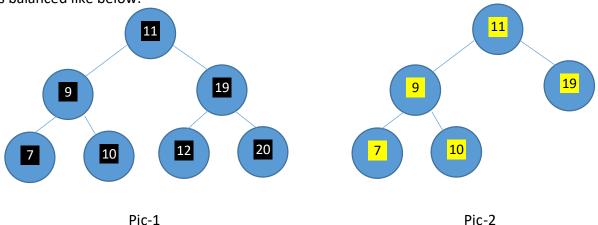
# AVL TREE

# V-1-(Intro):-

AVL TREE is named upon its investor (Adelson-Velsky & Landis) to reduce time complexity and it has self-balancing Algo.

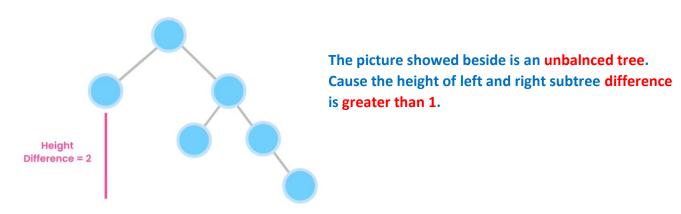
# V-2-(Balanced & UnBalanced Tree):-

Binary Search Tree(BST) has O (log n) time complexity in most of the operation but only if the tree is balanced like below:-

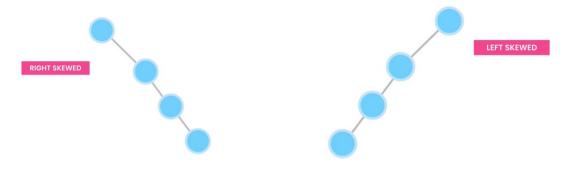


Looked at the picture-1 & picture-2 every node has two child or zero children for leaf node. In BST  $\underline{Height(left\ subtree)}$  -  $\underline{Height(right\ subtree)}$  <=  $\underline{I}$ . That mean we can't have a long branch.

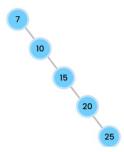
Picture -1 is an example of perfect tree(left and right subtree are full with children).



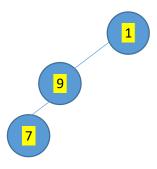
There are more trees. Look at the picture below(worst binary tree looked like linked list):-



If we insert items in ascending order then the tree become right skewed binary tree.



If we insert items in descending order then the tree become left skewed binary tree.



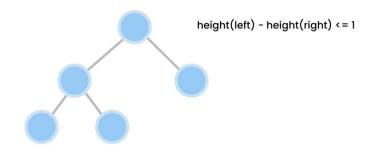
Even if BST (if not sorted) it became unbalanced.

## **SELF-BALANCING TREES**

- AVL Trees (Adelson-Velsky and Landis)
- Red-black Trees
- B-trees
- Splay Trees
- 2-3 Trees

Self-Balancing Trees has property of balancing itself. Actually , there are many types of trees but we don't need to learn them all.

## V-3-(Rotations):-



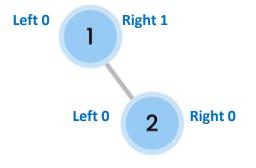
AVL tree is a special type of binary tree what automatically readjust themselves whenever we add nodes by considering that the left and right subtree height difference must be <= 1.

# (4 Types)

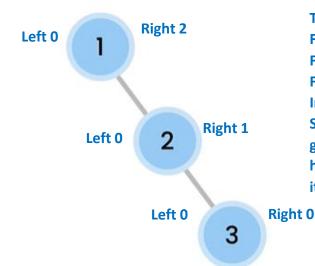
- Left (LL)
- Right (RR)
- Left-Right (LR)
- Right-Left (RL)

.

