOD - 1)) / n, MOD);
```

```
inv n = bigMod(n, MOD - 2, MOD), w[0] = w[n] =
for (int i = 1; i < n; ++i) w[i] = (ll) w[i -

→ 1] * r % MOD;
for (int i = 1; i < n; ++i) rev[i] = (rev[i >>
\rightarrow 1] >> 1) | ((i & 1) << (sz - 1));
void ntt (int *a, int n, int dir) {
for (int i = 1; i < n - 1; ++i) {
 if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
for (int m = 2; m <= n; m <<= 1) {
 for (int i = 0; i < n; i += m) {
  for (int j = 0; j < (m >> 1); ++j) {
   int \&u = a[i + j], \&v = a[i + j + (m >> 1)];
    int t = (ll) v * w[dir ? n - n / m * j : n
   / m * j] % MOD;
   v = u - t < 0 ? u - t + MOD : u - t;
    u = u + t >= MOD ? u + t - MOD : u + t;
if (dir) for (int i = 0; i < n; ++i) a[i] =
   (ll) a[i] * inv n % MOD;
int f a[N], f b[N];
vector <int> multiply (vector <int> a, vector
int sz = 1, n = a.size(), m = b.size();
while (sz < n + m - 1) sz <<= 1; prepare(sz);
for (int i = 0; i < sz; ++i) f a[i] = i < n ?
\rightarrow a[i] : 0;
for (int i = 0; i < sz; ++i) f b[i] = i < m?
\rightarrow b[i] : 0;
ntt(f a, sz, 0); ntt(f b, sz, 0);
for (int i = 0; i < sz; ++i) f a[i] = (ll)

    f a[i] * f b[i] % MOD;

ntt(f a, sz, 1); return vector <int> (f a, f a
  + n + m - 1);
// G = primitive root(MOD)
int primitive root (int p) {
vector <int> factor:
int tmp = p - 1;
for (int i = 2; i * i <= tmp; ++i) {
 if (tmp % i == 0) {
  factor.emplace back(i);
  while (tmp \% i == 0) tmp /= i;
if (tmp != 1) factor.emplace back(tmp);
for (int root = 1; ; ++root) {
 bool flag = true;
 for (int i = 0; i < (int) factor.size(); ++i)</pre>
→ {
```

```
if (bigMod(root, (p - 1) / factor[i], p) ==
    flag = false; break;
  if (flag) return root;
int main() {
//(x + 2)(x + 3) = x^2 + 5x + 6
 vector \langle int \rangle a = \{2, 1\};
 vector <int> b = {3, 1};
 vector <int> c = multiply(a, b);
 for (int x : c) cout \langle \langle x \rangle \rangle; cout \langle \langle end \rangle \rangle Euler Totient
 return 0;
```

#### 6.4 Catalan Number

```
unsigned long int binomialCoeff(unsigned int n,
unsigned int k) {
unsigned long int res = 1;
if (k > n - k)
 k = n - k:
for (int i = 0; i < k; ++i) {
 res *= (n - i); res /= (i + 1);
return res;
unsigned long int catalan(unsigned int n) {
unsigned long int c = binomialCoeff(2 * n, n);
return c / (n + 1);
```

### 6.5 Diophantine Equation

```
int gcd extend(int a, int b, int& x, int& y)
if (b == 0) {
 x = 1:
 v = 0;
 return a;
 else {
 int g = gcd extend(b, a % b, x, y);
 int x1 = x, y1 = y;
 x = y1;
  y = x1 - (a / b) * y1;
  return q;
void print solution(int a, int b, int c) {
int x, y;
if (a == 0 \& \& b == 0) {
 if (c == 0) {
   cout << "Infinite Solutions Exist" << endl;</pre>
  else {
   cout << "No Solution exists" << endl;</pre>
```

```
int gcd = gcd extend(a, b, x, y);
if (\check{c} \% gc\check{d} != 0) {
 cout << "No Solution exists" << endl;</pre>
else {
 cout^{-} << "x = " << x * (c / gcd) << ", y = "
\rightarrow << y * (c / qcd) << endl;
```

```
int phi(int n) {
int result = n;
for (int i = 2; i * i <= n; i++) {
 if (n % i == 0) {
  while (n \% i == 0) n /= i;
   result -= result / i:
if (n > 1)
 result -= result / n;
return result;
void phi 1 to n(int n) {
vector<int> phi(n + 1);
phi[0] = 0;
phi[1] = 1;
for (int i = 2; i \le n; i++)
 phi[i] = i;
for (int i = 2; i \le n; i++) {
 if (phi[i] == i) {
  for (int j = i; j <= n; j += i)
   phi[j] -= phi[j] / i;
```

## 6.7 FastSieve

