

## Linked List

### V-1) Introduction:-

▶ 1- Introduction

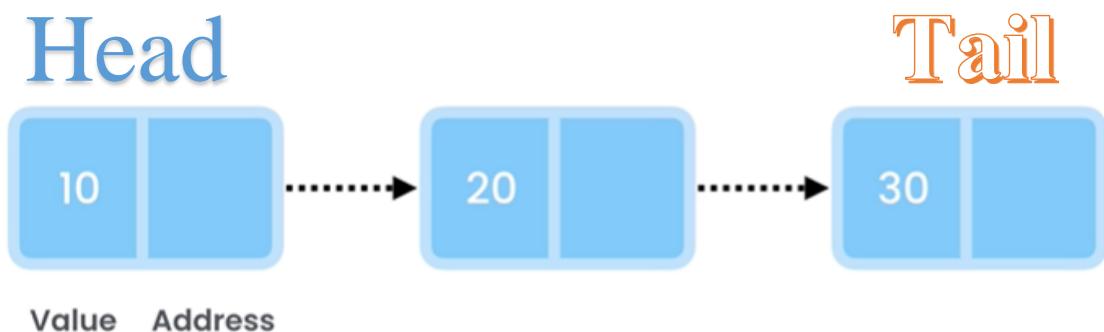
#### IN THIS SECTION

Overview of Linked Lists

Runtime Complexity

Build a Linked List

### V-2) What Are Linked List?



Here Data's 10, 20, 30 are linked with one next to each. Where One Node has value and address of next node that's goanna linked the two data's. That's why this structure of dataset is called Linked List.  
Above List is actually singly Linked List. There are total 8 types of linked lists.

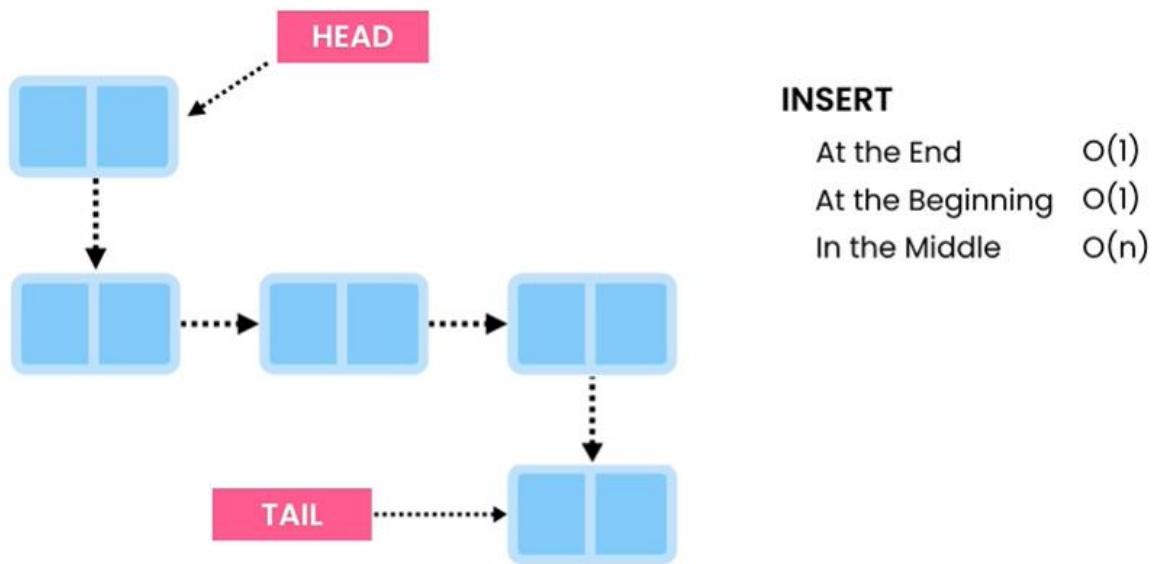
Time Complexities:-

## Scenario 1:-

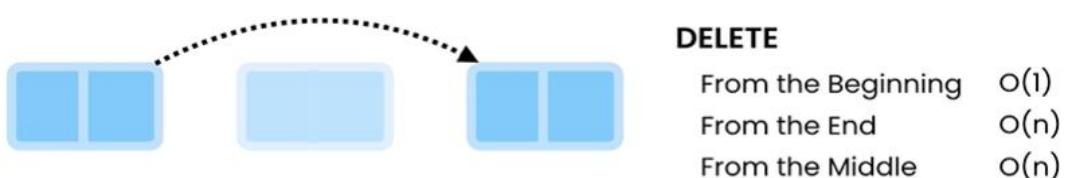


For scenario 1 this linked list cannot direct access like indices of Array. Each node has a reference to a node. So at the worst case scenario for both value and index search might be at the end of the list. So , that's why both cases has a complexity of  $O(n)$ .

## 2. Insertion



## 3. Deletion



## V- 3) Working with Linked\_List

```
import java.util.Arrays;
import java.util.LinkedList;

public class Main {
    public static void main(String[] args) {
        // LinkedList <int> = new LinkedList<int>; //this will give error cause it is primitive we have to write the whole class
        LinkedList list = new LinkedList(); // kisu nah khali create holo list
        list.addLast(10); // [ 10 ]
        list.add(20); // [ 10, 20 ]
        list.addLast(10); // [ 10, 20, 10 ]
        list.add(30); // [ 10, 20, 10, 30 ]
        list.addFirst(50); // [ 50, 10, 20, 10, 30 ]
        System.out.println(list);
        System.out.println(list.size()); // return size of the array means 5
        System.out.println(list.indexOf(10)); // find the first occurrence of the given element which is 1
        System.out.println(list.contains(10)); // return true if found the item in the list , false otherwise
        var array = list.toArray(); // eikhane variable er jaigai onno kisu dile kaj korbe nah
        System.out.println(Arrays.toString(array)); // array ta list na niye print korlam
    }
}
```

## V-(4-12)-working with LinkedList Descriptively:-

### Global Variable in the LinkedList Class:-

```
private Node first;// reference of head
private Node last;// reference of tail
private int size;// reference of the list items exist in the list
```

### Node class which is inside LinkedList class:-

```
private class Node {
    private int value;// Node er value store korbe
    private Node next;// Node er next element er reference rakbe
    private Node(int value){
        this.value = value;
    }
}
```

#### V-4. Add last Method (add new items at the end of the list)

```
public void addLast(int item){  
    // Node node = new Node(); eivabe likhe confused hoar dorkar nai  
    // kenona java compiler dan side dekhei buje nibe node kon type er object  
    //jodi var use kori  
    /*  
     *  
     *var node = new Node();  
     *node.value = item;  
     */  
    var node = new Node(item);  
    if(isEmpty())// jodi list empty thake tahole List er first & last node eki  
        first=last=node;  
    else{  
        last.next = node;//existing last node er pore ei notun node boslo  
        last = node;//last variable jeta node er tail reference kore seta update korlam  
    }  
    size++;  
}
```

#### main class:-

```
public class Main2 {  
    public static void main(String[] args) {  
        var list = new LinkedList();  
        list.addLast(10);  
        list.addLast(20);  
        list.addLast(30);  
    }  
}
```

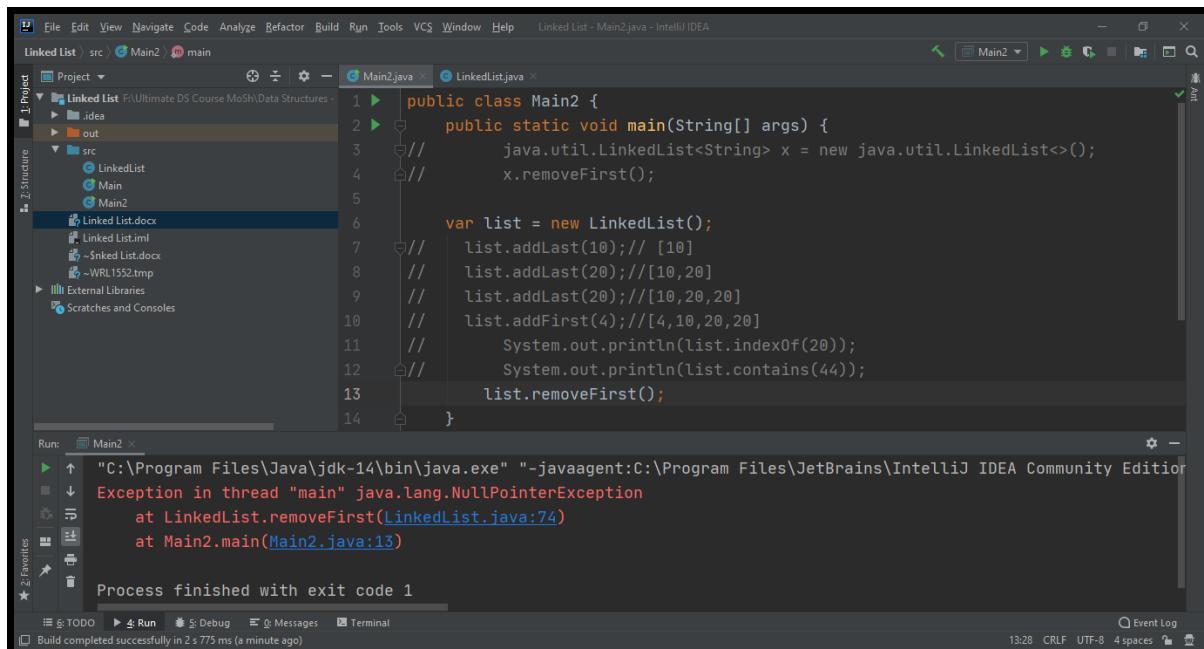
#### V-5-(addFirst method with listempty checking):-

```
public void addFirst(int item){  
    var node = new Node(item);  
    if(isEmpty())  
        first=last=node; //jodi list empty thake  
    else {  
        node.next = first;// first e notun node er sathe first ke link  
        first = node;// then update the first cause head of the list changes  
    }  
    size++;  
}
```

## V-6 -(indexof method)

```
public boolean contains(int item){// ei method e amra indexof method ke reuse  
korlam  
    return indexOf(item)!=-1;// eikhane item khuje paile true return korbe otherwise  
false return korbe  
// indexOf method integer return kore  
// se jodi uporer contains method er call e -1 return na kore mane list e  
// oi item ta ase taile contains method true return korbe  
// ar -1 return korle mane item ta list ei nai tai contains method  
// false return korbe  
}
```

## V-7-(RemoveFirst Method[little complicated])



```
Main2.java  
1 public class Main2 {  
2     public static void main(String[] args) {  
3         // java.util.LinkedList<String> x = new java.util.LinkedList<>();  
4         // x.removeFirst();  
5         var list = new LinkedList();  
6         list.addLast(10); // [10]  
7         list.addLast(20); // [10, 20]  
8         list.addLast(20); // [10, 20, 20]  
9         list.addFirst(4); // [4, 10, 20, 20]  
10        System.out.println(list.indexOf(20));  
11        System.out.println(list.contains(44));  
12        list.removeFirst();  
13    }  
14 }
```

Run: Main2  
Exception in thread "main" java.lang.NullPointerException  
at LinkedList.removeFirst(LinkedList.java:74)  
at Main2.main(Main2.java:13)

Process finished with exit code 1

As the list is empty so we got **null pointer exception**.

```

public class Main2 {
    public static void main(String[] args) {
        java.util.LinkedList<String> x = new java.util.LinkedList<>();
        x.removeFirst();

        var list = new LinkedList();
        // list.addLast(10); // [10]
        // list.addLast(20); // [10, 20]
        // list.addLast(20); // [10, 20, 20]
        // list.addFirst(4); // [4, 10, 20, 20]
        // System.out.println(list.indexOf(20));
        // System.out.println(list.contains(44));
        // list.removeFirst();
    }
}

Process finished with exit code 1

```

In java linkedlist built in class we got **NoSuchElementException** instead of **Nullpointer** exception. So, we need to implement our code in the way its built in java packages.