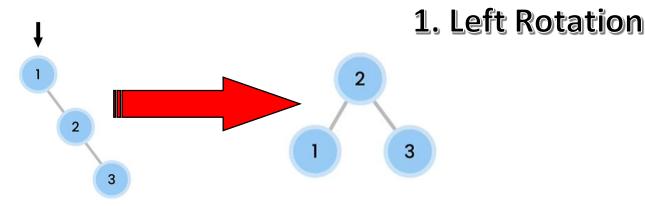
Illustrating imbalance in tree:-



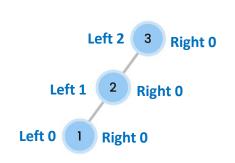
The beside picture tree is a balnce tree cause :For Node 2 , left subtree height = right subtree height = 0
For Node 1, left subtree height = 0 & right subtree height=1
In balance tree left subtree height - right subtree height <=1
So, that satisfy here.



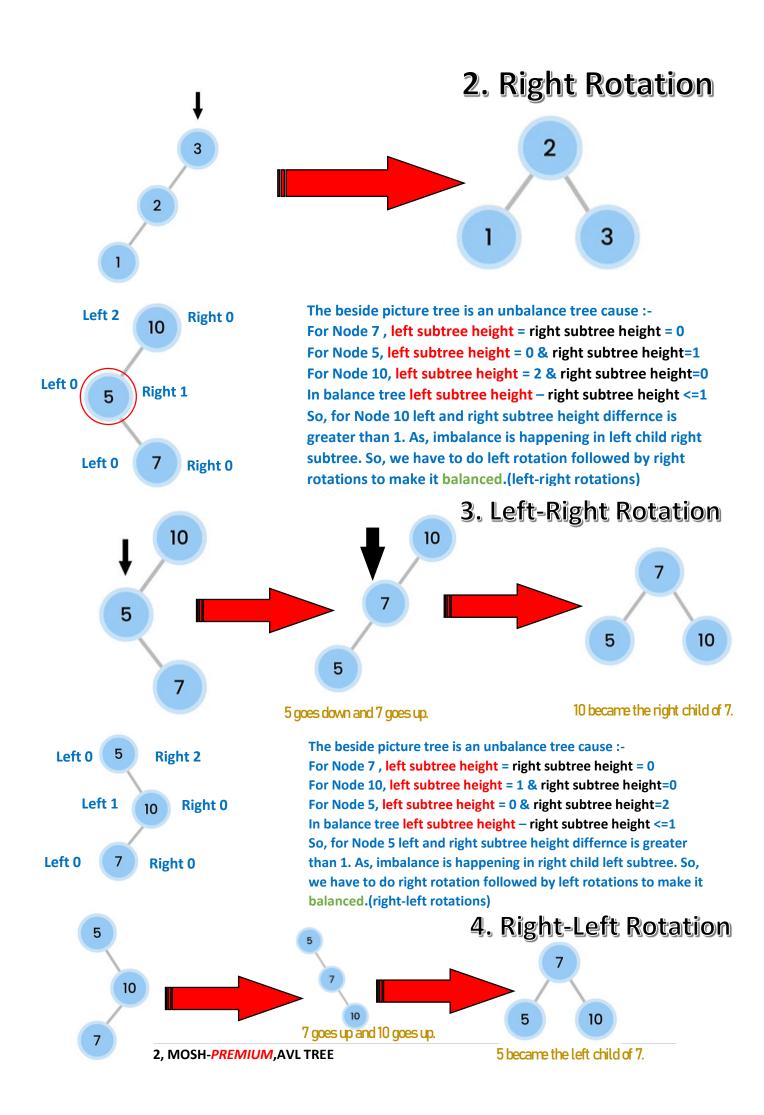
The beside picture tree is an unbalance tree cause:For Node 3, left subtree height = right subtree height = 0
For Node 2, left subtree height = 0 & right subtree height=1
For Node 1, left subtree height = 0 & right subtree height=2
In balance tree left subtree height — right subtree height <=1
So, for Node 1 left and right subtree height differnce is greater than 1. So, the whole tree is unbalance. As this tree is heavy on right side. So, we have to do left rotations to make it balanced.



After Rotations one side goes up other side came down. That means height of one side increases and height of other side decreases.



The beside picture tree is an unbalance tree cause:For Node 1, left subtree height = right subtree height = 0
For Node 2, left subtree height = 1 & right subtree height=0
For Node 3, left subtree height = 2 & right subtree height=0
In balance tree left subtree height — right subtree height <=1
So, for Node 3 left and right subtree height differnce is greater than 1. So, the whole tree is unbalance. As this tree is heavy on left side. So, we have to do right rotations to make it balanced.



