



**COMPUTER SCIENCE AND DATA ANALYTICS**

Course: CSCI 6511 Artificial Intelligence

## **Project 2**

Project title: **CSP: N-Queens**

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# N-Queens CSP Project Report

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## Project Code

The complete source code for this project:  
<https://github.com/NamigPlanov/P2-Team3-AI>

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## Introduction

The **N-Queens problem** is a classic puzzle where the goal is to place  $n$  queens on an  $n \times n$  chessboard such that no two queens attack each other.

This project implements a **Constraint Satisfaction Problem (CSP)** solution for N-Queens using both **Backtracking with heuristics** and **Min-Conflicts iterative repair**.

The program can read an initial board from a file or generate a random one. It automatically chooses the suitable algorithm depending on board size ( $n \leq 50 \Rightarrow$  CSP Backtracking,  $n > 50 \Rightarrow$  Min-Conflicts).

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## Problem Formulation

The N-Queens problem is modeled as a CSP:

- **Variables:** Each row on the chessboard (0 to  $n-1$ ).
- **Domains:** Columns where a queen can be placed (0 to  $n-1$ ).
- **Constraints:** No two queens can share the same column or diagonal.

This CSP model allows the use of backtracking with heuristics or iterative repair to find solutions efficiently.

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## Algorithms

### a. CSP Backtracking

For smaller boards ( $n \leq 50$ ), **CSP Backtracking** is used.

Key components:

- **Minimum Remaining Values (MRV):** Selects the row with the fewest remaining valid columns.
- **Least Constraining Value (LCV):** Orders columns that minimize conflicts with other queens.
- **AC-3 (Arc Consistency):** Propagates constraints to reduce domains before trying values.

Steps:

1. Initialize domains for each row.
2. Reduce domains using AC-3.
3. Select unassigned row using MRV.
4. Order possible columns with LCV.
5. Assign queen and propagate constraints recursively until solution is found.

$n = 11$ :

```
Number of lines (10-1000): 16
Solving N-Queens for n = 16

Algorithm          | Heuristics          | Conflicts | Time(ms) | Mem(KB) | Solved | Steps
-----|-----|-----|-----|-----|-----|-----
CSP Backtracking   | MRV + LCV + AC3     | 0         | 174.34   | 799.02  | Yes   | 18

Solution (column index per row):
[15, 13, 11, 3, 5, 12, 1, 9, 0, 2, 14, 8, 10, 7, 4, 6]

Board visualization:
. . . . . Q
. . . . . Q .
. . . . . Q .
. . Q . . . .
. . . . Q . .
. . . . . Q .
. Q . . . . .
. . . . . Q .
Q . . . . .
. Q . . . . .
. . . . . Q .
. . . . . Q .
. . . . . Q .
. . . . . Q .
. . . . . Q .
. . . . . Q .
```

## b. Min-Conflicts (Iterative Repair)

For larger boards ( $n > 50$ ), **Min-Conflicts** is used:

- Start with a random assignment of queens.
- Identify rows where queens are in conflict.

- Move conflicted queens to positions that minimize conflicts.
- Repeat up to a maximum number of steps (by default it is 200,000).

This algorithm scales well because it avoids full backtracking and focuses only on conflicts.

$n = 333$ :

```
Number of lines (10-1000): 333
Solving N-Queens for n = 333

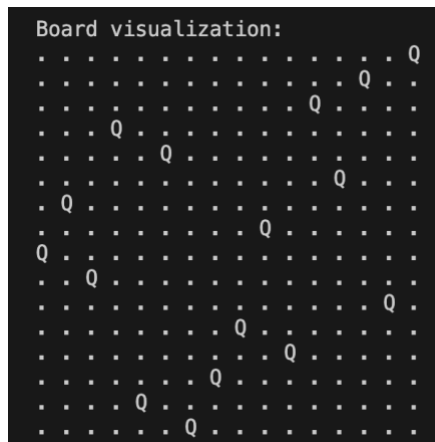
Algorithm      | Heuristics                  | Conflicts | Time(ms) | Mem(KB) | Solved | Steps
-----|-----|-----|-----|-----|-----|-----
Min-Conflicts  | Iterative Repair + Random TieBreak | 0         | 217.27   | 27.81   | Yes   | 230

Solution (column index per row):
[99, 76, 195, 235, 281, 197, 143, 43, 234, 332, 226, 156, 293, 125, 142, 112, 274, 278, 275, 78, 191, 81, 315, 128, 140, 18,
286, 185, 13, 220, 187, 229, 310, 55, 301, 105, 26, 158, 237, 212, 252, 122, 267, 228, 213, 146, 15, 254, 150, 41, 294, 103
, 263, 104, 2, 259, 232, 222, 86, 218, 246, 215, 147, 83, 170, 121, 320, 48, 163, 231, 202, 116, 89, 153, 6, 328, 186, 5, 16
5, 204, 127, 49, 131, 137, 50, 148, 32, 178, 92, 256, 180, 9, 74, 3, 151, 164, 66, 284, 262, 133, 37, 52, 248, 313, 182, 20,
73, 203, 244, 14, 230, 107, 31, 63, 12, 324, 209, 33, 65, 95, 11, 47, 108, 304, 139, 268, 270, 206, 130, 201, 249, 173, 287
, 241, 114, 132, 253, 311, 271, 61, 64, 319, 174, 16, 233, 327, 118, 261, 106, 302, 102, 314, 129, 211, 177, 290, 326, 240,
0, 123, 273, 285, 176, 266, 29, 59, 307, 17, 1, 330, 300, 57, 10, 62, 214, 184, 192, 79, 325, 205, 272, 117, 23, 331, 145, 3
16, 297, 309, 250, 21, 149, 70, 238, 243, 223, 42, 159, 264, 155, 56, 53, 305, 190, 22, 210, 82, 321, 162, 60, 247, 292, 152
, 27, 109, 7, 216, 289, 227, 189, 8, 24, 90, 19, 51, 288, 225, 93, 282, 115, 242, 296, 317, 91, 69, 322, 329, 255, 308, 58,
160, 279, 318, 217, 251, 45, 183, 80, 207, 46, 113, 199, 269, 75, 119, 100, 77, 124, 168, 299, 257, 141, 239, 35, 44, 157, 2
60, 200, 94, 303, 208, 291, 258, 283, 97, 136, 101, 306, 71, 30, 161, 154, 188, 194, 196, 40, 312, 85, 34, 144, 179, 68, 166
, 36, 193, 198, 54, 98, 181, 323, 67, 111, 280, 219, 236, 120, 175, 169, 221, 38, 298, 171, 295, 172, 134, 4, 277, 126, 87,
39, 72, 276, 84, 25, 265, 88, 245, 138, 96, 28, 224, 167, 110, 135]

Board not printed (n too large).
```

