

1. [Application of SMT solver (that uses DPLL algorithm)] Consider a grid workspace of size 10×10 with $(1, 1)$ being the left-bottom and $(10, 10)$ being the top-right cell. There are 10 robots:

- Initially, robots r_1, \dots, r_5 are placed at $(1, 1), (3, 1), (5, 1), (7, 1), (9, 1)$ cells respectively.
- Initially, robots r_6, \dots, r_{10} are placed at $(1, 10), (3, 10), (5, 10), (7, 10), (9, 10)$ cells respectively.

In each transition, a robot can move one cell in $\{\text{N}, \text{S}, \text{E}, \text{W}\}$ direction. Write a program for the following:

- The initial locations of r_1, \dots, r_5 should be swapped with the initial locations of r_6, \dots, r_{10} respectively.
The robots should move through the workspace to capture the new locations. At a time a cell can be occupied by a single robot. Find the path for each robot. Give a visual presentation of the same. What is the makespan ?
- Is it possible to further improve the makespan ?

Solve the problem using Microsoft's SMT solver **Z3**. The following link may be helpful – Z3-with-Python. 20

2. [Variants of A*]

- (a) For A* algorithm we use the $f(n) := g(n) + h(n)$ to estimate the cost of the path to goal via node n . Here, both $g(n)$ and $h(n)$ have equal weights. Now, if we want to assign more weight to the heuristic cost $h(n)$, i.e. $f(n) := g(n) + w \cdot h(n)$ where $w > 1$, what will be the effects in the number of explored nodes and optimality guarantee, and why ? 5
- (b) In conventional A*, we explore the nodes in a single direction – from start to goal node. Now, we modify the procedure and explore the nodes in two opposite directions – a) from the start toward the goal (forward search), and b) from the goal toward the start (backward search). Consider the branching factor of each node is b and the goal is d distance away from the start node. What effect will the new algorithm have in the time complexity of the search algorithm ? 5