[Built in java linkedlist removeFirst() call:-

```

java.util.LinkedList<String> x = new java.util.LinkedList<>();
x.removeFirst();
]

```

So, our final modified method for removefirst is:-

```

public void removeFirst(){
    if(isEmpty()) throw new NoSuchElementException();
    if(first==last) //list e jodi ekta element thake
        first = last = null;
    else{
        /*
        ekon amar onek sabdhane korte hobe remove
        ex [10->20->30]
        ami chailei 20 take first banai diye first.next null
        korte parbo nah ete pura list harai jabe
        */
        var second = first.next;//ex [10->20->30] eikhane 20 er reference save raklam
        first = null;// 10 er sathe 20 er bonding remove korlam
        first = second;// ekon first ke 20 e point korlam ekon list hobe [20->30]
    }
    size--;
}

```

## V-8-(RemoveLast() method)-[complex]

```

public void removeLast(){
    // list jodi empty hoi tahole remove korar kisui thakbe nah sekehetre exception
    dekhabo
    if(isEmpty()){
        throw new NoSuchElementException();
    }
    //list e only ekta element thakle jevabe remove korbo
    if(first==last)
        first=last=null;
    else {
        // [10,20,30]
        // the logic below written is only working for list with atleast two node
        var previous = getPrevious(last); // previous = 20
        last = previous; // last ekon 30 theke 20 ke point korlo
        last.next = null;// 20 er sathe 30 er bondhon chinno kore dilam (garbage
        collector)
    }
    size--;
}

```

## V-9-(Implementing Size)

```

public int size(){
    return size;
}

```

## V-9-(Cheat Sheet Gone Wrong If we memorize specially for Job Interview)

### Data Structures

Data Structure	Time Complexity								Space Complexity	
	Average				Worst					
	Indexing	Search	Insertion	Deletion	Indexing	Search	Insertion	Deletion		
Basic Array	O(1)	O(n)	-	-	O(n)	O(n)	-	-	O(n)	
Dynamic Array	O(1)	O(n)	O(n)	O(n)	O(1)	O(n)	O(n)	O(n)	O(n)	
Singly-Linked List	O(n)	O(n)	O(1)	O(n)	O(n)	O(n)	O(1)	O(1)	O(n)	
Doubly-Linked List	O(n)	O(n)	O(1)	O(n)	O(n)	O(n)	O(1)	O(1)	O(n)	
Skip List	O(log(n))	O(1..O(n))	O(1..O(n))	O(1..O(n))	O(1..n)	O(n)	O(1..n)	O(1..n)	O(n log(n))	
Hash Table	-	O(1)	O(1)	O(1)	-	O(n)	O(n)	O(n)	O(n)	
Binary Search Tree	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	O(n)	O(n)	O(n)	O(n)	
Cartesian Tree	-	O(1..O(n))	O(1..O(n))	O(1..O(n))	-	O(n)	O(n)	O(n)	O(n)	
B-Tree	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	
Red-Black Tree	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	
Splay Tree	-	O(log(n))	O(log(n))	O(log(n))	-	O(log(n))	O(log(n))	O(log(n))	O(n)	
AVL Tree	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	O(n)	

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## V-10-(Array & LinkedList Comparison)

### SPACE

- Static arrays have a fixed size
- Dynamic arrays grow by 50-100%
- Linked lists don't waste memory
- Use arrays if you know the number of items to store
- `new ArrayList(100)`

### ARRAYS

### LINKED LISTS (Singly List Only)

#### Lookup

By Index	$O(1)$	$O(n)$
By Value	$O(n)$	$O(n)$

#### Insert

For Shifting elements	Beginning/End	$O(n)$	$O(1)$
	Middle	$O(n)$	$O(n)$

#### Delete

For Shifting elements	Beginning	$O(n)$	$O(1)$
	Middle	$O(n)$	$O(n)$
	End	$O(n)$	$O(n)$

For traversing the whole list

For traversing the whole list

## V-15-(Types of Linked List)

### SINGLY

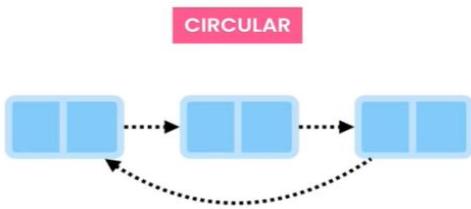


Delete from the End  $O(n)$

### DOUBLY



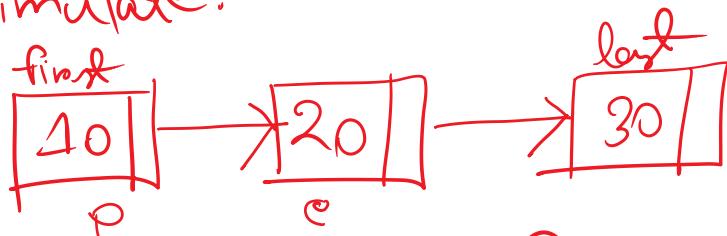
Delete from the End  $O(1)$



Circular Linked List needed when we build something called music player . Cause at the end of the playlist it needs to start from the beginning .

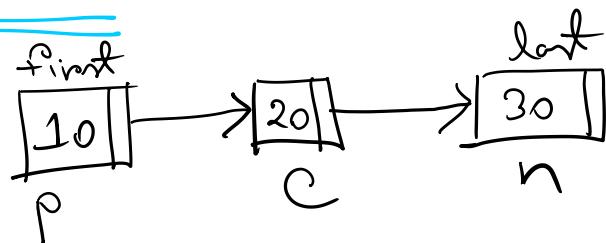
Ex-1 Reversing a singly linked list.  
 (Hottest problem for beginners)  
 even senior developers make it wrong sometimes.

# lets simulate:-



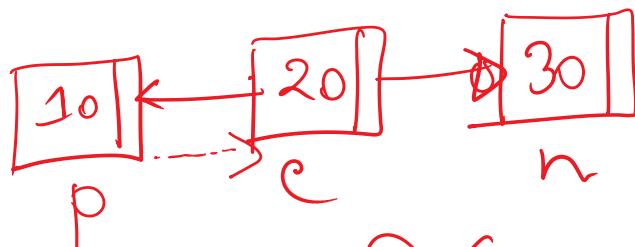
যদি ইলে reference দিতে হ্যাতে তার  
 গুরুত্ব এ কোনওভাবে পর্যবেক্ষণ। কোনো এক  
 স্থানে লিঙ্কগুলি প্রাপ্তি হাতে।

# মাঝে লিখা :-



এতে বলে singly linked list এর ফোন লিংক বা element হারাবে না।

// c.next = p করলে ট্যুপ্য



অর্থাত লিংক পরিষ্কার হবে।

কিন্তু যদি গুরুত্ব remove করে না এবং সেবা

জোরে first = first;

গুরুত্ব first.next = null;

p, c, n ও node first

পরিষ্কার করলে এস্ট

first change করাবে  
সম্ভব কী করা

```

public void reverseList(){
    //List ta inplace reverse korte hobe
    // [10->20->30]
    //hote hobe [30->20->10]
    //onek somossa
    // 10-> 20-> 30
    // p   c   n
    // n = c.next diye reference rakbo
    // link change korbo ekon c.next = p;
    if(isEmpty()) return;
    var previous = first;// p
    var current = first.next;// c
    while (current!=null){
        var next = current.next;// n
        current.next = previous;// link er direction change
        previous = current;
        current = next;
    }
    last = first;
    last.next = null;
    first = previous;
}

```

EX-2) Find kth node from the end. –[Popular interview question about **SLL**]

If List Has 10-> 20->30->40->50

And if we pass the value **3** for k it should return **30**. It should be solved through 1 call . That means we can't iterate through the whole list.

Ans:- We can solve this problem using 2 pointers.

**10-> 20 ->30 ->40 ->50**

**P1**                   **P2**

**10-> 20 ->30 ->40 ->50**

**P1**                   **P2**

*The idea is that when the 2<sup>nd</sup> pointer reached the end of the list it should find the expected value as in the 1<sup>st</sup> pointer. So, pointer should be selected properly and also handle the exceptions.*

Lets solve this:- [very much complicated]

```
public int getKthfromtheEnd(int k){  
    if(isEmpty()) throw new IllegalStateException();  
    //setting two pointers  
    // 10 -> 20 -> 30 -> 40-> 50  
    // a  
    // b  
    // if k = 3  
    var a = first;  
    var b = first;  
    //for loop simulation  
    // 10 -> 20 -> 30 -> 40-> 50  
    // a  
    //     b  
    // 10 -> 20 -> 30 -> 40-> 50  
    // a  
    //         b  
    for(int i=0;i<k-1;i++){// this loop will set the distance between two pointers  
        b= b.next;  
        if(b==null) throw new IllegalArgumentException();  
    }  
    // while loop simulation  
    // 10 -> 20 -> 30 -> 40-> 50  
    //     a  
    //             b  
    // 10 -> 20 -> 30 -> 40-> 50  
    //     a  
    //                 b  
    while (b!=last){  
        a=a.next;  
        b=b.next;  
    }  
    return a.value;  
}
```

## LINKED LISTS

- Second most used data structures
- Grow and shrink automatically
- Take a bit more memory

