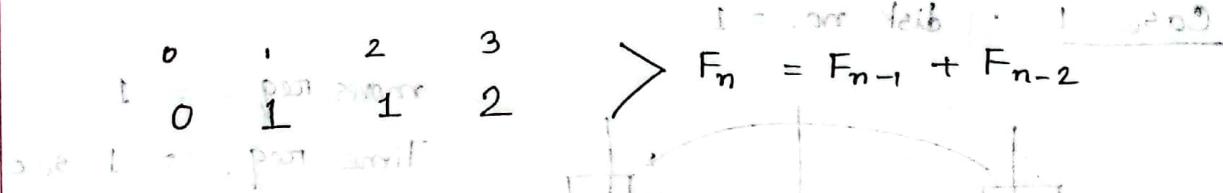


### Recurrent Problems

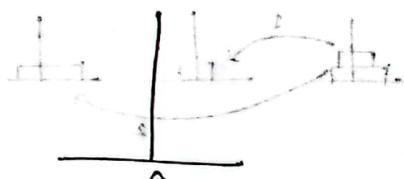
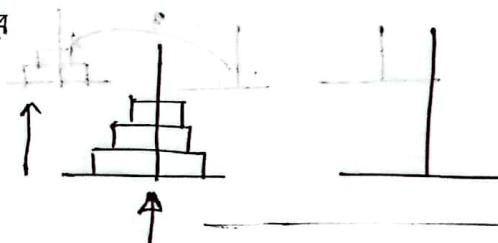
: ब्रह्मपति विद्या विद्यालय  
 → 3 niddle of diamond  
 → gold plate

→ Fibonacci series



Tower of Hanoi : Lucas, 'Tower of Brahma' का उपयोग करके बनाए गए [disk = 8]

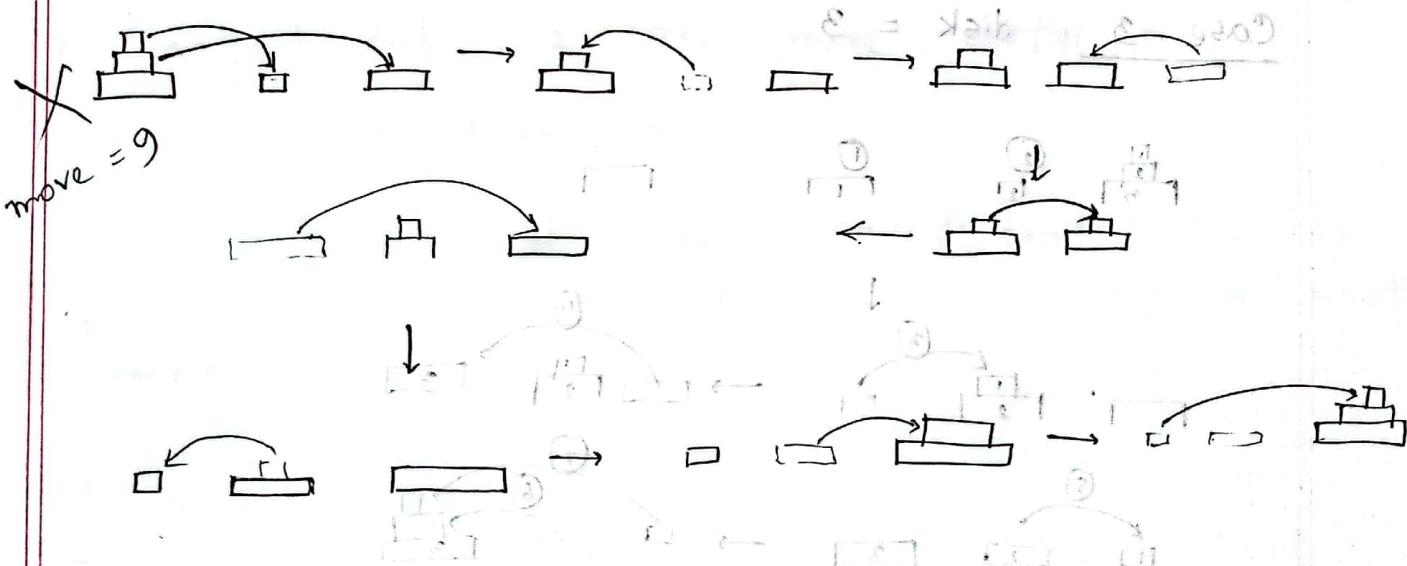
देख  
disk एवं  
उपकरण  
लाइट



- a) at a time एक ही  
 b) तूट रखे उपकरण वाले

बाथा यावे ना

move = 2<sup>n</sup> - 1

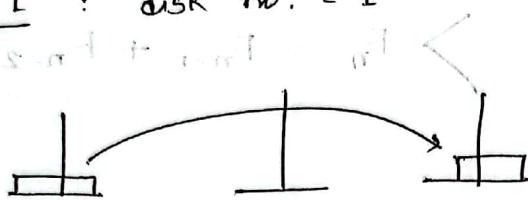


ECS - 309  
Algorithm of Tower of Hanoi

\* move required  
\* how much time required

Study block :-  
\* When disk no. = 0  $\rightarrow$  moves req. = 0 is required.

Case 1 : disk no. = 1

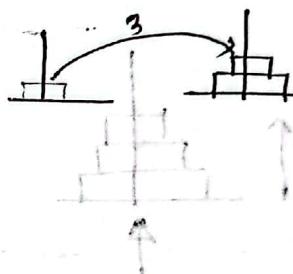
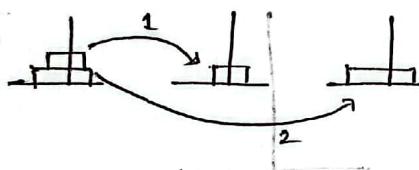


moves req. = 1

Time req. = 1 sec.

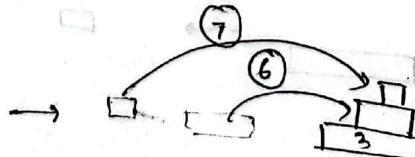
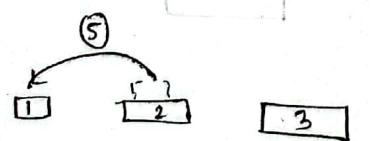
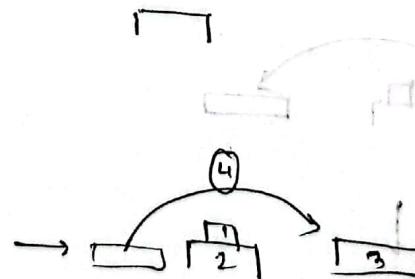
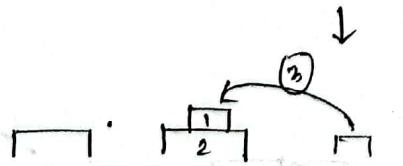
For even no. handling [to move] each : ionoff to mount  
for odd no. to go [2 - 4 sec] RST block

Case 2 : disk = 2



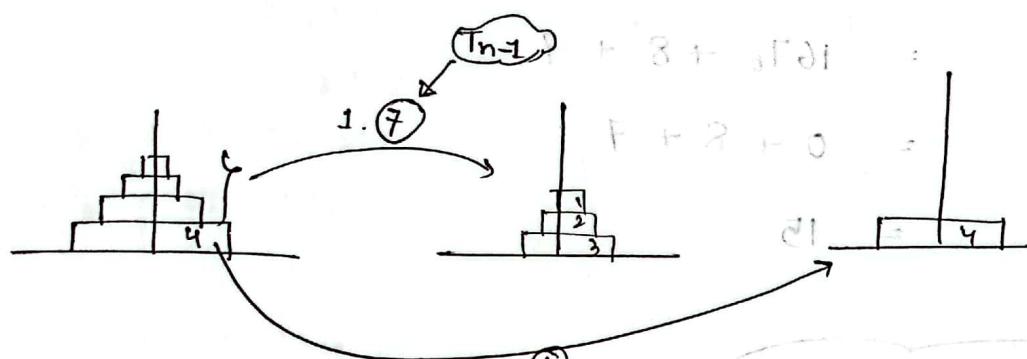
moves required = 3

Case 3 : disk = 3



Case 4 : disk = 4

moves = 15



$$L + \text{moves} \geq T$$
$$7 + 7 + 1 = 15 \text{ moves}$$



\* largest disk , 1 move

\* n no. of disks , 3 tower

$T_n$  = no. of moves required that'll transfer n disks from one tower/peg to another  
no. of moves (minimum)

$$T_0 = 0$$

$$T_4 = 7 + 1 + 7 = 15$$

$$T_n = 2T_{n-1} + 1$$

$$\begin{aligned} T_4 &= 2 \times T_3 + 1 \\ &= 2 \times (2T_2 + 1) + 1 \\ &= 4T_2 + 3 \\ &= 4(2T_1 + 1) + 3 \end{aligned}$$

$$P = 4 \times 16 + 1 = P = 65$$

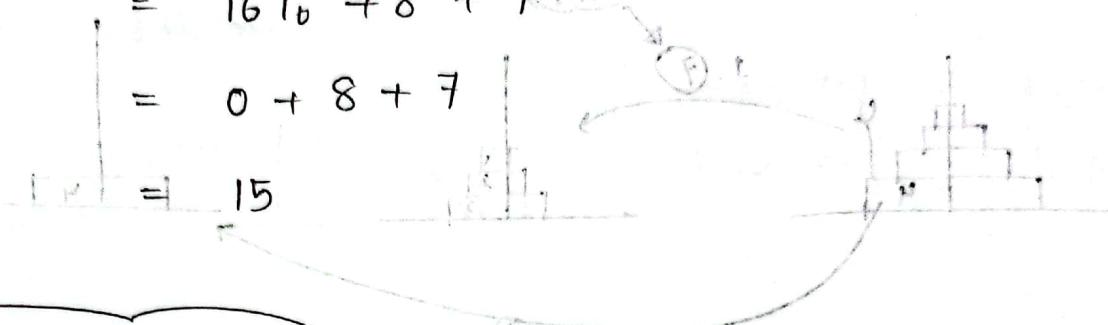
$$= 8T_1 + 7$$

$$= 8(2T_0 + 1) + 7 \quad \text{dt} = 2^{\log n}$$

$$= 16T_0 + 8 + 7$$

$$= 0 + 8 + 7$$

$$T_0 = 15$$



$$T_n \geq 2T_{n-1} + 1 \rightarrow T_{n(\min)} = 2T_{n-1} + 1$$

# Prove that, for  $n$  disks  $2T_{n-1} + 1$  more are sufficient and necessary for making to the destination peg.

From state 1, 42ib feasible

state E, 24ib to ...

statement if both beginning from state 1, no of moves are same then

(minimum)

$$T_n = T_{n-1} + 2T_{n-2} + \dots + 2T_1 + 1$$

$$T_n = 2(T_{n-1} + T_{n-2} + \dots + T_1) + 1$$

$$T_n = 2T_{n-1} + 1 + 2(T_{n-2} + T_{n-3} + \dots + T_1)$$

$$T_n = 2T_{n-1} + 1 + 2(T_{n-2} + T_{n-3} + \dots + T_1)$$

$$T_n = 2T_{n-1} + 1 + 2(T_{n-2} + T_{n-3} + \dots + T_1)$$

Closed Form :

$$n=0 \rightarrow T_0 = 0$$

$$n=1 \rightarrow T_1 = 1 (2T_0 + 1)$$

$$n=2 \rightarrow T_2 = 3 (2T_1 + 1)$$

$$n=3 \rightarrow T_3 = 7 (2T_2 + 1)$$

$$n=4 \rightarrow T_4 = 15 (2T_3 + 1)$$

$$n=5 \rightarrow T_5 = 31 (2T_4 + 1)$$

$$n=6 \rightarrow T_6 = 63 (2T_5 + 1)$$

$$\boxed{T_n = 2^n - 1} \rightarrow \text{Soln of closed form}$$

$$2^n - 1$$

$$\text{Base} \rightarrow T_0 = 2^0 - 1 = 0$$

$$\text{hypothesis} \rightarrow T_{n-1} = 2^{n-1} - 1$$

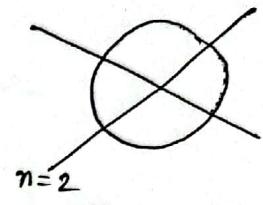
$$\text{Induction} \rightarrow T_n = \cancel{2^n} \cancel{+ \cancel{2^{n-1}} \cancel{+ \dots + 2^1}} + 2T_{n-1} + 1$$

$$= 2(2^{n-1} - 1) + 1$$

$$= 2^n - 1$$

Geometric problem :

How many slices of pizza can a person obtain by making  $n$  straight cut with a pizza knife?

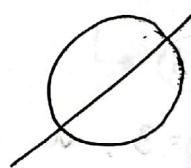


What's the maximum number ( $L_n$ ) of regions defined by  $n$  lines in the plane

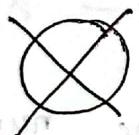
$$n=0, L_n = 1$$



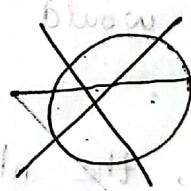
$$n=1, L_n = 2$$



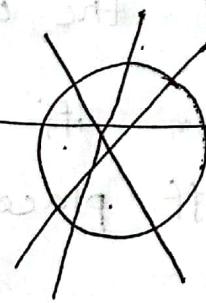
$$n=2, L_n = 4$$

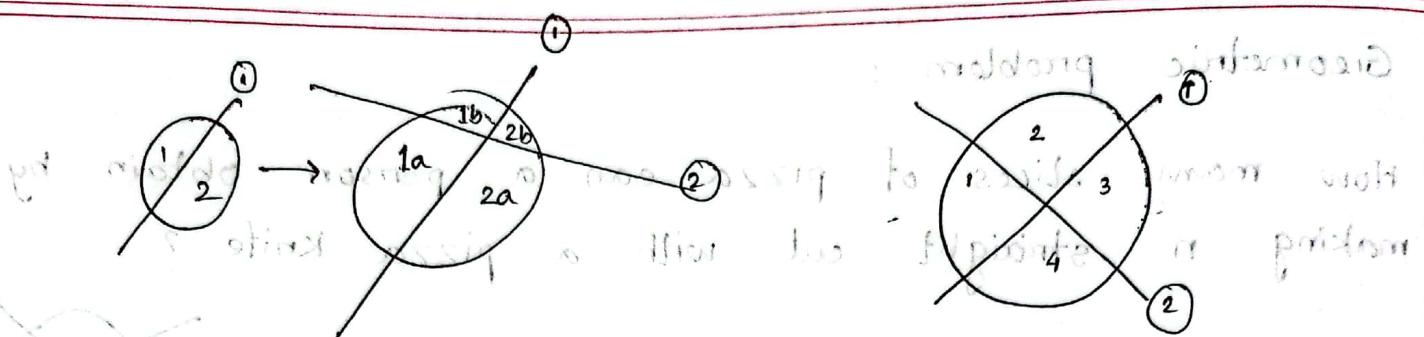


$$n=3, L_n = 7$$



$$n=4, L_n = 11$$



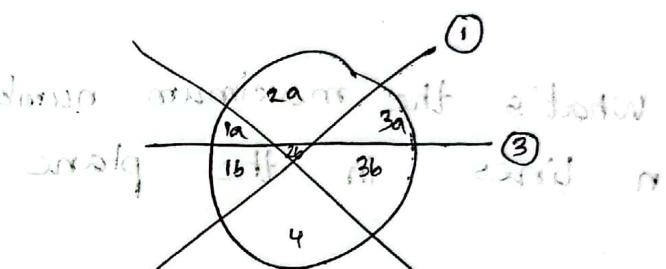


→ adding a line to two regions

→ passing to next region with  
line to 1st place in problem



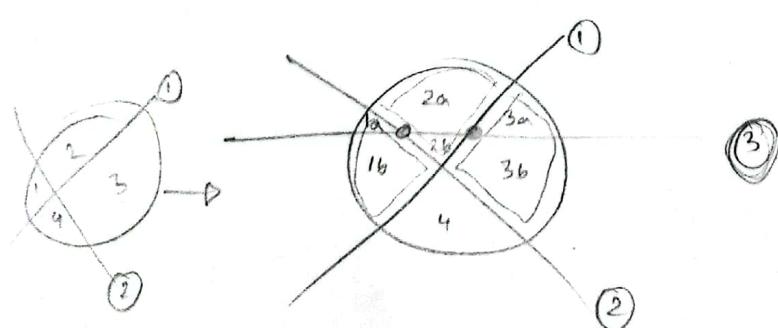
→ initial answer is (n!) regions when all lines



1. For  $n=3$ , a line can split at most three region

2.  $n$ th line would split an old regions at most  
 $\leq n-1$ ,  $O = n$

3. For  $n > 0$ , the  $n$ th line increases the number of regions by  $k$  iff it splits  $k$  of the old regions and it splits  $k$  old region iff it hits the previous lines in  $(k-1)$  diff places.



