ME 620: Fundamentals of Artificial Intelligence

Lecture 10: Searching AND/OR Graphs

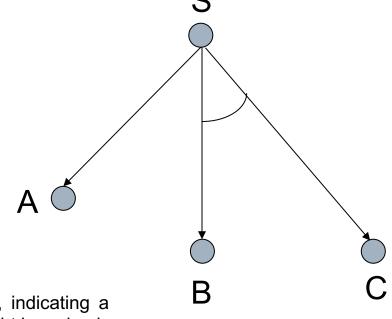


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- Represent the solution of problems that can be solved by decomposing them into a set of smaller problems, all of which must then be solved.
 - AND arc may point to any number of successor node, all of which must be solved in order for the arc to point to a solution.
 - AND arcs are indicated with a line connecting all the components.



Several arcs may emerge from a single node, indicating a variety of ways in which the original problem might be solved.

Hypergraph



- ☐ AND-OR Graphs are defined as hypergraphs.
 - Instead of arcs connecting pairs of nodes, there are hyperarcs connecting a parent node with a set of successor nodes.

These hyperarcs are called connectors.

Each k-connector is directed from a parent node to a set of k successor node.

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If all of the connectors are 1-connectors, we have the special case of an ordinary graph.

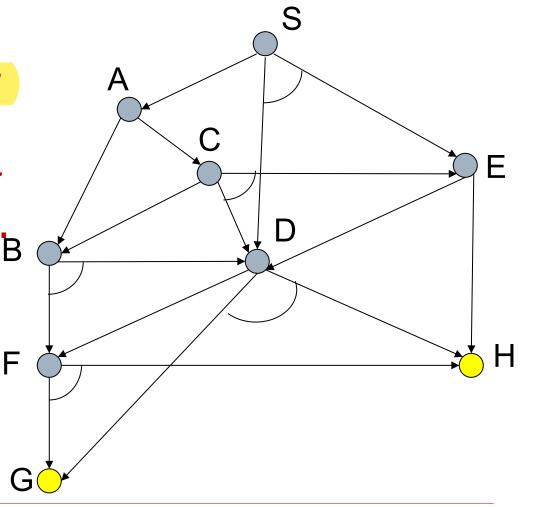


☐ Figure below shows an example AND-OR Graph.

Node S has a *one-connector* directed to successor A

Node S has a two-connector directed to successors (D, E)

For k> 1; each k-connector is denoted by a curved line joining the arcs from parent to elements of the successor set.

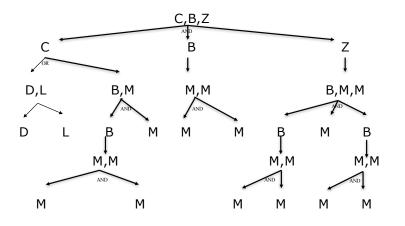


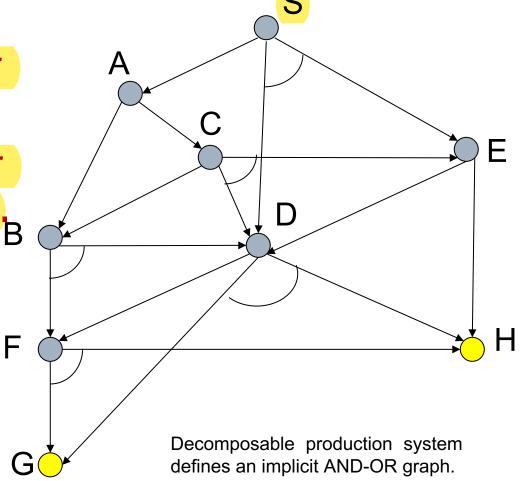


☐ Initial database corresponds to start node S.

■ Node S has a *one-connector* directed to successor A

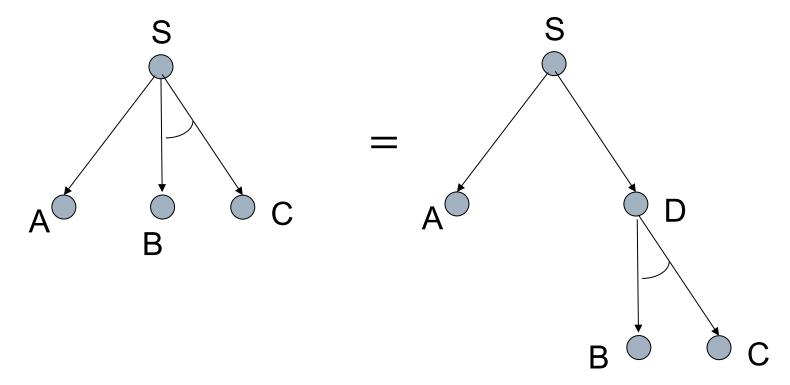
■ Node S has a *two-connector* directed to successors (D, E).







If required one can always transform the graph on the left to one of the right; of course, keeping track of the cost of the edges.



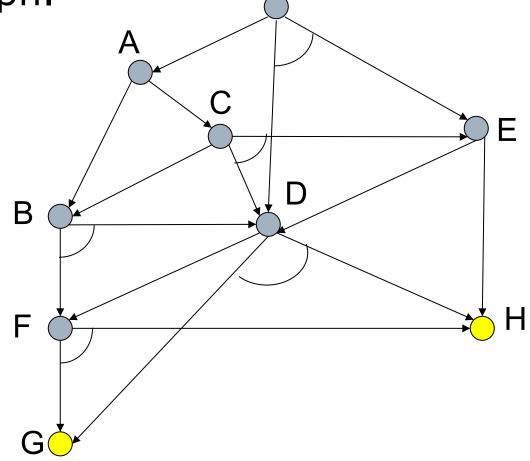
Recall that the AND-OR graphs used in our discussion of decomposable production systems had alternate AND and OR nodes.



☐ A solution graph of an AND-OR graph is analogous

to a path in an ordinary graph.

