

Homework 3

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Problem 1

a/

The register `%rdx` is increased by 8 every step in the inner loop. In C, the size of the data type `long` is 8 bytes. As a result, at each step of the inner loop, the register percent `%rdx` is increased by 8 to move to the next element of `A`. Therefore, `%rdx` holds the pointer to element `A[i][j]`.

b/

In Assembly, there are instructions `"movq (%rdx), %rcx"` and `"movq %rcx, (%rax)"`. These instructions swap the address of the register `%rdx` and `%rax`. The register `%rdx` holds the address of `A[i][j]`, so the register `%rax` must hold the address of `A[j][i]`.

c/

In Assembly, in line 7, there is an instruction `"addq $120, %rax"`. This instruction means changing the address of `A[j][i]` hold by the register `%rax` to the address of `A[j+1][i]`. So, the total size of `M` elements data types is 120. Hence $M = 120/8 = 15$.

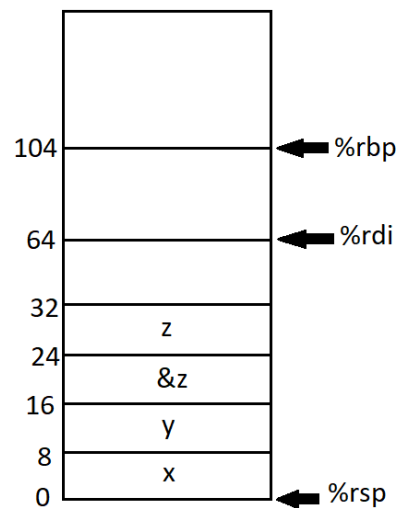
Problem 2

The instructions `"addq $1, %rdx"` and `"cmpq %rdi, %rdx"` are found on Assembly lines 13 and 15. In C, these instructions are equivalent to `"i++"` and `"i < NR(n)"`. As a result, the register `%rdi` keeps the result of `NR(n)`. There are instructions `"leaq (%rdi, %rdi, 2), %rax"` and `"movq %rax, %rdi"` on lines 3 and 4. Hence, $NR(n) = 3 \times \%rdi = 3n$.

The register `%rcx`, which stores the address of element `A[i][j]`, increases `%r8` bytes to move to `A[i+1][j]`, hence the register `%r8` stores the size of `NC(n)` elements. The value in the register `%r8` is shifted left by 3 or multiplied by 8 in Assembly line 7. Therefore, the value saved in the register `%r8` in line 2 is the definition of $NC(n) = \%rdi * 4 + 1 = 4n + 1$

Problem 3

a/



b/

In C, *eval* passes the function *process* with a *StrA* *s*, while in Assembly, there is an instruction "`leaq 64(%rsp),%rdi`". In Assembly, the register `%rdi` is used to pass the first parameter. Then, we know that `%rsp+64` is the pointer to the begin of the first parameter of function *process*. Therefore, *eval* passes `%rsp+64` to function *process*.

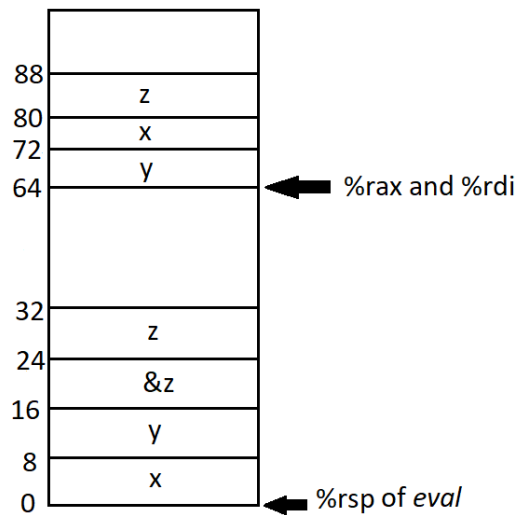
c/

Since *eval* passes `%rsp+64` to function *process*, the Assembly code adds the register `%rsp` with an offset to access the elements in structure argument *s*.

d/

Since *eval* passes `%rsp+64` to the register `%rdi`, the function *process* calculate `%rdi + offset` to access and set value of the elements of structure *r* in *eval*.

e/



f/

In Assembly, the caller allocates stack space and passes the address to the callee, which then stores data in the space and returns the address to the caller.

Problem 4

Cache	m	C	B	E	S	t	s	b
1.	32	1024	4	4	64	24	6	2
2.	32	1024	4	256	1	30	0	2
3.	32	1024	8	1	128	22	7	3
4.	32	1024	8	128	1	29	0	3
5.	32	1024	32	1	32	22	5	5
6.	32	1024	32	4	8	24	3	5

Problem 5

Cache	m	C	B	E	S	t	s	b
1.	32	2048	8	1	256	21	8	3
2.	32	2048	4	4	128	23	7	2
3.	32	1024	2	8	64	25	6	1
4.	32	1024	32	2	16	23	4	5

Problem 6

a/

CT	CT	CT	CT	CT	CT	CT	CT	CT	CI	CI	CO	CO
11	10	9	8	7	6	5	4	3	2	1	0	

b/

Operation	Address	Hit?	Read value (or unknow)
Read	0x834	Miss	Unknown
Write	0x836	Hit	
Read	0xFFD	Hit	0xC0

Problem 7

The size of (C) of this cache in bytes is 128

CT	CT	CT	CT	CT	CT	CT	CT	CI	CI	CI	CO	CO
12	11	10	9	8	7	6	5	4	3	2	1	0

Problem 8

Address format:

0	0	1	1	1	0	0	0	1	1	0	1	0
12	11	10	9	8	7	6	5	4	3	2	1	0

Memory reference:

Parameter	Value
Blockoffset (CO)	0x2
Index (CI)	0x6
Cache tag (CT)	0x38
Cache hit? (Y/N)	Y
Cache byte return	0xEB

Problem 9

a/

Address format:

1	0	1	1	0	1	1	1	0	1	0	0	0
12	11	10	9	8	7	6	5	4	3	2	1	0

Memory reference:

Parameter	Value
Blockoffset (CO)	0x0
Index (CI)	0x2
Cache tag (CT)	0xB7
Cache hit? (Y/N)	N
Cache byte return	—

b/

- 0x1140, 0x1141, 0x1142, 0x1143
- 0x16C8, 0x16C9, 0x16CA, 0x16CB
- 0x180C, 0x180D, 0x180E, 0x180F
- 0x1150, 0x1151, 0x1152, 0x1153
- 0x1798, 0x1799, 0x179A, 0x179B

Problem 10

dst array					src array				
	Col. 0	Col. 1	Col. 2	Col. 3		Col. 0	Col. 1	Col. 2	Col. 3
Row 0	m	m	m	m	Row 0 b	m	h	h	h
Row 1	m	m	m	m	Row 1	m	h	h	h
Row 2	m	m	m	m	Row 2	m	h	h	h
Row 3	m	m	m	m	Row 3	m	h	h	h