N:B:- Easy Thing to identify where to do left-right rotation is if we are doing a left rotation and after that the basic condition of a tree being balanced (left subtree height – right subtree height <=1) not satisfied then we have to do right rotation afterwards. On the other hand, if we are doing a right rotation and after that the basic condition of a tree being balanced (left subtree height – right subtree height <=1) not satisfied then we have to do left rotation afterwards.



# V-4-(AVL TREES):-

Avl Tree Working Simulations(How it works):-

#### 1. Insert 30.

Ans: As it is the first element to be inserted, It is considered as root.

30

#### 2. Insert 15.

Ans: As 30 is already inserted & it is the root node. And 15 is less than 30. So, it should be inserted as the left child of root node 30.

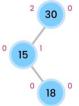


#### 3. Insert 18.

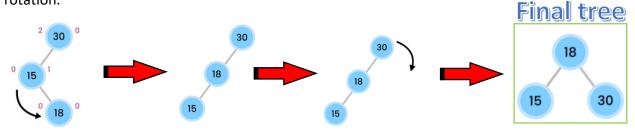
Ans: Now, 15>18<30. So, 18 should be right child of node 15.



But the tree is unbalanced here. Cause, for node 30 <u>left subtree height – right subtree</u> <u>height >1.</u>The picture below has the calculation of height for each node.



As the tree is imbalanced is left child right subtree. So, we have to implement left-right rotation.



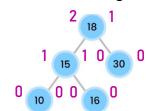
15 goes down and 18 goes up.

As the tree become left –skewd, we have to perform right rotation to make it balanced. So, 30 goes down and became the right child of node 18.



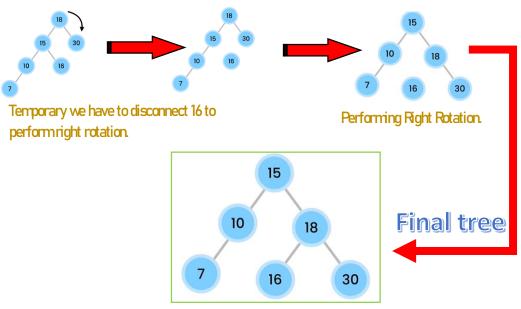
## 5. Insert 16.

Ans: As 16<18 & 16>15. So, 16 should be the right child of node 15. (It's still balanced)

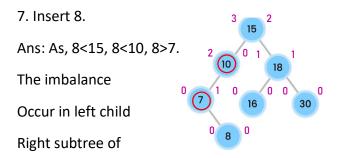


# 6. Insert 7. Ans: As 7<18, 7<15,7<10. Calculating Height and it breaks The balanced condition for node 18. 3-1 = 2 ×1.

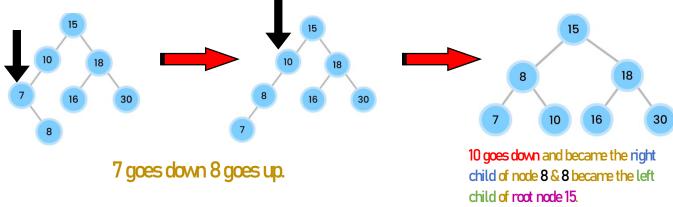
The imbalance occurs in left child, left subtree. So, we have to do right rotation. It is quite tricky.



As 16×15 but 16<18. So, 16 should be the left child of node 18.



root node 15. So, We have to perform left-right rotation.



# V-5-(AVL Rotation Exercise):-

# **AVL Rotations**

Add the following set of numbers to an AVL tree. For each set, start with

an empty tree. Draw the tree at each step and show the type of rotations

required to re-balance the tree. Solutions are on the next page.

