MPI: Microprocessors and Interfacing
Academic Year: 2022 - 23

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Ecole Centrale School of Engineering

Lab4- Due Date: September 20, 2022

You are given with n positive Integers. Write assembly language functions to compute the following
(use NASM assembler).

1. To find number of even numbers
2. To find number of odd numbers
3. To find number of prime numbers
4. To find GCD of n numbers
5. To find LCM of n numbers

Lab3- Due Date: September 10, 2022

Write assembly language functions to compute the following (use NASM assembler). Assume that all elements are integers.

1. Find the maximum and the minimum of n numbers.
2. Assume that you are given with a list of n elements and **key**, search the **key** using linear search. If the required element is found, return its position. Otherwise return -1.
3. Evaluate $1 + x + x^2 + x^3 + ... + x^n$
   (Assume that x and n are positive integers)
4. Assume that you are given with a list of n elements then find **mode** and **median**.

Lab2 - Due date: August 31, 2022.

Installations.

1. Install NASM in your laptop and run a sample code.
2. Install MASM in your laptop and run a sample code.

Lab1 - Due date: August 24, 2022.

Develop C-Programs for the following problem statements.

1. Print the internal representation of data stored in primary data types: **int, float, and double**.
2. **Bonus Problem** Perform addition and multiplication of two 32-bit numbers (Please remember that the input is 32-bit binary number). The numbers can either be integers or real numbers. Assume that the integers are represented in 2's complement form and the real numbers are represented using single-precision IEEE 754 standard.

Submission Guide Lines

- **Max. team size is 4.**
- Mail-ID: cs3106.mpi@gmail.com
- Sub:TEAM_NUM_LAB_NUM
- Attach.Name and Type: (Sub.).zip
- Late Submission<=3-Days:50%.
- Write a readme file to understand your solutions.

Number Systems

- Representation of Integer Numbers
  - Signed Magnitude Representation
  - 1's Complement Representation
  - 2's Complement Representation
- Representation of Real Numbers
  - Fixed Point Representation
  - Floating Point Representation

Resolution is difference between two successive numbers.

Representation of Integer Numbers

Let A $= a_{n-1}a_{n-2}a_{n-3}...a_0$ is an n-bit binary number
if A is an **unsigned integer**, then value of A is : $\sum_{i=0}^{n-1}(2^i \times a_i)$.
if A is a **signed integer**:

- Signed Magnitude Representation:
  - A $= \sum_{i=0}^{n-2}(2^i \times a_i)$, if $a_{n-1} = 0$
  - A $= -\sum_{i=0}^{n-2}(2^i \times a_i)$, if $a_{n-1} = 1$
- 1's Complement Rep.: A $= -(2^{n-1} - 1) \times a_{n-1} + \sum_{i=0}^{n-2}(2^i \times a_i)$
- 2's Complement Rep.: A $= -2^{n-1} \times a_{n-1} + \sum_{i=0}^{n-2}(2^i \times a_i)$

Resolution: 1

Range of Numbers

Let A $= a_{n-1}a_{n-2}a_{n-3}...a_0$ is an n bit binary number
if A is an **unsigned integer**, then range of A is : 0 to $(2^n - 1)$.
if A is a **signed integer**:

- Signed Magnitude Rep., range of A is : $-(2^{n-1} - 1)$ to $(2^{n-1} - 1)$.
- 1's Complement Rep., range of A is : $-(2^{n-1} - 1)$ to $(2^{n-1} - 1)$.
- 2's Complement Rep., range of A is : -2^{n-1} to $(2^{n-1} - 1)$.

Expansion of Bit Length

Add additional bit positions to the left and fill in with value of the sign bit.
Let A = **1 0 1 0** is a 4-bit binary number,
Representation of A using 8-bits (i.e. B): **1 1 1 1 1 0 1 0**.
is A=B ?

- In 2's Complement Rep.: **Yes**.
- In 1's Complement Rep.: **Yes**.
- In Signed Magnitude Rep.: **No**.