3 Br region too  
2 Bi point → touch  
→ split 316,

4. Two lines can intersect in at most one point.  
 Therefore the new line can intersect  $(n-1)$  old lines in at most  $(n-1)$  places.

\*  $L_n =$  maximum regions by  $n$  straight line

$$n=0 \rightarrow L_0 = 1$$

$$n=1 \rightarrow L_1 = 1 + 1 = 2$$

$$n=2 \rightarrow L_2 = 2 + 2 = 4$$

$$n=3 \rightarrow L_3 = 4 + 3 = 7$$

$$n=4 \rightarrow L_4 = 7 + 4 = 11$$

$$L_n = L_{n-1} + n$$

↳ Recurrence sol<sup>n</sup>

$$\rightarrow L_0 = 1$$

$$\rightarrow L_n = L_{n-1} + n$$

(for  $n > 0$ )

Closed Form:

An expression for a quantity  $f(n)$  is in closed form if we can compute it using at most a fixed number of well-known standard operations, independent of  $n$ .

Let's find the form of  $L_n$  in terms of  $n$  and  $L_{n-1}$ .

$$\begin{aligned} L_n &= L_{n-1} + n \quad \text{from } n \text{ is added at } n^{\text{th}} \text{ step} \\ &= L_{n-2} + (n-1) + n \quad \text{from } (n-1) \text{ from } n^{\text{th}} \text{ step} \\ &= L_{n-3} + (n-2) + (n-1) + n \end{aligned}$$

$$\text{and } L_4 = L_3 + (n-3) + (n-2) + (n-1) + n = \boxed{L_3 + \frac{n(n+1)}{2}}$$

When  $n=4$ ,  $L_4 = L_0 + 1 + 2 + 3 + 4 = \boxed{10}$

For  $n \rightarrow L_n = L_0 + [1 + 2 + 3 + \dots + (n-3) + (n-2) + (n-1) + n]$

$$= L_0 + \frac{n(n+1)}{2} \quad \boxed{L_n = n + \frac{n(n+1)}{2}}$$

So, total time taken for lifting up a set of stairs is  $n + \frac{n(n+1)}{2}$ .  
To reach the top of the stairs, we have to take  $n$  steps.  
So, total time taken for lifting up a set of stairs is  $n + \frac{n(n+1)}{2}$ .

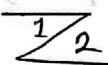
Self. study :  Straight line

If line isn't straight, line is bend



bend line  $\rightarrow$  max no. of regions determine,  
if 'n' no. of  $\times \rightarrow$  line is cut.

Zigzag line



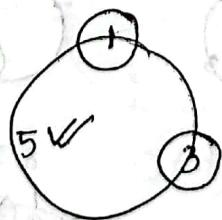
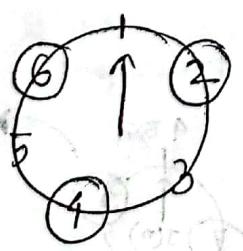
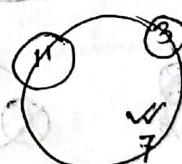
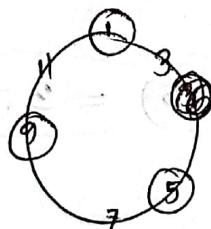
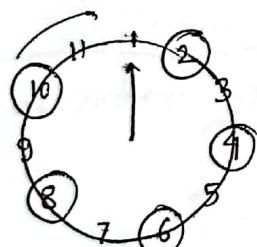
$$L = (S)T$$

$$C = (S)T$$

$$C = T$$

### JOSEPHUS PROBLEM

$$L = (A)T$$

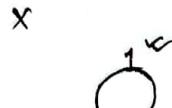


$$0 \leq \lambda \leq \frac{m}{3} - 1$$

$n$	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
$J(n)$	1	1	3	1	3	5	7	1	3	5	7	9	11	13	15	1	3	5	7	9

0 to 20 people in circle board

$$n = 0$$



$$n = 1$$

1 survive विजेता  $\rightarrow J(1) = 1$

$$n = 2$$



$$J(2) = 1$$

$$n = 3$$



$$J(3) = 3$$

$$n = 4$$



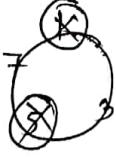
$$J(4) = 1$$

$$n = 5$$



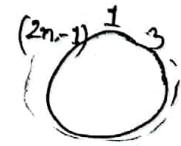
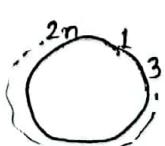
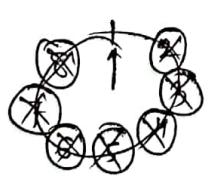
$$J(5) = 3$$

$$n = 7$$



$$J(7) = 7$$

$$n = 8$$



Visualization  
when  $n$  is  
even

2nd round

$$\begin{aligned} 1 &\rightarrow 1 \quad (2^{n-1}) \\ 3 &\rightarrow 2 \quad (2^{n-2}) \\ 5 &\rightarrow 3 \quad (2^{n-3}) \\ 7 &\rightarrow 4 \quad (2^{n-4}) \end{aligned}$$

$$J(2n) = 2^n J(n) - 1$$

## Observations :

1.  $J(n)$  is always odd. (at least for  $n=1, 2, 3, 4, 5, 6, 7$ )
2. Even positions are eliminated at first round.

Note:

•  $J(1) = 1$ , i.e., half of total odd numbers

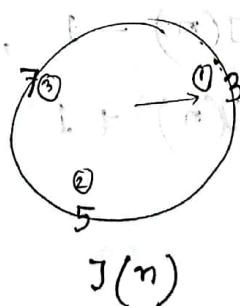
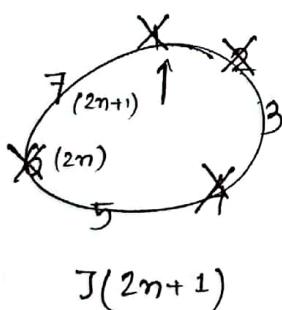
•  $J(2n) = 2^n J(n) - 1$

$\therefore J(10) = 2^5 J(5) - 1 = 5$

$J(20) = 2^5 J(10) - 1 = 2^5 \cdot 5 - 1 = 9$

$J(40) = 2^5 J(10) - 1 = 9^5 \cdot 2 - 1 = 17$

Visualization when  $n$  is odd



$$\begin{aligned} & J(n) \\ & 3 \rightarrow 1 \leftarrow (2^n n + 1) \\ & 5 \rightarrow 2 \leftarrow \\ & 7 \rightarrow 3 \leftarrow \\ & \boxed{J(2n+1) = 2^n J(n) + 1} \end{aligned}$$

•  $J(2n+1) = 2^n J(n) + 1$

$$J(18) = 2 J(9) + 1 = 2(2 J(4) + 1) + 1 = 2(2^3 \cdot 1 + 1) + 1 = 5$$

3. For even no. of people  
 Originally we have  $n$  no. of people,  
 after one round, it seems we have  $n$  no.  
 of people except each person's position  
 number has been doubled and decreased by 1.

4. For odd no,  $2n+1$   
 $\dots n^2 - 1 = (n^2)C + S = (n)C$   
 $\dots$  and increased by 1.  
 $C = 1 - 2^{n+1} + 1 = (n)C + S = (n+1)C$

### Recurrence problem:

$$J(1) = 1$$

$$J(2n) = 2^n J(n) - 1, n \geq 1$$

$$J(2n+1) = 2^n J(n) + 1, n \geq 1$$

(n)C

(n+1)C

$$1 + (n)C + (n+1)C = (n+1)C$$

$$1 + (n+1)C = 1 + (n+1)C + (n+1)C - (n+1)C = (n+1)C$$

Closed Form

The condition will be

1. When  $n$  is power of 2, i.e.  $J(n) = 1$  for every

2. When  $J(n)$  is 1, later  $J(n+1), J(n+2)$  are of consecutive odd numbers?  $n = 2^m$  is given

$$J(2^m) + J(2^m + 1)$$

$$J(n) = 1$$

$$J(2^m) = 1 \quad \text{here } 2^m = 2^0 \text{ as } (2 \times 0 + 1) = 1$$

$$J(2^m + 1) = 3 \quad (2 \times 1 + 1)$$

$$J(2^m + 3) = 5 \quad (2 \times 2 + 1) \quad \dots \quad (2 \times 2 + 1) = (2 \times 2)^2$$

$$J(2^m + 5) = 7 \quad (2 \times 3 + 1) \quad \dots \quad (2 \times 3 + 1) = (2 \times 3)^2$$

$$\text{when, } n = 2^m + l \quad (l \in \mathbb{N})$$

$$J(n) = J(2^m + l) = (2 \times l + 1)$$

$$J(n) = J(2^m + 1) = 2l + 1$$

$$\text{here, } 2^m \leq n < 2^{m+1}$$

$$0 \leq l \leq 2^m - 1$$

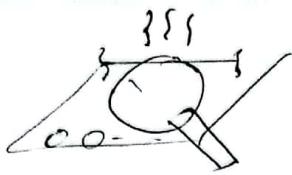
$$J(100) = J(2^6 + 36) = 2 \times 36 + 1 = 73$$

$$(2 \times 6 + 1) \in (2 \times 6)^2$$

$$1 + (2 \times 6)^2$$

$$1 + (2 \times 6)^2 \in \mathbb{Z}$$

$$1 + [1 + (2 \times 6)^2] \in \mathbb{Z} = \{1, 37, 73\}$$



Using the recurrence sol<sup>n</sup>, prove by induction that  $J(2^m + l) = 2l + 1$ .

Basis:  $m=0, l=0$

$$J(2^0 + 0) = J(1) = 1$$

Induction:  $l$  is even/odd

If  $l+m > 0$  and  $2^m + l = 2n$

$$J(2n) = 2J(n) - 1$$

$$J(2^m + l) = 2J\left(\frac{2^m}{2} + \frac{l}{2}\right) - 1$$

$$= 2J\left(2^{m-1} + \frac{l}{2}\right) - 1$$

$$= 2\left[2 \cdot \frac{l}{2} + 1\right] - 1 = l + 2 - 1$$

$$= 2l + 2 - 1$$

$$= 2l + 1$$

when  $l$  is odd,  $2^m + l = 2n + 1 \rightarrow n = \frac{2^m + l - 1}{2}$

$$J(2^m + l) = J(2n + 1)$$

$$= 2J(n) + 1$$

$$= 2J\left(\frac{2^m + l - 1}{2}\right) + 1$$

$$= 2\left[2 \cdot \frac{l-1}{2} + 1\right] + 1 = 2l + 1$$

The Josephus problem  $\rightarrow J(n) = J(2^m + l) \ (\Leftarrow 2l + 1)$

Examining the "closed form"

\* power of  $2^m$  played an important role.

if  $n$  can be represented in binary  $\rightarrow n \& J(n)$

$n$ 's binary?  $n = 2^m + l$

$$n = (b_m b_{m-1} \dots b_1 b_0)_2$$

$\hookrightarrow$  each  $b_i = \begin{cases} 0 \\ 1 \end{cases}$

loading bit,  $b_m = 1$ , so  $n = (1 b_{m-1} b_{m-2} \dots b_1 b_0)_2$

$$l = (0 b_{m-1} b_{m-2} \dots b_1 b_0)_2$$

$$2l = (b_{m-1} b_{m-2} \dots b_1 b_0 0)_2$$

$$2l + 1 = (b_{m-1} b_{m-2} \dots b_1 b_0 1)_2$$

$$\therefore J(n) = (b_{m-1} b_{m-2} \dots b_1 b_0 b_m)_2$$

$$\boxed{J((b_m b_{m-1} \dots b_1 b_0)_2) = (b_{m-1} b_{m-2} \dots b_1 b_0 b_m)_2}$$

$\hookrightarrow$  1 bit cyclic left shift

If,  $n = 10$ , then  $2l + 1 = (((01)_2)_2)_2 = 5$

$$J((10)_2) = 0101_2 = 5$$

If,  $n = 100$ ,

$$J((1100100)_2) = 1001001_2 = 73$$

$2 \rightarrow 10$
↓ double
$4 \rightarrow 100$
↓ double
$8 \rightarrow 1000$

#

$$J(n) \left( \approx n^m \right) \approx (n)^m \cdot J(n) \approx \frac{n^m}{2}$$

$n$	1	3	7	15	... $\Rightarrow$ অবশ্যই $m$ পরিমাণ,
$J(n)$	1	3	7	15	... এখন কোথায় হবে? কোথায় হবে?