- Starting from node S; select exactly one connector.
- From each successor node to which the connector is directed, select an outgoing connector and so on.
- Eventually every successor is an element of N

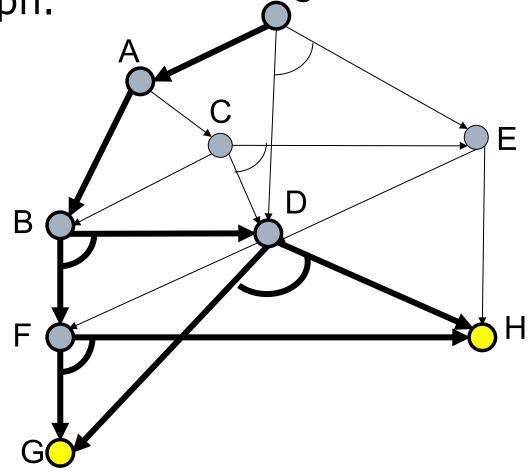




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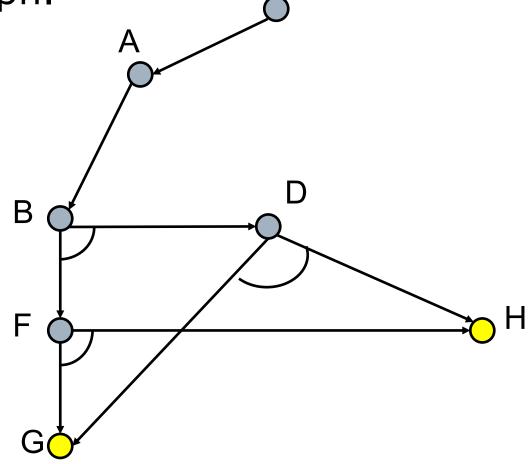




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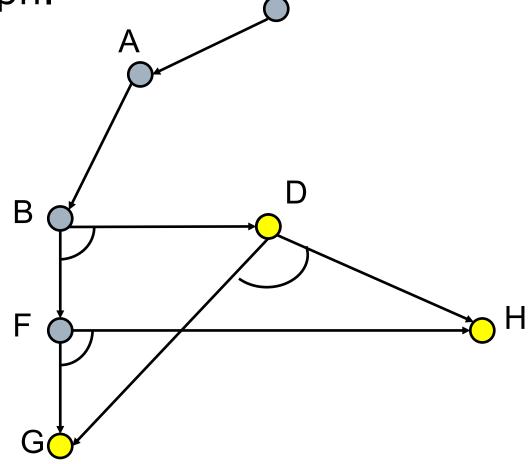




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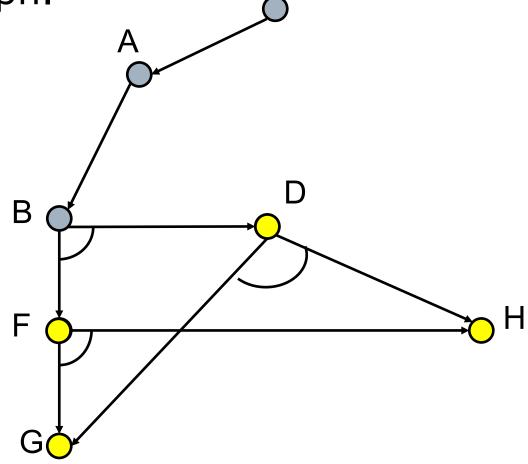




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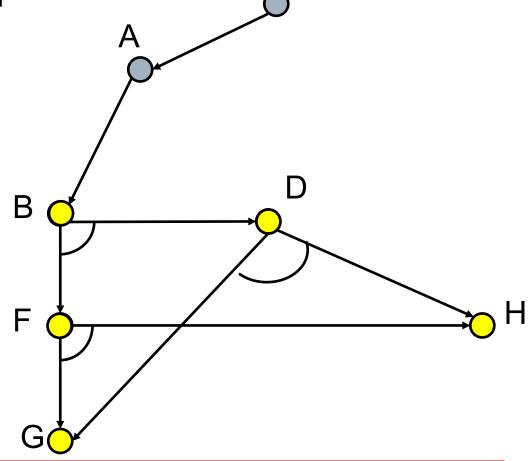




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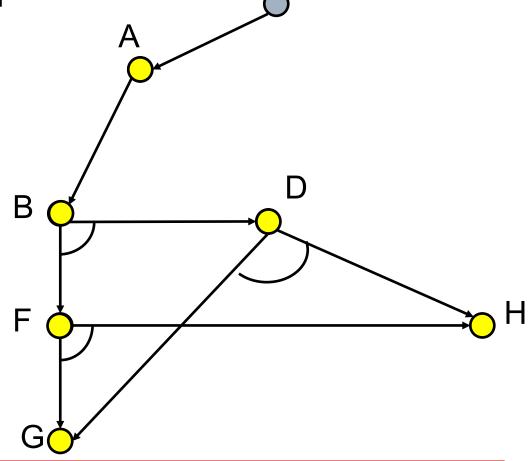




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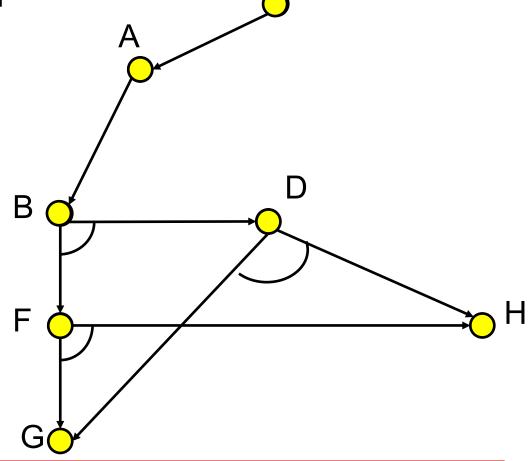