## RUNTIME COMPLEXITIES

### Lookup

By Index  $O(n)$   
By Value  $O(n)$

### Insert

Beginning/End  $O(1)$   
Middle  $O(n)$

### Delete

Beginning  $O(1)$   
Middle  $O(n)$   
End  $O(n) / O(1)$

Ex-3-exclusive) **Find the middle of a linked list in one pass.-**

```
public void printMiddle() {  
    if (isEmpty())  
        throw new IllegalStateException();  
    var a = first;  
    var b = first;  
    while (b != last && b.next != last) {  
        b = b.next.next;  
        a = a.next;  
    }  
    if (b == last)  
        System.out.println(a.value);  
    else  
        System.out.println(a.value + ", " + a.next.value);  
}
```

Ex-4) if There is any loop or not in a SLL?

```
public boolean hasLoop() {  
    var slow = first;  
    var fast = first;  
    while (fast != null && fast.next != null) {  
        slow = slow.next;  
        fast = fast.next.next;  
        if (slow == fast)  
            return true;  
    }  
    return false;  
}
```

Exercise 3 and 4 explanation is given below:-

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# Linked Lists

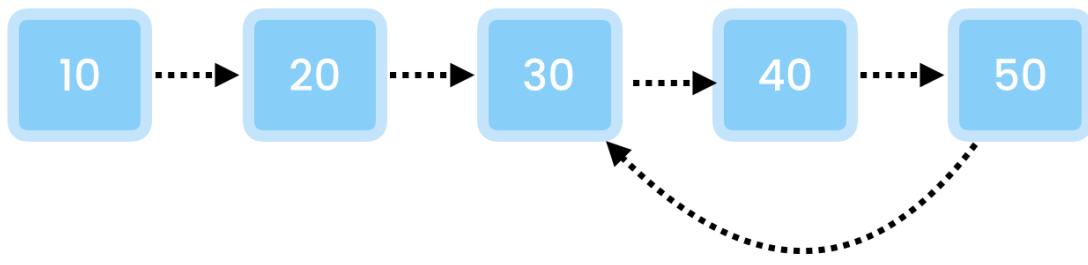
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## Exercises

1- Find the middle of a linked list in one pass. If the list has an even number of nodes, there would be two middle nodes. (Note: Assume that you don't know the size of the list ahead of time.)

**Solution: `LinkedList.printMiddle()`** - You can read my analysis of the solution in the next page.

2- Check to see if a linked list has a loop.



Hint: use two pointers (slow and fast) to traverse the list. Move the slow pointer one step forward and the fast pointer two steps forward. If there's a loop, at some point, the fast pointer will meet the slow pointer and overtake it. Draw this on a paper and see it for yourself. This algorithm is called *Floyd's Cycle-finding Algorithm*.

**Solution: `LinkedList.hasLoop()`** - You can read my analysis later in this document.

# Solutions

## 1- Find the middle of a linked list in one pass.

As always, we start by simplifying and narrowing the problem. Let's imagine the list has an odd number of nodes.

Since we have to find the middle node in one pass, we need two pointers. The tricky part here is that we should figure out how many nodes should the first and second pointers be apart. Let's throw a few numbers to find the relationship between these pointers.

Number of Nodes	Middle Node
1	1
3	2
5	3
7	4
9	5
11	6

Do you see a pattern here? In every row, the number of nodes is increasing by two where as the position of the middle node is increasing by one. Agreed?

So, we can define two pointers that reference the first node initially. Then, we use a loop to move these pointers forward. In every iteration, we move the first pointer one step and the second pointer two steps forward. The moment the second pointer hits the tail node, the first node is pointing the middle node.

Now, let's expand our problem. What if the list has an even number of nodes?

Number of Nodes	Middle Node
2	1,2
4	2,3
6	3,4
8	4,5
10	5,6

We see the same pattern. In every step, the number of nodes increases by two whereas the position of the middle node is increasing by one. The only difference is that here we need to return two middle nodes. That's easy. Once we find the left middle node, we'll also return the node next to it.

How do we know if the list has an even or odd number of items? We can declare a count variable and increment it in each step. But we don't really need this. If the list has an even number of nodes, at the end of the last iteration, the second pointer will reference the tail node; otherwise, it'll be null. (Try a few examples and you'll see this yourself.)

Here's how we can implement this algorithm:

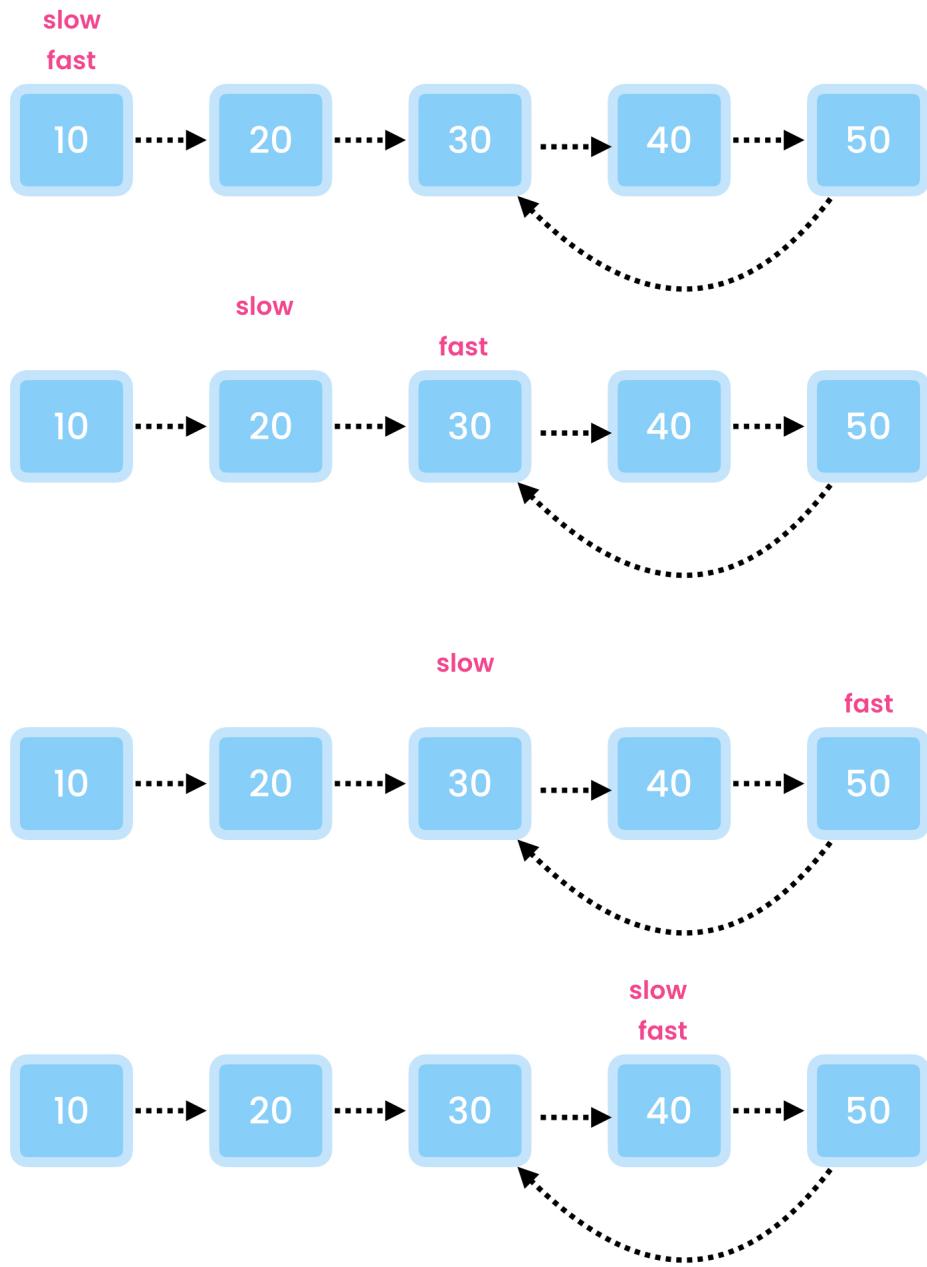
```

1:  var a = first;
2:  var b = first;
3:  while (b != last && b.next != last) {
4:      b = b.next.next;
5:      a = a.next;
6:  }
7:
8:  if (b == last)
9:      System.out.println(a.value);
10: else
11:     System.out.println(a.value + ", " + a.next.value);

```

This algorithm works for any lists with at least one element. For an empty list, we need to throw an **IllegalStateException**. You can see the complete solution in **LinkedList.printMiddle()**.

## 2- Check to see if a linked list has a loop.



You can find the implementation of this algorithm in **LinkedList.hasLoop()**. To test this:

```
var list = LinkedList.createWithLoop();
System.out.println(list.hasLoop());
```

Academic

Notes

CSE

220

(2)

## Linked List (Singly)

2) public class Node {

    Object element; // value to store

    Node next; // link to next node

    public Node (Object e, Node n) {

        element = e;

        next = n;

}

2) creating a list :- (Singly)

(type-1) Node head = null;

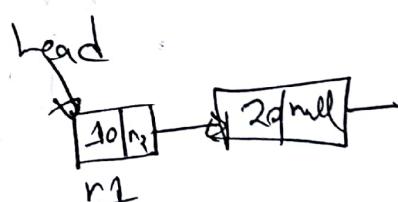
Node n1 = new Node ("10", null);

Node n2 = new Node ("20", null);

n1.next = n2;

~~n1 = head~~

head = n1;



③

### Type-3 (Singly list creation)

Node head = null;

[নোড তৈরি, লিস্ট তৈরি করা আগে দেখতে হব।  
node তৈরি করে ফিরুন্ত না।]

// শিখুন যদি যেভাবে last node তৈরি

create করতে গারুব।

Node n4 = new Node("40", null);

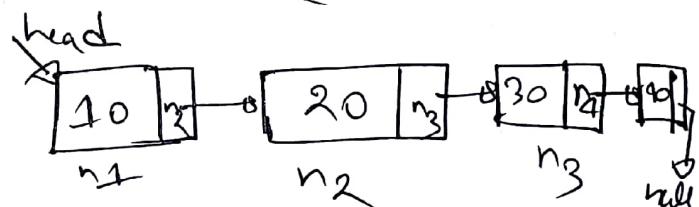
Node n3 = new Node("30", null);

Node n2 = new Node("20", n3);

Node n1 = new Node("10", n2);

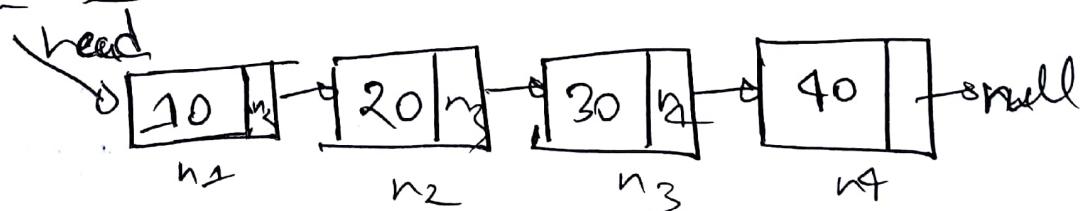
// Assigning the head reference

head = n1;