- repeatedly Josephus problem apply :

$$J(J(J(13)))$$

$$13 = 1101_2$$

$$J(1101_2) = 1011_2 \quad \text{এবং } 1011_2 = 11$$

$$J(1011_2) = 0111_2 = 7 \quad \text{একটি fixed point}$$

$$J(0111_2) = 1111_2 \quad (\text{এবং } 1111_2 = 7) \quad \text{এ চলে যাব।}$$

when  $J(n)$  will be  $n$ ?  $\Rightarrow$   $(\text{odd} + \text{odd}) = 12$

↳ binary representation এ যখন সঠিগুলো bit

$$\rightarrow 1 \quad (\text{odd} + \text{odd} + \text{odd} + \text{odd}) = (n)_{\text{E}}$$

13 → binary representation  $\Rightarrow$  3 টি 1 ( $1101$ )  
 $\rightarrow$  fix রয়ে  $\rightarrow 2^3 - 1 = 7$  এ,

$n$  প্রতি  
সঠিগুলো  
'1' bit  
লাগে

fix রয়ে

$\frac{2^x - 1}{2^x}$

$x = \text{no of '1'}$

$\therefore J(J(J(J(13)))) = 7$  (at least 2 টি জিনালে fix রয়ে)  
 $\Rightarrow$  repeat রয়ে 3 no রয়ে।

$$J(1101_2) = ((1101)_2)_E$$

$$J(J(J(J(J(2))))))$$

$$= 10101 \rightarrow \underbrace{3}_{(2^3 - 1)} \cdot 1 \xrightarrow{(2^3 - 1) + 1} 2^3 - 1 = 7$$

$$1 \rightarrow 01011$$

$$2 \rightarrow 0111 = 7$$

$$3 \rightarrow 111 = 7$$

$$4 \rightarrow 1111 = 7$$

$$J(J(J(J(J(J(10101110)))))))$$

$$\xrightarrow{2^6 - 1 = 63}$$

$$1 \rightarrow 01011101$$

$$2 \rightarrow 01111011$$

$$3 \rightarrow 111011$$

$$4 \rightarrow 110111$$

$$5 \rightarrow 101111$$

$$6 \rightarrow 111111$$

When  $J(n)$  will be  $\frac{n}{2}$ ?

$$J(n) = 2l + 1 = \frac{2^m + l}{2} \quad \text{longing } n = 2^m + l$$

$$\Rightarrow 4l + 2 = 2^m + 1$$

$$\Rightarrow 3l = 2^m - 2$$

$$\Rightarrow l = \frac{1}{3}(2^m - 2)$$

binary no. (or) half  $\rightarrow$  right shift

$$\begin{array}{rcl} n & \xrightarrow{\quad J(n) \quad} & \\ (2) 10 & \rightarrow & 01(1) \\ (10) 1010 & \rightarrow & 0101(5) \end{array}$$

$$n \text{ मध्यन } \underbrace{170}_{101010} \rightarrow J(n) = \frac{170}{2} = 85$$

\* The pattern is <sup>presence of</sup> consecutive '10' in in binary representation.

Generalization:

$$\left. \begin{array}{l} J(1) = 1 \\ J(2n) = 2 J(n) - 1 \\ J(2n+1) = 2 J(n) + 1 \end{array} \right\} \begin{array}{l} \text{constant values are} \\ \text{represented like } \alpha, \beta, \gamma. \end{array}$$

$$\therefore f(1) = \alpha$$

$$f(2n) = 2f(n) + \beta \quad (\text{for } n \geq 1)$$

$$f(2n+1) = 2f(n) + \gamma \quad (\text{for } n \geq 1)$$

According to the original eqn,

$$\alpha = 1$$

$$\beta = -1$$

$$\gamma = 1$$

## General method

$n$	$f(n)$
1	$\alpha$
2	$2\alpha + \beta$
3	$2\alpha + 0\beta + \gamma$
4	$4\alpha + 3\beta + 0\gamma$
5	$4\alpha + 2\beta + \gamma$
6	$4\alpha + 1\beta + 2\gamma$
7	$4\alpha + 0\beta + 3\gamma$
8	$8\alpha + 7\beta + \gamma$
9	$8\alpha + 6\beta + 1\gamma$
10	$8\alpha + 5\beta + 2\gamma$

↓      ↓      ↓      ↓  
 largest power of 2    decrease    increase    simple shifting rotation

$$\begin{aligned}
 f(2) &= 2f(1) + \beta \\
 &= 2\alpha + \beta \\
 f(3) &= 2f(1) + \gamma \\
 &= 2\alpha + \gamma \\
 f(4) &= 2f(2) + \beta \\
 &= 2(2\alpha + \beta) + \beta \\
 &= 4\alpha + 3\beta \\
 f(5) &= 2f(2) + \gamma \\
 &= 2(2\alpha + \beta) + \gamma \\
 &= 4\alpha + 2\beta + \gamma \\
 f(6) &= 2f(3) + \beta \\
 &= 2(2\alpha + \gamma) + \beta \\
 &= 4\alpha + \beta + 2\gamma \\
 f(7) &= 2f(3) + \gamma \\
 &= 2(2\alpha + \gamma) + \gamma \\
 &= 4\alpha + 3\gamma
 \end{aligned}$$

- ①  $\alpha$ 's coefficient is  $n$ 's largest power of 2
- ②  $\beta$ 's " decreasing by 2 down to 0
- ③  $\gamma$ 's " increasing by 1 up from 0

$$f(n) = \underbrace{A(n)\alpha}_{2^m} + \underbrace{B(n)\beta}_{2^m - 1} + \underbrace{C(n)\gamma}_l$$

## Chapter 2 : Sums

Notation

$$\textcircled{1} \quad 1 + 2 + 3 + \dots + n \quad \checkmark$$

$$\textcircled{2} \quad 1 + 2 + \dots + (n-1) + n \quad \checkmark$$

$$1 + (\textcircled{2}) 1 + \dots + n \quad (\text{not really clear})$$

$$\begin{aligned} \textcircled{3} \quad & * a_1 + a_2 + a_3 + \dots + a_n \rightarrow \text{summation of } n \text{-terms} \\ & \downarrow \\ & \text{term} \end{aligned}$$

$$\textcircled{4} \quad \text{Using } \sum : \sum_{k=1}^n a_k \rightarrow \begin{array}{l} \text{"sigma notation"} \\ (a_k \rightarrow \text{summand}) \end{array}$$

$$\begin{array}{l} \rightarrow \text{"delimited form"} \\ \text{Limit } i \text{ to } n \text{ (sum)} \end{array}$$

sigma notation (generalized)

$$\rightarrow \sum_{\substack{1 \leq k \leq n}} a_k \rightarrow \text{generalized sigma form}$$

relation under sigma

$$\textcircled{5} \quad \text{Ex: } \sum_{\substack{1 \leq k \leq 100 \\ k \text{ odd}}} k^2 \rightarrow 1^2 + 3^2 + 5^2 + \dots + 99^2$$

$$\sum_{k=0}^{49} (2k+1)^2$$

- Sum of reciprocal of the prime numbers upto 100 :

$$\sum_{p \text{ prime}} \frac{1}{p} = \frac{1}{2} + \frac{1}{3} + \frac{1}{5} + \frac{1}{7} + \frac{1}{11} + \dots + \frac{1}{23} + \frac{1}{29}$$

OR  $\sum_{p \text{ prime}} \frac{1}{p} = \sum_{n=1}^{\infty} \frac{1}{p_n}$

$$\sum_{\substack{1 \leq n \leq 100 \\ p_n \text{ prime}}} \frac{1}{n}$$

$$\text{OR, } \sum_{k=1}^{\pi(100)} \frac{1}{p_k}$$

$\pi$  grant function  
 $\sum_{k=1}^N \frac{1}{p_k}$  generate prime numbers  
 $p_k \rightarrow k^{\text{th}}$  prime number

$$\text{OR, } \sum_{\substack{p \text{ prime} \\ p \leq N}} \frac{1}{p}$$

$$\sum_{\substack{k=1 \\ (p \text{ upto } N)}}^{\pi(N)} \frac{1}{p}$$

$$\text{OR, } \sum_{k=1}^{\pi(N)} \frac{1}{p_k}$$

$\sum_{k=1}^n a_k \rightarrow$  we want to change limit from  $k$  into  $k+1$ :

$$\rightarrow \sum_{k=1}^{n-1} a_{k+1}$$

$$m^2 + m^2 \sum_{k=1}^{n-1} = m^2$$

$$m^2 + m^2 =$$

$\sum_{k=1}^n a_k \rightarrow$   $k$  into  $k+1$  followed up in case of  $\infty$

$$= \sum_{k=0}^{n-1} a_{k+1}$$

$m^2 + m^2 = m^2$  To write  $m^2 + m^2 = m^2$

Note:

→ problem statement represent → use delimited.

→ To manipulate sums → use relation under the  $\sum$ .

$$\rightarrow \sum_{k=1}^n a_k (k-1)$$

↓  
 1 first term → zero sum is not harmful; we don't have to handle it.

$$= \sum_{k=2}^n a_k (k-1)$$

\* Kenneth Iverson's conversion to KIC

$$[\text{p prime}] = \begin{cases} 1 & \text{if } p \text{ prime} \\ 0 & \text{if } p \text{ not prime} \end{cases}$$

↑  
p prime ता हले !  
p prime नहीं हले 1 o/p दिये ,

$$\sum_{\substack{p \leq N \\ p \text{ is prime}}} p \xrightarrow{\text{K.I.C}} \sum_{\substack{(n) \\ n \leq N}} p [\text{p prime}] [p \leq n]$$

↓  
that number  $\leq n$   
is prime

$p \leq N$  एवं एकें (इन्हें)  
o/p दिया ।

### Sums & Recurrence

$$\text{The sum, } S_n = \sum_{k=0}^n a_k$$

$S_0$  = 0.  $a_0$  दोनों गणितों में ज्ञान है।

$$\begin{aligned} S_n &= \sum_{k=0}^{n-1} a_k + a_n \\ &= S_{n-1} + a_n \end{aligned}$$

So, the sum is equivalent to,  $S_0 = a_0$

$$S_n = S_{n-1} + a_n ; n > 0$$

If,  $a_n = \underbrace{\text{constant}}_{\beta} + \underbrace{\text{multiple of } n}_{n\gamma}$ ,  $a_0 = \alpha$

$$R_0 = \alpha \quad \text{constant term}$$

$$R_n = R_{n-1} + (\beta + n\gamma) \quad \left. \begin{array}{l} \text{multiple of } n \\ \text{constant term} \end{array} \right\} \text{multiple of } n$$

Now it is observed that at every step it is constant but if added at each step.

$$(1+n)\beta + n\gamma$$

$$R_0 = 1$$

$$R_n = R_{n-1} + \beta + \gamma n$$

$n = 1, 2, 3, \dots$

$$R_0 = \alpha$$

$$R_1 = R_0 + \beta + \gamma$$

$$= \alpha + \beta + \gamma$$

$$R_2 = R_1 + \beta + 2\gamma$$

$$= \alpha + 2\beta + 3\gamma$$

$$R_3 = R_2 + \beta + 3\gamma$$

$$= \alpha + 3\beta + 6\gamma$$

generalised solution,  $R_n = A(n) \cdot \alpha + B(n) \cdot \beta + C(n) \cdot \gamma$

Set,  $R_n = 1$  :  $R_n = 1$

$$\begin{cases} R_n = 1 \\ R_{n-1} = 1 \\ R_0 = 1 \end{cases}$$

$$\begin{cases} 1 = 1 + (\beta + \gamma n) \\ \Rightarrow \beta + \gamma n = 0 \\ \text{so, } \beta = \gamma = 0 \end{cases}$$

$$\text{So, } \alpha = 1, \beta = 0, \gamma = 0$$

$$\text{So, eqn (iii), } 1 = A(n) \cdot 1 + B(n) \cdot 0 + C(n) \cdot 0$$

$$\therefore A(n) = 1$$

Set  $R_n = n$  :  $R_0 = \alpha \therefore R_0 = 0$

$$\begin{aligned} R_n &= R_{n-1} + \beta + \gamma n \\ \Rightarrow n &= n-1 + \beta + \gamma n \\ \Rightarrow \beta + \gamma n &= 1 \\ \Rightarrow \beta + \gamma n &= 1 + 0 \cdot n \end{aligned}$$

$$\text{and } \therefore \beta = 1, \gamma = 0$$

$$\text{So, from eq (iii), } R_n = A(n)\alpha + B(n)\beta + C(n)\gamma$$

$$\Rightarrow n = A(n) \cdot 0 + B(n) \cdot 1 + C(n) \cdot 0$$

$$\therefore B(n) = n$$

Set  $R_n = n^r$

$$R_0 = 0^r = 0 = \alpha$$

$$\therefore \alpha = 0$$

$$\therefore n^r = (n-1)^r + \beta + \gamma n$$

$$\Rightarrow n^r = n^r - 2n + 1 + \beta + \gamma n$$

$$\Rightarrow -1 + 2n = \beta + \gamma n$$

$$\therefore \beta = -1, \gamma = 2$$

$$\text{From eqn } \textcircled{iii}, n^r = A(n) \cdot 0 + B(n) \cdot (-1) + C(n) \cdot 2$$

$$\Rightarrow n^r = -B(n) + 2 \cdot C(n)$$

$$\Rightarrow C(n) = \frac{B(n) + n^r}{2}$$

$$\begin{aligned} & (0+1)+1=1 \\ & 0+0+0=0 \\ & 2+2=4 \end{aligned}$$

So, from eqn  $\textcircled{ii}$ ,

$$R(n) = 1 \cdot \alpha + n \cdot \beta + \frac{n+n^r}{2} \cdot \gamma \quad \text{--- } \textcircled{iv}$$

$\rightarrow$  closed form

$$S_n = \sum_{k=0}^n (a+bk)$$

$$\text{when, } k=0 \rightarrow a+b \cdot 0$$

$$k=1 \rightarrow a+b \cdot 1$$

$$k=2 \rightarrow a+b \cdot 2$$

$$k=n \rightarrow a+b \cdot n$$

$$\sum_{k=0}^{n-1} (a+bk) + a+bn$$

$$a \cdot (n+1) + 1 \cdot (n+1) + b \cdot R_n = R_{n+1} \quad \text{so, } \beta = \alpha$$

$$\text{and } R_0 = \alpha = \alpha$$

$$\therefore \alpha = a, \beta = b$$

half solution H

$$\therefore R(n) = \alpha + n\beta + \frac{n+n}{2}\gamma \quad (\text{from eqn } ⑦)$$

$$= a + na + \frac{n+n}{2}b$$

$$= a(n+1) + b \cdot \frac{n(n+1)}{2} \quad (\rightarrow \text{soln from } \sum_{k=0}^n (a+bk))$$

## TOWER OF HANOI :

$$\text{Recurrence soln} \rightarrow T_0 = 0$$

$$\rightarrow T_n = 2T_{n-1} + 1 \quad (\text{for } n > 0) \quad \rightarrow ②$$

divide eqn ② by  $2^n$ ,

$$\frac{T_0}{2^0} = \frac{0}{2^0} = 0 \quad (n=0) \quad \rightarrow (S_0)$$

$$\frac{T_n}{2^n} = \frac{2T_{n-1} + 1}{2^n} \quad (\text{for } n > 0) \quad \text{not yet in form of eqn ③}$$

$$\Rightarrow \frac{T_n}{2^n} = \frac{T_{n-1}}{2^{n-1}} + \frac{1}{2^n} \quad (\text{for } n > 0) \quad \text{eqn ③ in form of}$$

assume,  $\frac{T_n}{2^n} = S_n$ , right hand side of ③ must

$$\text{L.H.S.} - S_0 = 0 \quad \text{and L.H.S. of ③ is also 0}$$

$$S_n = S_{n-1} + \frac{1}{2^n}$$

It follows that,

$$S_n = \sum_{k=0}^n 2^{-k}$$

$$\frac{a(1-r^n)}{1-r}$$

(for insert)

$$S_n = 2^{-1} + 2^{-2} + 2^{-3} + \dots + 2^{-n} + 0$$

$$= \left(\frac{1}{2}\right)^1 + \left(\frac{1}{2}\right)^2 + \left(\frac{1}{2}\right)^3 + \dots + \left(\frac{1}{2}\right)^n$$