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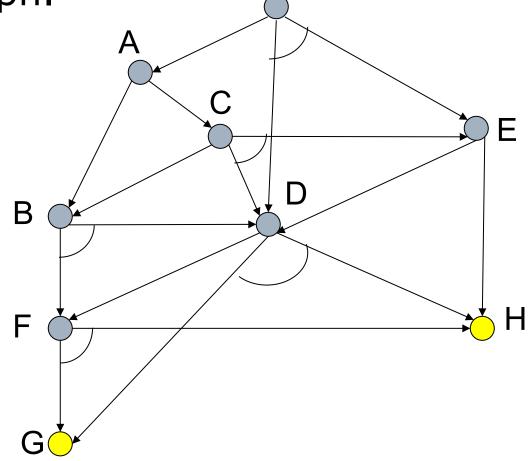




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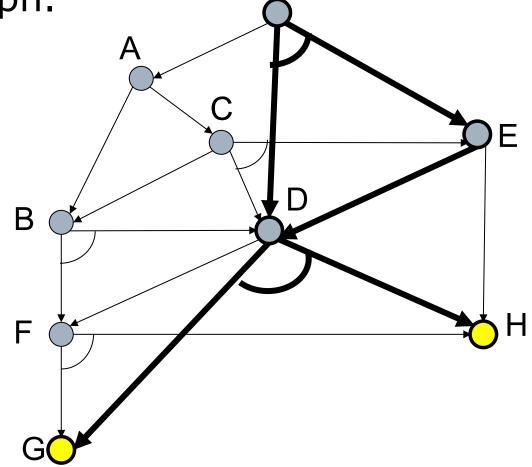




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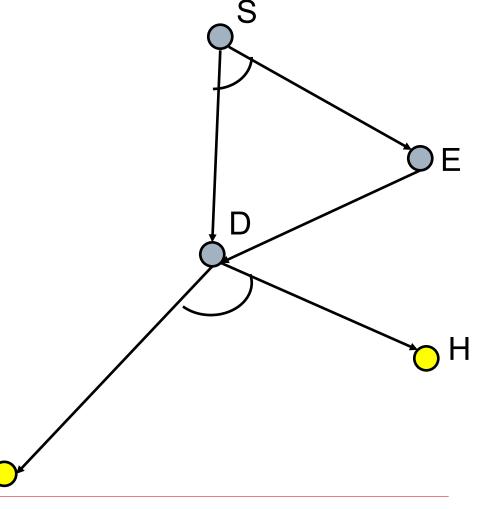




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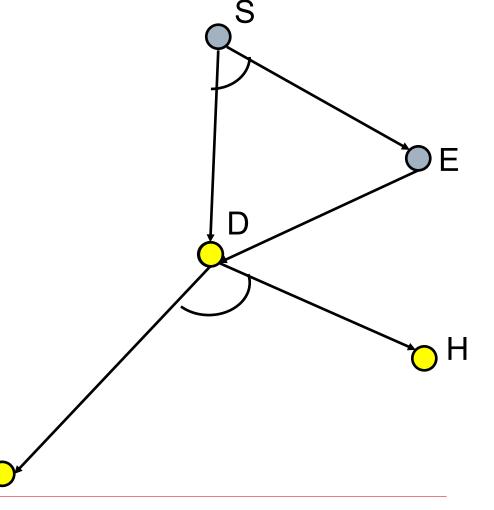




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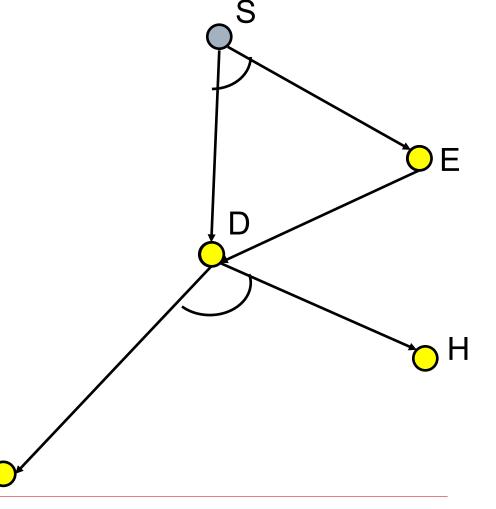




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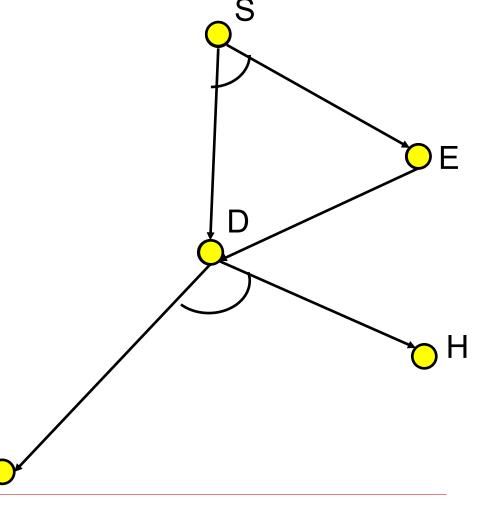




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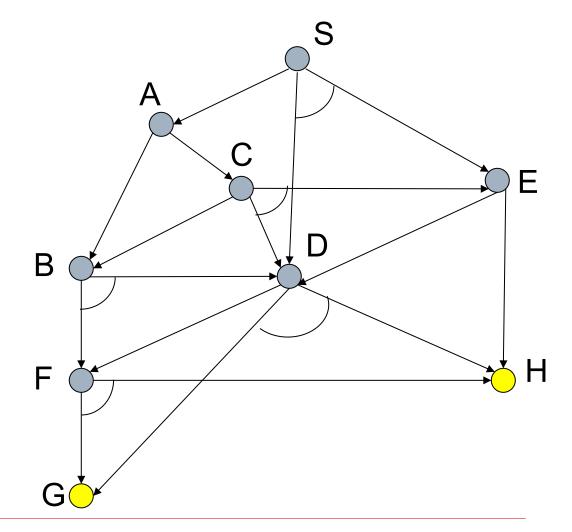
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Searching AND-OR Graph



- To find solutions, we need an algorithm similar to Best-First Search
 - ☐ But with the ability to handle the AND arcs appropriately.
- Algorithm should find a path from the start node to a set of nodes representing the solution states.
 - May be required to reach to more than one solution state;
 Each arm of an AND arc must lead to a solution node.



