

## CS323 Assignment 3

### 1 Requirements

You are expected to complete all required homework exercises and encouraged to complete the optional ones. For submission, please put all your answers in a single PDF file and submit it via the assignment channel on SAKAI. The name of the file should follow the format “studentID\_A#” (e.g., 30003554\_A3). **The submission deadline is 11:55 PM, November 1, 2022.** Late submissions are allowed within one week after the deadline (grace period). If you submit your assignment during the grace period, your score will be 80% of the score you could get if the submission was made in time. Assignment submitted after the grace period will not be graded, meaning that you will get a zero for the assignment.

### 2 Required Exercises (100 points)

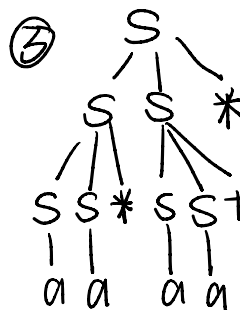
**Exercise 1 (Grammar Basics):** Consider the following context-free grammar  $G$ :

$$S \rightarrow SS + \mid SS * \mid a$$

1. Give a leftmost derivation for the string  $aa * aa + *$ . [10 points]
2. Give a rightmost derivation for the string  $aa * aa + *$ . [10 points]
3. Give a parse tree for the string  $aa * aa + *$ . [10 points]
4. Give an equivalent grammar without immediate left recursions. [10 points]
5. Is the grammar ambiguous? [10 points]

①  $S \Rightarrow SS*$   
 $\Rightarrow SS*SS*$   
 $\Rightarrow aS*SS*$   
 $\Rightarrow aa*S*$   
 $\Rightarrow aa*SS+$   
 $\Rightarrow aa*aS+$   
 $\Rightarrow aa*aa+$

②  $S \Rightarrow SS*$   
 $\Rightarrow SSS+*$   
 $\Rightarrow SSa+*$   
 $\Rightarrow Saa+*$   
 $\Rightarrow SS*aa+*$   
 $\Rightarrow Sa*aa+*$   
 $\Rightarrow aa*aa+*$



④  $S \rightarrow \alpha S'$   
 $S' \rightarrow S + S' \mid S * S' \mid \epsilon$

⑤ It is unambiguous

EX2: ① First:  $\text{First}(S) = \text{First}(aB) = \text{First}(SB) = \{a\}$   
 $S \rightarrow aB$   $\text{First}(B) = \{a, \epsilon\}$

Follow:  $\text{Follow}(S) = \{\$, *\}$   
 $B \rightarrow S * B \mid \epsilon$   $\text{Follow}(B) = \{\$, *\}$

For  $S \rightarrow aB$ ,  $a \in \text{First}(aB)$ ,  $\therefore$  put  $S \rightarrow aB$  into  $M(S, a)$

For  $B \rightarrow S * B$ ,  $a \in \text{First}(SB)$ ,  $\therefore$  put  $B \rightarrow SB$  into  $M(B, a)$ .

For  $B \rightarrow \epsilon$ ,  $\epsilon$  in  $\text{First}(B)$ , for each terminal in  $\text{Follow}(B)$ , that is

put  $B \rightarrow \epsilon$  into  $M(B, \$)$ ,  $M(B, *)$

	a	*	\$
S	$S \rightarrow aB$		
B	$B \rightarrow S * B$	$B \rightarrow \epsilon$	$B \rightarrow \epsilon$

② Yes

③

3. Can an LL(1) parser accept the input string  $aaaa***$ ? If yes, please list the moves made by the parser; otherwise, state the reason. Before parsing, please resolve conflicts in the parsing table if any. [20 points]

Matched	Stack	Input	action
	S\$	aaaa***\$	$S \rightarrow aB$
a	aB	aaaa***\$	match a
a	B\$	aaa***\$	$B \rightarrow S * B$
a	S*B\$	aaa***\$	$S \rightarrow aB$
a	aB*B\$	aaa***\$	match a
aa	B*B\$	aa***\$	$B \rightarrow S * B$
aa	S*B*B\$	aa***\$	$S \rightarrow aB$
aa	aB*B*B\$	aa***\$	match a
aaa	B*B*B\$	a***\$	$B \rightarrow S * B$
aaa	S*B*B*B\$	a***\$	$S \rightarrow aB$
aaa	aB*B*B*B\$	a***\$	match a
aaaa	B*B*B*B\$	***\$	$B \rightarrow \epsilon$
aaaa	*B*B*B\$	***\$	match *
aaaa*	B*B*B\$	**\$	$B \rightarrow \epsilon$
aaaa*	*B*B\$	**\$	match *
aaaa**	B*B\$	*\$	$B \rightarrow \epsilon$
aaaa**	*B\$	*\$	match *
aaaa***	B\$	\$	$B \rightarrow \epsilon$
aaaa***	\$	\$	

**Exercise 2 (Top-Down Parsing):** Consider the following grammar  $G$ :

$$S \rightarrow aB$$

$$B \rightarrow S * B \mid \epsilon$$

1. Construct the predictive parsing table for  $G$ . Please put down the detailed steps, including the calculation of FIRST and FOLLOW sets. [25 points]
2. Is the grammar LL(1)? [5 points]
3. Can an LL(1) parser accept the input string  $aaaa^{***}$ ? If yes, please list the moves made by the parser; otherwise, state the reason. Before parsing, please resolve conflicts in the parsing table if any. [20 points]

### 3 Optional Exercise (10 bonus points)

1. Justify your answer to Question 5 of Exercise 1. You do not need to provide a very rigorous proof. Informal explanations/examples are also acceptable.

For an ambiguous grammar, we can find two leftmost/rightmost derivation for it. Let  $w$  be any string that can be derived by the language. Suppose we want to find a rightmost derivation. When we choose production in each step, we should always consider the terminal that is currently unmatched and on the far right, if it is  $a$ , we choose  $S \rightarrow a$ ; if it is  $+$  we choose  $S \rightarrow SS+$  if it is  $*$ , we choose  $S \rightarrow aa*$ , so there are only one choice each step, so the parse tree constructed should be unique.