(40)

$$= \left[1 - \left(\frac{1}{2}\right)^n\right]$$

$$a = \frac{1}{2}, r = \frac{1}{2}$$

$$= 1 - \frac{1}{2^n}$$

$$S_n = \frac{2^n - 1}{2^n}$$

$$\therefore 2^n S_n = 2^n - 1$$

$$\therefore T_n = 2^n - 1$$

$$\frac{1}{2}(1 - (\frac{1}{2})^n)$$

(1-1/2)

$$a_n T_n = b_n T_{n-1} + c_n \quad (3)$$

↓

↓

↓

Summation factor

$$T_n = 2 T_{n-1} + 1$$

let the summation factor is  $S_n$

from (3),  $S_n$  multiply,

$$S_n \cdot a_n T_n = S_n b_n \cdot T_{n-1} + S_n c_n \quad 3:2$$

$$\text{let, } \boxed{S_n b_n = S_{n-1} \cdot a_{n-1}}$$

$$S_n = S_n a_n T_n$$

\* Recurrence  $\rightarrow$  sol<sup>n</sup> এবং কর্ণত ২<sup>n</sup> sum

$$\begin{aligned} \cdot 3.2 \rightarrow S_n &= S_{n-1} b_{n-1} T_{n-1} + s_n c_n \text{ at } n \\ S_n &= S_{n-1} + s_n c_n \rightarrow S_{n-2} + S_{n-1} c_{n-1} + s_n c_n \end{aligned}$$

Hence,

$$S_n = s_0 a_0 T_0 + \sum_{k=1}^n s_k c_k \cdot$$

$$= s_1 b_1 T_0 + \sum_{k=1}^n c_k s_k$$

$$S_n a_n T_n = s_1 b_1 T_0 + \sum_{k=1}^n c_k s_k$$

$$\Rightarrow \boxed{T_n = \frac{1}{s_n a_n} \left[ s_1 b_1 T_0 + \sum_{k=1}^n c_k s_k \right]} \rightarrow \text{sol } n$$

but this is same with to see, multiple breaking of

$$n=1 \rightarrow T_1 = \frac{1}{s_1 a_1} [s_1 b_1 T_0 + c_1 s_1]$$

$$\text{if we take } T_0 = \frac{1}{a_1} [b_1 T_0 + c_1]$$

$$S_n b_n = S_{n-1} a_{n-1}$$

$$S_n = \frac{S_{n-1} a_{n-1}}{b_n}$$

$$= \frac{S_{n-2} \cdot a_{n-2}}{b_{n-1}} \cdot \frac{a_{n-1}}{b_n}$$

$$S_n = \frac{a_{n-1} a_{n-2} a_{n-3} \dots a_1}{b_n b_{n-1} b_{n-2} \dots b_2}$$

$$a_n T_n = b_n T_{n-1} + c_n \quad \text{and} \quad T_{n-1} = 2T_{n-2} + 1$$

$$\downarrow \quad \downarrow \quad \downarrow$$

$$a_n T_n = 2 \cdot T_{n-1} + 1$$

$$\therefore s_n = \frac{1 \cdot 1 \cdot 1 \cdots 2 \cdot 3}{2 \cdot 2 \cdots 2} + \alpha T_{n-2}$$

$$\boxed{s_n = \frac{1}{2^{n-1}}} \rightarrow \text{summation factor}$$

### Problem

In quicksort algorithm, one of the mostly used method for sorting data, data's need to compare with other element. The average number of comparison steps made by quicksort when it's applied to  $n$  items in random order which satisfies the following recurrence.

$$c_0 = 0$$

$$c_n = n+1 + \frac{2}{n} \sum_{k=0}^{n-1} c_k, \text{ for } n > 0$$

1. Form the above recurrence in  $a_n T_n = b_n T_{n-1} + c_n$
2. Find the summation factor ( $s_n$ )
3. Prove that i)  $c_n = 2(n+1) \sum_{k=1}^n \frac{1}{k+1}$

ii)  $c_n = 2(n+1)H_n - 2n$ , where  $H_n$

is the harmonic sum

iii) Find closed form for  $c_n$

iv) Find  $c_1, c_2$  and  $c_3$ .

$$a_n T_n = b_n T_{n-1} + c_n$$

$$n c_n = (n+1) c_{n-1} + 2n$$

$\uparrow \text{with } a_n \text{ to obtain } c_n \leftarrow (-1)^n c_{n-1}$

$\uparrow b_n \qquad \uparrow \text{from } n c_n \leftarrow (-1)^n 2n-1$

$\uparrow b_{n-1} \qquad \uparrow c_{n-1}$

$$S_n = \frac{a_{n-1} a_{n-2} a_{n-3} \dots a_1}{b_n \cdot b_{n-1} \cdot b_{n-2} \dots b_2}$$

$$= \frac{(n-1)(n-2)(n-3) \dots 1}{(n+1) n \cdot (n-1) \dots 3}$$

$$= \frac{(n-1)!}{(n+1) n \cdot (n-1)!} = \frac{2}{n(n+1)}$$

$$= \frac{2}{n(n+1)} = \frac{2}{n^2 + n}$$

summation factor

remove  $n^2 + n$

cancel and split a partial sum

## Manipulation of Sums:

1. Distributive law.

2. Associative law.

3.

$$\sum_{k \in K} a_k = \sum_{P(k) \in K} a_{P(k)}$$

→ we can reorder the sum in any way

→  $P(k)$  → any permutation of the set.

Ex :

$$\sum_{k=1,2,3} a_k = a_1 + a_2 + a_3$$

↳ natural order → trans-

pose

- For comparing different set of indices,

$$\sum_{k \in K} a_k + \sum_{k \in K'} a_k = \sum_{k \in K \cup K'} a_k$$

*Kenette  
iteration*

$$\sum_{k=1}^m a_k + \sum_{k=m}^n a_k = a_m + \sum_{k=1}^{m-1} a_k$$



only  $k = m$  common

↳ Perturbation method

→ splitting a single term from a sum

constant

$$\sum_{k \in K} c a_k = c \sum_{k \in K} a_k$$

$$\sum_{k \in K} (a_k + b_k) = \sum_{k \in K} a_k + \sum_{k \in K} b_k$$

The operation of splitting of a term is the basis of perturbation method. It can be used to determine the closed form of any sum.

$$\begin{aligned}
 S_n &= \sum_{0 \leq k \leq n} a_k \\
 \Rightarrow S_n + a_{n+1} &= \sum_{0 \leq k \leq n} a_k + a_{n+1} \\
 &= \sum_{0 \leq k \leq n} a_k \\
 &\stackrel{=} {=} a_0 + \sum_{1 \leq k \leq n+1} a_k \\
 &= a_0 + \sum_{0 \leq k-1 \leq n} a_k \\
 &= a_0 + \sum_{0 \leq k \leq n} a_{k+1}
 \end{aligned}$$

Problem :

Find the sum of a general geometric progression using perturbation method, i.e.

$$S_n = \sum_{0 \leq k \leq n} a x^k$$

$$\rightarrow S_n = a x^0 + a x^1 + a x^2 + a x^3 + \dots + a x^n$$

$$S_0 = a$$

apply perturbation method, to minimize  $\alpha x$   
 b/w  $S_n$  &  $S_{n+1}$ , b/w  $x$  &  $x^{n+1}$

$$S_n + \alpha x = \sum_{0 \leq k \leq n} \alpha x^k + \alpha x^{n+1}$$

$$\begin{aligned} &= \alpha x^0 + \sum_{0 \leq k \leq n} \alpha x^{k+1} \\ &= \alpha x^0 + \sum_{0 \leq k \leq n} \alpha x^k \cdot x \\ &= \alpha x^0 + x \cdot \sum_{0 \leq k \leq n} \alpha x^k \end{aligned}$$

$$\Rightarrow S_n - x S_n = \alpha x^0 - \alpha x^{n+1}$$

$$\Rightarrow \cancel{(1-x)} S_n = \frac{\alpha(1 - x^{n+1})}{1-x}$$

∴  $S_n = \frac{\alpha(1 - x^{n+1})}{1-x}$

$$S_n = \frac{\alpha(1 - x^{n+1})}{1-x}$$

## Multiple of sums :

$$* \sum_{1 \leq j, k \leq 3} a_j b_k = a_1 b_1 + a_1 b_2 + a_1 b_3 + \\ a_2 b_1 + a_2 b_2 + a_2 b_3 + \\ a_3 b_1 + a_3 b_2 + a_3 b_3$$

$$* \sum_{P(j,k)} a_{j,k} = \sum_{j,k} a_{j,k} [P(j,k)] \rightarrow \text{Kronecker Delta Form}$$

$$* \sum_j \sum_k a_{j,k} [P(j,k)]$$

• Law of interchanging the order of summation.

$$\sum_j \sum_k a_{j,k} [P(j,k)] = \sum_{P(j,k)} a_{j,k} = \sum_k \sum_j a_{j,k} [P(j,k)]$$

$$\sum_{1 \leq j, k \leq 3} a_j b_k = \sum_{j,k} a_j b_k [1 \leq j, k \leq 3]$$

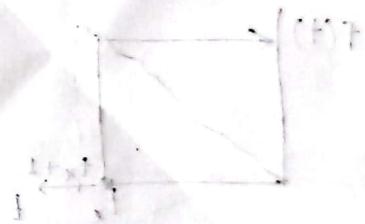
$$= \sum_{j,k} a_j b_k [1 \leq j \leq 3] [1 \leq k \leq 3]$$

$$= \sum_j \sum_k a_j b_k [1 \leq j \leq 3] [1 \leq k \leq 3]$$

$$\begin{aligned}
 &= \left( \sum_j a_j [1 \leq j \leq 3] \right) \left( \sum_k b_k [1 \leq k \leq 3] \right) \\
 &= \left( \sum_{j=1}^3 a_j \right) \left( \sum_{k=1}^3 b_k \right)
 \end{aligned}$$

General distributive law:

$$\sum_{\substack{j \in J \\ k \in K}} a_j b_k = \left( \sum_{j \in J} a_j \right) \left( \sum_{k \in K} b_k \right)$$



analogous to addition of two vectors. Similarly, A

is to obtain  $(A+B)^2$  using distributive law  
and  $(A+B)$  denotes multiplying after applying  
both vectors to have a single unique vector  
whose  $(x,y)$  coordinates sum to zero  
with two square signs will give the negative  
of the previous one.

2

\* Find the closed form for the following,

$$\begin{aligned}
 & 2 \sum_{\substack{1 \leq j, k \leq n \\ i > j, k}} a_k b_k \\
 & = 2 \sum_{1 \leq j \leq n} \sum_{1 \leq k \leq n} a_k b_k \\
 & = 2 \sum_{1 \leq k \leq n} \left( \sum_{1 \leq j \leq n} a_k b_k \right) \\
 & = 2 \sum_{1 \leq k \leq n} \left[ a_k b_k \sum_{1 \leq j \leq n} 1 \right] \\
 & = 2 \sum_{1 \leq k \leq n} [a_k b_k \cdot n] \\
 & = 2n \sum_{1 \leq k \leq n} a_k b_k
 \end{aligned}$$

\* Find closed form of the sum,  $S_n = \sum_{1 \leq j < k \leq n} \frac{1}{k-j}$

i) First summing on  $j$ ,

$$S_n = \sum_{1 \leq j < k \leq n} \frac{1}{k-j}$$

$$= \sum_{1 \leq k \leq n} \sum_{1 \leq j < k} \frac{1}{k-j}$$

$$= \sum_{1 \leq k \leq n} \sum_{1 \leq k-j < k} \frac{1}{k-(k-j)}$$

$$= \sum_{1 \leq k \leq n} \left( \sum_{1 \leq j < k-1} \frac{1}{j} \right)$$

$$= \sum_{1 \leq k \leq n} H_{k-1} = \sum_{0 \leq k \leq n} H_k$$

(Ans)

2) summing on  $k$ , not most helpful with  $j$

$$S_n = \sum_{1 \leq j < k \leq n} \frac{1}{k-j}$$

$$\begin{aligned} & j < k+j \leq n \\ \Rightarrow & j-j < k+j & j \leq n-j \\ \Rightarrow & 0 < k \leq n-j \end{aligned}$$

$$= \sum_{1 \leq j \leq n} \sum_{j < k \leq n} \frac{1}{k-j}$$

$$1 \leq k \leq n-j$$

$$= \sum_{1 \leq j \leq n} \sum_{j < k \leq n} \frac{1}{k+j-j} \quad \text{(adding)} \rightarrow \text{replacing } k \text{ with } k+j$$

$$= \sum_{1 \leq j \leq n} \left( \sum_{j < k+j \leq n} \frac{1}{k} \right)$$

$$= \sum_{1 \leq j \leq n} \left( \sum_{0 < k \leq n-j} \frac{1}{k} \right)$$

$$= \sum_{1 \leq j \leq n} H_{n-j}$$

$$1 \leq n-j \leq n$$

$$\Rightarrow 1-n \leq n-j-k \leq n-n$$

$$= \sum_{1 \leq n-j \leq n} H_{n-(n-j)}$$

$$1 \leq n-j \leq 0 \Rightarrow 1-n \leq -j \leq 0$$

$$= \sum_{1 \leq n-j \leq n} H_j$$

$$n-1 \geq j \geq 0$$

$$= \sum_{0 \leq j \leq n-1} H_j$$

$$0 \leq j \leq n-1$$

$$= \sum_{0 \leq j \leq n} H_j$$

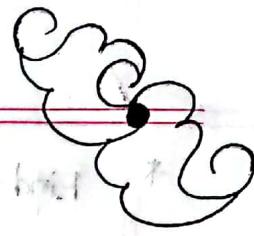
$$0 \leq j \leq n$$

$$H_{\infty} = \lim_{n \rightarrow \infty} H_n$$

$$(cont)$$

$$\sum_{1 \leq k \leq n} \frac{n}{k} - \sum_{1 \leq k \leq n} 1$$

$$= nH_n - n$$



3)  $S_n = \sum_{1 \leq j < k \leq n} \frac{1}{k-j}$  not much help here

$$= \sum_{1 \leq j < k \leq n} \frac{1}{k}$$

[replacing  $k$  by  $k+j$   
 $\because k+j-j = k$ ]

$$= \sum_{1 \leq k \leq n} \left( \sum_{j \leq k \leq n} \frac{1}{k} \right)$$

$$= \sum_{1 \leq k \leq n} \left( \sum_{\substack{j \leq k \\ 0 \leq k}} \frac{1}{k} \right)$$

