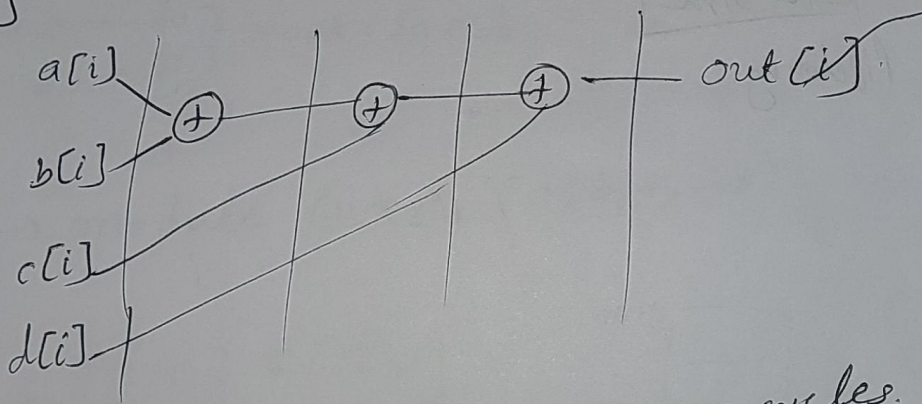


⇒ HLS for loop

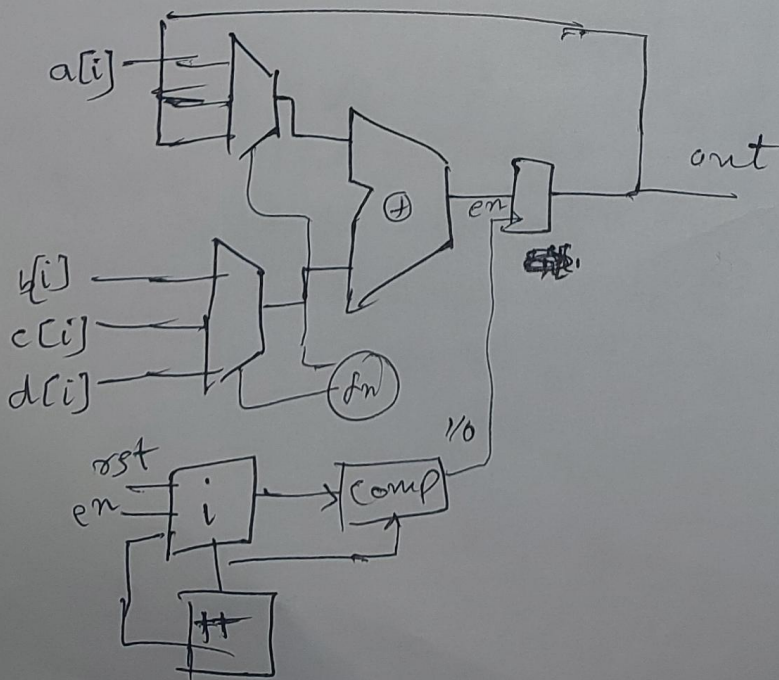
```
void fn (a[10], b[10], c[10], d[10])
{
    for (i=0; i<10; i++) { #pragma HLS pipeline
        out[i] = a[i] + b[i] + c[i] + d[i];
    }
    return out;
}
```

3x10 cycles

clk to execute iter: K
iterations: n



1) iterative implementation: 30 cycles. ($K \times n$)



2) Pipelined implementation: 12 cycles $(K + (n-1) \times 1)$
 $= n + K - 1$

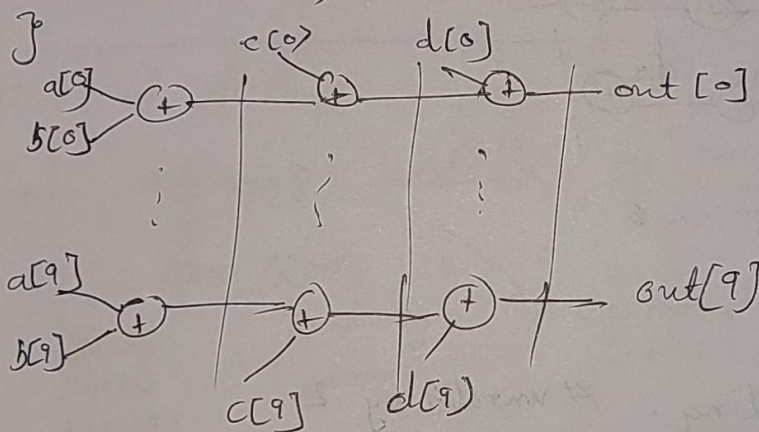
3) Unrolling:

void fn(a[0], b[0], c[0], d[0]) {

out[0] = a[0] + b[0] + c[0] + d[0]

out[9] = a[0] + b[0] + c[0] + d[0]

return out;

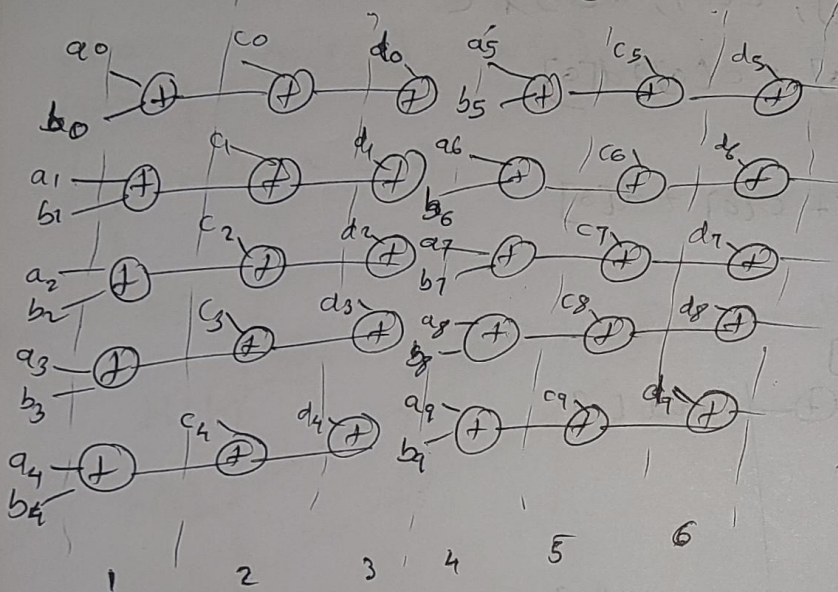


	time	Resources
Iterative	30	1
★ Pipelined	12	3 → best sol ⁿ w.r.t. resource/cost
Unrolled	3	10
Partial Unroll	15	2

3a) Unroll with resource constrained (6 cycles)

RC = 5 adders

obj. = ~~minimize~~ schedule with min time stamp
w/o violating RC



4) Partial Unrolling: # unrolls by 2

```
void fn(a[10], b[10], c[10], d[10]) {
```

```
    for (i=0; i<10; i+=2) {
```

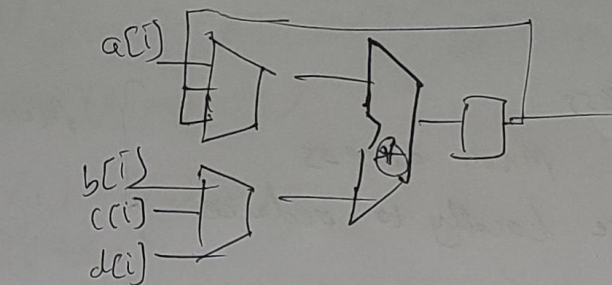
```
        out[i] = a[i] + b[i] + c[i] + d[i];
```

```
        out[i+1] = a[i+1] + b[i+1] + c[i+1] + d[i+1];
```

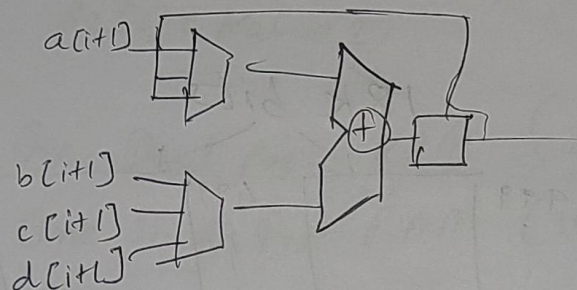
```
    }
```

```
    return out;
```

```
}
```