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Real Numbers

1. $(4.5)_{10} = (100.1)_2$
2. $(8.25)_{10} = (1000.01)_2$
3. $(16.125)_{10} = (10000.001)_2$
4. $(0.875)_{10} = (0.111)_2$
5. $(4.5)_{10} = (1.001)_2 \times 2^2$
6. $(8.25)_{10} = (1.00001)_2 \times 2^3$
7. $(16.125)_{10} = (1.0000001)_2 \times 2^4$
8. $(0.875)_{10} = (1.11)_2 \times 2^{-1}$

Representations 5 to 8 are called Normalized Representations.

Normalized Rep.: $(\pm 1.xxxxxx)_2 \times 2^E$, Where 'E' is a**True Exponent**, 'xxxxxx' is a **Fraction/Mantissa**.

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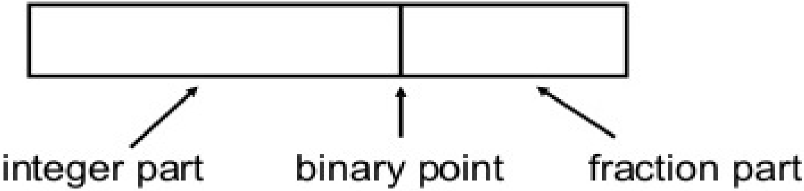
Fixed Point (FP) Representation



Find the decimal equivalent of the following binary numbers. Assume that the binary numbers are represented using FP representation (6,2), i.e., 6 bits for integer part and 2 bits for fractional part.

- $(00101011)_2 =$ **?**
- $(11111011)_2 =$ **?**

Smallest $+ve$ number that can be represented using FP (6,2) rep.: **?**
Biggest $+ve$ number that can be represented using FP (6,2) rep.: **?**

Fixed Point Arithmetic

Consider a 16-bit binary representation, in which least significant 8 bits are used for precision then give the range of values that can be represented using the 16-bit representation.

- $+ve$ Values: 2^{-8} to $2^{7} - 2^{-8}$
- $-ve$ Values: -2^{7} to -2^{-8}
- Zero
- Resolution is **?**

Advantages of FP Arithmetic:

- Easy to implement and occupies less space.
- If performance is important than precision.
- Once can choose a trade off between range and precision.

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IEEE 754 Representation of Real Numbers

- IEEE 754 format for Real Numbers.

Sign	Biased Exponent	Mantissa/Fraction

Single Precision N=32	1 bit	8 bits	23 bits	Bias Value:+127
Double Precision N=64	1 bit	11 bits	52 bits	Bias Value :+1023

Biased Exponent=True Exponent + Bias Value

- Rep. of $(4.5)_{10}$ using Single Precision
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IEEE 754

- Biased Exponent=True Exponent + Bias Value, where $1 \leq$ Biased Exponent $\leq (2^{Length\ of\ Biased\ Exponent} - 2)$.
- Single Precision (N=32), $1 \leq$ Biased Exponent ≤ 254.
- Biased Exponent $= 0$,
  - Mantissa $= \pm 0$, then Value is ± 0.
  - Mantissa $\neq 0$, then Value is **not a normalized number**.
- Biased Exponent $= 255$,
  - Mantissa $= \pm 0$, then Value is $\pm\infty$.
  - Mantissa $\neq 0$, then Value is **NAN**.
- Range of positive values: $[1.0 \times 2^{-126}, (2 - 2^{-23}) \times 2^{127}]$
- Range of negative values: $[-(2 - 2^{-23}) \times 2^{127}, -1.0 \times 2^{-126}]$
- Single Precision Number Resolution: $2^{-23} \times 2^{TrueExponent}$

BCD Representation

BCD: Binary Coded Decimal

It uses a 4-bit binary number to represent each decimal digit.

Decimal Number	BCD Rep.
0	0000
1	0001
2	0010
3	0011
9	1001
10	0001 0000
25	0010 0101
99	1001 1001

Table 1: BCD equivalent of a decimal number.

All the best ☺