Best-First Algorithm?

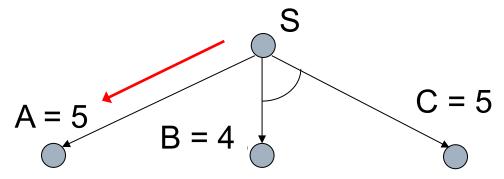


□ To find solution in an AND-OR graph we need an algorithm similar to best-first search; but with the ability to handle the AND arc appropriately.

Edge cost = 1 unit.
Best-first Search? Node B

Involving
$$B = 4 + 5 + 2 = 11$$

Path S-A = 5 +1 = 6; Better



B is part of an AND arc; If we choose to use B, we must also use C.

Assumption: Each k-connector has a cost of 1 for each of its k successors.

AND arc with TWO successors have a cost of 2.

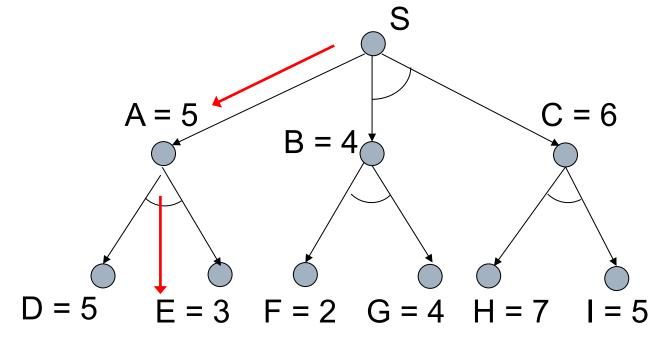
Best-First Algorithm?



□ To find solution in an AND-OR graph we need an algorithm similar to best-first search; but with the ability to handle the AND arc appropriately.

Edge cost = 1 unit.
Best-first Search? Edge B-FG

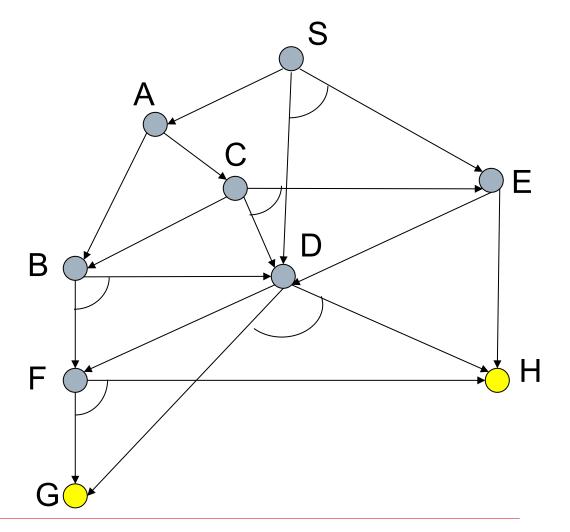
Involving B-FG; Also C-HI Path S-A; A-DE Better



Searching AND-OR Graph



- To find solutions, we need an algorithm similar to Best-First Search
 - ☐ But with the <u>ability to handle</u> the <u>AND arcs</u> appropriately.
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Use a heuristic function h(n); estimate of the cost of an optimal solution from node n to set of terminal nodes.

Each node will have an associated h value; serve as measure of goodness of the node.

Rather than two lists – OPEN and CLOSED - used for A* algorithm, the AO* algorithm use a single structure - a graph - representing part of the search graph that has been explicitly generated so far.

Two major operations:

- 1. Top-down, graph-growing
- 2. Bottom-up, cost-revising, connnector-marking, SOLVE-labelling.

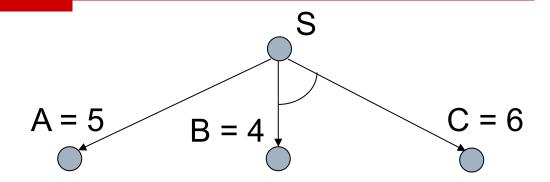
Not store a value for the cost to the current node; such a value is not necessary because of the top down traversal of the best-known path.



After graph G is initialized to S

Follow which hyperedge?

Alternatives – Path S-A; Path S-BC

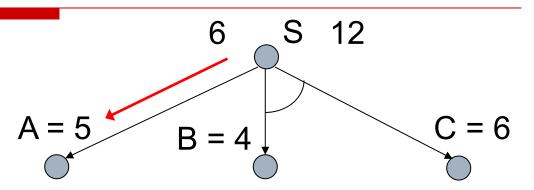




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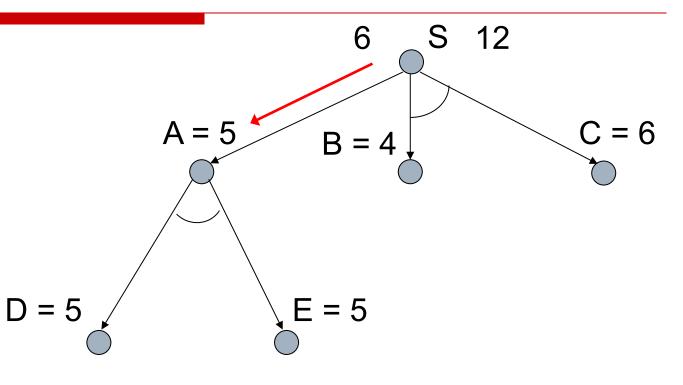


After graph A is expanded;

Nodes – D and E

How promising?