$\sum_{j \leq k+j \leq n}$   
 $j-k < j \leq n-k$   
 $1 \leq j \leq n-k$

$$= \sum_{1 \leq k \leq n} \frac{1}{k} \cdot n$$

$$L.S. = \sum_{1 \leq k \leq n} \sum_{1 \leq j \leq n-k} \frac{1}{k}$$

$j \leq k+j$   
 $0 \leq k$   
 $j \leq n-k$

$$= \sum_{1 \leq k \leq n} \frac{n-k}{k}$$

$$= \sum_{1 \leq k \leq n} \frac{n}{k} - \sum_{1 \leq k \leq n} 1$$

$$= nH_n - n$$

\* Find closed form for the sum of squares, i.e  $D_n$

$$S_n = D_n = \sum_{0 \leq k \leq n} k^2 = 0^2 + 1^2 + 2^2 + \dots + n^2$$

→ Perturbation method :

$$D_n + (n+1)^2 = \sum_{0 \leq k \leq n} k^2 + (n+1)^2$$

$$= \sum_{0 \leq k \leq n+1} k^2 = \sum_{0 \leq k \leq n} (k+1)^2$$

$$= \sum_{0 \leq k \leq n} (k^2 + 2k + 1)$$

$$= \sum_{0 \leq k \leq n} k^2 + \sum_{0 \leq k \leq n} 2k + \sum_{0 \leq k \leq n} 1$$

$$\Rightarrow D_n + (n+1)^2 = D_n + 2 \sum_{0 \leq k \leq n} k + n+1$$

$$\Rightarrow 2 \sum_{0 \leq k \leq n} k = (n+1)^2 - (n+1)$$

$$\Rightarrow \sum_{0 \leq k \leq n} k = \frac{(n+1)(n+1-1)}{2}$$

$$\therefore \boxed{\sum_{0 \leq k \leq n} k = \frac{n(n+1)}{2}}$$

$$\square_n = \sum_{0 \leq k \leq n} k^3$$

$$\begin{aligned}
 \square_n + (n+1)^3 &= \sum_{0 \leq k \leq n} k^3 + (n+1)^3 \\
 &= \sum_{0 \leq k \leq n} (k+1)^3 = \sum_{0 \leq k \leq n} (k^3 + 3k^2 + 3k + 1) \\
 &= \sum_{0 \leq k \leq n} k^3 + 3 \sum_{0 \leq k \leq n} k^2 + 3 \sum_{0 \leq k \leq n} k + \sum_{0 \leq k \leq n} 1 \\
 &= \square_n + 3 \sum_{0 \leq k \leq n} k^2 + 3 \sum_{0 \leq k \leq n} k + \sum_{0 \leq k \leq n} 1 \\
 &= 3 \square_n + 3 \frac{n(n+1)}{2} + (n+1)
 \end{aligned}$$

$$\begin{aligned}
 \Rightarrow \square_n &= \frac{(n+1)^3}{3} - \frac{n(n+1)}{2} + \frac{n+1}{3} \\
 &= (n+1) \left\{ \frac{2n^2+4n+2}{6} - \frac{3n+2}{6} \right\} \\
 &= \frac{(n+1)(-n^2-n+6)}{6} \\
 &= \frac{n(n+1)(2n+1)}{6}
 \end{aligned}$$

$$1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{\sum_{k=1}^n k^2}{2} = \frac{2 + 3k}{2} = 15$$

Reperatoire method:

$$R_0 = \alpha$$

$$R_n = R_{n-1} + \beta + \gamma n + \delta n^2$$

$$R_n = A(n)\alpha + B(n)\beta + C(n)\gamma + D(n)\delta$$

$$S_n = \sum_{k=1}^n k^2$$

$$S_n = S_{n-1} + n^2$$

$$R_0 = \alpha$$

$$R_n = R_{n-1} + \beta + \gamma n + S_n^2$$

$$R_0 = \alpha = 1$$

$$n \cdot \alpha = n \cdot 1 = n(1 + 1 + \dots + n)$$

$$n^2 \cdot \alpha = \frac{n^2 + n}{2} \cdot 6(1 + 2 + 3 + \dots + n)$$

Set,

$$R(n) = n^3$$

$$R_0 = 0^3 = 0 \rightarrow \boxed{\alpha = 0}$$

$$R_n = R_{n-1} + \beta + n\gamma + n^2\delta$$

$$\Rightarrow n^3 = (n-1)^3 + \beta + n\gamma + n^2\delta$$

$$\Rightarrow \beta + n\gamma + n^2\delta = n^3 - (n-1)^3$$

$$(n+1)^3 - n^3 = (n^3 - 3n^2 + 3n - 1)$$

$$= 3n^2 - 3n + 1$$

$$\Rightarrow \beta + n\gamma + n^2\delta = (3n^2 - 3n + 1)$$

$$\alpha = 0, \beta = 1, \gamma = -3, \delta = 3$$

$$R(n) = A(n)\alpha + B(n)\beta + C(n)\gamma + D(n)\delta$$

$$\Rightarrow n^3 = 1 \cdot 0 + n \cdot 1 + \frac{n^2 + n}{2} \cdot (-3) + 3D(n)$$

$$\Rightarrow 3D(n) = 2n^3 + 3(n^2 + n) - 2n$$

$$= \frac{1}{3}(2n^3 + 3n^2 + n)$$

$$\underline{\underline{n(n+1)(2n+1)}}$$

## Chapter 4

### NUMBER THEORY

Divisibility :

We say that,  $m$  divides  $n$  or  $n$  is divisible by  $m$ .

$$\hookrightarrow m \mid n$$

if  $m > 0$  then,  $m \mid n$  is an integer

$m \mid n \rightarrow m > 0$  and  $n = mk$  for some integer  $k$ .

GCD & LCM :

- $\gcd(m, n) = \max \{k \mid k \mid m \text{ and } k \mid n\}$
- $\text{lcm}(m, n) = \min \{k \mid m \mid k \text{ and } n \mid k, k > 0\}$

\*  $\gcd(5, 14)$  using Euclid's algo :

$$\begin{array}{r} 5 \mid 14 \mid 2 \\ \hline 10 \\ \hline 4 \mid 5 \mid 1 \\ \hline 4 \mid 4 \mid 1 \\ \hline 0 \end{array}$$

Euclid's algorithm :

$$\gcd(m, n) = \gcd(n \bmod m, m)$$

$\hookrightarrow$  here,  $m < n$

$$\gcd(0, n) = n$$

$$\text{So, } \gcd(0, n) = n$$

$$\gcd(m, n) = \gcd(n \bmod m, m) \text{ for } m > 0$$

Euclid's algo (another)

$$\gcd(m, n) = m'm + n'n$$

$$\text{Ex: } \gcd(4, 14) = 2 = (-3) \cdot 4 + (1) \cdot 14$$

$$\gcd(0, 2) = 2 = m'm + n'n$$

$$\therefore \gcd(0, 2) = 0 + 1 \cdot 2 = (0, 2) \text{ bsp}$$

determining  $m'$  &  $n'$ :

$$\text{base case: if } m = 0, \text{ take } m' = 0, n' = 1 \text{ bsp}$$

$$\gcd(m, n) = 0 \cdot m + 1 \cdot n = n$$

$$\therefore \gcd(0, n) = n$$

Second case:

remainder,  $r = n \bmod m$

$$\gcd(m, n) = \gcd(n \bmod m, m)$$

$$(\Leftrightarrow \gcd(r, m))$$

$$= \gcd(m \bmod r, r) \quad ((r, m) \text{ bsp})$$

$$= \gcd(s, r) \quad r = (r, s) \text{ bsp}$$

$$\begin{array}{r} 4 | 14 | 3 \\ 12 | 2 | 0 \\ \hline 4 | 2 | 0 \end{array}$$

$$\begin{array}{r} 3 | 18 | 3 \\ 9 | 6 | 0 \\ \hline 3 | 6 | 0 \end{array}$$

$$\begin{array}{r} 2 | 10 | 2 \\ 5 | 5 | 0 \\ \hline 2 | 5 | 0 \end{array}$$

$$\gcd(r, m) = \overline{r}r + \overline{m}m$$

If  $r = n \bmod m$ , recursively apply  $r$  &  $m$  in place of  $m$  and  $n$ , computing  $\overline{r}$  and  $\overline{m}$ , such that,

$$\overline{r}r + \overline{m}m = \gcd(r, m)$$

here,  $r = n \bmod m$ ,

$$= n - \left\lfloor \frac{n}{m} \right\rfloor \cdot m$$

$$\gcd(m, n) = \gcd(r, m)$$

$$\begin{aligned} \Rightarrow m'm + n'n &= \overline{r}r + \overline{m}m \\ &= \overline{r} \left( n - \left\lfloor \frac{n}{m} \right\rfloor \cdot m \right) + \overline{m}m \\ &= \overline{r}n - \left\lfloor \frac{n}{m} \right\rfloor m \overline{r} + \overline{m}m \\ &= \overline{r}n + m \left( \overline{m} - \left\lfloor \frac{n}{m} \right\rfloor \overline{r} \right) \end{aligned}$$

equating co-efficients of  $m$  &  $n$ ,

$m' = \overline{m} - \left\lfloor \frac{n}{m} \right\rfloor \overline{r}$
$n' = \overline{r}$

$$\begin{array}{c} m \\ | \\ 3 \\ | \\ 14 \\ | \\ 12 \\ | \\ 2 \end{array}$$

$$14 = 3 \times 4$$

$$14 = 3 \times \left\lfloor \frac{14}{3} \right\rfloor$$

$$\downarrow \quad \downarrow \\ n = m \cdot \left\lfloor \frac{n}{m} \right\rfloor$$

$$m' = \bar{m} - \lfloor \frac{n}{m} \rfloor \bar{n}$$

$$n' = \bar{n}$$

$$\text{Ex: } \gcd(12, 18) = (-1) \cdot 12 + 1 \cdot 18$$

$$= \gcd(6, 12) = 1 \cdot 6 + 0 \cdot 12$$

$$= \gcd(\underbrace{\cancel{m}}_1, \underbrace{\cancel{n}}_0) = \underbrace{m'}_m \underbrace{m}_b + \underbrace{n'}_0 \underbrace{n}_1$$

$$\frac{m'}{m} = \frac{1}{6}$$

$$\begin{array}{r} 12 \\ | \\ 18 \end{array} \quad \begin{array}{r} 12 \\ | \\ 12 \end{array} \quad \begin{array}{r} 2 \\ | \\ 0 \end{array}$$

$$m' = 0 \cdot 1 - \left\lfloor \frac{12}{6} \right\rfloor \cdot 0 = 1$$

$$\gcd(6, 12) \quad n' = \bar{n} = 0$$

$$m' = 0 - \left\lfloor \frac{18}{12} \right\rfloor \cdot 1 = -1$$

$$n' = \bar{n} = 0 \quad \Rightarrow \left( m \cdot \left\lfloor \frac{n}{m} \right\rfloor - n \right) \bar{n} =$$

$$m \bar{n} + \bar{n} \left\lfloor \frac{n}{m} \right\rfloor - n \bar{n} =$$

$$\left( \bar{n} \left\lfloor \frac{n}{m} \right\rfloor - \bar{n} \right) m + n \bar{n} =$$

$$(m - \bar{n}^2 m) \bar{n} \quad (\text{durch } \bar{n} \text{ teilen}) \rightarrow \text{rest } \bar{n}$$

$$\boxed{\left. \begin{array}{l} \bar{n} \left\lfloor \frac{n}{m} \right\rfloor - \bar{n} = \bar{n} \\ \bar{n} = m \end{array} \right\}}$$

4.2 : Primes :

A positive integer  $n$  is called prime if it has just two divisors, 1 &  $n$  itself.  
1 is not prime.

So, prime sequence is 2, 3, 5, 7, 11, ...  
composite number,  $15 = 1 \times 15 = 3 \times 5$

A number can be either prime or composite.

- Prove that, any positive integer ( $n$ ) can be written as a product of primes and it is unique.

$\rightarrow$  *Mathematical representation*  $n = P_1 P_2 P_3 \dots P_m = \prod_{i=1}^m P_i$  — (4.1)

where,  $P_1 \leq P_2 \leq P_3 \leq \dots \leq P_m$

Part 1 :

Product of prime numbers is called

$$n = P_1 \cdot P_2 \cdot \dots \cdot P_m$$

if  $m=0$ ,  $n=1$  [empty product]; which implies  $n=1$ .

$\rightarrow$  Note: यहाँ यहाँ product का नाम है और यहाँ यहाँ का नाम है

if  $n > 1$  and, it is not prime, then there's some divisor  $n_1$ , such that,  $n_1 < n$

$$n = n_1 \times n_2$$

$\downarrow$

$$n_1' \times n_2'$$

$$1 < n_1 < n$$

with some examples

$n_1$  and  $n_2$  can be rewritten as products of prime otherwise both is prime factorizing  $n$ .  
So,  $n$  can be always represented as product of primes.

Ex :  $36 = 4 \times 9$        $n_1 = 2 \times 2$        $n_2 = 3 \times 3$   
 $\downarrow$        $\downarrow$        $\downarrow$   
 $n$        $n_1$        $n_2$

$36 = 2 \times 2 \times 3 \times 3$

Part II : Product of primes is unique.

There is only one way to write  $n$  as a product of primes in non-decreasing order.

Proof : when  $n = 1$ , there's only one possibility since the product must be empty.  
let's assume,  $n > 1$ , and we have two factorizations, where  $p$ 's and  $q$ 's are all prime.

$$n = p_1 p_2 p_3 \cdots p_m = q_1 q_2 q_3 \cdots q_k$$

where,  $p_1 \leq p_2 \leq p_3 \leq \dots \leq p_m = q_1 \leq q_2 \leq q_3 \leq \dots \leq q_k$

We'll prove,  $p_1 = q_1$

$$(p_1 - q_1) \mid n$$

$$p_1 \times p_2 \times \dots \times p_m = n$$

$$q_1 \times q_2 \times \dots \times q_k = n$$

if  $p_i \neq q_1$ , then we can assume that  $q_1 < p_i$   
making  $q_1$  smaller than  $p_i$ .  
As,  $p_i$  and  $q_1$  is prime,

$$\gcd(p_i, q_1) = 1$$

According to the euclid's self certifying algorithm.

$$\gcd(p_i, q_1) = ap_i + bq_1 \quad [\text{where } a \text{ and } b \text{ are integers}]$$

$$\Rightarrow ap_i + bq_1 = 1$$

$$\Rightarrow ap_i(p_2 p_3 \cdots p_m) + bq_1(p_2 p_3 \cdots p_m) = 1 \cdot (p_2 p_3 \cdots p_m)$$

$$\therefore \underbrace{ap_i}_{\text{1st term}} + \underbrace{bq_1 \cdot p_2 p_3 \cdots p_m}_{\text{2nd term}} = \underbrace{p_2 p_3 \cdots p_m}_{\text{R.H.S.}}$$

L.S  $\rightarrow q_1$  first divide first term

but, R.S  $\rightarrow q_1$  term  $\cancel{\text{not}}$

In eqn 4.2,  $q_1$  divides the R.H.S i.e  $p_2 p_3 \cdots p_m$

Thus,  $p_2 p_3 \cdots p_m / q_1$  will be an integer and  
 $p_2 p_3 \cdots p_m$  has a prime factorization  $q_1$  which  
contradicts our assumption.