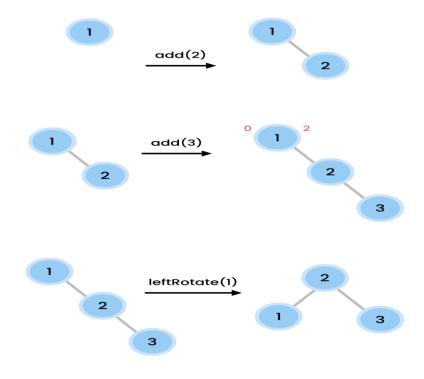
Set 1: (1, 2, 3, 4, 5)

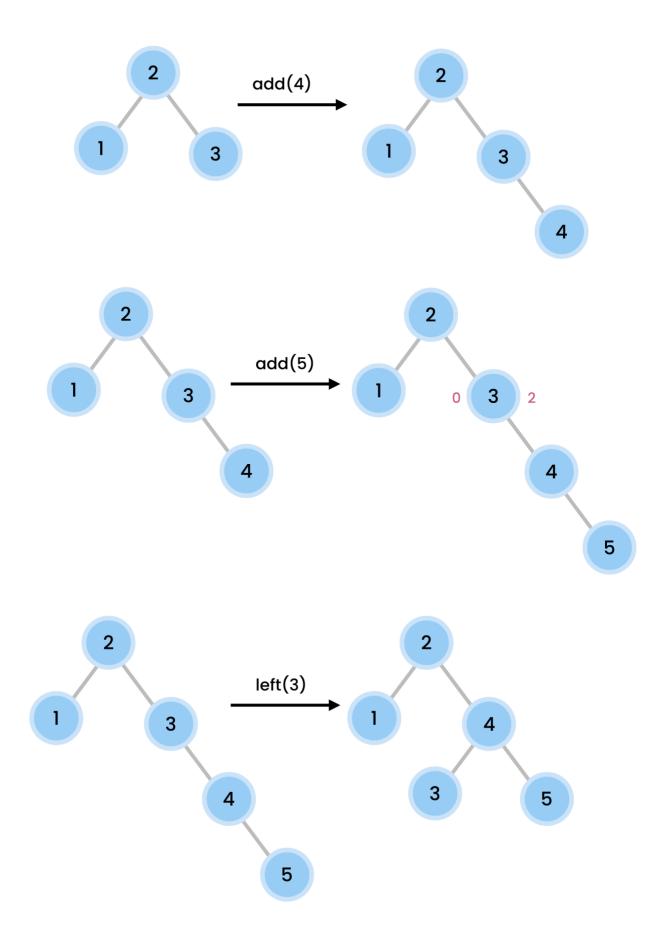
Set 2: (5, 10, 3, 12, 15, 14)

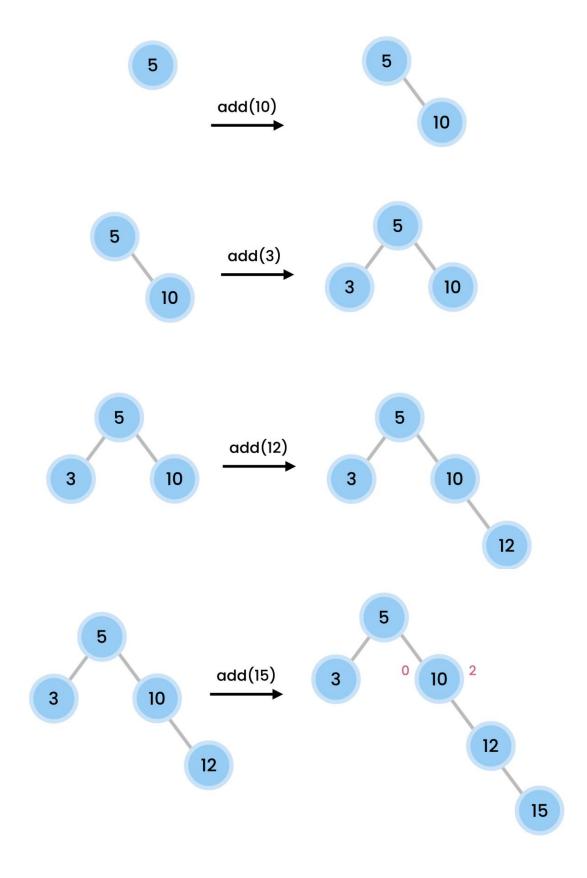
Set 3: (12, 3, 9, 4, 6, 2)