```
primes up to 5e8 within 0.35 seconds
// primes up to 1e9 within 1 second
vector <int> fastSieve (const int N, const int
\rightarrow Q = 17, const int L = 1 << 15) {
 const int M = (N + 29) / 30;
 const int two = sqrt(N), four = sqrt(two);
  static const int r[] = \{1, 7, 11, 13, 17, 19,

→ 23, 29};

 struct P {
    P (int p) : p(p) {}
   int p, pos[8];
 auto approxPrimeCount = [] (const int N) ->

    int {
    return N > 60184 ? N / (log(N) - 1.1) :
   \max(1.0, N / (\log(N) - 1.11)) + 1
```

```
vector <bool> isPrime(two + 1, true);
 for (int i = 2; i <= four; ++i) if
  (isPrime[i]) {
     for (int j = i * i; j <= two; j += i)
  isPrime[i] = false:
 const int r size = approxPrimeCount(N + 30);
 int p size = 3;
 vector <P> s primes;
 vector \langle int \rangle primes = \{2, 3, 5\};
 int p beg = 0, prod = 1;
 primes.resize(r size);
 for (int p = 7; p <= two; ++p) {
   if (!isPrime[p]) continue;
   if (p <= Q) prod *= p, ++p beg,

¬ primes[p size++] = p;

   auto cur = P(p);
   for (int t = 0; t < 8; ++t) {
     int j = (p \le Q) ? p : p * p;
     while (j % 30 != r[t]) j += p << 1;
     cur.pos[t] = j / 30;
   s primes.push back(cur);
 vector <unsigned char> pre(prod, 0xFF);
 for (size t it = 0; it < p beg; ++it) {
   auto cur = s primes[it];
   const int p = cur.p;
   for (int t = 0; t < 8; ++t) {
     const unsigned char m = \sim (1 \ll t);
     for (int i = cur.pos[t]; i < prod; i +=
  p) pre[i] \&= m;
 const int block size = (L + prod - 1) / prod
→ * prod;
vector <unsigned char> block(block size);
 unsigned char *p block = block.data();
 for (int beg = 0; beg < M; beg += block_size,</pre>
→ p block -= block size) {
   int end = min(M, beg + block size);
   for (int i = beg; i < end; i + prod) {
     copy(pre.begin(), pre.end(), p_block + i);
   if (beg == 0) p block[0] &= 0xFE;
   for (size t it = p beg; it <

¬ s primes.size(); ++it) {

     auto &cur = s primes[it];
     const int p = cur.p;
     for (int t = 0; t < 8; ++t) {
       int i = cur.pos[t];
       const unsigned char m = \sim (1 \ll t);
       for (; i < end; i += p) p block[i] &= m;
       cur.pos[t] = i;
```

```
for (int i = beg; i < end; ++i) {
     for (int m = p block[i]; m > 0; m &= m -
→ 1) {
        primes[p_size++] = i * 30 +
→ r[__builtin_ctz(m)];
 assert(p size <= r size);</pre>
 while (p size > 0 and primes[p size - 1] > N)

→ --p size;

 primes.resize(p size); return primes;
int main() {
 int LIM; cin >> LIM;
 auto primes = fastSieve(LIM);
```

```
6.8 Primes
const int N = 10000000 + 6;
vector<long long>primes;
bitset<N>flag;
vector<long long>v;
void siv() {
flag[1] = 1;
for ( int i = 2; i * i <= N; i++ ) {
 if ( flag[i] == 0 ) {
  for ( int j = i * i; j < N; j += i ) flag[j]
for ( int i = 2; i < N; i++ ) {
 if ( flag[i] == 0 ) primes.push_back(i);
long long mul(long long a, long long b, long
→ long mod) {
long long res = 0;
a \approx mod;
while (b) {
 if (b \& 1) res = (res + a) % mod;
 a = (2 * a) % mod;
 b >>= 1; // b = b / 2
return res;
long long mod inverse( long long n, long long p
long long x, y, g;
g = gcd extended( n, p, x, y );
if ( g < 0 ) x = -x;
return (x % p + p) % p;
long long mpow( long long x, long long y, long
→ long mod ) {
long long ret = 1;
while ( v )
```

