Example 9 (Worksheet) — RSA Decryption (work the process, not just the answer)

Problem. We receive the ciphertext blocks 0981–0461 produced by the RSA cryptosystem from Example 8. The public key was (n, e) = (2537, 13) with $n = 43 \cdot 59$. Decrypt the message.

The big idea (the why)

RSA works in blocks. If a block c was encrypted with $c \equiv m^e \pmod{n}$, then anyone who knows the *private* exponent d (the inverse of $e \mod \phi(n)$) can recover the plaintext block via

$$m \equiv c^d \pmod{n}$$
.

This is fast thanks to repeated squaring. After recovering each numeric block m, translate back to letters using two digits per letter: $A = 00, \ldots, Z = 25$. Leading zeros matter!

Step 1 — Compute the private exponent d

$$\phi(n) = (p-1)(q-1) = 42 \cdot 58 = 2436,$$
 find d with $13d \equiv 1 \pmod{2436}$.

Extended Euclid gives d = 937 (indeed $13 \cdot 937 = 12181 = 1 + 5 \cdot 2436$).

Step 2 — Decrypt each block with repeated squaring

Block 1: $c = 0981 \Rightarrow c = 981$.

$$m \equiv 981^{937} \pmod{2537}, \qquad 937 = 512 + 256 + 128 + 32 + 8 + 1.$$

Squares mod 2537:

Multiply only the needed entries (powers 1, 8, 32, 128, 256, 512):

$$r \leftarrow 1$$

 $r \cdot 981 \equiv 981$
 $r \cdot 1325 \equiv 981 \cdot 1325 \equiv 1717$
 $r \cdot 441 \equiv 1717 \cdot 441 \equiv 1251$
 $r \cdot 2472 \equiv 1251 \cdot 2472 \equiv 282$
 $r \cdot 1688 \equiv 282 \cdot 1688 \equiv 1292$
 $r \cdot 293 \equiv 1292 \cdot 293 \equiv \boxed{704}$

So $m_1 = 0704 \Rightarrow 07 \text{ 04} = \text{H E}.$

Block 2: $c = 0461 \Rightarrow c = 461$.

$$m \equiv 461^{937} \pmod{2537}, \qquad 937 = 512 + 256 + 128 + 32 + 8 + 1.$$

Squares mod 2537:

Multiply the needed entries (powers 1, 8, 32, 128, 256, 512):

$$r \leftarrow 1$$

 $r \cdot 461 \equiv 461$
 $r \cdot 1261 \equiv 461 \cdot 1261 \equiv 1327$
 $r \cdot 1737 \equiv 1327 \cdot 1737 \equiv 1559$
 $r \cdot 316 \equiv 1559 \cdot 316 \equiv 1122$
 $r \cdot 913 \equiv 1122 \cdot 913 \equiv 82$
 $r \cdot 1433 \equiv 82 \cdot 1433 \equiv \boxed{1115}$

So $m_2 = 1115 \Rightarrow 11 \ 15 = L \ P.$

Step 3 — Read the plaintext

Blocks 0704 1115 translate to HELP.

Tips, tricks, and common pitfalls

• Keep the two-digit mapping straight: $A = 00, \ldots, J = 09, \ldots, Z = 25$. Leading zeros are part of the block!

- Choose the block size so that each four-digit block m is < n. Here 2N = 4 works because 2525 < 2537 < 252525.
- When doing repeated squaring, build a small table of c^1, c^2, c^4, \ldots and then multiply only the powers that add up to d.
- Arithmetic gets easier if you reduce *often*. Every product should be brought back modulo n immediately.

Practice — **Your Turn** (use n = 2537, e = 13, d = 937)

Use the same key as above. Show your exponentiation steps and *keep* leading zeros when converting back to letters.

Problem A (easier). Decrypt the single block 2081. What two letters do you get? Hint:

compute 2081^{937} mod 2537 and then split the result as __ __.

Problem B (similar). Decrypt the two-block ciphertext 2081 2182. Reminder: convert

each four–digit block separately, then map back to letters.

Problem C (harder). The ciphertext 0981 0724 1774 was made with the same key.

- Decrypt all three blocks.
- Translate to letters. If the last block ends with a padding letter X, circle it.

Reflection. In one or two sentences, explain why knowing e and n does not make decryption

easy, but knowing d does.