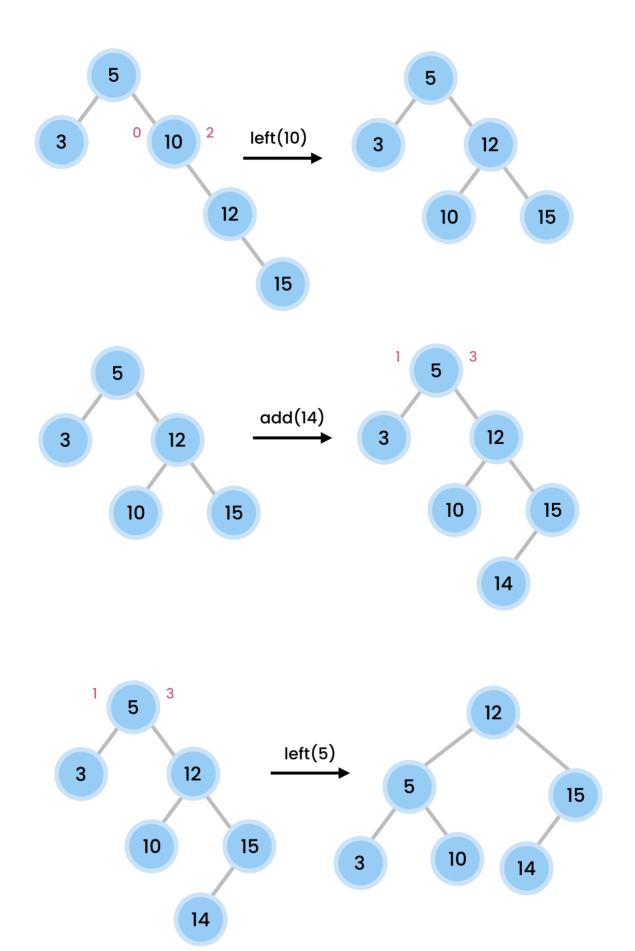
You can use the following tool to visualize an AVL tree: <a href="https://www.cs.usfca.edu/~galles/visualization/AVLtree.html">https://www.cs.usfca.edu/~galles/visualization/AVLtree.html</a>

Solution to Set 1 (1, 2, 3, 4, 5)