```
if ( y & 1 ) ret = mul(ret, x, mod);
 y >>= 1, x = mul(x, x, mod);
return ret % mod;
int isPrime( long long p ) {
if (p < 2 | | !(p \& 1)) return 0;
if (p == 2) return 1;
long long q = p - 1, a, t;
int k = 0, b = 0;
while (!(q \& 1)) q >= 1, k++;
 for ( int it = 0; it < 2; it++ ) {
 a = rand() % (p - 4) + 2;
 t = mpow( a, q, p );
b = (t == 1) || (t == p - 1);
 for ( int i = 1; i < k && !b; i++ ) {
  t = mul(t, t, p);
  if ( t == p - \bar{1} ) b = 1;
 if ( b == 0 ) return 0:
 return 1:
long long pollard rho( long long n, long long c
long long x = 2, y = 2, i = 1, k = 2, d;
while ( 1 ) {
 x = ( mul(x, x, n) + c );
 if (x >= n) x -= n;
 d = gcd(x - y, n);
  if (d > 1) return d;
 if ( ++i == k ) y = x, k <<= 1;
return n;
map<long long, int>mp;
void factorize( long long n ) {
int l = primes.size();
for ( int i = 0; primes[i]*primes[i] <= n && i</pre>
if ( n % primes[i] == 0 ) {
  mp[primes[i]] = 1;
  while ( n % primes[i] == 0 ) n /= primes[i];
if (n != 1) mp[n] = 1:
void lfactorize( long long n ) {
if ( n == 1 ) return;
if ( n < 1e9 ) {
 factorize(n);
  return;
if ( isPrime(n) ) {
 mp[n] = 1:
 return;
 long long d = n;
```

```
for ( int i = 2; d == n; i++ ) d =
→ pollard rho(n, i);
lfactorize(d);
lfactorize(n / d);
long long f(long long r, vector<long long> v1) {
int sz = v1.size();
long long res = 0;
for (long long i = 1; i < (1 << sz); i++) {
 int ct = 0;
 long long mul = 1;
 for (int j = 0; j < sz; j++) {
  if (i & (1 << j)) {
   ct++;
   mul *= v1[j];
 long long sign = -1;
 if (ct \& 1) sign = 1;
 res += sign * (r / mul);
return r - res;
```

#### 6.9 Striling Number of 2nd kind

```
long long p = 1e9 + 7;
long long fact[1000005];
int n, m, k;
long long s( long long N, long long R )
if ( N == 0 && R == 0 ) return 1;
if ( N == 0 \mid \mid R == 0 ) return 0:
long long ans = 0;
for (int i = 1; i \le R; i++) {
 long long par;
 if ( (R - 1) \% 2 == 0 ) par = 1;
 else par = -1;
 par = (par + p) % p;
 long long temp = (ncr(R, i) * bm(i, N)) % p;
 temp = (temp \% p * par \% p) \% p;
 ans = (ans \% p + temp \% p) \% p;
return (ans * bm( fact[R], p - 2 )) % p;
```

#### 7 Misc

#### 7.1 Build (Nafi)

```
"cmd" : ["g++ -std=c++14 $file_name -o
$file_base name && timeout 6s
- ./$file_base name<in>out"],
"selector" : "source.c, source.cpp, source.Cc",
"shell": true,
"working dir" : "$file path"
```

# 7.2 Build files //pragma #pragma GCC optimize("03") #pragma GCC optimize("unroll-loops") compile: q++-std = c++17 - I. - Dakifpathan -→ '0 "%e" "%f" build: q++-std = c++17DHFTF - Wshadow - o "%e" "%f" fsanitize = address - fsanitize = undefined D GLIBCXX DEBUG run: - "./%e" //for sublime "cmd" : ["g++ -std=c++14 \$file name -o \$file base name && timeout 6s ./\$file base name<in>out"], "selector" : "source.c, source.cpp, source.Cc", "shell": true, "working dir" : "\$file path" //windows "cmd": ["g++.exe", "-std=c++14", "\${file}", "\${file base name}.exe", "&&", "\${file base name}.exe<in>out"], "shell": true, "working dir": "\$file path", "selector": "source.cpp, source.c, source.c++, → source.cc"

# 7.3 Ternary Search

```
while (hi >= lo)
int mid1 = lo + (hi - lo) / 3; int mid2 = hi -
\rightarrow (hi - lo) / 3;
if (f(mid1) > f(mid2)) { } //change
else //change
}//ittehad
double x1, why1, z1, x2, y2, z2, x, y , z;
double f( double t )
double xt = x1 + (x2 - x1)t;
double yt = why1 + (y2 - why1)t;
double zt = z1 + (z2 - z1)t;
return ((xt - x)(xt - x) + (yt - y)(yt - y) +
\rightarrow (zt - z)(zt - z));
double Tsearch()
double low = 0, high = 1, mid;
int step = 64;
while ( step-- ) {
 double t1 = (2low + high) / 3;
```