Manipulation :-

Q) List iteration :- (Singly)



for(Node n= head; n!=null; n=n.next){

// Do whatever we want to do

}

8

## 8/ countNode :- (Singly list)

```
int count = 0;
for(Node n = head; n != null; n = n.next) {
    count++;
}
return count;
```

## 9/ Getting a specific Node :- (singly list)

- যদি কোন specific index এর node এর value return করা।  
 public Object get (Node head, int index) {
 // এইটি random access list এ নথি। তাই অব্যবহৃত  
 // iteration use করে করে যেতে হবে। [index valid  
 // শর্ত নিয়ে]

```
int c = 0;
```

```
for(Node n = head; n != null; n = n.next) {
    if (c == index) // index এর পাস  

        return n.element(); // node এর value  

    c++;
}
```

return -1; // কোন অস্থিতি index দিলে  
 -1 return করে দিব।

৬// NodeAt(); // কোনো প্রকার নোডের কোনো মাধ্যমে

তা দুটি এর একটা।

//random access use a iteration.

public Node NodeAt(Node head, int size, int index)  
//index এর validation check করব।

if (index < 0 || index >= size) {

System.out.println("invalid index");

return null;

int c = 0;

for (Node n = head; n != null; n = n.next) {

if (c == index) {

return n; // এভাবে "iterate করে  
দ্বিতীয় গুরুত্ব

c++;

}

return null; // যদিও index ফির

null return করব।

↓

123456789

## 9// set function(): (singl)

(যোগ করে নিয়ে)  $\rightarrow$  specific node  $\sqsubset$  user থেকে প্রেরণ

element insert করব।

```
public void setNewElement(Node head, int index,  
                           //validation of index  
                           if (index < 0 || index >= size) {  
                               System.out.println("Invalid index");  
                           }  
                           int c = 0;  
                           for (Node n = head; n != null; n = n.next) {  
                               if (c == index) {  
                                   n.element = elem; // element ব্যাপার  
                                   break; // insert করে ফিরব  
                               }  
                               c++;  
                           }  
                       }
```

b) search function (Singh) পাঠ একটি element

একটি ট্রেন ক্ষেত্র লিস্ট (LRU)

//এই অস্থির index এলেমেন্টটি আছে নেটা এবং এক্ষয়

public int indexOf(Node head, Object elem) {

int c = 0;

```
for (Node n = head; n != null; n = n.next) {  
    if (n.element.equals(elem)) {  
        return c; // এই node এর node  
    } // এই এলেমেন্ট  
    c++; // আছে গুরুত্বপূর্ণ  
}
```

return -1;

// কাউন্ট মোড়ে লিস্ট তে আছে কি-না দেখো

public boolean contains(Node head, Object elem) {

boolean contain = false;

```
for (Node n = head; n != null; n = n.next) {  
    if (n.element.equals(elem)) {
```

contain = true; // পাই গোল  
break;

}

return contain; // true বা false দিবিবি

অন্তর্যামী গুরুত্বপূর্ণ ফাংশন

## Q/ insertion in Linked List :- (Single)

9

(1) List এর প্রথম element insert করা।

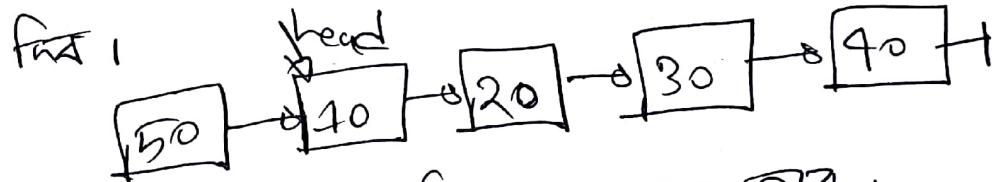
(2) List middle/end

Q. first  
point  
simulation

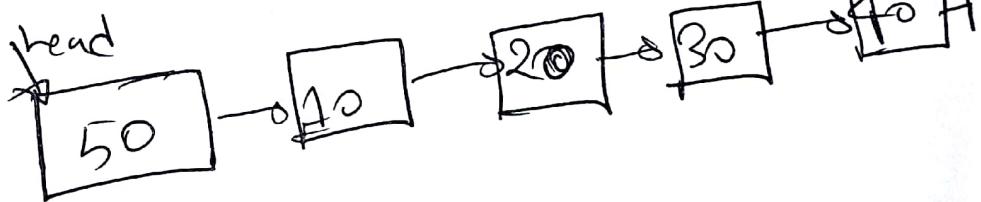
insert 50 0th index

(1) node create [50] +

(II) এই নতুন node next & head  $\Rightarrow$



(III) head গোলাটে হবে।



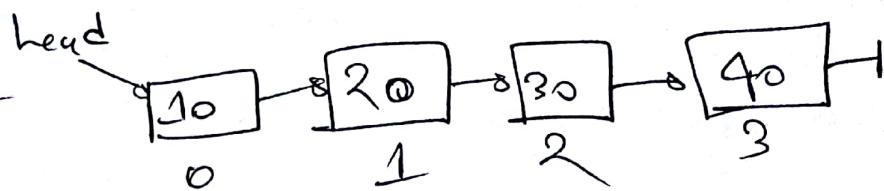
# code logic :-

Node n = new Node()

P.T.O.

(5)

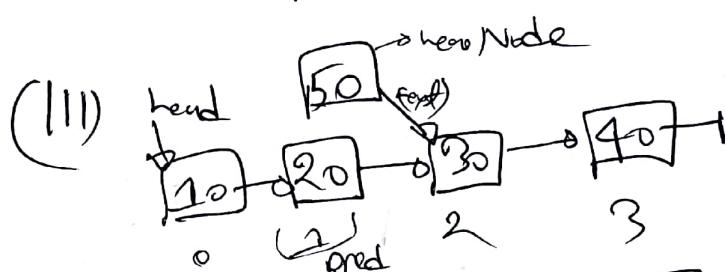
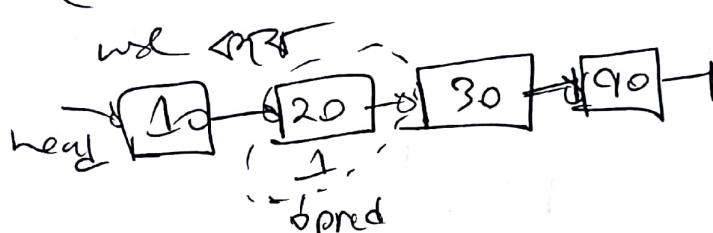
12:

~~insert node~~~~in the middle~~~~on end  
of the list~~~~simulation~~

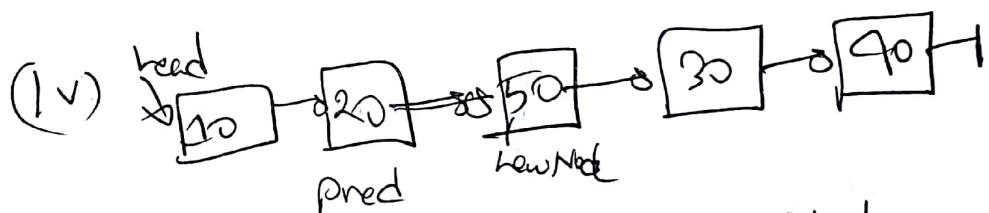
insert 50 @ 2 :-

(i) create node [50]

(II) আবেগ নোড একে পরে সম্ভাল



50 এর next [20] পরে [20] এর next  
এসলাম, 20 টি list হাবে না



pred Node এর next [20] nextNode এর  
প্রদর্শন করা হবে।

ফলস্বরূপ কাঠ ঘোষ।

P.T.O

# code of inserting new nodes in last (Singly) (2)

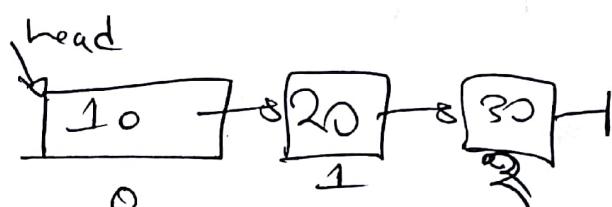
```
public Node insert (Node head, int size; Object elem,  
                    int index)  
{  
    if (index < 0 || index >= size){  
        System S.O.P ("Invalid index");  
    }  
    Node newNode = new Node (elem, null);  
  
    if (index == 0) {  
        newNode.next = head;  
        head = newNode;  
    }  
    else {  
        Node pred = NodeAt (index - 1);  
        newNode.next = pred.next;  
        pred.next = newNode;  
    }  
    return head;  
}
```

## // Removing Nodes (Singly)

20

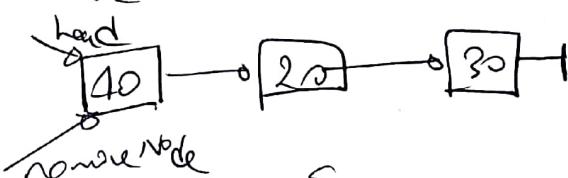
- (1) 0 index থেকে remove করত এটি (remove)
- (2) যদি এই কোড রিভিউ remove,

(i) simulation  
removeNode

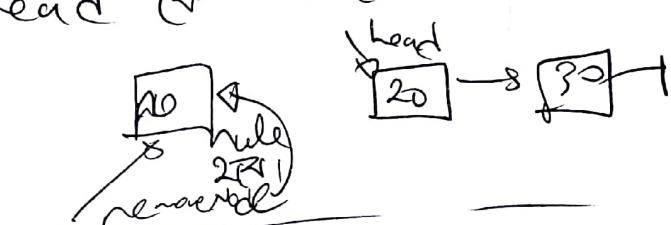


(a) এখন removeNode = null কোড কৈমনি

(2) removeNode হল head নোডের সঙ্গে



(3) Head কে পুনরাবৃত্ত করা।



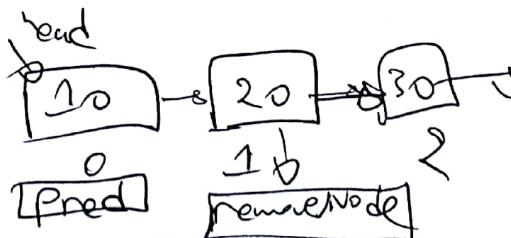
(ii) simulation

removeNode



predecessor Node (পূর্ব নোড) হিসেবে

(2)



p.7.0

(3)



[ Object garbage collection ]

... .