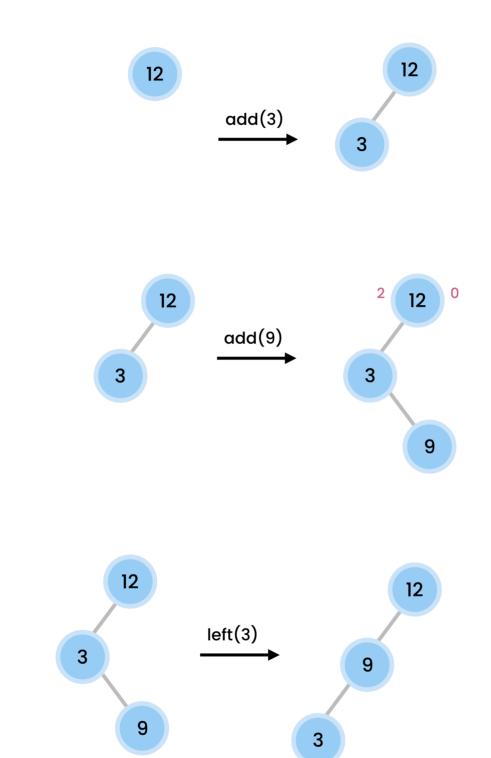


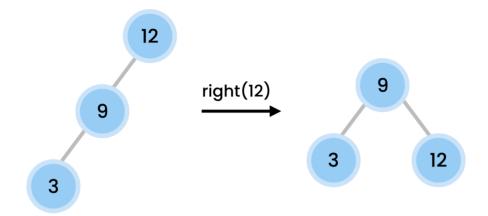


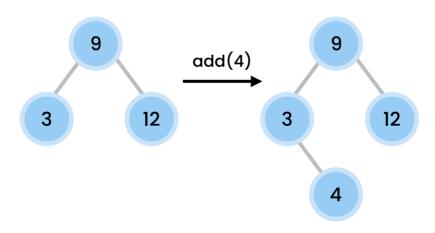


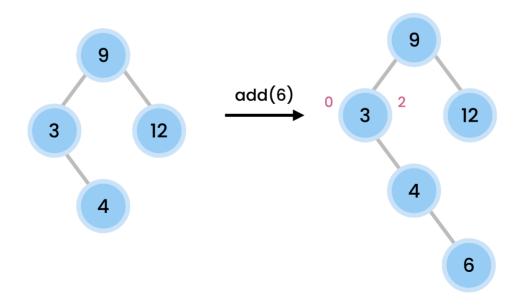


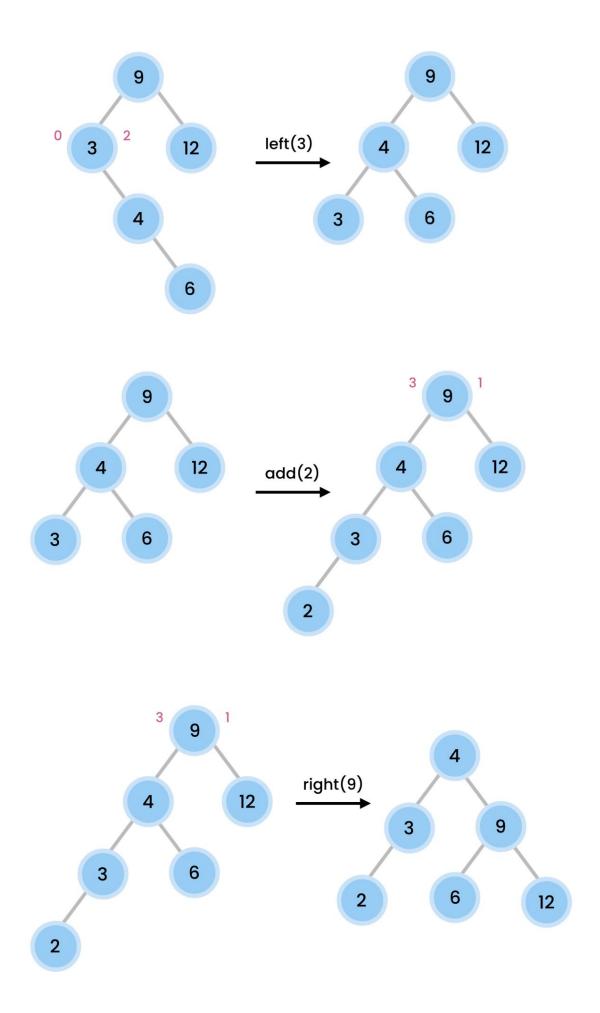
# Solution to Set 3 (12, 3, 9, 4, 6, 2)











# V-6-(AVL Tree Building):-

```
// AVLTree
// AVLNode
// insert() - recursion
```

# V-(7+8)-(Insert Method+height+Basic Class):-

```
public class AVLTree {
    private class AV1Node{
      private int value;
      private AVINode leftchild;
      private AVlNode rightchild;
      public AVlNode(int value) {
          this.value=value;
      @Override
      public String toString() {
          return "Value= "+this.value;
    private AVlNode root;
    public void insert(int value) {
        root = insert(root, value);
    public AV1Node insert(AV1Node root, int value) {
        if(root==null)
            return new AVlNode(value);
        if(value<root.value)</pre>
root.leftchild=insert(root.leftchild, value);
root.rightchild=insert(root.rightchild, value);
        root.height =
Math.max(height(root.leftchild), height(root.rightchild)
        return root;
    private int height(AV1Node node) {
    return (node==null) ? -1: node.height;
```

## V-(10+11)-(Balance tree checking+Updated Codes of AVLTree Class):-

```
public void insert(int value) {
    root = insert(root,value);
}
public AVINode insert(AVINode root,int value) {
    if(root==null)
        return new AVINode(value);
    if(value<root.value)
        root.leftchild=insert(root.leftchild,value);
    else
        root.rightchild=insert(root.leftchild), height(root.rightchild))+1;
    var balancefactor = balancefactor(root);
    if(isLeftHeavy(root))
        System.out.println(root.value+" is leftHeavy");
    if(isRightHeavy(root))
        System.out.println(root.value+" is rightHeavy");
    return root;
}
private boolean isLeftHeavy(AVINode node) {
    return balancefactor(node)>1;
}
private boolean isRightHeavy(AVINode node) {
    return balancefactor(node)<-1;
}
private int balancefactor (AVINode node) {
    // for empty tree it is balanced already
    // otherwise calculate the differences of left and right subtree and return
    return (node==null)? 0:height(node.leftchild)-height(node.rightchild);
}
private int height(AVINode node) {
    return (node==null)? -1: node.height;
}</pre>
```