```
double t2 = (low + 2high) / 3;
  double d1 = \dot{f}(t1);
  double d2 = f(t2);
  if ( d1 < d2 ) high = t2;
  else low = t1;
 return low;
7.4 fastIO
ios base::sync with stdio(false);

    cin.tie(NULL); cout.tie(NULL)

8 String
8.1 Aho-Corasick
struct vartex {
 int next[30], endmark, link;
 vector<int>dlink;
 vartex() {
  memset(next, -1, sizeof(next));
  endmark = -1;
  link = 0;
void addstring(string& s, vector<vartex>&trie) {
 int v = 0:
 for (auto x : s) {
  if (trie[v].next[x - 'a'] == -1) {
   trie[v].next[x - 'a'] = trie.size();
   trie.emplace back();
  v = trie[v].next[x - 'a'];
 trie[v].endmark = 0;
void fail(vector<vartex>&trie) {
 int v = 0;
 trie[v].link = 0;
 queue<int>q;
 q.push(0);
 while (!q.empty()) {
  v = q.front();
  q.pop();
  for (int i = 0; i < 26; i++) {
   if (trie[v].next[i] != -1) {
    if (v == 0) {
     trie[trie[v].next[i]].link = 0;
     else {
     int x = trie[v].link;
     while (x != 0 \& \& trie[x].next[i] == -1) {
      x = trie[x].link;
     if (trie[x].next[i] == -1) {
      trie[trie[v].next[i]].link = 0;
```

```
else -
     trie[trie[v].next[i]].link =
  trie[x].next[i];
   q.push(trie[v].next[i]);
void dictionary link(vector<vartex>&trie) {
queue<int>q;
q.push(0);
while (!q.empty()) {
 int u = q.front();
 q.pop();
 for (int i = 0; i < 26; i++) {
  if (trie[u].next[i] != -1) {
   q.push(trie[u].next[i]);
 int k = u;
 while (k != 0) {
  if (trie[k].endmark != -1 && k != u) {
   trie[u].dlink.push back(k);
  k = trie[k].link;
 debug(u, trie[u].dlink);
int search(string& s, vector<vartex>&trie) {
int v = 0:
for (auto x : s) {
 v = trie[v].next[x - 'a'];
return trie[v].endmark;
```

### 8.2 Aho-Corasick\_New

```
struct node {
    bool f = false;
    char c;
    int p = -1, link = -1, ex = -1;
    int to[26], go[26];
    vector <int> id;
    node() {
        memset(to, -1, sizeof(to));
        memset(go, -1, sizeof(go));
    }
};
int siz = 1;
vector <node> trie(1);
void insert(string &s, int len, int j) {
    int u, v = 0;
    for (int i = 0; i < len; i++) {</pre>
```

```
int c = s[i] - 97;
        if (trie[v].to[c] == -1) {
            trie.emplace back();
            trie[v].to[c] = siz++;
        \ddot{u} = v;
        v = trie[v].to[c];
        trie[v].p = u, trie[v].c = s[i];
    trie[v].f = true;
    trie[v].id.push back(j);
int go(int v, char c);
int get link(int v) {
    if (trie[v].link == -1) {
        if (v == 0 \mid | trie[v].p == 0)

    trie[v].link = 0;

        else trie[v].link =
    go(get link(trie[v].p), trie[v].c);
    return trie[v].link;
int get exit link(int v) {
    if (trie[v].ex == -1) {
        int u = get link(v);
        if (u == 0 | | trie[u].f) trie[v].ex = u; long long h[400005];
        else trie[v].ex = get exit link(u);
    return trie[v].ex;
int go(int v, char c) {
    int x = c - 97;
    if (trie[v].qo[x] == -1) {
        if (trie[v].to[x] != -1) trie[v].go[x]
else trie[v].go[x] = v ?
    qo(qet link(v), c) : 0;
    return trie[v].go[x];
string s, t;
vector <int> a[100005];
int n, k[100005], len, ln[100005];
void get id(int v, int i) {
    int sz = trie[v].id.size();
    for (int j = 0; j < sz; j++) {
        int p = trie[v].id[j];
        a[p].push back(i - ln[p]);
void fun(int i, int v) {
    if (trie[v].f) get id(v, i);
    int u = get exit link(v);
    while (u > 0) {
        if (trie[u].f) get id(u, i);
        u = get exit link(\overline{u});
    if (i < s.size()) fun(i + 1, go(v, s[i]));
```