## 22 Code of deletion of nodes (Singly) :-

```
public Node remove (Node head, int size, int index)
{
    // index validation check
    if (index < 0 || index > size)
        S.O.P (" Invalid index ");
    else
        return null;
}

Node removeNode = null;

if (index == 0)
    removeNode = head;
    head = head.next;
}

else
    Node pred = nodeAt (index - 1);
    removeNode = pred.next;
    pred.next = removeNode.next;
    removeNode.next = null;
    removeNode.element = null;
}

return head;
```

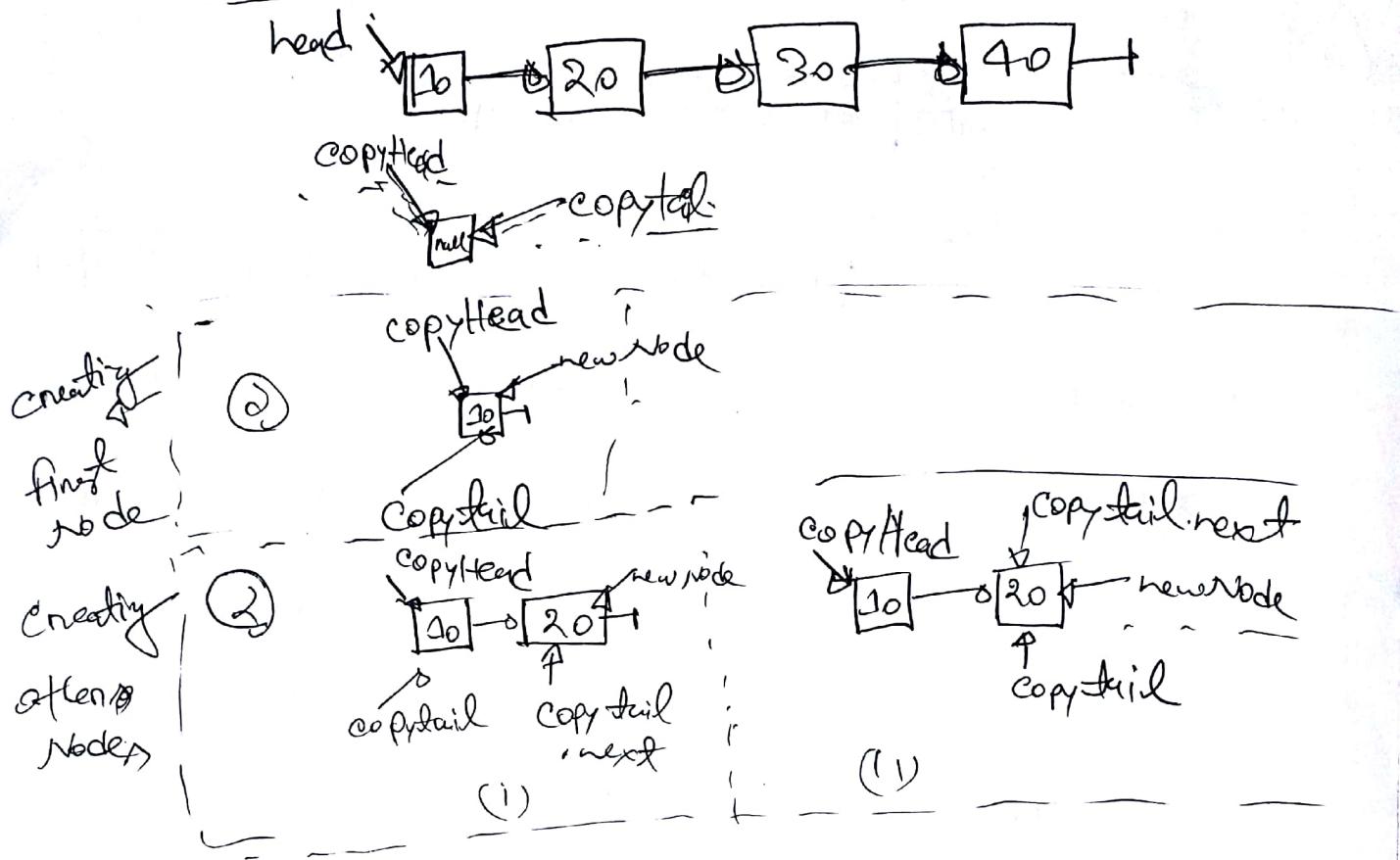
(23)

22

22

22/

## Copying a list



public Node copyList (Node head) {

    Node copyHead = null;

    Node copyTail = null;

    for (Node n = head; n != null; n = n.next)

        Node newNode = new Node(n.element, null);

        if (copyHead == null) {

            copyHead = newNode;

            copyTail = newNode;

        } else {

            copyTail.next = newNode;

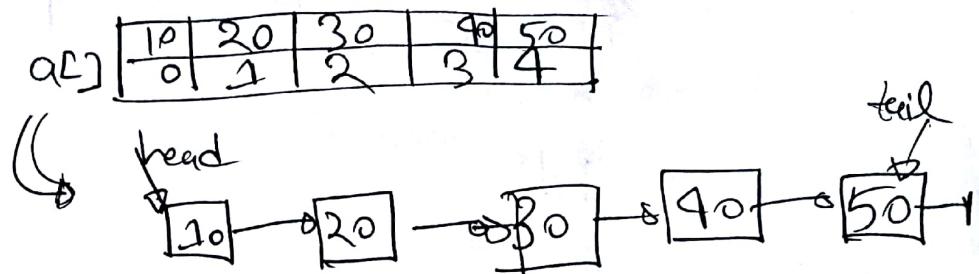
            copyTail = newNode;



# 20/1 Creating a <sup>Singly</sup> List from an array :-

(28)

(copylist  
ক্ষেত্র মাত্র)



(1) 0th index থেকে শুরু করে List এর creation

শুরু মুখ্য এবং কোন index এ গোলে copy করা হবে না।

public Node createList (int [] a) {

Node head = null;

Node tail = null;

for (int i=0; i < a.length; i++) {

Node newNode = new Node (a[i], null);

if (head == null) {

head = newNode;

tail = newNode;

}

else {

tail.next = newNode;

tail = newNode;

}

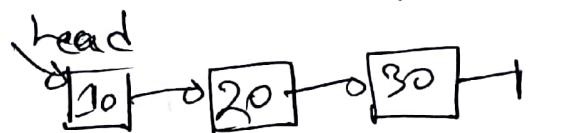
return head;

## ~~28~~ # Reversing a List :-

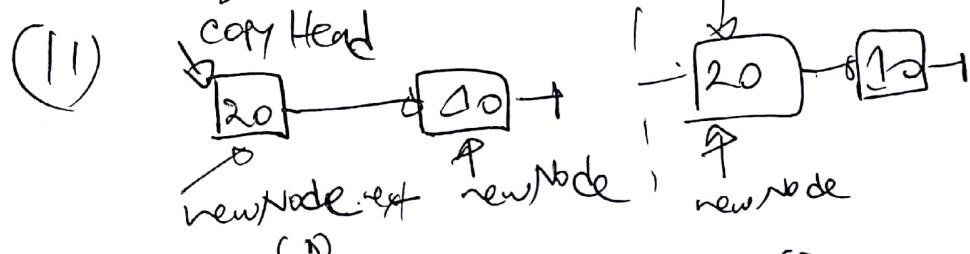
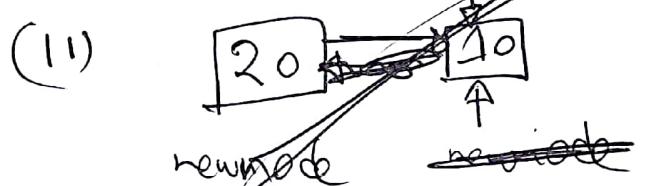
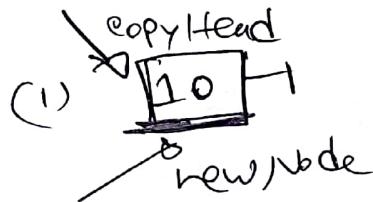
↳ (I) ~~inplace~~ → given list we to reverse

(II) ~~outplace~~ → creating a new list  
in reverse.  
and reverse.

(II) Outplace :- (নতুন বাইকে reverse way  
list create করলি)



(I)  
→ reverse ~~বাইকে~~



public Node reverseList (Node head) {

    Node copyHead = null;

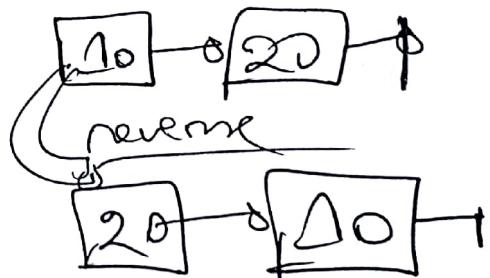
    for (Node n = head; n != null; n = n.next) {

        Node newNode = ~~new~~ Node;

P. 70

~~#Code~~ (reversing a list (outplace) : (try) ②)

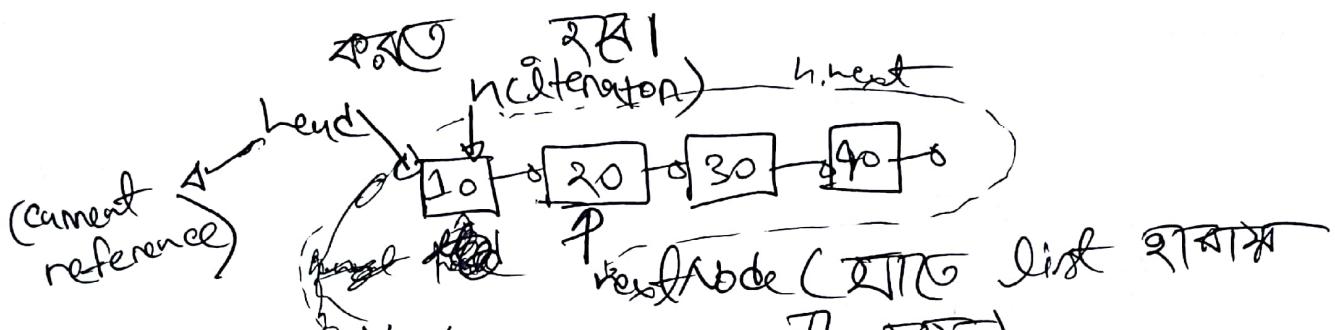
```
public Node reverseList (Node head) {  
    Node copyHead = null;  
    for (Node n = head; n != null; n = n.next) {  
        Node newNode = new Node (n.element, null);  
        if (copyHead == null) {  
            copyHead = newNode;  
        } else { // if Node is not copyHead  
            newNode.next = copyHead; // newNode is now the head  
            copyHead = newNode; // copyHead update  
        }  
    }  
    return copyHead;
```