Evaluate h values – back-up



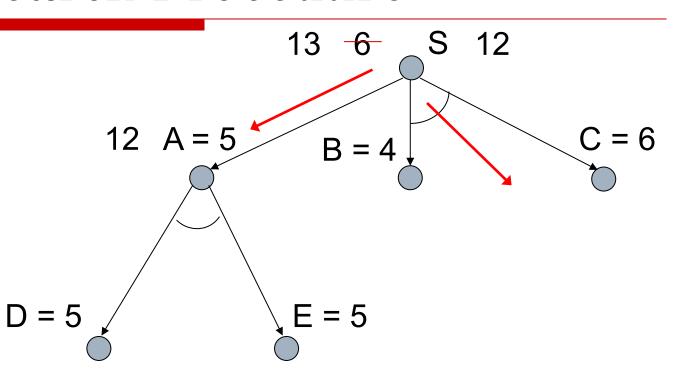


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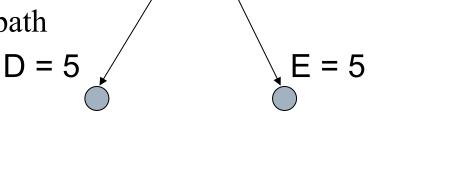
More promising?

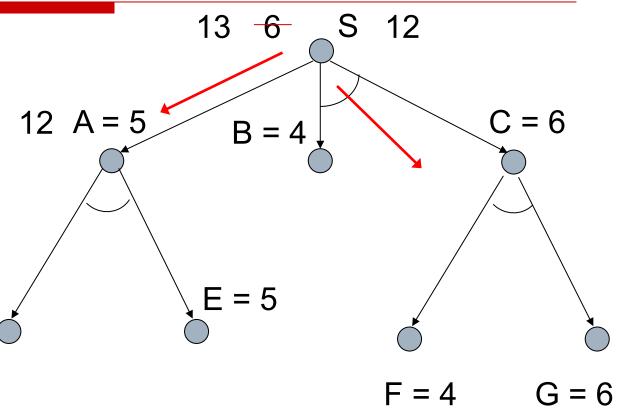
After node C is expanded;

Nodes – F and G

How promising?

Evaluate h values – revise on path







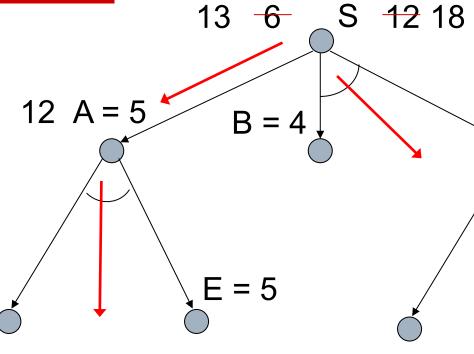
More promising?

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Nodes – F and G

How promising?

Evaluate h values – revise on path



F = 4

G = 6

How promising?

Nodes – D and E

Evaluate h values – back-up

Basic Idea of A0*



□ Top-down graph growing

- Picks out best available partial solution sub-graph from explicit graph by tracing down marked connectors.
 - Here *marked* indicates the current best partial solution graph from each node in the search graph.
- One of the nonterminal leaf nodes of this best partial solution graph is expanded.
 - cost is assigned to its successors.

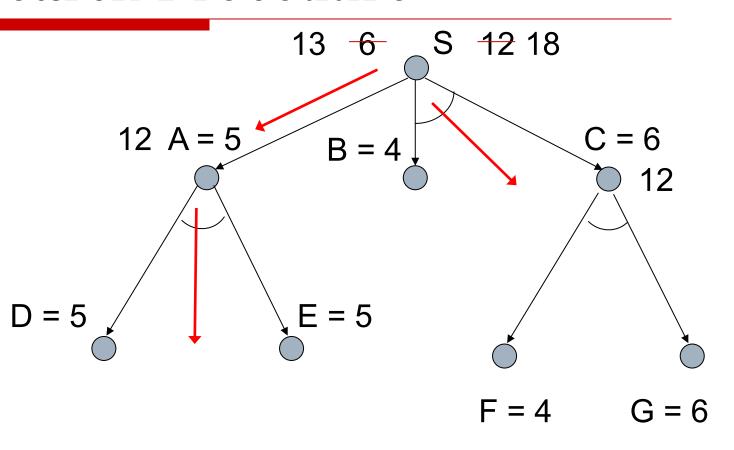


Node E is expanded;

Nodes – D and E

How promising?

Evaluate h values – back-up



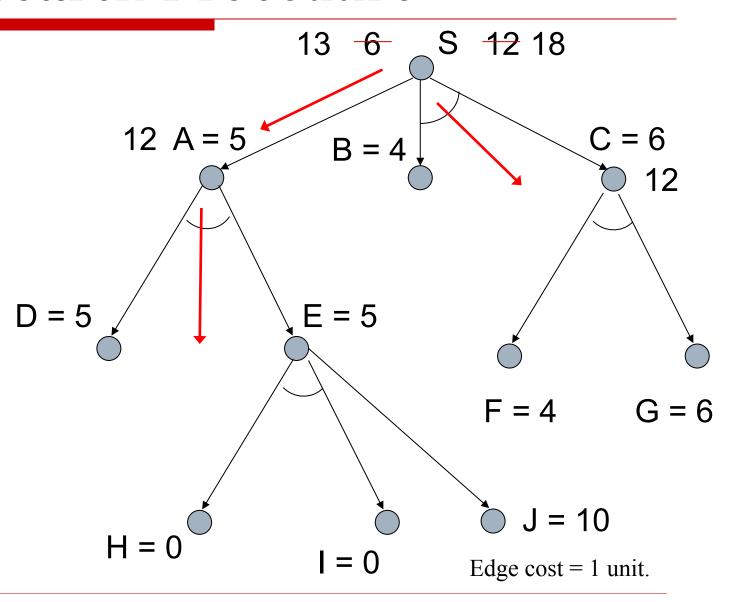


Node E is expanded;

Nodes – H-I; J

How promising?