15 cycles



HLS of Arrays

int a[1000]

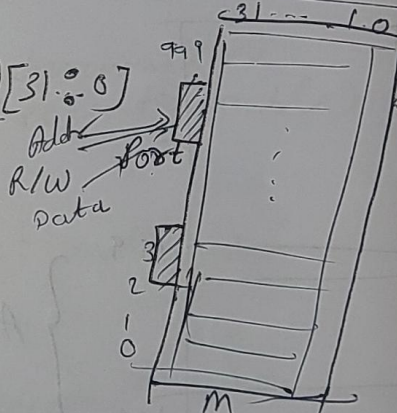
MAP

999:0
m[1000][31:0]
RAM

Random
Access
memory

Readonly
memory

Access through
Index



Memory: 1/2 ports

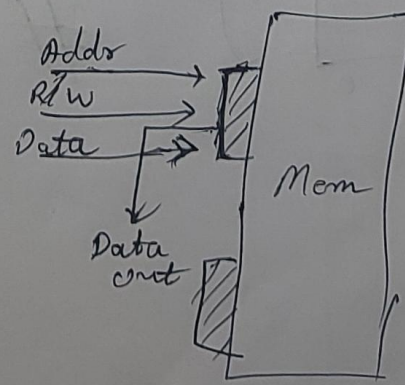
```
void fun(int A[20], int* S){
```

```
    for(i=3; i<20; i++){
        S[i] = A[i] + A[i-1] + A[i-2]
              + A[i-3]
```

```
}
```

⇓

mem access = 68



```
void fun(int A[20], int* S){
```

```
    t1 = A[0], t2 = A[1], t3 = A[2]
```

```
    for(i=3; i<20; i++){ x = A[i]
```

```
        S[i] = A[i] + t1 + t2 + t3;
```

```
        t3 = t2
```

```
        t2 = t3
```

```
        t3 = A[i] * x
```

```
}
```

Mem Access = 20

Array

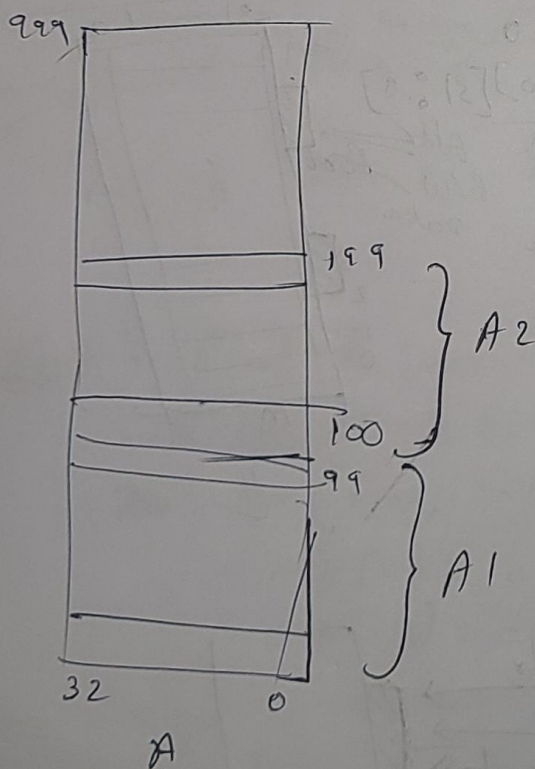
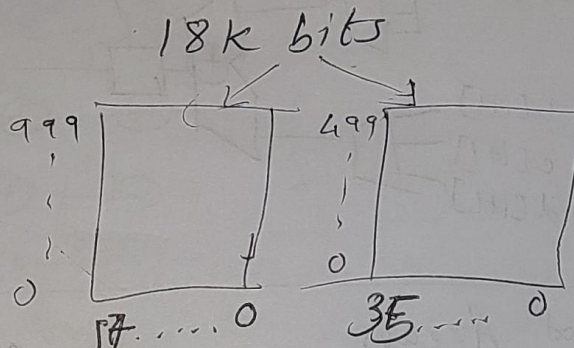
- 1) Reduce Mem Access
- 2) Remove Redundant Mem access
- 3) Read ^{once} and store locally to reduce