## #(1) Reversing a List (Singly)

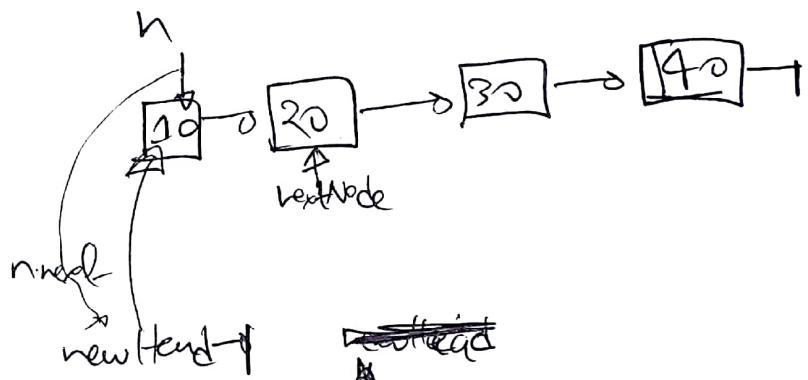
In place approach :-

এখন list পরিষ্কার করে তার মাধ্যমে reverse

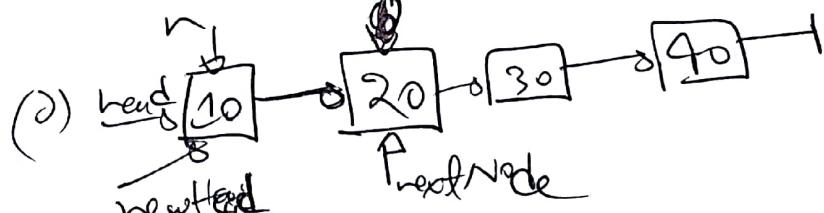


simulation

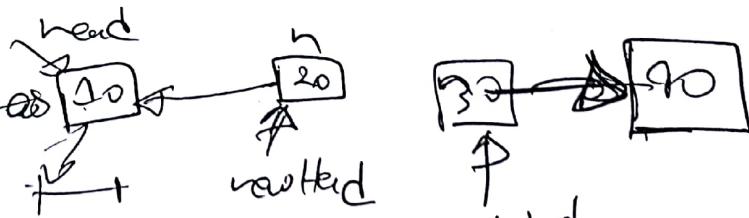
(1)



(2)



(3)



(4) newHead



#code 1

(26)

Publie Node reverseList (Node head) {

Node newHead = ~~new~~ null;

Node n = head;

while (n != null) {

// এখনকার প্রথম node nextnode এর reference  
পাইব।

Node nextnode = n.next;

// ~~newHead~~ newHead কিছু n এর স্থানে  
পাওয়া যাবে।

n.next = newHead; // 

newHead = n; // 

n = nextNode; // n কে এখন ব্যবহার করে

}

এখন ফির আসো

ব্যবহার করে

[n.next কি হবে  
last node]

return newHead;

}

~~#Code-3~~ (Reverse a singly list) &  
~~Alternative~~ (in-place Approach)

public Node reverseList (Node head) {

    Node current = head;

    Node previous = null;

    Node next = null;

    while (current != null) {

        next = current.next;

        current.next = previous;

        previous = current;

        current = next;

}

    return previous;

}

Simulation

Current

previous ~~→ null~~, next → null,

(1)

Current



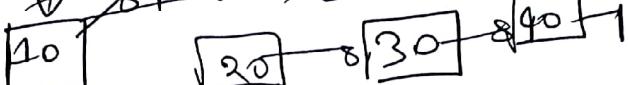
next

$$\text{next} = \text{current.next}$$

→ ΔP  
loop

(11)

previous

previous ~~→ null~~

current

current

previous

null

(111)

next

null

→ ΔP  
loop

(1)

previous

current

null

(11)

previous

current

null

(111)

current

next

null

(14)

current

next

null

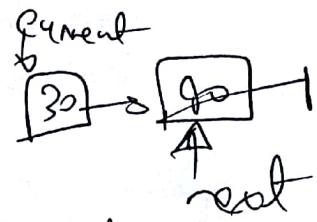
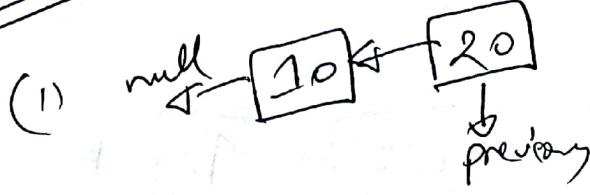
(14)

previous

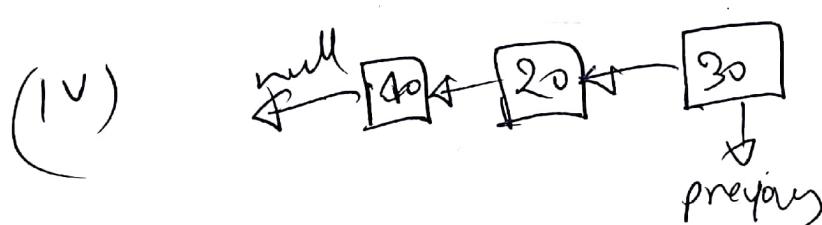
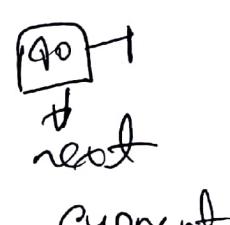
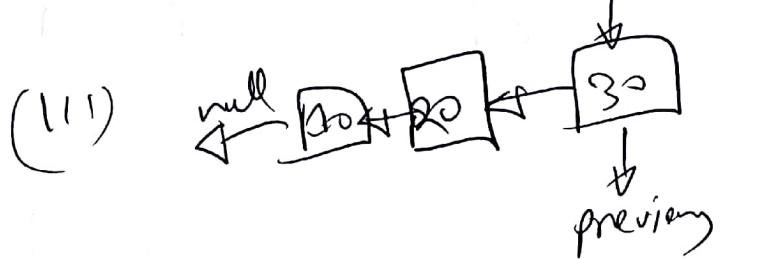
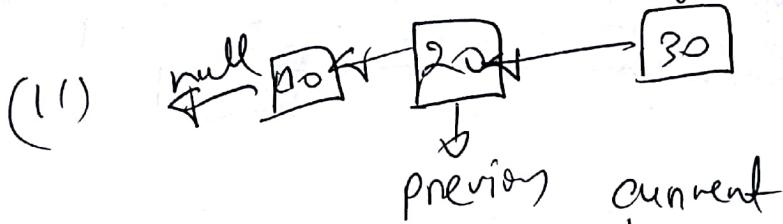
current

null

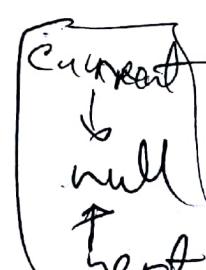
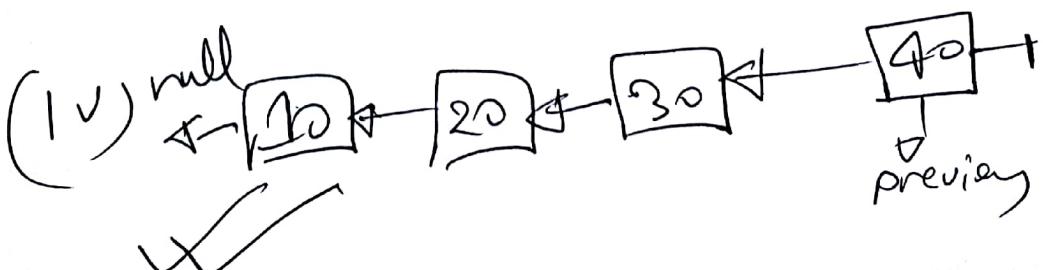
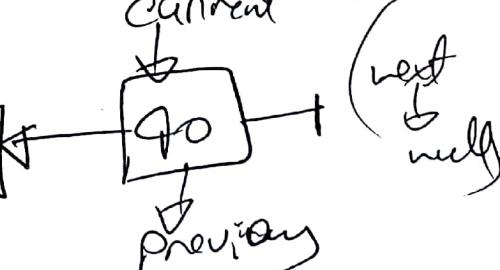
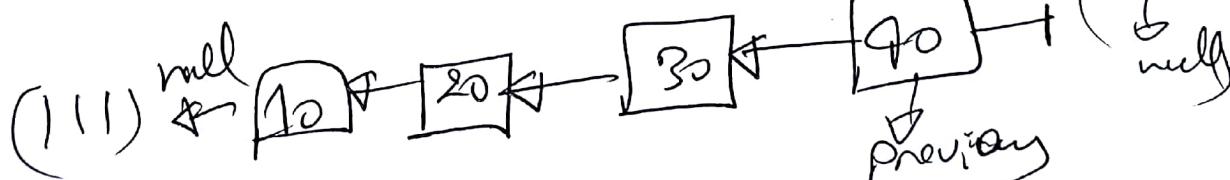
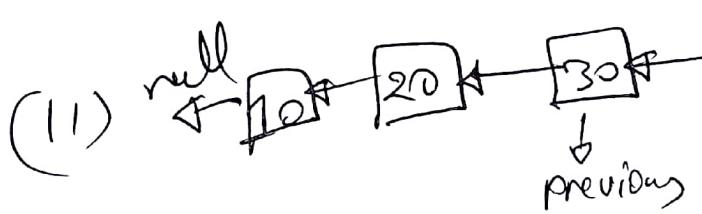
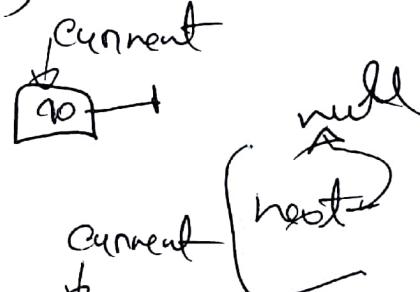
~~921 573~~