```
int query(int i) {
    int s = a[i].size(), ans = -1;
    for (int j = k[i] - 1, p = 0; j < s; j++,
→ p++) {
        int x = a[i][j] + ln[i] - a[i][p];
        ans = (ans == -1) ? x : min(ans, x);
    return ans;
|int main() {
    cin >> s >> n;
    len = s.length();
    for (int i = 0; i < n; i++) {
    cin >> k[i] >> t;
         ln[i] = t.length();
        insert(t, ln[i], i);
    fun(0, 0);
    for (int i = 0; i < n; i++) cout <<
    querv(i) << endl:</pre>
```

### 8.3 Hashing without inv

```
long long MOD[400005];
int L;
void pre hash( string s ) {
long long p = 31 , m = 1e9 + 9, power = 1,
\rightarrow hash = 0;
int z = 0:
for ( int i = s.size() - 1; i >= 0; i-- ) {
 hash = ( hash * p + (s[i] - 'A' + 1) ) % m;
 h[i] = hash;
 MOD[z] = power;
 power = (power * p) % m;
long long f( int l, int r ) {
long long val = h[r], m = 1e9 + 9;
if ( l != L - 1 ) {
 long long val2 = (h[l + 1] % m * MOD[l - r +
→ 1] % m ) % m;
 val -= val2;
 val += m;
 val %= m;
if ( val < 0 ) val = (val + m) % m;
return val;
```

## 8.4 KMP

```
#define pii pair<int,int>
vector<int> prefix function (string Z) {
int n = (int) Z.length();
vector<int> F (n);
```

```
F[0] = 0;
for (int i = 1; i < n; ++i) {
int i = F[i - 1];
while (j > 0 && Z[i] != Z[j])
j = F[j - 1];
 if (Z[i] == Z[j]) ++j;
F[i] = j;
return F;
```

#### 8.5 Manacher

```
/// When i is even, pal[i] = largest
    palindromic substring centered from str[i /
/// When i is odd, pal[i] = largest palindromic
   substring centered between str[i / 2] and
\rightarrow str[i / 2] + 1
vector <int> manacher(char *str) {
int i, j, k, l = strlen(str), n = l << 1;</pre>
vector <int> pal(n);
for (i = 0, j = 0, k = 0; i < n; j = max(0, j)
\rightarrow - k), i += k) {
 while (j \le i \&\& (i + j + 1) < n \&\& str[(i - i)]
\rightarrow j) >> 1] == str[(i + j + 1) >> 1])
  1++;
 for (k = 1, pal[i] = j; k \le i \&\& k \le pal[i]
\sim && (pal[i] - k) != pal[i - k]; k++) {
  pal[i + k] = min(pal[i - k], pal[i] - k);
pal.pop back();
return pal;
int main() {
char str[100];
while (scanf("%s", str)) {
 auto v = manacher(str);
 for (auto it : v) printf("%d ", it);
 puts("");
return 0:
```

#### 8.6 Palindromic Tree

```
#define CLR(a) memset(a,0,sizeof(a))
* str is 1 based
Each node in the palindromic tree denotes a

→ STRING

Node 1 denotes an imaginary string of size -1
Node 2 denotes a string of size 0
They are the two roots
There can be maximum of (string length + 2)
    nodes in total
```

```
It's a directed tree. If we reverse the
direction of the suffix links, we get a dag. In
this DAG, if node v is reachable from node u
iff, u is a substring of v.
* if ( tree[A].next[x] == B )
then, B = xAx
* if ( tree[A].suffixLink == B )
Then B is the longest possible palindrome which

→ is a proper suffix of A

(node 1 is an exception)
* occ[i] contains the number of occurrences of
the corresponding palindrome
* st[i] denotes starting index of the first

→ occurrence of the corresponding palindrome

* st[] or occ[] or both can be ignored if not

    needed

* If memory limit is compact, a map has to be

    used instead of

ed[MAXN][MAXC]. Swapping row and column of the

→ matrix will

save more memory.
Example:
map <int,int> ed[MAXC];
ed[c][u] = v means, there is an edge from node
node v that is labeled character c.
namespace pt {
const int MAXN = 100010; /// maximum possible