Manually

→ Array Merging:

int A[100]

int B[100]



$$A1 + A2 \Rightarrow A$$

$$A2[99:0] \leftrightarrow A[199:100]$$

$$A1[99:0] \leftrightarrow A[99:0]$$

#pragma HLS ARRAY_MAP

Variable A1 instance A
horizontal

#pragma HLS ARRAY_MAP

variable A2 instance A horizontal

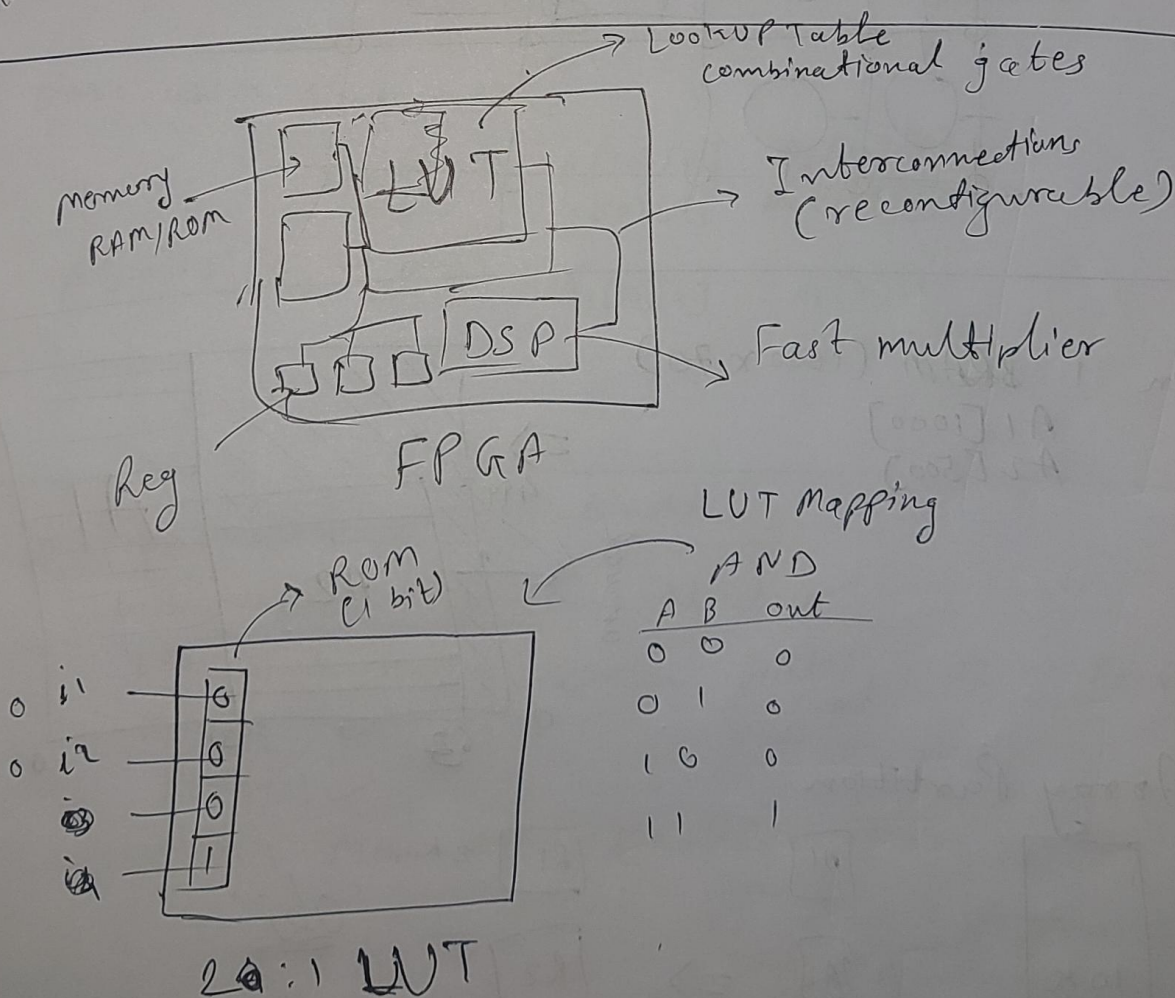

```

int A1[1000], A2[500]; // Map A1 → A; map A2 → A
for (i2=0; i2<1000; i2++)
{
    A1[i2] = value1;
    if (i2 < 500)
        A2[i2] = value2;
}

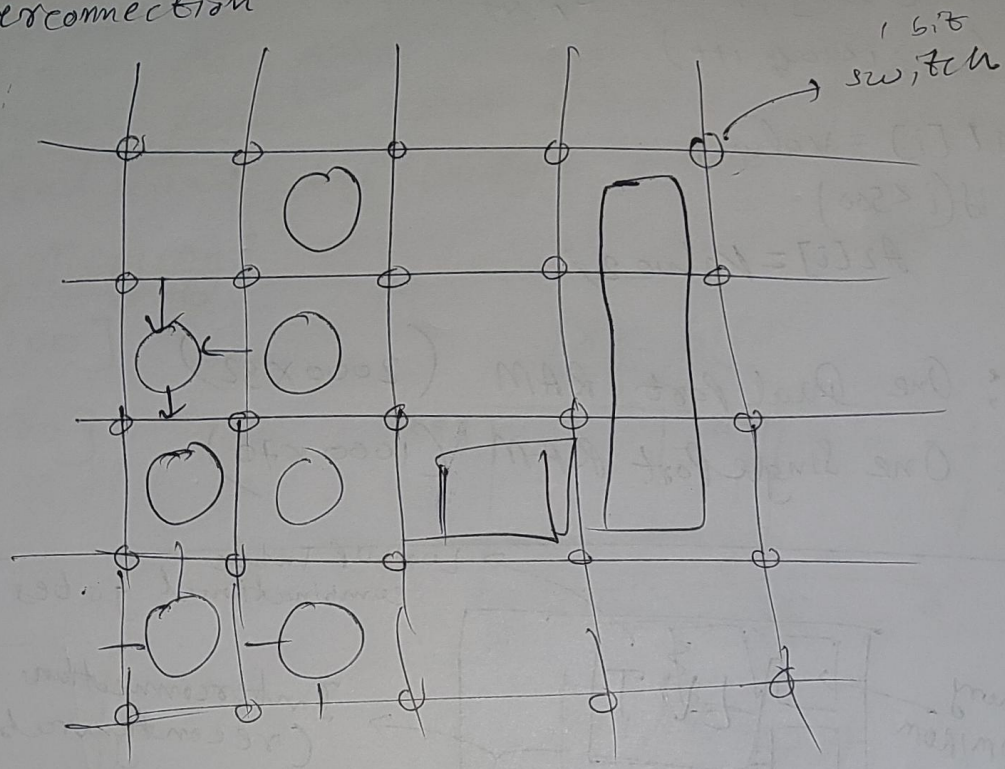
```