# V-(12+13)-(Detect Rotation on tree+Updated Codes of AVLTree Class):-

```
public class AVLTree {
    private class AVlNode(
        private int height;
    private int value;
    private AVlNode leftchild;
    private AVlNode rightchild;
    public AVlNode(int value) {
        this.value=value;
    }
    @Override
    public String toString() {
        return "Value= "+this.value;
    }
}

private AVlNode root;
public void insert(int value) {
    root = insert(root, value);
}

public AVlNode insert(AVlNode root, int value) {
    if(root==null)
        return new AVlNode(value);
    if(value<root.value)</pre>
```

```
balance(root);
private boolean isLeftHeavy(AV1Node node){
private boolean isRightHeavy(AV1Node node){
```

# V-(12+13)-(Detect Rotation on tree+Updated Codes of AVLTree Class):-

```
10 root
                                                                 newRoot = root.right
           20 newRoot
                                      10 goes down and became the right
                                                                 newRoot.left = root
              30
                                      child of node 20.
        leftRotate(10)
                                                                10
                                     But if 20 has a leftchild then what?
                                                                      20 newRoot
newRoot = root.right
root.right = newRoot.left
newRoot.left = root
                                                   20
                                              10
                                                       30
                                                            Node 10 (root) right child should be node 20
                                                            (new root) leftchild, then node 20 left child = node 10.
```

## V-(14+15)-(Performing Rotation on tree+Updated Codes of AVLTree Class):-

```
public AV1Node(int value){
  public String toString(){
    root = insert(root, value);
public AVINode insert(AVINode root, int value) {
        return new AVlNode (value);
        root.rightchild=insert(root.rightchild, value);
    return balance(root);
private AV1Node balance(AV1Node root) {
        if (balancefactor(root.leftchild) < 0) {</pre>
       if (balancefactor(root.rightchild)>0) {
        return rotateLeft(root);
```

```
setHeight(newRoot);
private void setHeight(AV1Node node) {
private boolean isLeftHeavy(AV1Node node) {
private boolean isRightHeavy(AV1Node node) {
    return balancefactor(node)<-1;</pre>
private int balancefactor(AV1Node node) {
private int height(AVlNode node){
```

# V-(16)-( AVLTree Exercise):-

Ex-1- Check to see if a binary tree is balanced.

Ans:-

```
public boolean isBalanced() {
    return isBalanced(root);
}

private boolean isBalanced(AV1Node root) {
    if (root == null)
        return true;

    var balanceFactor = height(root.leftchild) - height(root.rightchild);

    return Math.abs(balanceFactor) <= 1 &&
        isBalanced(root.leftchild) &&
        isBalanced(root.rightchild);
}</pre>
```

Ex-2. Is the tree perfect or not?

[hint :- If all parents had exactly two child in a tree then we called it a perfect binary Tree]

[Resource(for more):- <a href="https://www.programiz.com/dsa/perfect-binary-tree">https://www.programiz.com/dsa/perfect-binary-tree</a> ]
Ans:-

```
public boolean isPerfect() {
    return size() == (Math.pow(2, height(root) + 1) - 1);
}
public int size() {
    return size(root);
}

private int size(AVINode root) {
    if (root == null)
        return 0;

    if (isLeaf(root))
        return 1;

    return 1 + size(root.leftchild) + size(root.rightchild);
}
private boolean isLeaf(AVINode node) {
    return node.leftchild == null && node.rightchild == null;
}
private int height(AVINode node) {
    return (node==null) ? -1: node.height;
}
```

# **AVL TREES**

- Balanced and unbalanced
- BST
  - Average: O(log n)
  - Worst: O(n)
- Self-balancing trees
- Rotations: Left, Right, Left-Right and Right-Left