CSP: N-Queens

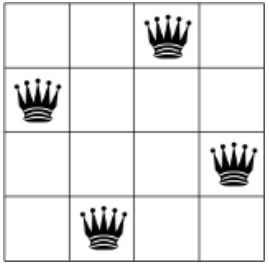
Team 3

AI Class – Project 2

February 2026

Problem Description

- Ensure exactly one queen per row and column.
- Avoid conflicts along diagonals.
- Apply MRV and LCV heuristics for efficient search.
- Use AC-3 algorithm for constraint propagation to reduce possibilities and speed up solving.



Algorithms

CSP Backtracking

For smaller boards (n <= 50)

Minimum Remaining Values (MRV)

Selects the row with the fewest remaining valid columns.

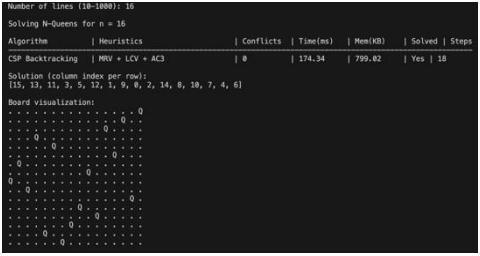
Least Constraining Value (LCV)

Orders columns that minimize conflicts with other queens.

AC-3 (Arc Consistency)

Propagates constraints to reduce domains before trying values.

N = 11



Min-Conflicts (Iterative Repair)

For larger boards (n > 50)

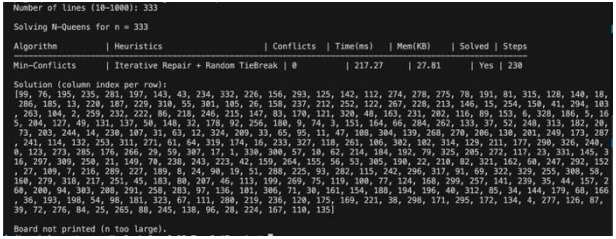
Start with a random assignment of queens.

Identify rows where queens are in conflict.

Move conflicted queens to positions that minimize conflicts.

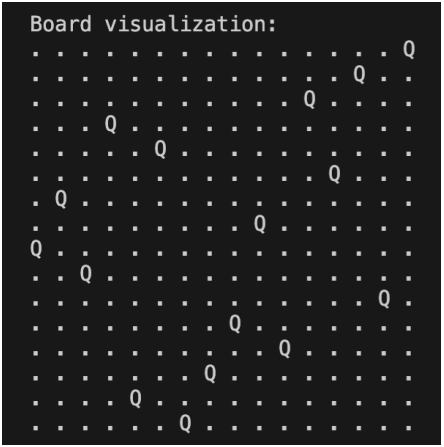
Repeat up to a maximum number of steps

N = 333



Implementation Notes

- Initial Board Input
- Conflict Counting
- Memory & Performance Tracking
- Board Visualization



Testing and Validation

Example result table for n = 23:

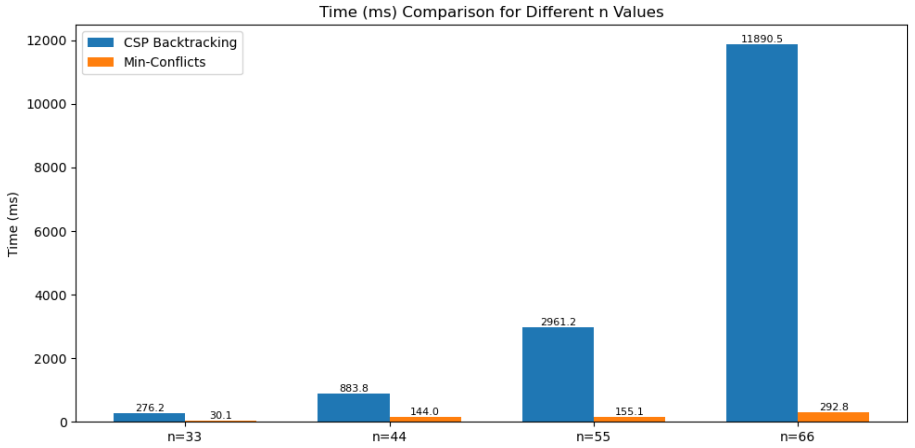
Algorithm	Conflicts	Time(ms)	Mem(KB)	Solved	Steps
CSP Backtracking	0	912.59	2995.83	Yes	27
Min-Conflicts	0	2.83	1.97	Yes	168

Example result table for n = 200:

Algorithm	Conflicts	Time(ms)	Mem(KB)	Solved	Steps
CSP Backtracking	Skipped	0.00	0.00	Skipped	–
Min-Conflicts	0	86.55	12.57	Yes	232

Testing and Validation

4 different test cases in one graph



Conclusion: Observations and Notes

CSP Backtracking is slower on large boards due to exponential growth in possibilities.

Min-Conflicts performs well even for n = 1000 but may not guarantee a solution if max steps are reached.

Memory usage scales linearly with n for both algorithms.

AC-3 significantly reduces the search space for small boards.

Thank You!