Given: One Dual Port RAM (2000 × 32)

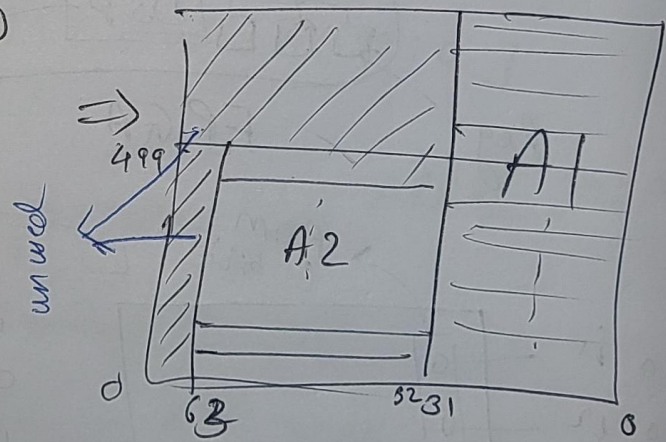
Given: One Single Port RAM (1000 × 70)



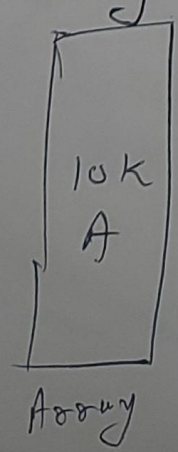
Interconnection



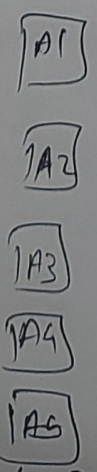
given 1 BRAM (1000 x 70)
 A1 [1000]
 A2 [500]



⇒ Array Partition

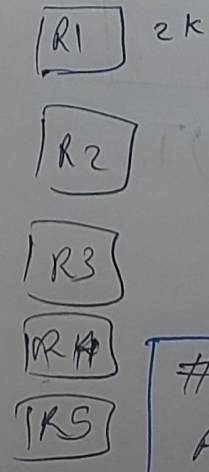


⇒



each 2K

⇒

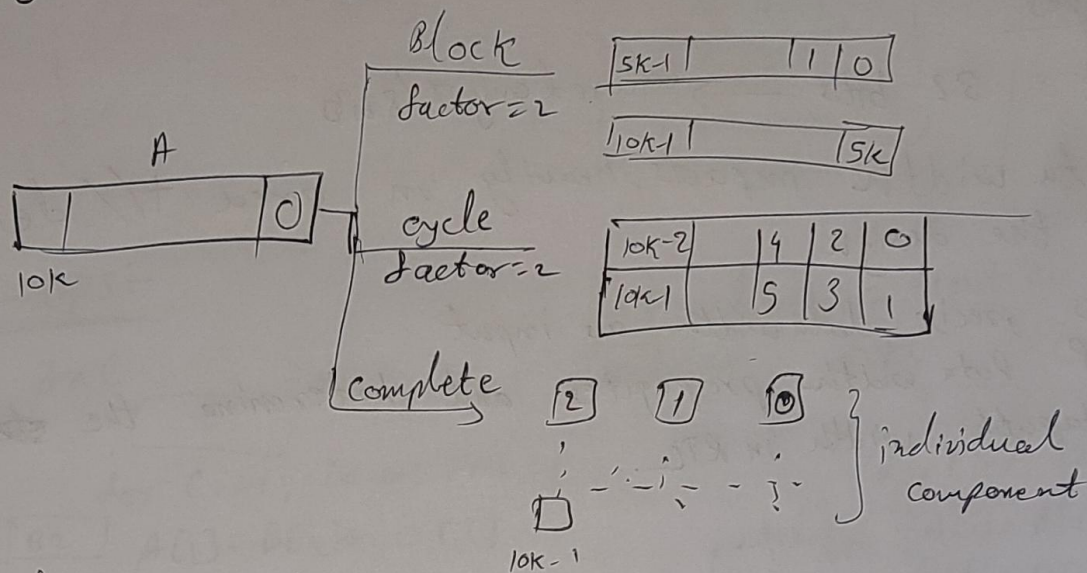


obj:

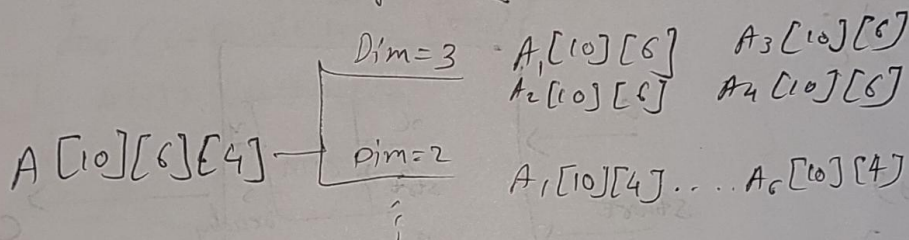
Keep parallel access memory in separate RAMs.

