CS528 SIMD/GCC/FP

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A Sahu 1

Outline

- Intro to Code Optimization
- Machine independent/dependent optimization
- Common sense of Optimization
 - Do less work, avoid expensive Ops, shrink working set
- Simple measure Large impact : simd, branch, comm sub expre
- C++ Optimization
- Scalar Profiling
 - Manual Instrumentation (get_wall_time, clock_t)
 - Function and line based profiling (gprof, gcov)
 - Memory Profiling (valgrind, callgraph)
 - Hardware Performance Counter (oprofile, likwid)

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Further Optimizations for Serial Code

- Simple measure Large impact : simd, branch, comm sub expre
- C++ Optimization

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Simple measures, large impact

- Elimination of Common Sub-expressions
- Avoid Branches:
 - Code Can be SIMdized by compiler/gcc
 - Effective use of pipeline for loop code
- Use of SIMD Instruction sets
 - 512 bit AVX SIMD in modern processor
 - ML/Al app use 8 bit Ops, can be speed up
 512/8=64 time by simply SIMD-AVX

Elimination of Common Subexpressions

```
//value of s, r, x don't change in this loop
for (i=0; i<ALargeN; i++) {
    A[i]=A[i]+s+r+sinx(x);
}</pre>
```

```
//value of s, r, x don't change in this loop
Tmp=s+r+sinx(x);
for (i=0; i<ALargeN; i++) {
    A[i]=A[i]+Tmp;
}</pre>
```

Avoid Branches

```
for (i=0; i<N; i++)
  for(j=0; j<N; j++) {
    if(i<j) S=1; else S=-1;
    C[i] =C[i]+S*A[i][j]*B[i];
}</pre>
```

```
for (i=0; i<N; i++) {
    for (j=0;j<i;j++)
        C[i] =C[i] -A[i][j]*B[i];
    for (j=i;j<N;j++)
        C[i] =C[i] +A[i][j]*B[i];
}</pre>
```

```
for (i=0; i<ALargeN; i++) {
    A[i]=A[i]+B[i]*D[i];
}
All iterations in this loop are
independent : gcc SIMD utilize very
nicely
//ML application uses 8 bit OPS, 512
bit AVX SIMD 512/8=64 OPS can be done
in parallel.</pre>
```

The ith iteration access: A[i], B[i], D[i]

$$S = \sum_{i=0}^{N} w_i x_i$$

Vector dot product : is the most common and

frequent kernel in

- Matrix multiplication,
- Neuron calculation (neural network NN)
 - Conv NN, Deep NN, //ML domain
- Digital Signal Processing, Image Sig Processing, etc.
- Media Applications: audio, video, JPG/MPG, DCT.

$$S = \sum_{i=0}^{N} w_i x_i$$

```
for (i=0; i<ALargeN; i++){
    S=S+A[i]+B[i];
}
All iterations in this loop are
independent : gcc SIMD utilize very
nicely
//ML application uses 8 bit OPS, 512 bit
AVX SIMD 512/8=64 OPS can be done in
parallel.</pre>
```

The ith iteration access: A[i], B[i]

```
for (i=0; i<N; i++) {
    A[i]=A[i]+B[i]; //S1
    B[i+1]=C[i]+D[i]; //S2
}</pre>
```

The ith iteration access: A[i], B[i], B[i+1], C[i] D[i]

Dependent loop iteration: i and i+1

```
A[0]=A[0]+B[0];
for (i=0; i<N; i++){
   B[i+1]=C[i]+D[i]; //S2
   A[i+1]=A[i+1]+B[i+1];//S1
}
B[N]=C[N-1]+D[N-1];</pre>
```

The ith iteration access: A[i+1], B[i+1], C[i], D[i]

```
for (i=0; i<N; i++) {
    A[i]=A[i]+B[i]; //S1
    B[i+1]=C[i]+D[i]; //S2
}</pre>
```



```
A[0]=A[0]+B[0];
for (i=0; i<N; i++){
   B[i+1]=C[i]+D[i]; //S2
   A[i+1]=A[i+1]+B[i+1];//S1
}
B[N]=C[N-1]+D[N-1];</pre>
```

Affine access: index a.x+b form

```
for (i=0; i<N; i++) {
    X[a*i+b]=X[c*i+d];
    //where a,b,c,d are integer
}</pre>
```

- GCD(c,a) divides (d-b) for loop dependence
- Ref Book : Hennesy Paterson, Advanced Computer Architecture, 5th Edition Book,

GCD test Example

```
for (i=0; i<N; i++) {
    X[2*i+3]=X[2*i]+5.0;
    //X[a*i+b]=X[c*i+d];
}</pre>
```

- GCD(c,a) must divides (d-b) for loop dependence
- Value of a=2, b=3, c=2, d=0;
- GCD(a,c)=2, d-b=-3
- 2 does not divide -3 → No dependence Possible

Role of Compilers

- General Compiler Optimization Options
- Inlining
- Aliasing
- Computational Accuracy

General Compiler Optimization Options

- GCC optimization: -00, -01, -02,-03
- \$man gcc
- At -00 level:
 - Compiler refrain from most of the opt.
 - It is correct choice for analyzing the code with debugger
- At high level
 - Mixed up source lines, eliminate redundant variable, rearrange arithmetic expressions
 - Debugger has a hard time to give user a consistent view on code and data

General Compiler Optimization Options

• Level 1

 fauto-inc-dec, -fmove-loop-invarient, -fmergeconstants, -ftree-copy-prop, -finline-fun-called-once

Level 2

 -falign-functions, -falign-loops, level, -finlining-smallfun, -finling-indirect-fun, -freorder-fun, -fstrictaliasing

Level 3

– -ftree-slp-vectorize, -fvect-cost-model

Inlining

- Inlining
 - Tries to save overhead by inserting the complete code of function
 - At the place where it called
- Saves time and resources by
 - not using function call, stack
 - All complier to use registers
 - Allows compiler to views a larger portion of code and employ OPTimization
- Auto inline or hint in program to function to be inlined

Aliasing

- Assuming a and b don't overlap
 - double __restrict *a, double __restrict *b
 - __restrict say no overlap
- Load and stores in the loop can be rearranged by compiler
- Apply software-pipeling, unrollling, group load/store, SIMD, etc

Computational Accuracy

- Compiler some time refrain from rearranging arithmetic expression
- FP domain associative rule a+(b+c)≠(a+b)+c
- If accuracy need to be maintained
 - Compared to non optimized code
 - Associative rules must not be used by compiler
 - Should be left to programmer to regroup safely
- FP underflow are push to zero

Computational Accuracy

- FP domain associative rule a+(b+c)≠(a+b)+c
 - -Let $a=1.0x10^{38}$, $b=-1.0x10^{38}$, c=1
 - -Result of a+(b+c)
 - $= 1.0 \times 10^{38} + (-1.0 \times 10^{38} + 1)$
 - $=1.0x10^{38} + (-1.0x10^{38}) = 0 //Big+Small=Big$
 - -Result of (a+b)+c
 - $= (1.0x10^{38} + -1.0x10^{38}) + 1$
 - =0+1 = 1

Computational Accuracy

- Why it happens for FP?
 - FP format use 32 bit represent number up to $\pm 2^{127}$
 - Int use 32 bit represent up to $\pm 2^{31}$
 - Used same 32 bit for large numbers, numbers are not equal-spaced
 - From 36000ft, both IITG and Amingoan are not distinguishable [Resolution:]
 - Going by Air: Delhi, Noida, Gurgaon use the same Airport