→ string size

const int MAXC = 26; /// Size of the character

→ set

int n; /// length of str
char str[MAXN];
int len[MAXN], link[MAXN], ed[MAXN][MAXC],

    occ[MAXN], st[MAXN];

int nc, suff, pos;
/// nc -> node count
/// suff -> Index of the node denoting the
   longest palindromic proper suffix of the
- current prefix
void init() {
str[0] = -1;
nc = 2; suff = 2;
len[1] = -1, link[1] = 1;
len[2] = 0, link[2] = 1;
CLR(ed[1]), CLR(ed[2]);
occ[1] = occ[2] = 0;
inline int scale(char c) { return c - 'a'; }
inline int nextLink(int cur) {
while (str[pos - 1 - len[cur]] != str[pos])

    cur = link[cur];

return cur;
inline bool addLetter(int p) {
pos = p;
int let = scale(str[pos]);
int cur = nextLink(suff);
```

```
if (ed[cur][let]) {
  suff = ed[cur][let];
  occ[suff]++;
  return false;
 suff = ++nc:
 CLR(ed[ncl):
 len[nc] = len[cur] + 2;
 ed[cur][let] = nc;
 occ[nc] = 1;
 if (len[nc] == 1) {
  st[nc] = pos;
  link[nc] = 2;
  return true;
 link[nc] = ed[nextLink(link[cur])][let];
 st[nc] = pos - len[nc] + 1;
 return true;
void build(int n) {
n = n;
init();
 for (int i = 1; i <= n; i++) addLetter(i);</pre>
 for (int i = nc; i >= 3; i--) occ[link[i]] +=
 → occ[i];
occ[2] = occ[1] = 0;
void printTree() {
 puts(str);
 cout << "Node\tStart\tLength\t0cc\n";</pre>
 for (int i = 3; i <= nc; i++) {
  cout << i << "\t" << st[i] << "\t" << len[i]
 - << "\t" << occ[i] << "\n";</pre>
int main() {
scanf("%s", pt::str + 1);
 pt::build(strlen(pt::str + 1));
 return 0:
8.7 String Hashing
ll bigmod(ll x, ll p, ll md) {
 ll res = 1;
 while (p) {
  if (p \& 1) res = res * x % md;
  x = x * x % md;
  p >>= 1;
 return res;
ll modinv(ll x, ll md) {
return bigmod(x, md - 2, md);
namespace Hash {
|ll pw[M][2];
ll invpw[M][2];
const int pr[] = {37, 53};
```

```
const int md[] = {1000000007, 1000000009};
void precalc() {
 pw[0][0] = pw[0][1] = 1;
 for (int i = 1; i < M; i++) {
 pw[i][0] = pw[i - 1][0] * pr[0] % md[0];
 pw[i][1] = pw[i - 1][1] * pr[1] % md[1];
 invpw[M - 1][0] = modinv(pw[M - 1][0], md[0]);
 invpw[M - 1][1] = modinv(pw[M - 1][1], md[1]);
 for (int i = M - 2; i >= 0; i--) {
 invpw[i][0] = invpw[i + 1][0] * pr[0] % md[0];
 invpw[i][1] = invpw[i + 1][1] * pr[1] % md[1];
pii get hash(const string &s) {
pii re\overline{t} = \{0, 0\};
for (int i = 0; i < s.size(); i++) {</pre>
  ret.first += (s[i] - 'a' + 1) * pw[i][0] %
 ret.second += (s[i] - 'a' + 1) * pw[i][1] %
\rightarrow md[1]:
 if (ret.first >= md[0]) ret.first -= md[0];
 if (ret.second >= md[1]) ret.second -= md[1];
return ret:
void prefix(const string &s, pii *H) {
H[0] = \{0, 0\};
 for (int i = 1; i <= s.size(); i++) {
 H[i].first = H[i - 1].first + (s[i - 1] - 'a')
\rightarrow + 1) * pw[i - 1][0] % md[0];
 H[i].second = H[i - 1].second + (s[i - 1] -
\rightarrow 'a' + 1) * pw[i - 1][1] % md[1];
 if (H[i].first >= md[0]) H[i].first -= md[0];
 if (H[i].second >= md[1]) H[i].second -=
\rightarrow md[1];
void reverse prefix(const string &s, pii *H) {
int n = s.size();
for (int i = 1; i <= s.size(); i++) {
 H[i].first = H[i - 1].first + (s[i - 1] - 'a')
\rightarrow + 1) * pw[n - i][0] % md[0];
 H[i].second = H[i - 1].second + (s[i - 1] -
\rightarrow 'a' + 1) * pw[n - i][1] % md[1];
 if (H[i].first >= md[0]) H[i].first -= md[0];
 if (H[i].second >= md[1]) H[i].second -=
\rightarrow md[1];
pii range hash(int L, int R, pii H[]) {
ret.first = (H[R].first - H[L - 1].first +

