## Design Patterns

## Singletons, Pools and Factories

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https://github.com/NagyAttila/DesignPatterns

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## Content

- Object Oriented Design Principles
  - Singleton
  - Meyers Singleton
  - Phoenix Singleton

- Object Pool
- Connection Pool
- Thread Pool

- Factory Method
- Abstract Factory
- Object Factory

## Open Close Principle

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## Interface Segregation Principle

Fat interface avoidance.

Clients should not be forced to depend on interface that they don't use.

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## Single Responsibility Principle

One class should be responsible for one thing.

## Singleton

### Purpose:

- class has at most one instance,
- that instance is accessible in a global scope.

### Application Areas:

- logging,
- Thread and Connection Pools,
- Factories
- system clock,
- etc.

## Singleton - How?

private constructors

•static \_pInstance\_ and \_Instance\_

```
1 class Singleton {
 public:
 static Singleton& Instance() {
    if (nullptr == pInstance) {
4
      pInstance = new Singleton;
6
7
    return *pInstance;
8
9
 private:
11 Singleton () = default;
12 Singleton (const Singleton&) = delete;
13 Singleton operator= (const Singleton&) = delete;
14 \sim \text{Singleton} () = delete;
16 static Singleton* pInstance;
17 // More Functions and Data
19 Singleton* Singleton::pInstance = nullptr;
```

# Singleton - How?

```
private constructors
                         •static _pInstance_ and _Instance_
 class Singleton {
 public:
 static Singlet
   if (null:
                    Destructor was
                      never called!
 priva
                    (memory leak)
 Single
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13 Singleton
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# Meyers Singleton

•local static \_instance\_

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⇒ private destructor is called at process termination



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#### How?

 Two Singleton classes, Logger and Keyboard, in different compilation units



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## Why?

order of deallocation of static objects is not deterministic



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#### Solution?

- on-demand reallocation of Singleton after destruction
  - ⇒ Phoenix Singleton



## Phoenix Singleton

### Combination of the 3 approaches:

- Simple: \_pInstance\_ pointer
- Meyers: static life-time
- Dead Reference: detection

### Plus:

- on-demand reallocation
- destructor called using \_atexit\_
- multiple time "reborn"

```
class Singleton {
  public:
  static Singleton& Instance() {
    if (destroyed) {
      new(pInstance) Singleton;
      atexit(KillSingleton);
    }
    if (nullptr == pInstance) {
      static Singleton instance;
      pInstance = &instance;
    }
    return *pInstance;
}

static void KillSingleton(void) {
    pInstance ->~Singleton();
}
```

```
private:

Singleton () {
    destroyed = false;
}
Singleton (const Singleton&) = delete;
Singleton operator= (const Singleton&) =
    delete;

Singleton () {
    destroyed = true;
    pInstance = nullptr;
}
Static bool destroyed;
static Singleton* pInstance;
};
bool Singleton::destroyed = false;
Singleton* Singleton::pInstance = nullptr;
}
```

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• functions registered on stack



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#### Solution?



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- registering from a registered function
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#### Solution?

- Read the Manual!
- Use the newest compilers!



- general lazy initialization problem
- race condition

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### Problem:

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### Solution:

A local static instance (Meyers approach)

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#### Problem:

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### Solution:

- A local static instance (Meyers approach)
- **B** locking

```
static Singleton& Instance() {
   Lock guard(mutex);
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C double-check locking

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## **Notes**

## Why not global?

- uniqueness
- lack of laziness
- pollute global scope

- not always possible
  - dependency
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### Possible extensions and alternative solutions:

- longevity control for dead reference problem
- registry high number of singletons
- inheritance

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## Final thought:

should be used sparingly

- if needs a lot:
  - ⇒ use registry, or
  - ⇒ change design

# Object Pool

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- report error
- increase pool size
- blocking request

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# Similar patterns

- Connection Pool
- Thread Pool

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### Used as

Singleton



### Connection Pool

- cache of DB connections
- creates new connection if empty pool (2nd approach)

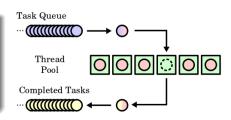


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#### Thread Pool

- asynchronous task processing
- more tasks than threads
- fix or dynamic number of threads



#### Threads

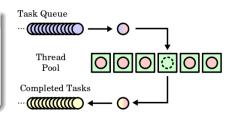
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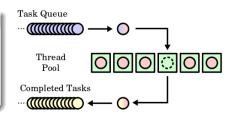
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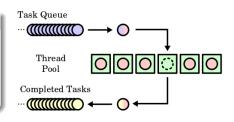
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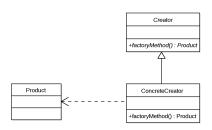
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# Factory Method

# A.k.a: Virtual Constructor

- uses an abstract class for interface
- instantiation determined by subclasses

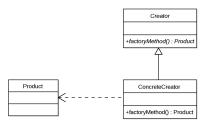


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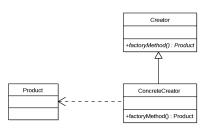
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#### Features:

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# Application areas:

- unit testing
- Abstract Factory



# Features:

• create product from families of products



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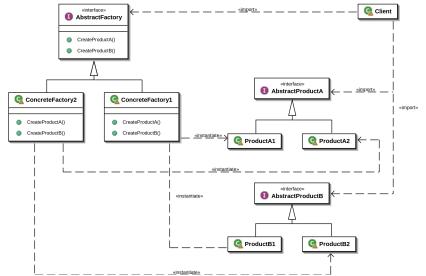


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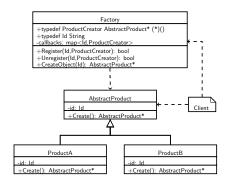
### Drawbacks:

- adding new product-family is cumbersome
  - ⇒ abstract factory class must change
    - $\Rightarrow$  concrete factory classes must follow the change



# **Object Factory**

- each product has to register
- uses the priority inversion principle
- one family of product
- create callbacks stored in map



# Resources

- Design Patterns: Elements of Reusable Object-Oriented Software "Gang of Four"
- Modern C++ Design: Generic Programming and Design Patterns Applied Andrei Alexandrescu
- Head First: Design Patterns
   Eric Freeman & Elisabeth Freeman



# Questions?



# Thank You!