Evaluate h values – back-up



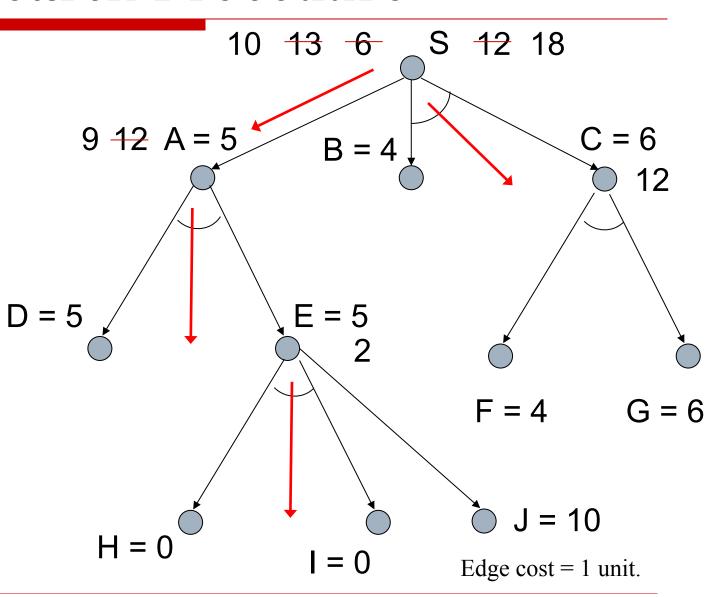


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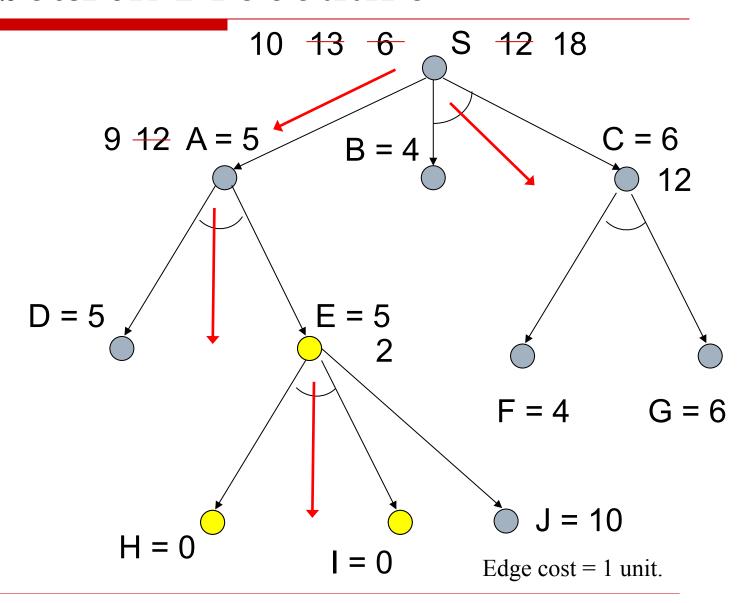
Evaluate h values – back-up





Node E is expanded; Nodes – H-I; J

SOLVE - Labelling



AO*: A Heuristic Search Procedure

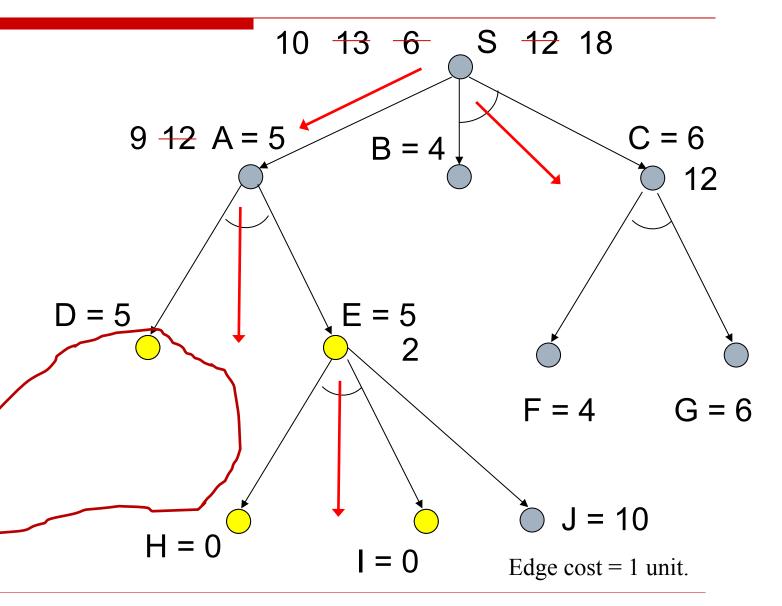


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SOLVE - Labelling

Sub-tree at Node D SOLVED!



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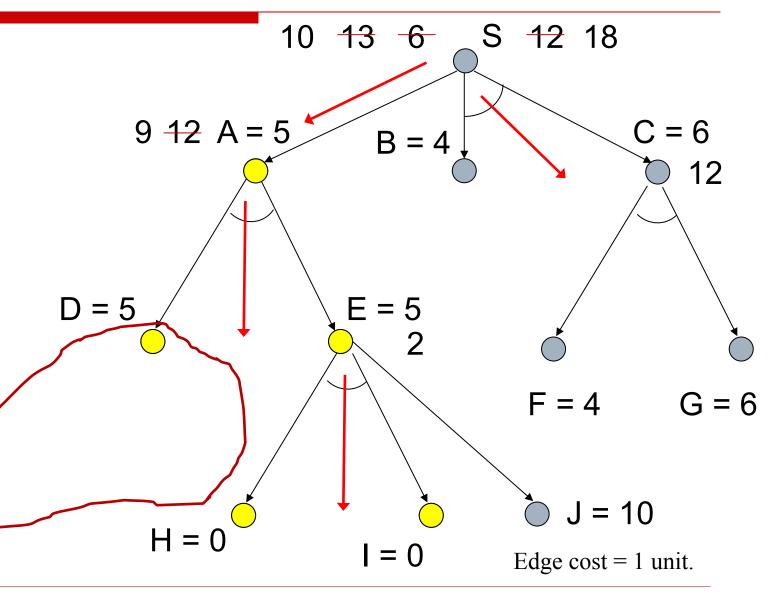


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AO*: A Heuristic Search Procedure