②

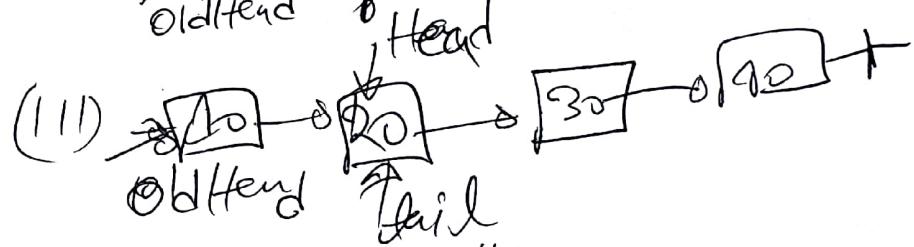
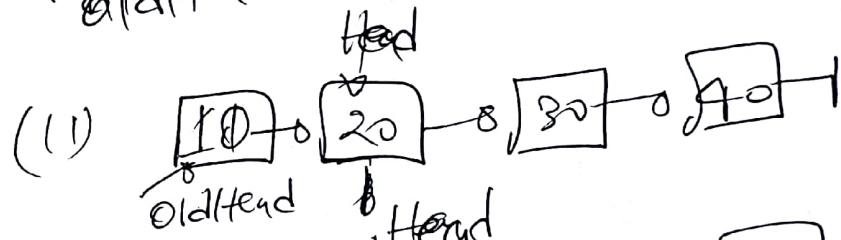
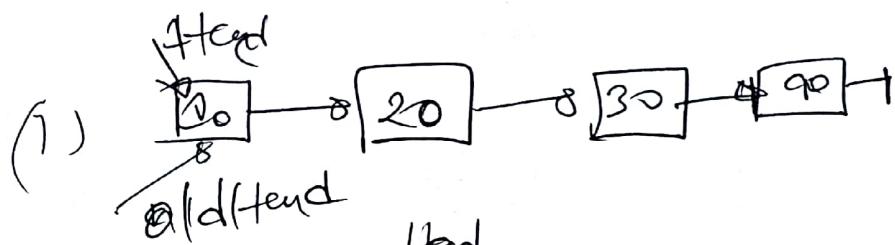
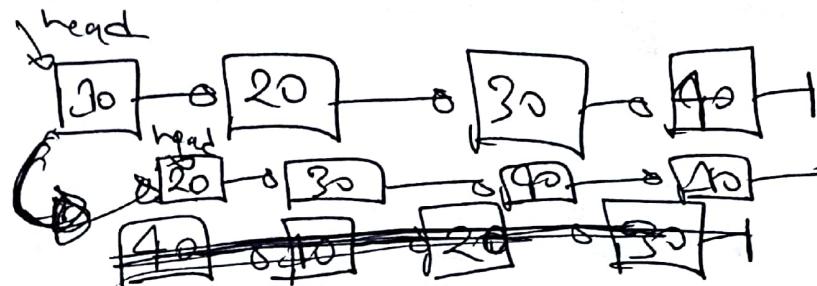


~~921 573~~

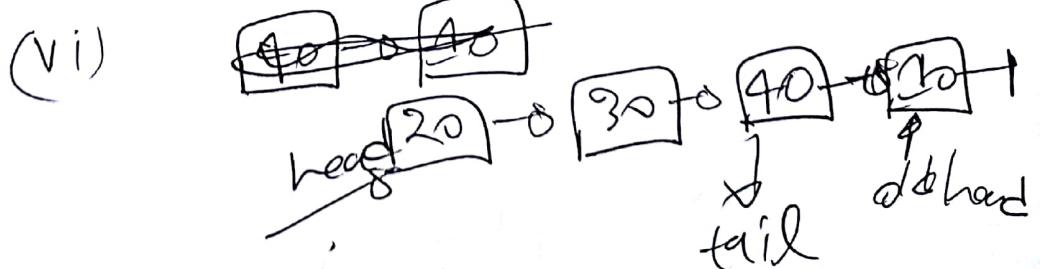
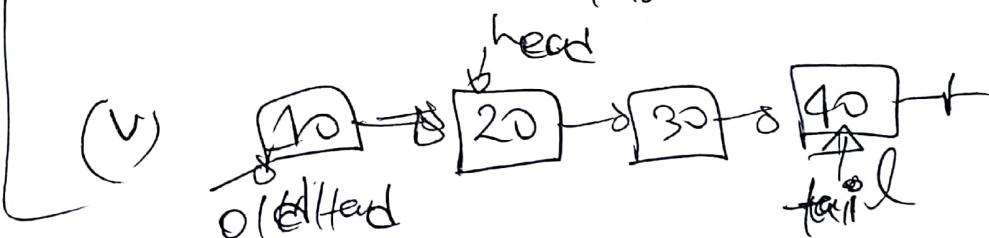
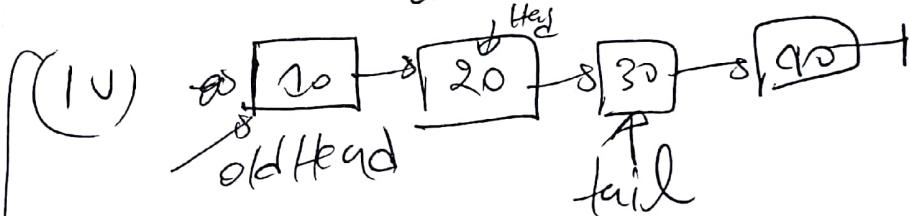


## 20 note left (singy) :-

(2)



while  
loop



public Node rotateLeft (Node ~~head~~  
int size) {

if (size == 0) {

s.o.p ("Nothing do. ~~Node~~");  
return null;

}

Node oldHead = head;

head = head.next;

Node tail = head;

while (tail.next != null) {

tail = tail.next;

}

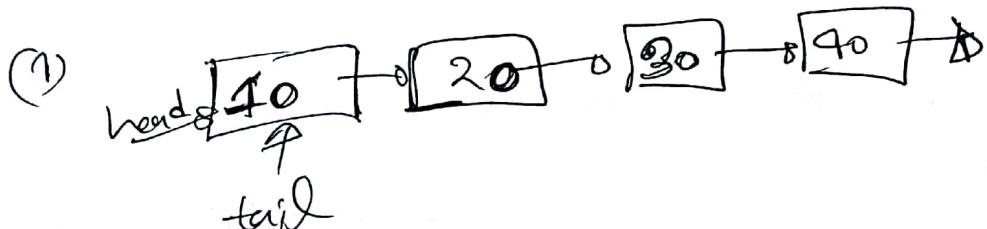
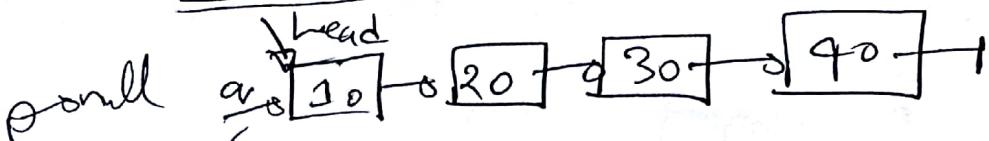
tail.next = oldHead;

oldHead.next = null;

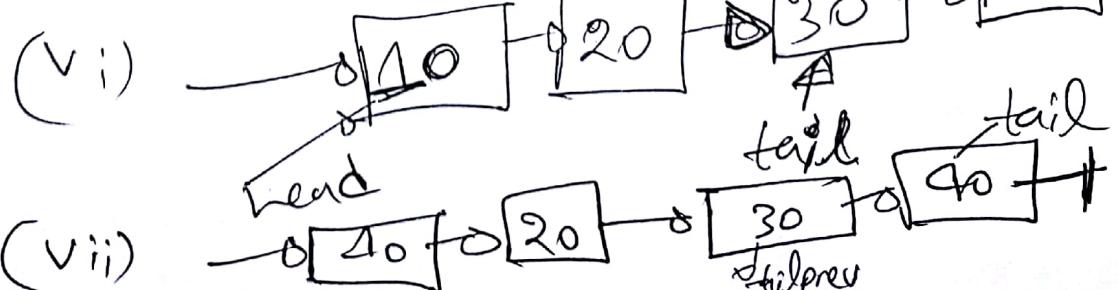
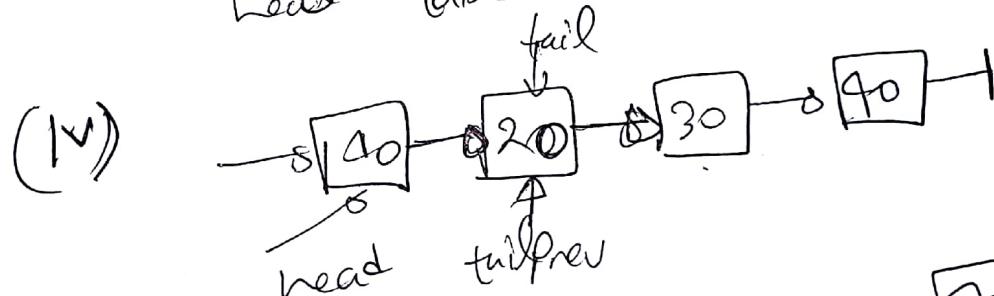
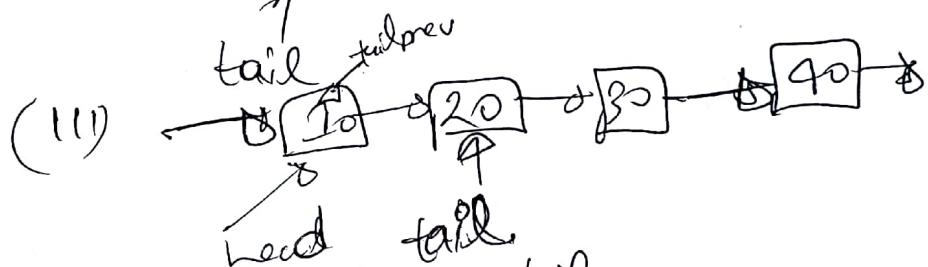
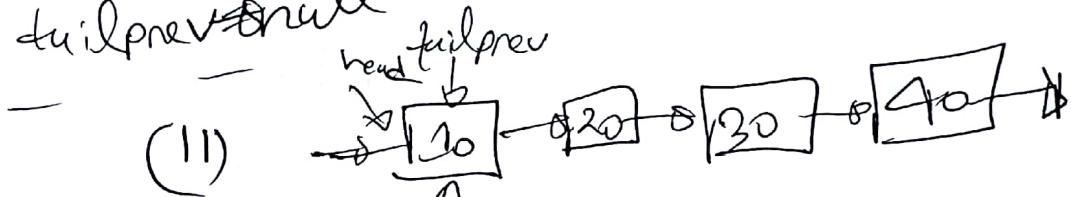
return head;

}

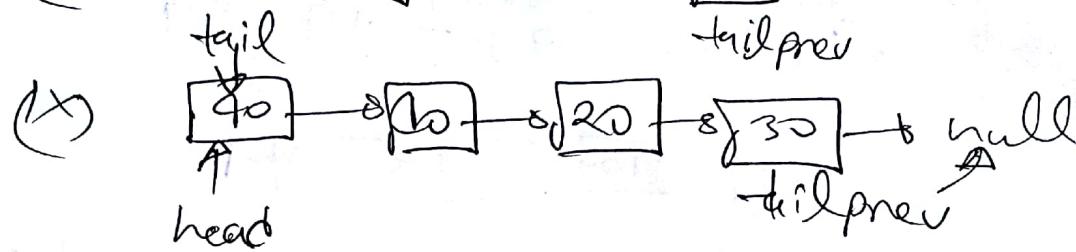
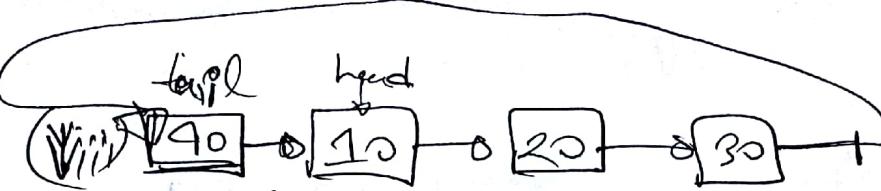
25 rotate Right (Singly) :-



wt  
 0/1  
 -  
 tailprev null



(22)



