Hands on - 3

Nagandla Sai Sreeja
1002269499

```
function x = f(n)
    x = 1;
  for i = 1:n
    for j = 1:n
      x = x+1;
```

1) find the runtime of the algorithm mathematically.

Initialize x=1 (constant time, O(1))

Executes a loop:

*) the outer loop runs from i=1 to n, and it iterates n times for each

*) the inner loop runs from j=1 to n, and it iterates n times for each outer loop iteration

$$T(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} 1$$

$$= \sum_{i=1}^{n} n \Rightarrow n * n$$

$$\Rightarrow n^2$$

The total Runtime complexity $T(n) = O(n^2)$

3) the upper and lower bounds on curve, from this specify. big-O, big-Omega, big-theta is

Big-O (upper Bound) $\Rightarrow$
the f(n) doesnot grow faster than $O(n^2)$.

Big-Omega → Lower bound:-

The $f(n)$ does not grow slower than $\Omega(n^2)$.

Big-Theta → tight Bound:-

$\Theta(n^2)$ this function grows exactly at the rate of n^2.

4) If I modified the function to be

```
x = f(n)
x = 1;
y = 1;
for i = 1:n
    for j = 1:n
    x = x + 1;
    y = y + i + j;
```

will this increase how long it takes the algorithm to run?

By adding an extra operation $y = i + j$; means each iteration takes slightly longer in real execution time

$\therefore O(n^2)$ is the time complexity.

5) It increase execution time but does not effect and change the overall $O(n^2)$ time complexity.