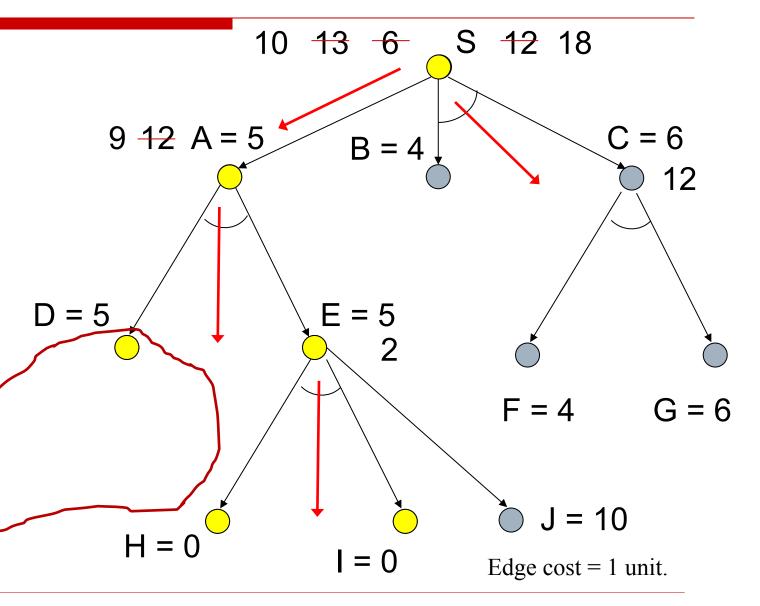


Node E is expanded;

Nodes – H-I; J

SOLVE - Labelling

Sub-tree at Node D SOLVED!



Basic Idea of AO*



- Bottom-up cost-revising, connector-marking
 - SOLVE-labeling.
 - If a direction has all successors SOLVED then n is marked SOLVED.
 - Starting with the node just expanded, the procedure revises its cost (using the newly computed cost of its successors).
 - Marks the outgoing connector on the estimated best path to terminal nodes.
 - Propagated upward in the graph.

AO* Algorithm



Create a search graph $G = \langle s \rangle$; q(s) = h(s); If s is a terminal node, mark s SOLVED Until s labeled SOLVED do:

begin

Top-down graph growing - picks out best available partial solution sub-graph

- a. Compute G' partial solution graph in G by tracing down marked connectors in G from s.
- b. Select any nonterminal leaf node n of G'.
- **c. Expand n**, place successors in G, For each successor not already in G let q(successor) = h(successor). Label SOLVED all successors that are terminal nodes.
- d. Let $\hat{S} := \{n\}$.

Until S is empty do:

begin

- a. Remove a node, m, from S which has no descendent in G also in S (minimal node).
- b. Revise cost for m q(m)=min [c + q(n1)+...+q(nk)].

Mark chosen connector.

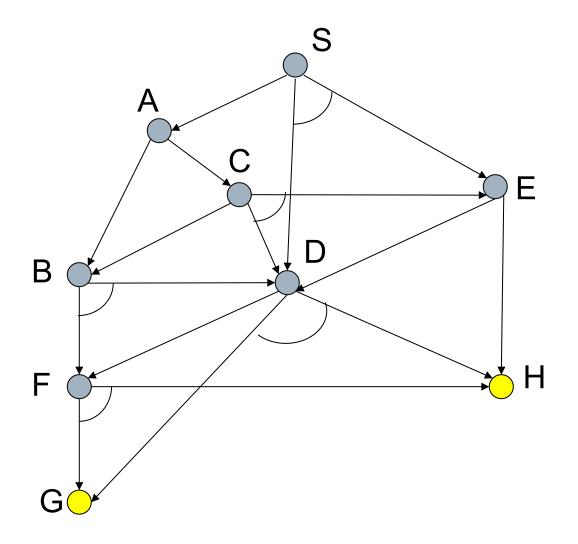
If all successors through the connectors are SOLVED then mark m SOLVED.

c. If m SOLVED or changed q(m) then add to S all "preferred" parents of m.

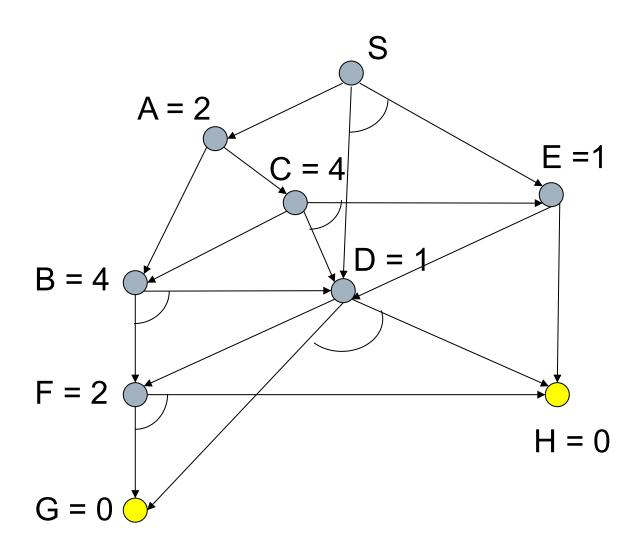
End.

Bottom-up, cost-revising, connector marking; SOLVED labelling procedure



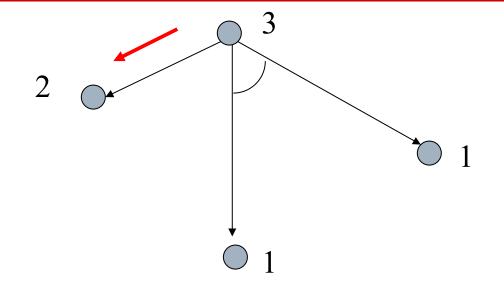


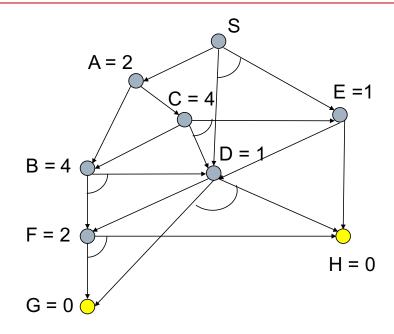




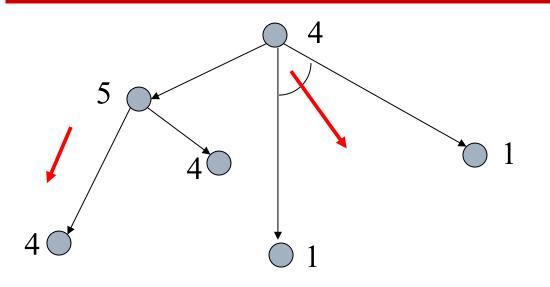
Nils J. Nilsson, Principles of Artificial Intelligence, Narosa. Chapter 3, Pages 107-108.