So,  $p_1$  must be equal to  $q_1$ ,  $p_2$  to  $q_2$ , ...,  $p_m$  to  $q_m$ .  
 Thus, we can prove for all  $p_2, \dots, p_m$  are equal to  $q_2, \dots, q_m$ . So, the product of primes is unique.

[proved].

Prime exponent representation:

$$18 = 2^1 \times 3^2$$

$$18 = (1 \cdot 2^1 \cdot 3^2)_{\text{base } b} = (1 \cdot 2^1 \cdot 3^2)_{\text{base } 10}$$

$$n = \prod_p p^{n_p} \text{ where } n_p \geq 0$$

$$1) \gcd(m, n) \rightarrow ? \text{ with } \text{exponent } p \leq 2 \text{ if } p \mid m \text{ and } p \nmid n$$

$$2) \text{lcm}(m, n) \rightarrow ? \text{ with } \text{exponent } p \leq 2 \text{ if } p \mid m \text{ and } p \nmid n$$

$$n_p \text{ } 18 \rightarrow \langle n_2, n_3, n_5, n_7 \rangle \rightarrow 2 \times 3^5 \text{ (odd, both small)}$$

$$m_p \text{ } 16 \rightarrow \langle 4, 0, 0, 0 \rangle \rightarrow 2^4 \text{ (big, even)}$$

$$\gcd \rightarrow \min(n_p, m_p)$$

$$\text{lcm} \rightarrow \max(n_p, m_p)$$

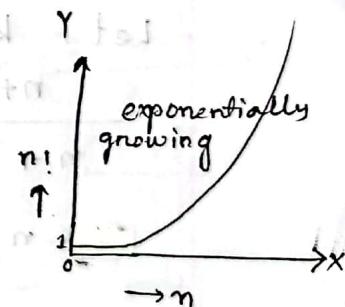
#### 4.4 factorial's factors

$$n! = 1 \cdot 2 \cdot 3 \cdots (n-1) \cdot n = \prod_{k=1}^n k, \text{ if } n \geq 0$$

for  $n=0$ , empty product. ie.  $0! = 1$

$$\begin{aligned} n! &= n(n-1)! \\ &= n(n-1)(n-2)! \end{aligned}$$

$n$	0	1	2	3	4	5	6	7	8
$n!$	1	1	2	6	24	120	720	5040	40320



$$\left(\frac{d}{dx}\right)^n \left(\frac{1}{x}\right) =$$

$$\left(\frac{d}{dx}\right)^n \left(\frac{1}{x+1}\right) =$$

$$\left(\frac{d}{dx}\right)^n \left(\frac{1}{(x+1)^2}\right) =$$

Prove that,  $n!$  increases exponentially using Gauss Trick or,  $n!$  is very big

$$n! = 1 \cdot 2 \cdot 3 \cdots n$$

$$(n!)^r = (1 \cdot 2 \cdot 3 \cdots n) \left( \frac{1 \cdot 2 \cdot 3 \cdots n}{1 \cdot 2 \cdot 3 \cdots n} \right)^{r-1}$$

$$= (1 \cdot 2 \cdot 3 \cdots n) (n \cdot (n-1) \cdots 1)$$

$$= 1 \cdot n \cdot 2(n-1) \cdot 3(n-2) \cdots n \cdot 1$$

$$= \underbrace{1(n-0)}_{1}, \underbrace{2(n-1)}_{2}, \underbrace{3(n-2)}_{3}, \cdots, \underbrace{(n-0) \cdot 1}_{1}$$

$$(n!)^r = \prod_{k=1}^n k(n+1-k)$$

$$\text{Let, } k = a$$

$$n+1-k = b$$

$$n+1 = a+b$$

$$2k - n - 1 = a - b$$

$$\text{From } 4.4.1, k(n+1-k) = ab$$

$$= \left(\frac{a+b}{2}\right)^r - \left(\frac{a-b}{2}\right)^r$$

$$= \left(\frac{n+1}{2}\right)^r - \left(\frac{2k-n-1}{2}\right)^r$$

$$= \frac{(n+1)^r}{4} - \frac{\{2k-(n+1)\}^r}{4}$$

$k \rightarrow 1 \text{ to } n$

$$k(n+1-k) = \frac{1}{4}(n+1)^2 - \left(k - \frac{n+1}{2}\right)^2 \quad \text{--- (4.4.2)}$$

$\left(\frac{n+1}{2}\right)^2 \geq 0 \rightarrow k(n+1-k) \text{ max}$   
 $\left(\frac{n+1}{2}\right)^2 \geq 1 \rightarrow k(n+1-k) \text{ min}$

In 4.4.2 has largest value, when  $\left(k - \frac{n+1}{2}\right)^2 = 0$ .

i.e.  $k = \frac{n+1}{2}$  without loss of generality.

and has smallest value when  $k=1$   $[1 \leq k \leq n]$

when,  $k = \frac{n+1}{2}$  then  $k(n+1-k) = \frac{1}{4}(n+1)^2 - 0 = \frac{(n+1)^2}{4}$   
(largest val)

$$\begin{aligned}
 \text{when } k=1, \text{ then, } k(n+1-k) &= \frac{1}{4}(n+1)^2 - \left(1 - \frac{n+1}{2}\right)^2 \\
 &= \frac{1}{4}(n+1)^2 - \left(\frac{2-n-1}{2}\right)^2 \\
 &= \frac{1}{4}(n+1)^2 - \frac{(n-1)^2}{4} \\
 &= \left(\frac{n+1}{2}\right)^2 - \left(\frac{n-1}{2}\right)^2 \\
 &= n \cdot 1 \\
 &= n \quad (\text{smallest val})
 \end{aligned}$$

$$\text{So, } n \leq k(n+1-k) \leq \frac{1}{4}(n+1)^2$$

$$\Rightarrow \sum_{k=1}^n n \leq \sum_{k=1}^n k(n+1-k) \leq \sum_{k=1}^n \frac{1}{4}(n+1)^2$$

$$\begin{aligned} & \Rightarrow n^n \leq (n!)^r \leq \left\{ \frac{1}{4} (n+1)^n \right\}^r \\ & \Rightarrow n^n \leq (n!)^r \leq \left( \frac{(n+1)^n}{2} \right)^r \\ & \Rightarrow \boxed{n^{\frac{n}{2}} \leq n! \leq \left( \frac{n+1}{2} \right)^n} \end{aligned}$$

↳ exponentially growing

From this relation, we conclude that, the factorial function grows exponentially.

$$\begin{aligned} & (1+\alpha)^n \cdot (1+\alpha)^{\frac{1}{2}} = (1+\alpha)^{n+\frac{1}{2}} \\ & (1+\alpha)^n \cdot (1+\alpha)^{\frac{1}{4}} = \\ & (1+\alpha) \cdot (1+\alpha)^{\frac{1}{2}} \\ & (1+\alpha) \cdot (1+\alpha)^{\frac{1}{4}} \end{aligned}$$

(for example)  $\alpha =$

$$\begin{aligned} & (1+\alpha)^{\frac{1}{2}} \geq (1+\alpha)^{\frac{1}{4}} \geq \dots \text{etc} \\ & (1+\alpha)^{\frac{1}{2}} \geq (1+\alpha)^{\frac{1}{4}} \geq \dots \geq \dots \end{aligned}$$

Chapter 5  
 "Binomial Co-efficient"

Symbol :  $\binom{n}{k}$  read as  $n$  choose  $k$  (or)  $nC_k$ .

Ex :  $\{1, 2, 3, 4\} \rightarrow$  choose two element  $\binom{4}{2} = 6$

$\{1, 2\}, \{2, 3\}, \{3, 4\}, \{1, 3\}, \{1, 4\}, \{2, 4\}$

If I have to choose  $k$  things from  $n$  things :

$$\frac{n}{k} \times \frac{n-1}{k-1} \times \dots \times \frac{n-(k-1)}{1}$$

$$\boxed{\binom{n}{k} = \frac{n(n-1)(n-2)\dots(n-k+1)}{k(k-1)(k-2)\dots1}} \rightarrow \text{when } k=2, \quad \binom{n}{2} = \frac{n(n-1)}{2}$$

$$\text{For real numbers} \rightarrow \binom{n}{k} = \binom{-1}{2} = \frac{(-1)(-1-1)}{2(2-1)}$$

$$\hookrightarrow \binom{\frac{1}{2}}{2} = \frac{\frac{1}{2}(\frac{1}{2}-1)}{2(2-1)}$$

Note : whole numbers are real numbers

$\binom{3}{2}$  can't be negative

$$\text{Ex 11. } \frac{n(n-1)(n-2)\cdots(n-k+1)(n-k)\cdots 1}{k! (n-k)\cdots 1}$$

$$= \frac{n!}{k!(n-k)!} \rightarrow \binom{n}{k} \quad \text{using } \frac{n!}{k!(n-k)!} = \frac{n(n-1)(n-2)\cdots(n-k+1)(n-k)\cdots 1}{k! (n-k)\cdots 1}$$

(\*) For real numbers  $n > 0$  then  $\binom{n}{k}$  : Inductive

$$\binom{n}{k} = \begin{cases} \frac{n(n-1)(n-2)\cdots(n-k+1)}{k(k-1)(k-2)\cdots 1} & ; k \geq 0 \\ 0 & ; k < 0 \end{cases}$$

$n^m$  is write as if  $n$  has  $m$  falling factors of  $n$

$$= n(n-1)(n-2) \cdots (n-m+1)$$

$$\binom{n}{k} = \frac{(n)(n-1)(n-2)\cdots(n-k+1)}{k!(n-k)!} \quad (n, k \text{ integers}, 0 \leq k \leq n)$$

# Pascal's Triangle :

$n$	$\binom{n}{0}$	$\binom{n}{1}$	$\binom{n}{2}$	$\binom{n}{3}$	$\binom{n}{4}$	$\binom{n}{5}$
0	1	0	0	0	0	0
1	1	1	0	0	0	0
2	1	2	1	0	0	0
3	1	3	3	1	0	0
4	1	4	6	4	1	0
5	1	5	10	10	5	1

symmetry identity

### Symmetry Identity :

From  
table

$$\rightarrow \begin{aligned} \binom{3}{0} &= \binom{3}{3}, & \binom{5}{0} &= \binom{5}{5} \\ \binom{3}{1} &= \binom{3}{2}, & \binom{5}{1} &= \binom{5}{4} \\ && \binom{5}{2} &= \binom{5}{3} \end{aligned} \quad \left\{ \boxed{\binom{n}{k} = \binom{n}{n-k}} \right.$$

So,  $\binom{n}{k} = \binom{n}{n-k}$  where  $n \geq 0$ , integer  $k$ . this formula is not for real numbers

(9) why  $n$  cannot be negative in symmetry identity ?

$$\rightarrow \text{let } n = -1, \quad \binom{-1}{k} = \binom{-1}{-1+k} \text{ ac.t symmetry rule}$$

$$= \binom{-1}{-1+k} \rightarrow k < 0 \text{ hence}$$

proof: L.S.  $= \binom{-1}{k}$  or  $\binom{-1}{k} \cdot \binom{0}{-1+k}$  Not equal

if  $k = 0$ ,  $\binom{-1}{0} = \frac{(-1)!}{0!(-1-0)!} = \frac{1}{0!} = 1$

[with example (d)] So,  $n$  cannot be negative. as.  $\binom{-1}{k} \neq \binom{-1}{-1-k}$

Formally:  $\binom{-1}{k} = \frac{(-1)(-2)(-3)\dots(-k)}{k!} = \frac{(-1)^k \cdot k!}{k!} = (-1)^k$

+1 or -1

[with example (d)]  $\binom{-1}{1}$  or  $\binom{-1}{2}$

$$\binom{-1}{-1-k} = \binom{-1}{-(1+k)} = 0$$

$\nwarrow (1+k) < 0$

So, as  $\binom{-1}{k} \neq \binom{-1}{-1-k}$ ,  $n$  cannot be negative.

### Absorption identity:

$$\binom{n}{k} = n \frac{n}{k} \binom{n-1}{k-1}$$

$$\binom{n}{k} = \frac{n(n-1) \cdots (n-k+1)}{k(k-1) \cdots (1)} = \binom{n-1}{k-1} \cdot \frac{n}{k}$$

Q) Prove that  $(n-k) \cdot \binom{n}{k} = n \binom{n-1}{k-1}$  using symmetry & absorption identity:

$$(n-k) \cdot \binom{n}{k} = (n-k) \binom{n}{n-k} \quad [\text{by symmetry}]$$

$$= (n-k) \cdot \binom{n}{n-k} \binom{n-1}{n-k-1} \quad [\text{by absorption}]$$

$$= n \binom{n-1}{n-1-k}$$

$$= n \binom{n-1}{k} \quad [\text{by symmetry}]$$

## Formula :

- Addition Formula: no. of ways -

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}, \text{ integer } k.$$

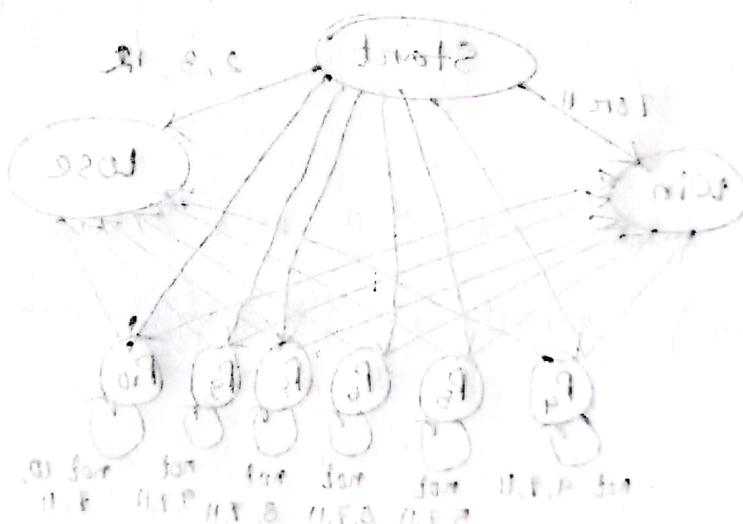
maps: first part, taking all -

$$\binom{4}{2} = \binom{3}{2} \binom{3}{1}$$

Note :  
\* recursive

## Home Study :

Prove this addition formula



Proof must

Addition formula :

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} + \dots + \binom{n-1}{0}$$

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k-1} + \dots + \binom{n-1}{0}$$

