

Explanation

Hashing

Int -

Table - $a[n][2]$

1 2 3 2 4 1 4 5 count sort -- c

$C[3] = 1$

$C[key] = value$

`cout<<c[key];`

keys	values
1	2
2	2
3	1
4	2
5	1

Stl - `unordered_map< string , vector<int> > arr;`

`Arr[1] = 2;`

`Arr[3] =2;`

...

`arr.insert(make_pair(3,2));`

Keys - int, long int, char, string

Values - int, long int, char, string, vector, pair

Size - `table_size`

Abc bcd abcd abcd asd

Keys must be unique

Hashmaps -

Search

Insert

Delete -

$O(1)$ best , avg

Worst - $O(n)$ $O(1)$

Insertion

$H[key] = value$

`h.insert(make_pair(key,value))`

Deletion

`h.erase(key)`

Search

$H[key]$

`h.count(key)` 1 or 0

`h.find(key)` `h.end()` - if not present

How to find the index of key

Hash function - input key - return index

Table size = $n = 10$

Int,int

19,20

23,56

34,54

73, 76

$key \% n = 19 \% 10 = 9$

$23 == 3$

$34 == 4$

$73 ==$

Collision -

Open hashing, chaining, open addressing, linear probing, double hashing

Hash function

Modulo with prime number - 11

2 primes - p_1, p_2

P_2 = nearest to table size

P_1 = nearest to data/input size

Basic implementation

Hash function, double hashing example

Hashmaps - unordered_map array O(1)
Map - bst O(logn)

2 sum

Find the pairs which has their sum as target

Basic = $O(n^2)$

Sorting = $O(n \log n)$

Hashmaps = $O(n)$

Sum = 8

1 3 5 4 2 1 6 7

Int c=0;

unordered_map<int,int> h;

for(int i=0;i<n;i++)

{

 Int x = sum - a[i];

 if(h.count(x) > 0)

 C += h[x];

 H[a[i]] += 1;

}

C = 4 i = 7

1 2

3 1

5 1

4 1

2 1

Count subarrays with 0 sum

2 -4 2 4 -6 -3 2

2 -2 0 4 -2 -5 -3

Intersection of 2 arrays

Basic = $O(n^2)$

Sorting = $O(n \log n)$ 2 pointers = i and j

Hashmap = $O(n)$

Heaps -

Complete binary tree

0 based indexing

Curr = i

Parent = floor((i-1)/2)

Child = 2*i+1, 2*i+2

1 based indexing

Curr = i

Parent = floor(i/2)

Child = 2*i, 2*i+1

10 15 30 40 50 100 40

0 1 2 3 4 5 6

priority_queue = priority - max element

max-heap

priority_queue = priority - min element

min-heap

Explanation :

Heaps

Insert - add new node to end and up heapify

O(logn)

Delete - delete top node and down heapify

O(logn)

Space - O(n)

```
Priority-queue <int> pq;                // max heap
pq.push(0);
pq.pop();
Pq.top;
```

Class comp

```
{
```

```
Public:
```

```
    Bool operator()(int a,int b)
```

```
    {
```

```
        //comp
```

```
    }
```

```
};
```

```
Priority-queue <int, vector<int>, greater<int>> pq;
```

```
// min heap
```

```
Priority-queue <int, vector<int>, comp> pq;
```

```
// custom heap
```

Priority-queue <pair<int,int>, vector<pair<int,int>>, comp> pq; // custom heap
sort(a,a+n, greater<int>())

Heap sort

Push all elements in min heap
Pop one by one and add to array

K th max min element in array

Time - $O(n \log n + k \log n)$

Space - $O(n)$

Time - $O(n \log k + (n-k+1) \log k)$

Space - $O(k)$

Max element - min heap

K elements

1 pop - kth max

2nd element - k-1 th max

1, 2, 3, ..., n-k-1, n-k, n-k+1, ... n

Running stream

Input size - very large

Kth Max element in running stream

Merge k arrays in sorted arrays of same size

Brute force - add all to single array and sort

Min heap of k size

Value, ind, array ind

pair<int, pair<int, int>>

Struct