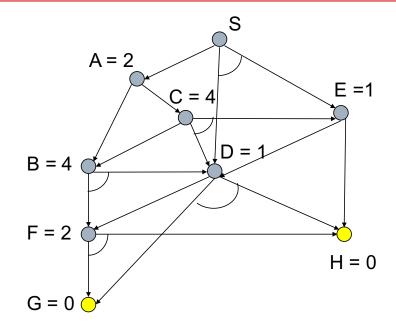




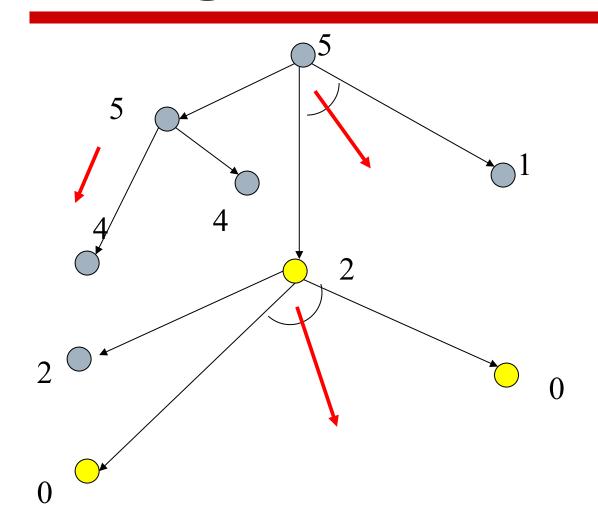


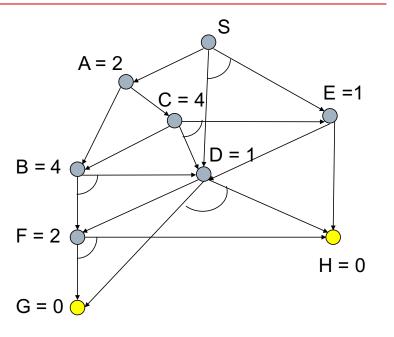




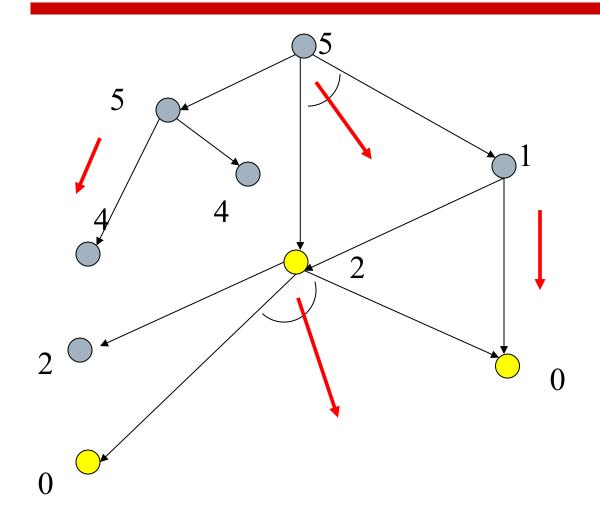


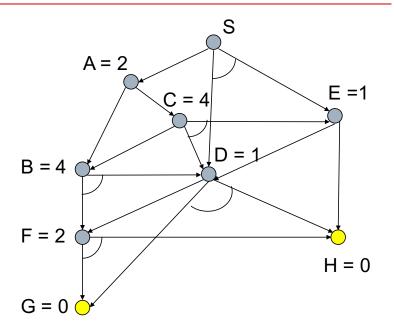




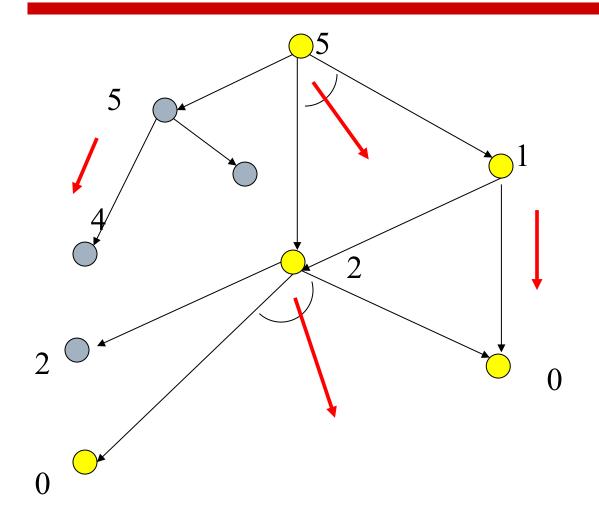


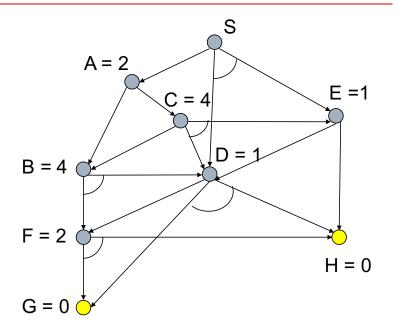












AO* Modifications



- □ AO* may be modified in a variety of ways to render it more practical.
 - Rather than recompute a new estimated partial solution graph after every node expansion; expand one or more leaf nodes and some number of their descendants all at once. Recompute an estimated best partial solution graph.
 - This would reduce the overhead of frequent bottom-up operations.
 - Staged-search strategy may be used for AND-OR graphs
 - ☐ Identify few partial solution graphs having largest estimated cost; these can be discarded periodically.
 - Risk of discarding one that might turn out to the top of an optimal solution graph.