# Example :  $\binom{4}{3} + \binom{3}{2} = \binom{3}{2} + \binom{2}{1}$

$$\binom{5}{3} = \binom{4}{3} + \binom{4}{2}$$

$$= \binom{4}{3} + \binom{3}{2} + \binom{1}{1}$$

$$\binom{4}{3} + \binom{3}{2} = \binom{4}{3} + \binom{3}{2} + \binom{2}{1} + \binom{2}{0}$$

$$= \binom{4}{3} + \binom{3}{2} + \binom{2}{1} + \binom{1}{0} + \binom{1}{-1} \quad \text{for } k < 0, \text{ so } 0.$$

$$\Rightarrow a_1 \quad a_2 \quad a_3 \quad a_4$$

$$= \sum_{k=1}^4 a_k$$

$$\boxed{\binom{4}{3} = \binom{1+1+1}{1+1+1}}$$

$$\binom{2+3}{3} = \binom{1+3}{3} + \binom{1+2}{2} + \binom{1+1}{1} + \binom{1+0}{0}$$

$$\binom{1+1+3}{3} = \binom{1+3}{3} + \binom{1+2}{2} + \binom{1+1}{1} + \binom{1+0}{0}$$

$$= \binom{x+k}{k} + \binom{x+(k-1)}{k-1} + \dots + \binom{x+0}{0}$$

$$\binom{x+n+1}{n} = \sum_{k=0}^n \binom{x+k}{k}$$

$$\boxed{\binom{x+n+1}{n} = \sum_{k=0}^n \binom{x+k}{k}}$$

Closed form

Note:  
→ x constant term

~~1st Option  
1st element  
frac~~

$$\begin{aligned}
 \binom{5}{3} &= \binom{4}{3} + \binom{4}{2} \\
 &= \binom{3}{3} + \binom{3}{2} + \binom{4}{2} = \binom{7}{4} \\
 &= \binom{2}{3} + \binom{2}{2} + \binom{3}{2} + \binom{4}{2} : \text{approx} \\
 &= \underbrace{\binom{1}{3}}_{1 < 3} + \underbrace{\binom{1}{2}}_{1} + \underbrace{\binom{2}{2}}_{1} + \underbrace{\binom{3}{2}}_{1} + \underbrace{\binom{4}{2}}_{1} \\
 &= \underbrace{\left( \binom{0}{3} + \binom{0}{2} \right)}_{\text{pattern}} + \underbrace{\left( \binom{1}{2} + \binom{2}{2} + \binom{3}{2} \right)}_{\text{pattern}} + \binom{4}{2}
 \end{aligned}$$

$$\binom{4+1}{2+1} = \sum_{k=0}^m \binom{k}{2} \quad \text{pattern}$$

$$\Rightarrow \boxed{\binom{n+1}{m+1} = \sum_{k=0}^n \binom{k}{m}}$$

(closed form for the sum)  $\binom{1}{0} + \binom{2}{1} + \binom{3}{2} + \dots + \binom{n}{m}$

\* Prove it  $\sum_{k=0}^r \binom{r}{m+k} \binom{s}{n-k} = \binom{r+s}{m+n} \binom{r+s}{n}$

$$\binom{r+s}{0} + \dots + \binom{r+s}{m+n} + \binom{r+s}{n}$$

$$\binom{r+s}{0} + \dots + \binom{r+s}{m+n}$$

$$\left\{ \binom{r+s}{m+n} \right\} = \binom{r+s}{n}$$

cancel

Table 174  
formulas

- Q) Ans
- $$\sum_{k=0}^n \binom{n}{k} x^k y^{n-k}$$
- HOME STUDY:  $\sum$  of something  $= \sum_{k \leq n} \binom{n}{k} x^k \cdot y^{n-k}$
- Find  $(x+y)^n$  in terms of summand binomial co-efficient
  - Find  $2^n$  when  $x=1$  and  $y=1$ ,  $n=r$
  - Find  $0^n$  when  $x=-1$  and  $y=1$ ,  $n=r$
  - Prove that,  $\binom{n}{k} = (-1)^{k+1} \binom{k-r-1}{k}$
  - Table 174.

1.  $\{P, S, S, P\}$

2.  $\{P, S, S, P, S\}$

3.  $\{P, S, S, P, S, P\}$

4.  $\{P, S, S, P, S, P, S\}$

5.  $\{P, S, S, P, S, P, S, P\}$

6.  $\{P, S, S, P, S, P, S, P, S\}$

7.  $\{P, S, S, P, S, P, S, P, S, P\}$

$F = \{P, S\}$

$S = \{P, S\}$

## Chapter 6

### Special Numbers

- Stirling Numbers is amount of  $\binom{[n]}{k}$  having  $n$  elements  
→ First kind  
→ Second kind

Second kind :  $\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \# \text{ways to partition a } (n \text{ element}) \text{ set into } k \text{ non-empty subsets.}$

Example :

$\left\{ \begin{matrix} 4 \\ 2 \end{matrix} \right\} \rightarrow 4 \text{ element}$   
 $\rightarrow 2 \text{ non-empty subset}$

$$\{1, 2, 3, 4\} \longrightarrow \begin{array}{l} \textcircled{1} \quad \{1, 2, 3\}, \{4\} \\ \textcircled{2} \quad \{1, 2\}, \{3, 4\} \\ \textcircled{3} \quad \{1, 3\}, \{2, 4\} \\ \textcircled{4} \quad \{2, 3\}, \{1, 4\} \\ \textcircled{5} \quad \{1\}, \{2, 3, 4\} \\ \textcircled{6} \quad \{2\}, \{1, 3, 4\} \\ \textcircled{7} \quad \{3\}, \{1, 2, 4\} \end{array}$$

$$\left\{ \begin{matrix} 4 \\ 2 \end{matrix} \right\} = 7$$

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = ?$$

$$2^{\binom{n}{k}} = 2$$

In  $\binom{n}{k}$ , if  $k=1$ . i.e., 1 non-empty subset

$\binom{n}{1} = 1$  choice elements have to be a pair.

$\binom{0}{1} = 0$  if no element then nothing is of

when  $k=0$  i.e., 0 non-empty subset.

If there is no element, then it is empty set

$\binom{n}{0} = 0$

$\binom{0}{0} = 1 \rightarrow \{ \}$

below part shows some idea behind it

when  $k=2$ , there are 2 non-empty subsets,

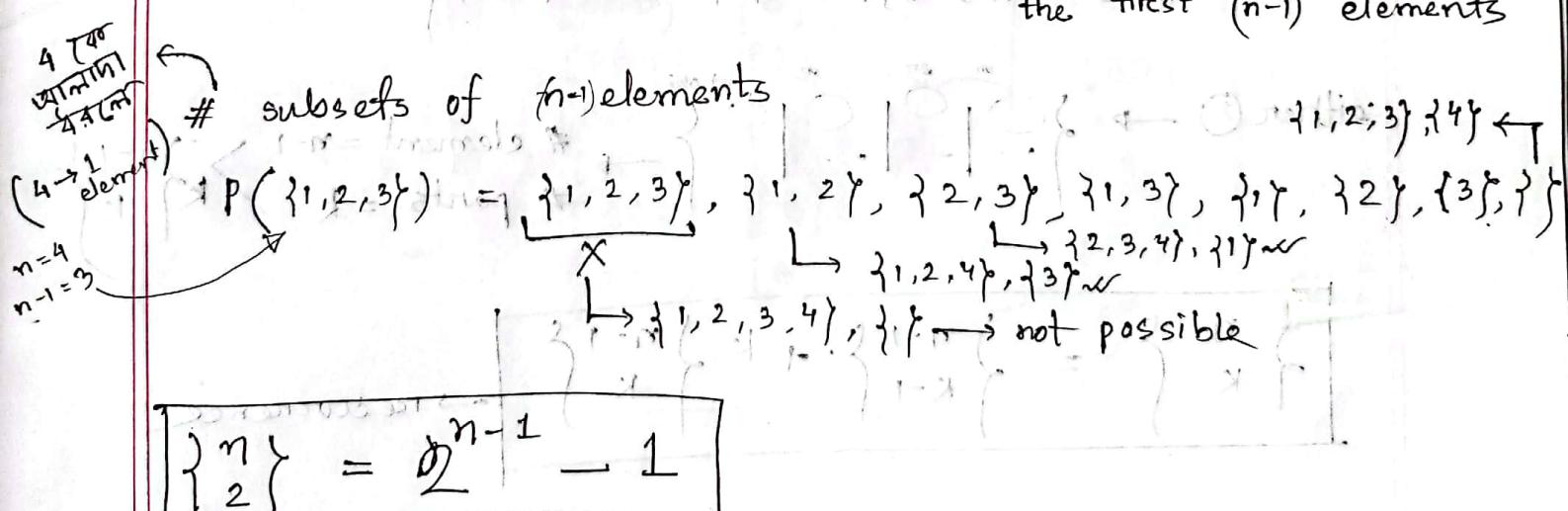
$\binom{n}{2} = ?$

$\binom{0}{2} = 0$ , if  $n > 1$  &  $k=2$

↳ if,  $n=0$

↳ 2 pairs one of these part contains the last element

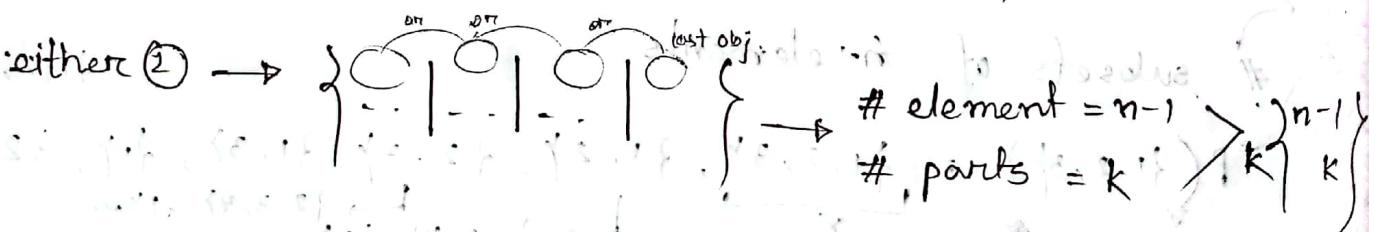
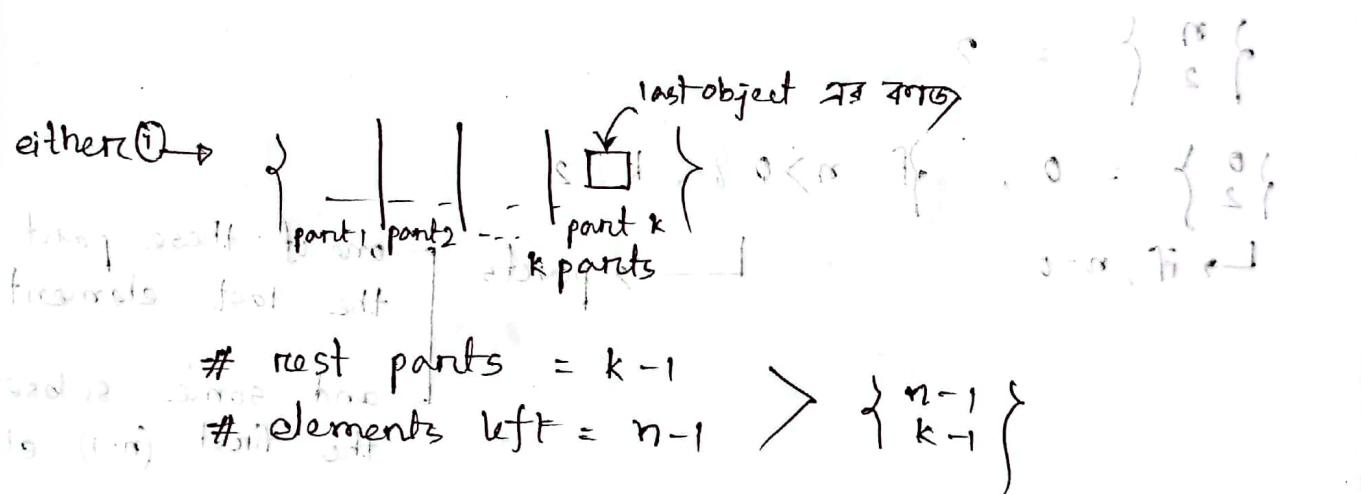
$\binom{n}{2} < \text{and some subset of } \{1, 2, 3, 4\}$  from the first  $(n-1)$  elements



$$\boxed{\binom{n}{2} = 2^{n-1} - 1}$$

## Generalization of $\{n\}_k$

- Given a set of  $n > 0$  elements / objects / things
- to be partition into  $k$  non-empty subsets
- # ways
  - either - last object is in a class/part by itself
  - last object is with some non-empty subset of the first  $(n-1)$  objects



$$\boxed{\{n\}_k = \{n-1\}_{k-1} + k \{n-1\}_k} \rightarrow \text{recurrence}$$

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} \rightarrow \text{recurrence}$$

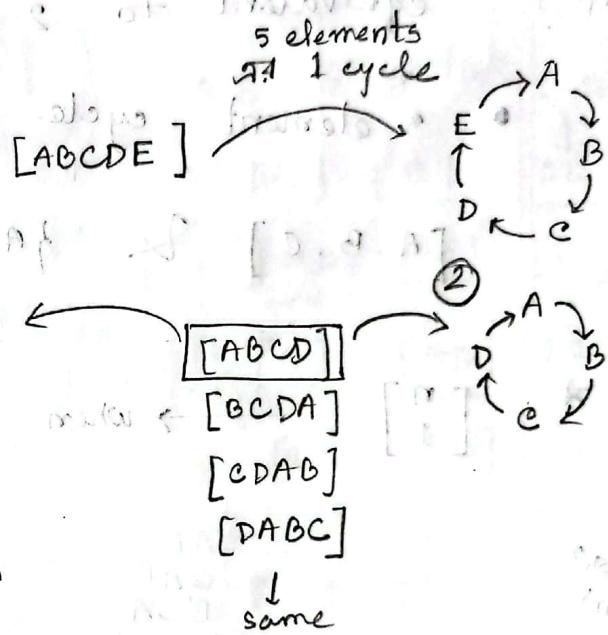
Next class → first kind

## First kind of the stirling numbers

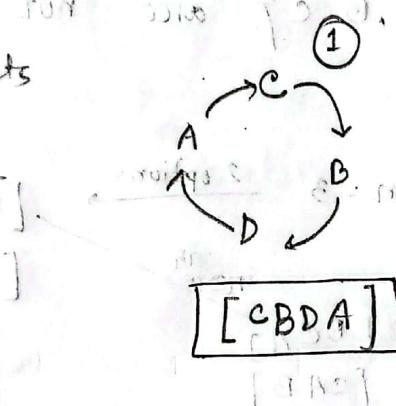
$\begin{bmatrix} n \\ k \end{bmatrix} \rightarrow$  counts #ways to arrange  $n$  elements / objects into  $k$  cycle

→  $n$  cycle  $k$ . ways of arranging

→ cyclic arrangement like necklace.



$$\begin{bmatrix} 4 \\ 1 \end{bmatrix} \xrightarrow{4 \text{ elements}} \begin{bmatrix} 1 \text{ cycle} \\ [B \circ A \circ] \end{bmatrix}$$



$$\begin{bmatrix} 4 \\ 2 \end{bmatrix} \xrightarrow{4 \text{ elements}} \begin{bmatrix} 2 \text{ cycle} \\ [A \circ B \circ C \circ D] \end{bmatrix}$$

[ABC]	[D]
[CBA]	[D]
[ABD]	[C]
[DBA]	[C]
[BCD]	[A]
[DCB]	[A]
[ACD]	[B]
[DCA]	[B]

3 | 1

2 | 2

$$\begin{array}{ll} [AB] & [CD] \\ [BC] & [AD] \\ [BD] & [AC] \end{array}$$

$$\begin{array}{l} [ ] \\ [ ] \end{array}$$

→ 11 options.

