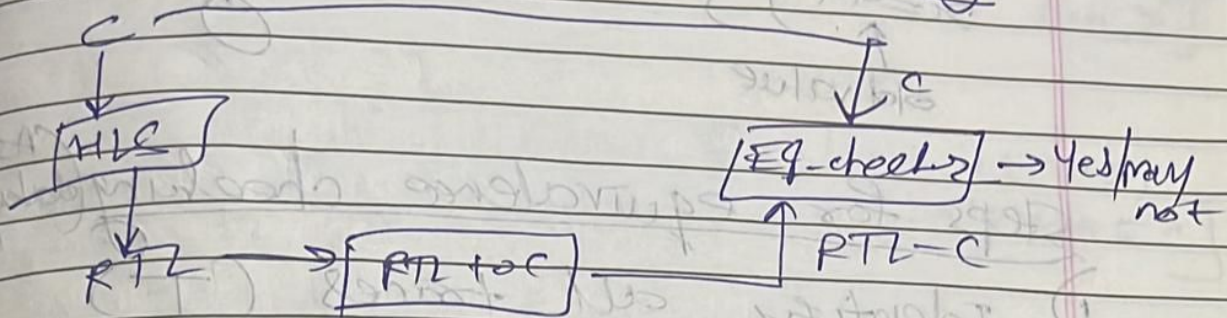


## C to RTL Equivalence checking



FSM-D (Finite state machine with datapath)

fn (i1, i2, i3) {

    a = i1 + i2;

    if (a > 5)

        x = a + 5

    else

        x = a - 5

    out = x + i3;

    return out;

}

Program → FSM-P

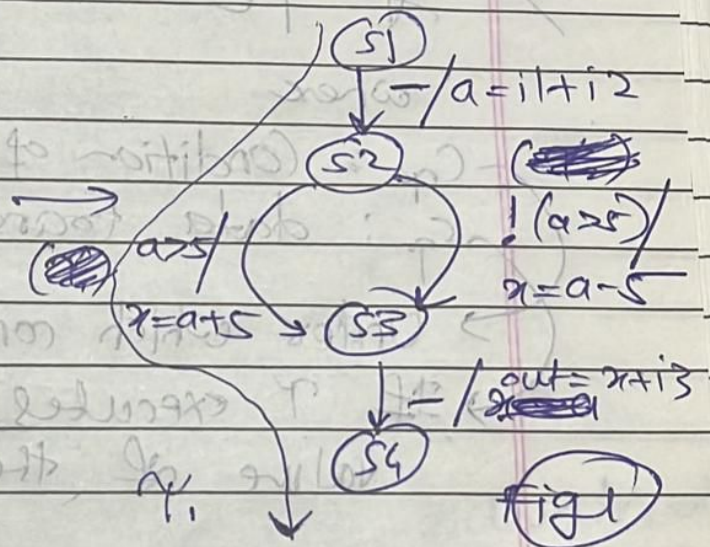
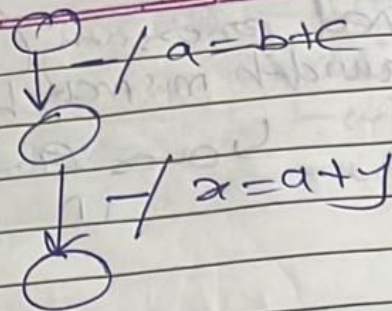


Fig 1

Note: 2 opns in same state ⇒ opns will be executed in parallel



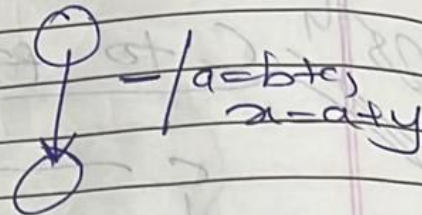
\*  $a = b + c$   
 $x = a + y$



(Fig 2)