→ md[0]) % md[0];

 ret.second = (H[R].second - H[L - 1].second +
\rightarrow md[1]) % md[1];
 ret.first = ret.first * invpw[L - 1][0] %
\rightarrow md[0];
```

```
ret.second = ret.second * invpw[L - 1][1] %
\rightarrow md[1];
return ret;
pii reverse hash(int L, int R, pii H[], int n) {
 ret.first = (H[R].first - H[L - 1].first +
\rightarrow md[0]) % md[0]
 ret.second = (H[R].second - H[L - 1].second +
\rightarrow md[1]) % md[1];
 ret.first = ret.first * invpw[n - R][0] %
\rightarrow md[0];
 ret.second = ret.second * invpw[n - R][1] %
\rightarrow md[1]:
 return ret;
8.8 Suffix Array
  O(|S| + |alphabet|) Suffix Array
  LIM := max{s[i]} + 2
void inducedSort (const vector <int> &vec, int
   val range, vector <int> &SA, const vector
- <int> &sl, const vector <int> &lms idx) {
vector <int> l(val range, 0), r(val range, 0);
 for (int c : vec) \overline{\{}
 ++r[c]; if (c + 1 < val range) ++l[c + 1];
 partial sum(l.begin(), l.end(), l.begin());
 partial sum(r.begin(), r.end(), r.begin());
 fill(SA.begin(), SA.end(), -1);
for (int i = lms idx.size() - 1; i >= 0; --i)
\rightarrow SA[--r[vec[lms idx[i]]] = lms idx[i];
 for (int i : SA) if (i > 0 and sl[i - 1])
\rightarrow SA[l[vec[i - 1]]++] = i - 1;
fill(r.begin(), r.end(), 0);
 for (int c : vec) ++r[c];
partial sum(r.begin(), r.end(), r.begin());
 for (int k = SA.size() - 1, i = SA[k]; k; --k,
\rightarrow i = SA[k]) \rightarrow
  if (i and !sl[i - 1]) SA[--r[vec[i - 1]]] = i
 vector <int> suffixArray (const vector <int>
const int n = vec.size();
 vector <int> sl(n), SA(n), lms idx;
 for (int i = n - 2; i >= 0; --\overline{i}) {
 sl[i] = vec[i] > vec[i + 1] or (vec[i] ==
 \rightarrow vec[i + 1] and sl[i + 1]);
  if (sl[i] and !sl[i + 1])
→ lms idx.emplace back(i + 1);
```

```
reverse(lms idx.begin(), lms idx.end());
inducedSort(vec, val range, SA, sl, lms_idx);
vector <int> new lms idx(lms_idx.size()),
\rightarrow lms vec(lms i\overline{d}x.s\overline{i}ze());
for (int i = 0, k = 0; i < n; ++i) {
 if (SA[i] > 0 and !sl[SA[i]] and sl[SA[i] -
\rightarrow 1]) new lms idx[k++] = SA[i];
int cur = 0; SA[n - 1] = 0;
for (int k = 1; k < new lms idx.size(); ++k) {
 int i = new lms idx[k - 1], j =

→ new lms idx[k]:

 if (vec[i] ^ vec[i]) {
  SA[i] = ++cur; continue;
  bool flag = 0;
  for (int a = i + 1, b = j + 1; ; ++a, ++b) {
  if (vec[a] ^ vec[b]) {
   flaq = 1: break:
  if ((!sl[a] and sl[a - 1]) or (!sl[b] and
  sl[b - 1])) {
   flag = !(!sl[a] \text{ and } sl[a - 1] \text{ and } !sl[b]
→ and sl[b - 1]); break;
 SA[j] = flag ? ++cur : cur;
for (int i = 0; i < lms idx.size(); ++i)</pre>
→ lms vec[i] = SA[lms idx[i]];
if (cur + 1 < lms idx.\overline{size}) {
 auto lms SA = suffixArray(lms vec, cur + 1);
 for (int i = 0; i < lms idx.size(); ++i)
→ new lms idx[i] = lms idx[lms SA[i]];
inducedSort(vec, val range, SA, sl,
   new lms idx); return SA;
vector <int> getSuffixArray (const string &s,
vector <int> vec(s.size() + 1);
copy(begin(s), end(s), begin(vec)); vec.back()
\Rightarrow = '$';
auto ret = suffixArrav(vec, LIM);
ret.erase(ret.begin()); return ret;
// build RMQ on it to get LCP of any two suffix
vector <int> getLCParray (const string &s,