- $\{A\}$  or  $\{B\}$
- $\{B, C\}$  or  $\{C, A\}$
- 1 element cycle or singleton (1-cycle) is equivalent to 1-set (1 element set)
  - 2 " " " " or doubleton (2-cycle) is equivalent to 2-set.
  - 3 element cycle (3-cycle) has several options.  
 $[A, B, C]$  &  $\{A, B, C\}$  are not equal.

\*

$$\begin{bmatrix} n \\ 1 \end{bmatrix}$$

→ when  $n=3$   $\xrightarrow{2 \text{ options}}$

$$\begin{array}{c} ABC \\ CAB \\ BCA \end{array}$$

3 element  
↳ 3! way

$$\begin{array}{c} BCA \\ CAB \end{array}$$

पर्याप्त

$$\begin{array}{c} ABC \\ BAC \end{array}$$

पर्याप्त

$$\begin{array}{c} ACB \\ CAB \end{array}$$

पर्याप्त

$$\begin{array}{c} CBA \\ BAC \end{array}$$

पर्याप्त

$$\begin{array}{c} A \\ B \\ C \end{array}$$

$$\text{So, } \frac{3!}{3} = 2$$

$\begin{bmatrix} n \\ 1 \end{bmatrix}$	$= \frac{n!}{n}$	$= (n-1)!$
$\{1\}$		
$\begin{bmatrix} n \\ 1 \end{bmatrix}$	$= 1$	

$$\begin{bmatrix} n \\ 1 \end{bmatrix} \geq \{1\}$$

$$\begin{array}{c} A \\ B \\ C \\ D \\ E \\ F \\ G \\ H \\ I \end{array}$$

समिक्षा न हो

When  $k = n$ ,  $\left[ \begin{smallmatrix} n \\ n \end{smallmatrix} \right]_{\text{part}} = \left[ \begin{smallmatrix} n \\ 1 \end{smallmatrix} \right]$  periodic partition

$$\left[ \begin{smallmatrix} n \\ n \end{smallmatrix} \right]_{\text{part}} = 1 \text{ since only } \left[ \begin{smallmatrix} n \\ n \end{smallmatrix} \right] \text{ is a partition of } n \text{ into } n \text{ parts}$$

$$\text{Hence } \left\{ \begin{smallmatrix} n \\ n \end{smallmatrix} \right\}_{\text{part}} = 1 \text{ since } \left\{ \begin{smallmatrix} n \\ n \end{smallmatrix} \right\} \text{ is a partition of } n \text{ into } n \text{ parts}$$

When,  $k = n-1$

$$\text{Since } \left[ \begin{smallmatrix} n \\ n-1 \end{smallmatrix} \right]_{\text{part}} = \left( \begin{smallmatrix} n \\ 2 \end{smallmatrix} \right) \text{ number of ways}$$

$$\left\{ \begin{smallmatrix} n \\ n-1 \end{smallmatrix} \right\} = \left( \begin{smallmatrix} n \\ 2 \end{smallmatrix} \right) \text{ number of ways}$$

$A B C D$  partition  
2 elements

[A] [B] [CD]

[AD] [B] [C]

[A] [BD] [C].

[Ac] [C] [D]

The number of ways to arrange  $n$  objects into

$(n-1)$  cycles or subsets is the number of ways to choose the two objects that will be in the same cycle or subset.

## Generalizing $\binom{n}{k}$ :

- $n > 0$
- # ways to choose  $k$  cycles from  $n$  elements
- Put the last object into a cycle by itself.

either -  $\binom{n-1}{k-1} \leftarrow \text{last obj } \cancel{\text{not}}$

or insert the last object into one of the cycle arrangements of the first  $(n-1)$  objects and there are  $(n-1)$  different options of insertions.

example:

$[A_1, B_1, C_1] \rightarrow D$  insert वर्षम्

$(n-1)$  options

$$\boxed{\binom{n}{k} = \binom{n-1}{k-1} + (n-1) \binom{n-1}{k}}$$

Eulerian Number : $\langle \cdot \cdot \rangle$  : stuff

Eulerian

 $\langle \cdot \cdot \rangle$  with,  $n < k$  no point

- represented by  $\langle \cdot \cdot \rangle \curvearrowright n$  ascent,  $k$
- $n$  ascent  $k$
- Angle bracket suggests less, than, greater than

 $\langle \cdot \cdot \rangle$  - is the # permutations  $\pi_1, \pi_2, \pi_3, \dots, \pi_n$ 

of  $\{n_1, n_2, n_3, \dots, n\}$  elements, that, have  $k$  ascents:

Example : 4 element  
 $\{1, 2, 3, 4\} \rightarrow \langle \cdot \cdot \rangle = ?$   
 permutation =  $4! = 24$

 $\langle \cdot \cdot \rangle \rightarrow$  $1 < 2 < 3$  $\downarrow \pi_2$ permutation  
 $\geq 2$  st ascent $\langle \cdot \cdot \rangle \rightarrow$ ①  $1 < 2 < 4 > 3$ ②  $1 < 3 > 2 < 4$ 3 element  $\rightarrow 3! = 6$ 

$\pi_1$	$\pi_2$	$\pi_3$	# < or number of ascents	$\langle \cdot \cdot \rangle = ?$
1	2	3	2 ( $1 < 2 < 3$ )	$\langle \cdot \cdot \rangle = 1$
1	3	2	1 ( $1 < 3 > 2$ )	$\langle \cdot \cdot \rangle = 4$
2	1	3	1 ( $2 > 1 < 3$ )	
2	3	1	1 ( $2 < 3 > 1$ )	$\langle \cdot \cdot \rangle = 1$
3	1	2	1 ( $3 > 1 < 2$ )	
3	2	1	6 ( $3 > 2 > 1$ )	$\langle \cdot \cdot \rangle = 0$

Note :  $\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \rangle$

- If  $k = n$  or  $k > n$ , then  $\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \rangle = 0$

Table - 268 : Eulerian Triangle

$n$	$\langle \begin{smallmatrix} n \\ 0 \end{smallmatrix} \rangle$	$\langle \begin{smallmatrix} n \\ 1 \end{smallmatrix} \rangle$	$\langle \begin{smallmatrix} n \\ 2 \end{smallmatrix} \rangle$	$\langle \begin{smallmatrix} n \\ 3 \end{smallmatrix} \rangle$	$\langle \begin{smallmatrix} n \\ 4 \end{smallmatrix} \rangle$	$\langle \begin{smallmatrix} n \\ 5 \end{smallmatrix} \rangle$
0	1					
1	1	1				
2	1	1				
3	1	4	1			
4	1	11	11	1		
5	1	26	66	26	1	

AC = IP = positive integer

$$\langle \begin{smallmatrix} 2 \\ 0 \end{smallmatrix} \rangle = \langle \begin{smallmatrix} 2 \\ 1 \end{smallmatrix} \rangle$$

$$\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \rangle = \langle \begin{smallmatrix} n \\ n-1-k \end{smallmatrix} \rangle$$

$$\langle \begin{smallmatrix} 3 \\ 0 \end{smallmatrix} \rangle = \langle \begin{smallmatrix} 3 \\ 2 \end{smallmatrix} \rangle$$

2 = 18 e - formula e

$$\langle \begin{smallmatrix} 4 \\ 0 \end{smallmatrix} \rangle = \langle \begin{smallmatrix} 4 \\ 3 \end{smallmatrix} \rangle$$

$$\langle \begin{smallmatrix} 4 \\ 1 \end{smallmatrix} \rangle = \langle \begin{smallmatrix} 4 \\ 2 \end{smallmatrix} \rangle$$

$$\langle \begin{smallmatrix} 5 \\ 0 \end{smallmatrix} \rangle = \langle \begin{smallmatrix} 5 \\ 4 \end{smallmatrix} \rangle$$

$$\langle \begin{smallmatrix} 5 \\ 1 \end{smallmatrix} \rangle = \langle \begin{smallmatrix} 5 \\ 3 \end{smallmatrix} \rangle$$

	X	X	X
	8	8	8
	8	8	8
	8	8	8
	8	8	8

### Observations:

- There can be  $(n-1)$  positions for measuring ascent or not ascent.
- at most,  $(n-1)$  ascents possible.
- There is symmetry between left & right part

So, symmetry  $\langle \frac{n}{k} \rangle = \langle \frac{n}{n-k} \rangle$

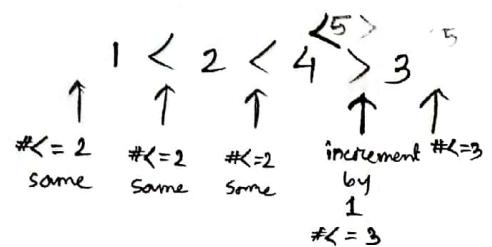
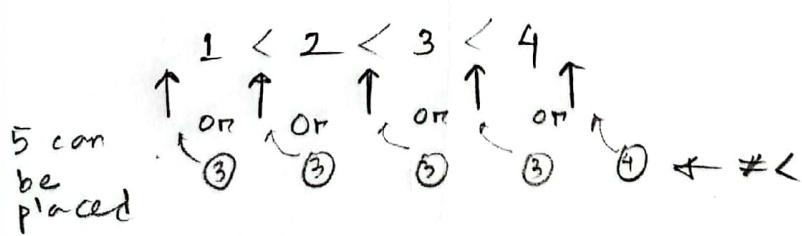
The permutation  $\pi_1 \pi_2 \pi_3 \dots \pi_n$  has,  $n-1-k$  ascents if its reflection  $\pi_n \pi_{n-1} \dots \pi_1$  has  $k$  ascents.

$$\text{reflecting in } \rightarrow \text{top row}$$

$$\langle \frac{3}{2} \rangle = 1 \quad \langle \frac{3}{0} \rangle = 1$$

General soln of  $\langle \frac{n}{k} \rangle$ : to between  $n-1$

Each permutation  $P = P_1 P_2 \dots P_{n-1}$  of elements  $\{1, 2, 3, \dots, n-1\}$  leads to  $n$  permutations of  $\{1, 2, 3, \dots, n\}$  if we insert the new element  $n$  in all possible way i.e.,  $n=5 \quad \{1, 2, 3, 4, 5\}$



first  $\pi_1 \pi_2 \pi_3 \pi_4$  last position  
 ascent  $\pi_1 \pi_2 \pi_3 \pi_4$   $\pi_n$   $\rightarrow$  ক্রমে always increment by 1  
 last element  $\pi_1 \pi_2 \pi_3 \pi_4$   $\pi_n$   $\rightarrow$  (1-n)st position  $\rightarrow$   
 মাঝের position  $\rightarrow$  ascent নামে no change  
 descent  $\rightarrow$   $\pi_1 \pi_2 \pi_3 \pi_4$   $\pi_n$   $\rightarrow$  increment by 1.  
 length of the resulting permutation is  $n+1$

Two case

the permutation has  $k$  ascents

the next  $\pi_1 \pi_2 \dots \pi_{k-1} \pi_n$  without  $n$  will  
starts at  $\pi_1, \pi_2, \dots, \pi_{k-1}$  itself or if  $\pi_i$

If we put  $n$  in position  $\pi_j$ ,

we get the permutation,  $\pi = \rho_1 \rho_2 \dots \rho_{j-1} n \rho_j \dots \rho_{n-1}$

# ascents in  $\pi$  same as  $\rho = \rho_1 \rho_2 \dots \rho_{n-1}$

if  $n$  inserted at:  $\pi$  first position

or middle of any ascent positions last

if there are  $k$  ascents

length will remain same if  $n$  is inserted  
for example  $\pi = 1 2 3 4 5$  and  $n = 6$  then  $\pi$  will be

1 2 3 4 6 5

1 2 3 5 6 4

1 2 4 3 6 5

1 3 2 4 6 5

1 3 4 2 6 5

1 3 4 5 6 2

ascent  $\rightarrow k$   
not ascent  $\rightarrow \underbrace{n-1}_{\text{composition}} - k$

# ascents in  $\pi$  increased by 1 from  $P$ ,

| n inserted at the last position

| or middle of any not ascent position

| there are  $(n-1-k)$  not ascent position

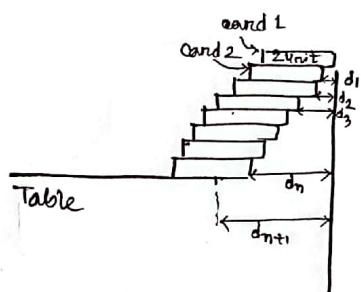
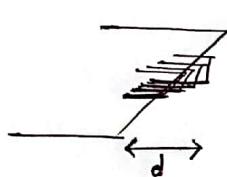
$$\rightarrow \underbrace{\binom{n-1}{k} + k \binom{n-1}{k}}_{\text{first ascent}} + \underbrace{\binom{n-1}{k-1}}_{\text{rest}} + (n-1-k) \binom{n-1}{k-1} = \binom{n}{k}$$
$$\Rightarrow \binom{n}{k} = (k+1) \binom{n-1}{k} + (n-k) \binom{n-1}{k-1}$$

base cond<sup>n</sup>:  $\binom{0}{0} = 1$

### Harmonic Number ( $H_n$ ) :

$$H_n = \frac{1}{1} + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n}$$

- (৭) 52 টা card এবের উপর এক table এর এক বৈণবিক্ষয় রাখলে বলুন দূর সর্বক- extend কোথা যাবে ?



$$d_{n+1} = ?$$

### Worm on the rubber band :

w → অতি sec  $\rightarrow$  1 cm

rubber  $\rightarrow$  100 cm

w 1 cm  $\xrightarrow{\text{stretch}}$  rubber stretch ২গুণ doubled হওয়া যায়

$$\begin{array}{ccc} 1\% & & 99\% \\ 1 \text{ sec} \rightarrow 1 & \rightarrow 100 \\ 2 \text{ sec} \rightarrow 2 & \rightarrow 200 \\ 3 \text{ sec} \rightarrow 3 & \rightarrow 300 \end{array}$$

$$\frac{H_n}{100} = \frac{1}{100} + \frac{1}{200} + \frac{1}{300} + \dots = \frac{1}{100} \left( 1 + \frac{1}{2} + \frac{1}{3} + \dots \right)$$

CSE-219 → Sec A

chap-(1-6)

compulsory 1 & 6

previous question

exercise → related problem - 23

self study

Sec A - CSE-213

chap-(1-7)

1, 5 full, 6 partial, 7 self study

compulsory 2 & 3

1st final  
2nd

\*1 Maths

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