~~\*  $a = b + c$   
 $x = a + y$~~

$a = b + c$   
 $x = a + y$  } parallel  
 ↑  
 old value



# Steps for equivalence checking (Assuming data indep. 100% length)

1) Identify all traces (T)  
 Trace is one execution path of the program.  
 Eg In Fig 1  $\rightarrow T = 2$

2)  $\forall \gamma \in T, C_\gamma \neq S_\gamma$

where

$C_\gamma$ : Condition of execution of  $\gamma$

$S_\gamma$ : data transformation of  $\gamma$

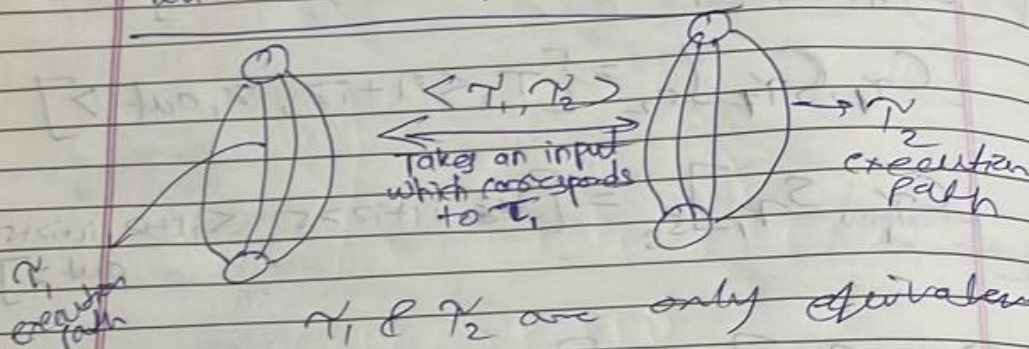
$\rightarrow$  Under which condition  $\gamma$  will execute

$\rightarrow$  If  $\gamma$  executes, what will be the final value of the variables - outputs.



How to do step 4 in d(nk) →

Data driven approach



Coverage Driven Testing

corresponding traces b/w two programmes using coverage driven testing.

Modified algo after scanning step 4 is →

5) Merge compatible traces in  $T_1$

6) " " in  $T_2$

7) Apply data driven approach to identify corresponding traces

8) for each corresponding trace pair,  $\langle \gamma_1, \gamma_2 \rangle$  ( $\gamma_1 \in T_1, \gamma_2 \in T_2$ ),

check  $\gamma_1 \subseteq \gamma_2$

↓  
Binary decision program



substitution method / symbolic substitution

$$[C_{\gamma_1}, S_{\gamma_1}]_{s_1} = [T, \text{state}, \text{only these vars are updating}, \text{out}]$$

$$[C_{\gamma_2}, S_{\gamma_2}]_{s_2} = [T, \langle i1+i2, \text{out} \rangle]$$

$$[C_{\gamma_1}, S_{\gamma_1}]_{s_3} = [\text{state}, \langle i1+i2 > 5, \text{out} \rangle]$$

$$[C_{\gamma_1}, S_{\gamma_1}]_{s_4} = [i1+i2 > 5, \langle i1+i2, i1+i2+5, i1+i2+i3+5 \rangle]$$

3) Passform step 1 & 2 in  $P_1$  &  $P_2$  ( $T_1, T_2$ )

4) For each  $\gamma_1 \in T_1$ , identify  $\gamma_2 \in T_2$  s.t.

$$O(n^2 \times \overset{\text{time for SMT solver}}{C}) \quad \gamma_1 \equiv \gamma_2$$

where  $\gamma_1 \equiv \gamma_2$  iff  $C_{\gamma_1} = C_{\gamma_2}$  and  $S_{\gamma_1} = S_{\gamma_2} \rightarrow ?$

Two expressions

check  $C_{\gamma_1} \neq C_{\gamma_2}$

check SAT  $\rightarrow$  SAT  $\Rightarrow$  not equivalent

UNSAT  $\Rightarrow$  equivalent