# Code RotateRight :-

```
public Node RotateRight(Node head){
```

```
    Node tail = head;
```

```
    Node tailprev = null;
```

```
    while(tail.next != null) {
```

```
        tailprev = tail;
```

```
        tail = tail.next;
```

```
}
```

```
tail.next = head;
```

```
head = tail; // tail is head
```

```
tail.prev.next = null; // tail
```

```
return head;
```

```
}
```

30 গুরু  
root head =  
null  
tail  
prev

## Linked Lists - (2)

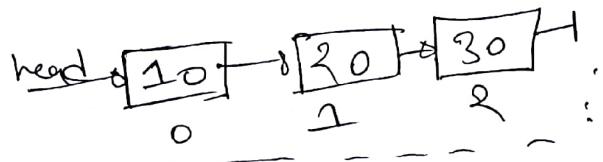
4

### Various types

total 8 types (eight)

Singly

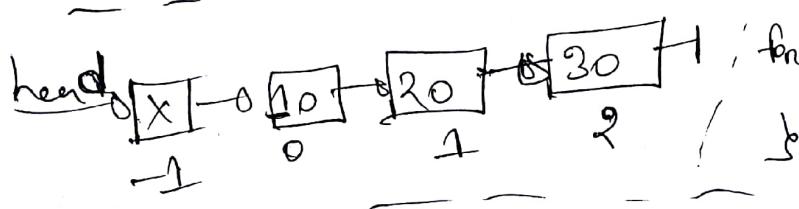
1// Singly Linked List :-



; iteration

for (Node n = head; n != null; n = n.next)

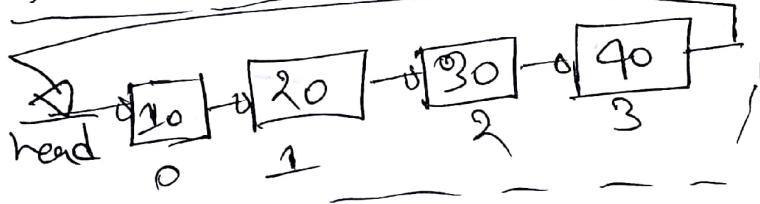
2// Dummy Headed singly List :-



; iteration :-

for (Node n = head.next; n != null; n = n.next)

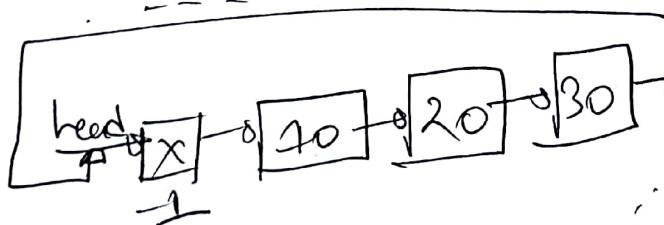
3// Singly Circular Linked List :-



; iteration

for (Node n = head; n.next != head; n = n.next)

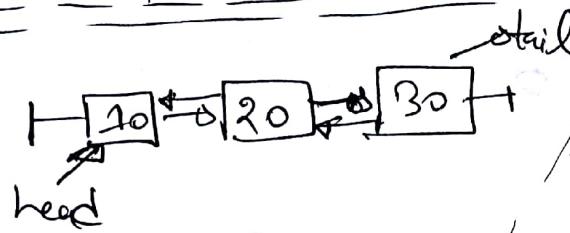
4// Dummy Headed Circular Singly Linked List :-



; iteration :-

for (Node n = head.next; n != head; n = n.next)

### Q/ Doubly Linked List :-



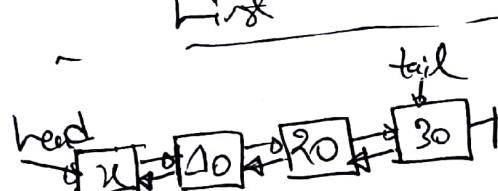
#### (i) Forward iteration:-

```
for(Node n = head; n!=null; n=n.next){}
```

#### (ii) Backward iteration:-

```
for(Node n = tail; n!=null; n=n.prev){}
```

### W/ Dummy Headed Doubly Linked List



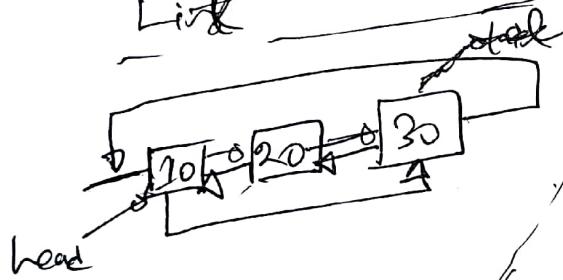
#### (i) Forward iteration:-

```
for(Node n = head.next; n!=null; n=n.next){}
```

#### (ii) Backward iteration:-

```
for(Node n = tail; n.prev!=null; n=n.prev){}
```

### 9/ Doubly Circular Linked List



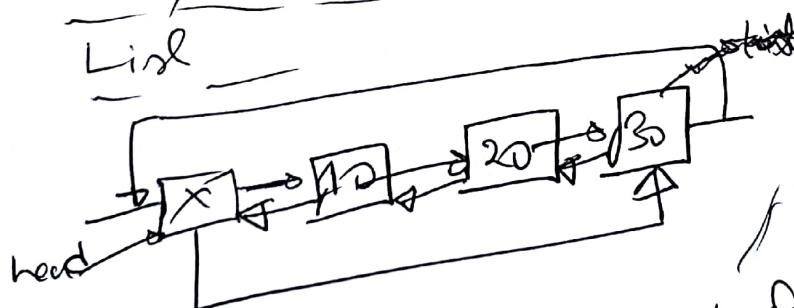
#### (i) Forward iteration:-

```
for(Node n = head; n.next!=head; n=n.next){}
```

#### (ii) Backward iteration:-

```
for(Node n = tail; n!=head; n=n.prev){}
```

### W/ Dummy Headed Doubly Circular Linked List



#### (i) Forward iteration:-

```
for(Node n = head.next; n!=head; n=n.next){}
```

#### (ii) Backward iteration:-

```
for(Node n = head.prev; n!=head; n=n.prev){}
```

# Experimenting with Dummy Headed Doubly circular

Linked List:

public class Node {

    public Object element;

    public Node next;

    public Node prev;

    public Node(Object e, Node n, Node p) {

        element = e;

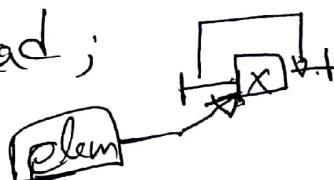
        next = n;

        prev = p;

    } //Creating a list.

    Node head = new Node(null, null, null);

    head.next = head.prev = head;



// Adding first node,

    Node n = new Node("elem", null, null);

    n.next = head.next;

    n.prev = head;

    head.next = n;

    head.prev = n;

(2) node insertion :-

public Node insertAfter(Node p, Object elem)

Node n = new Node (elem, null, null);

Node q = p.next;

n.next = q;

n.prev = p;

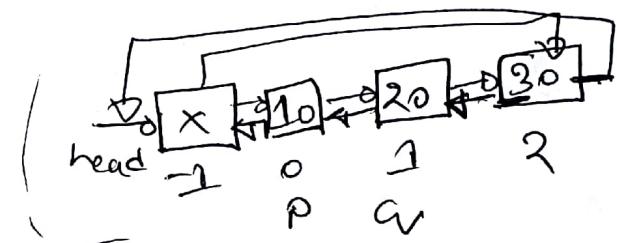
p.next = n;

q.prev = n;

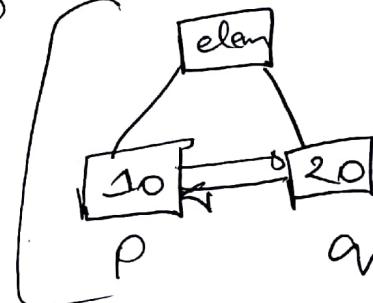
return n;

}

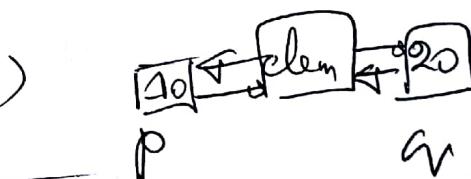
(1) 



(2) 



(3) 



(1) node insertions // inserting specific elements in list @

public ~~Node~~ insertBefore(Node head, Object elem, Object newElement)

Node p = new Node(newElement, null, null);

Node q = null;

for (Node n = head; n !=

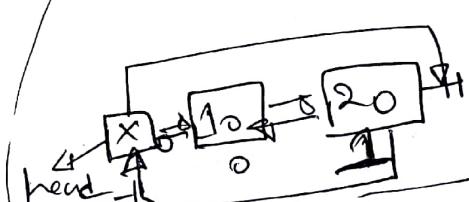
for (Node n = head.next; n != head; n = n.next) {

if (n.element.equals(elem))

q = n;

break;

50  
r



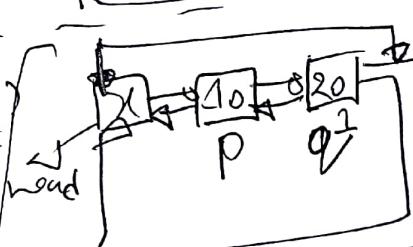
q.prev = p;

n.next = q;

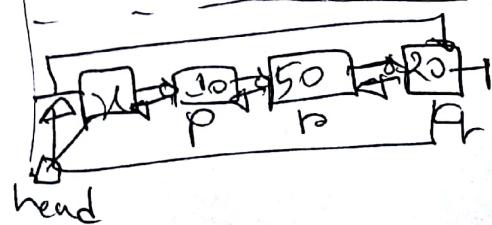
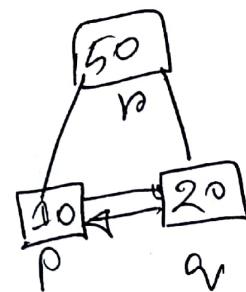
p.prev = p;

p.next = n;

p.prev = r;



(II)



# Remove :-  
public static void removeNode(Node n) {

Node p = n.prev;

Node q = n.next;

p.next = q;

q.prev = p

n.next = n.prev = null;

n.element = null;

}