→ const vector <int> &SA) {
int n = s.size(), k = 0;
vector <int> lcp(n), rank(n);
for (int i = 0; i < n; ++i) rank[SA[i]] = i;
for (int i = 0; i < n; ++i, k? --k : 0) {
 if (rank[i] == n - 1) {
  k = 0; continue;
```

```
8.9 Suffix Automaton
// collected from cp algorithm
struct state {
int len, link, cnt, firstpos; // cnt -> endpos
- set size, link -> suffix link
map <char, int> next;
const int MAXLEN = 100002;
state st[MAXLEN * 2];
struct SuffixAutomata { // 0-based
int sz, last;
SuffixAutomata() { // init
 st[0].cnt = st[0].len = 0;
 st[0].link = -1;
 sz = 1, last = 0;
void add(char c) { // add new char in automata
 int cur = sz++;
 st[cur].len = st[last].len + 1;
 st[cur].firstpos = st[cur].len - 1;
 st[cur].cnt = 1;
 int p = last;
 while (p != -1 \&\& !st[p].next.count(c)) {
  st[p].next[c] = cur;
   p = st[p].link;
 if (p == -1) {
  st[cur].link = 0;
 else {
  int q = st[p].next[c];
   if (st[p].len + 1 == st[q].len) {
    st[cur].link = q;
   else { // clone state
    int clone = sz++;
    st[clone].len = st[p].len + 1;
    st[clone].next = st[q].next;
    st[clone].link = st[q].link;
    st[clone].firstpos = st[q].firstpos;
    st[clone].cnt = 0;
    while (p != -1 \&\& st[p].next[c] == q) {
    st[p].next[c] = clone;
    p = st[p].link;
    st[q].link = st[cur].link = clone;
```

```
last = cur:
void occurrence() { // calculate number of
→ occurrences of all possible substring
 vector <int> rank(sz);
 iota(all(rank), 0);
 sort(all(rank), [&](int i, int j) {
  return st[i].len > st[j].len;
 for (int ii : rank) if (st[ii].link != -1)
   st[st[ii].link].cnt += st[ii].cnt;
int count(string s) { // number of occurrences
   of string s. #prerequisite -> call
→ occurrence()
 int node = 0;
 for (char ch : s) {
 if (!st[node].next.count(ch)) return 0;
  node = st[node].next[ch];
 return st[node].cnt;
int firstOcc(string s) { // first
→ position(occurence) of string s
 int node = 0;
 for (char ch : s) {
 if (!st[node].next.count(ch)) return -1;
  node = st[node].next[ch];
 return st[node].firstpos + 2 - (int)s.size();
void build(string S) { // build suffix automata
for (char ch : S) add(ch);
bool find(string s) { // find string s in

→ automata
 int node = 0;
 for (char ch : s) {
```

```
if (!st[node].next.count(ch)) return false;
   node = st[node].next[ch];
  return true;
8.10 Trie
//define\ M,\ K = alphabet\ size
int trie[M][K], word[M * K + 3], cnt[M * K +
→ 3], SZ;
void Insert(string s) {
int node = 0:
for (int i = 0; i < s.size(); i++) {
  int c = s[i] - 'a';
  if (!trie[node][c]) {
  trie[node][c] = ++sz;
  node = trie[node][c];
  cnt[node]++;
word[node]++;
bool Search(string s) {
int node = 0, ret = 0;
for (int i = 0; i < s.size(); i++) {
  int c = s[i] - 'a';
  if (!trie[node][c]) return false;
 node = trie[node][c];
return (word[node] > 0);
void Delete(string s) {
int node = 0;
vector < int > v(1, 0);
for (int i = 0; i < s.size(); i++) {</pre>
 int c = s[i] - 'a';
```

```
node = trie[node][c];
  cnt[node]--:
 v.push back(node);
word[node]--;
 for (int i = 1; i < v.size(); i++) {
 int c = s[i - 1] - 'a';
 if (!cnt[v[i]]) {
  trie[v[i - 1]][c] = 0;
8.11 Z algo
```

```
/* z[i] denotes the maximum length of substring
* starting from position(i) which is also a
→ prefix
* of the string
* call with Z \tilde{z}f(x) where x is the desired

    strina*/

struct Z {
int n; string s;
vector<int>z:
Z(const string \&a) {
 n = a.size(); s = a; z.assign(n, \theta);
void z function() {
 for (int i = 1, l = 0, r = 0; i < n; ++i) {
  if (i \le r) z[i] = min (r - i + 1, z[i - l]);
  while (i + z[i] < n \&\& s[z[i]] == s[i +
  z[i]]) ++z[i];
   if (i + z[i] - 1 > r) l = i, r = i + z[i] - r
   1;
```