#Pragma HLS array-partition
 Array-name <type> factor=
 Dim =

type: block/cycle/complete



Dim: which dim to partition



Unsupported C constructs

Data types: int, char, float ✓

Structure, Union: ✓

pointer: ✓

Array: ✓

Dynamic mem Alloc (malloc/calloc): ✗

Function: ✓

Recursive Function: ✗

Standard template Library: ✗

System Calls: ✗

printf/scanf: ✗ → Remove

fsd: Rewrite program by statically allocating memory

Rewrite non recursive version as the fn

✗ → write it yourself.
→ Remove them

Coding Style for HLS

Data types

Int: 32 bits, \rightarrow required only 7/8 bits

- 1) Data width impacts heavily on area +// delay of the design
- 2) Uses precise datawidth as input
- 3) I/P Data width propagates and determine the datapath width in RTL.

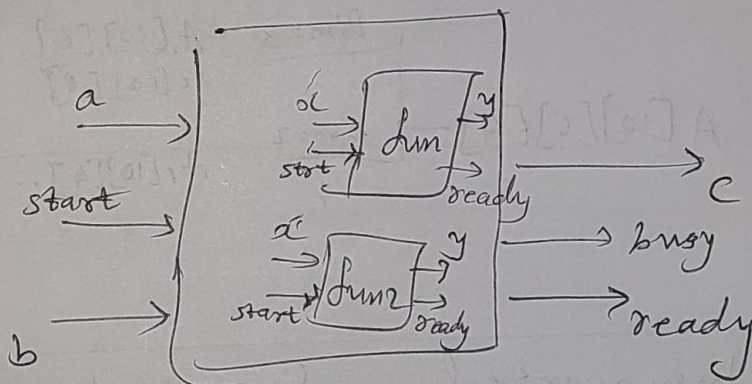
Function

```
top(a, b, &c) {
```

```
    fun(x, y)
```

```
    fun2(x, y)
```

```
}
```



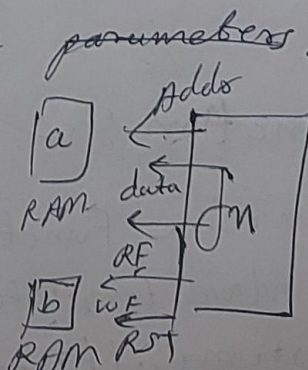
- 1) small function should be inlined to reduce area overhead.
- 2) Each fn is allocated a dedicated HW module
- 3) Functions can be scheduled in parallel.
 \Rightarrow improves performance.
- 4) Must be careful about the ~~parameters~~ arguments

```
top(---) {  
    int a[100], b[100];
```

```
    fun(a[], b[]);
```

5) declare inside to avoid

```
}
```

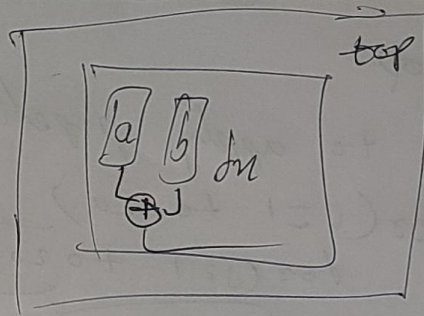


- 6) Array may be duplicated to avoid this.


```

fnl( ) {
  // at [100, 4100]
}

```



Loops:-

```

fnl( ) {
  B1
  for (i=1; i<100; i++) {
    B2 A[i] = B[i] + C[i]
  }
  for (i=1; i<100; i++) {
    B3 D[i] = E[i] + F[i]
  }
  B4
}

```

Loop limit 1 = limit 2

```

parallel for ( ) {
  B1
  for (i=1; i<100; i++) {
    parallel B2 A[i] = B[i] + C[i]
    D[i] = E[i] + F[i]
  }
  B3
}

```

limit 1 \neq limit 2

```

for ( ) {
  for ( ), fnl( ) // loop 1
  fnl( ) // loop 2
} both scheduled
  parallel
}

```


Nested loop:

① when to apply pipeline

```
for (i = 1 to 20) _____ → pipeline  
  for (j = 1 to 20) { _____ → pipeline
```

B1

}

}

⇒ outer loop:

→ Inner loop will be unrolled.

→ Run Faster but with lot of area overhead.

→ pipeline for 20 times

* ⇒ Inner loop:
suggested

→ Pipeline inner loop

→ Require less resource

→ Pipeline for 400 times.