## Implementation Notes

- **Initial Board Input:** p2\_n-queen.txt (optional). If missing, the program asks for n and generates a random board.
- **Conflict Counting:** Function conflicts(board, row, col) checks column and diagonal attacks.
- **Memory & Performance Tracking:** Uses tracemalloc and time.perf\_counter() to measure memory and execution time.
- **Board Visualization:** Only displayed for  $n \leq 50$  to avoid huge output.



## Testing and Validation

- test.py validates algorithms for different board sizes.
- Measures:
  - Total conflicts
  - Execution time
  - Peak memory usage
  - Steps taken
- Results for small boards include full board output, but large boards skip visualization for readability.

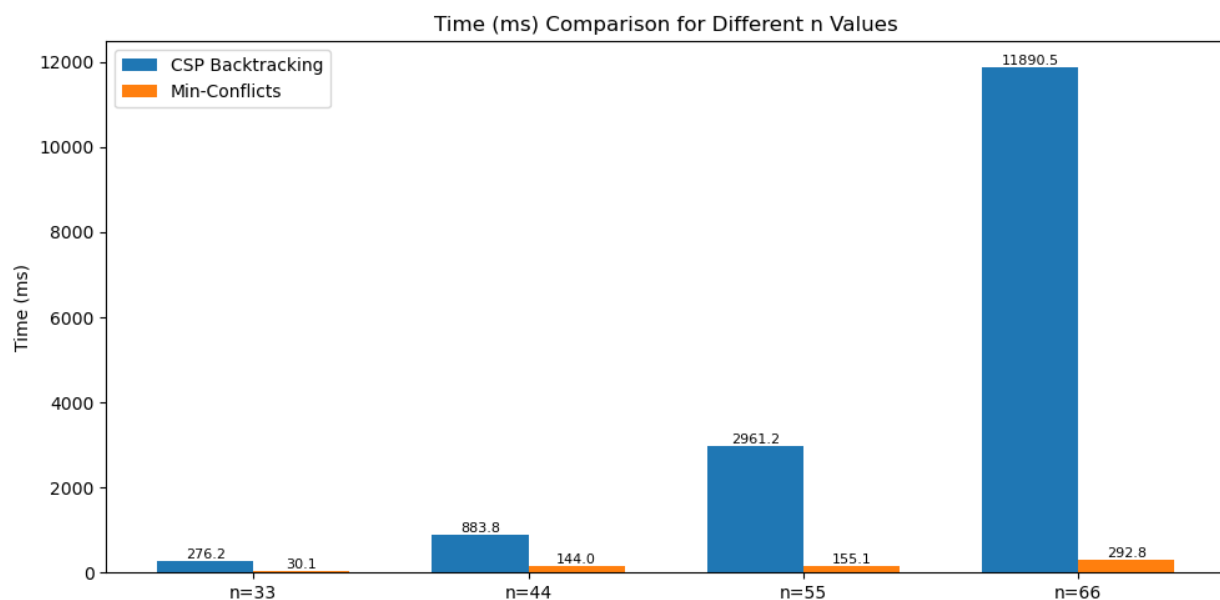
Example result table for  $n = 23$ :

Algorithm	Conflicts	Time(ms)	Mem(KB)	Solved	Steps
CSP Backtracking	0	912.59	2995.83	Yes	27
Min-Conflicts	0	2.83	1.97	Yes	168

Example result table for  $n = 200$ :

Algorithm	Conflicts	Time(ms)	Mem(KB)	Solved	Steps
CSP Backtracking	Skipped	0.00	0.00	Skipped	–
Min-Conflicts	0	86.55	12.57	Yes	232

Visualization of 4 different test cases:



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## Observations and Notes

- **CSP Backtracking** is slower on large boards due to exponential growth in possibilities.
- **Min-Conflicts** performs well even for  $n = 1000$  but may not guarantee a solution if max steps are reached.
- Memory usage scales linearly with  $n$  for both algorithms.
- AC-3 significantly reduces the search space for small boards.

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## Conclusion

This project successfully implements two CSP approaches for the N-Queens problem:

- Backtracking with heuristics and AC-3 for small boards.
  - Min-Conflicts iterative repair for large boards.
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