



# CS201 DISCRETE MATHEMATICS FOR COMPUTER SCIENCE

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# Review

- |                              |                            |
|------------------------------|----------------------------|
| 01. Propositional Logic      | 08. Cryptography           |
| 02. Predicate Logic          | 09. Mathematical Induction |
| 03. Mathematical Proofs      | 10. Recursion              |
| 04. Sets                     | 11. Counting               |
| 05. Functions                | 12. Relation               |
| 06. Complexity of Algorithms | 13. Graphs                 |
| 07. Number Theory            | 14. Tree                   |



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Discrete Probability

Groups, Rings and Fields



# Logic

- Logical connectives



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$$\neg p, p \vee q, p \wedge q, p \oplus q, p \rightarrow q, p \leftrightarrow q$$



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De Morgan's laws, commutative laws, distributive laws, ...

- Predicate logic

contains variables





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De Morgan's laws, commutative laws, distributive laws, ...

- Predicate logic

contains variables

- Quantified statements

universal, existential, equivalence



# Methods of Proving Theorems

## ■ Basic methods to prove theorems:

### ◇ *direct proof*

- $p \rightarrow q$  is proved by showing that if  $p$  is true then  $q$  follows

### ◇ *proof by contrapositive*

- show the contrapositive  $\neg q \rightarrow \neg p$

### ◇ *proof by contradiction*

- show that  $(p \wedge \neg q)$  contradicts the assumptions

### ◇ *proof by cases*

- give proofs for all possible cases

### ◇ *proof of equivalence*

- $p \leftrightarrow q$  is replaced with  $(p \rightarrow q) \wedge (q \rightarrow p)$



# Function

- function?



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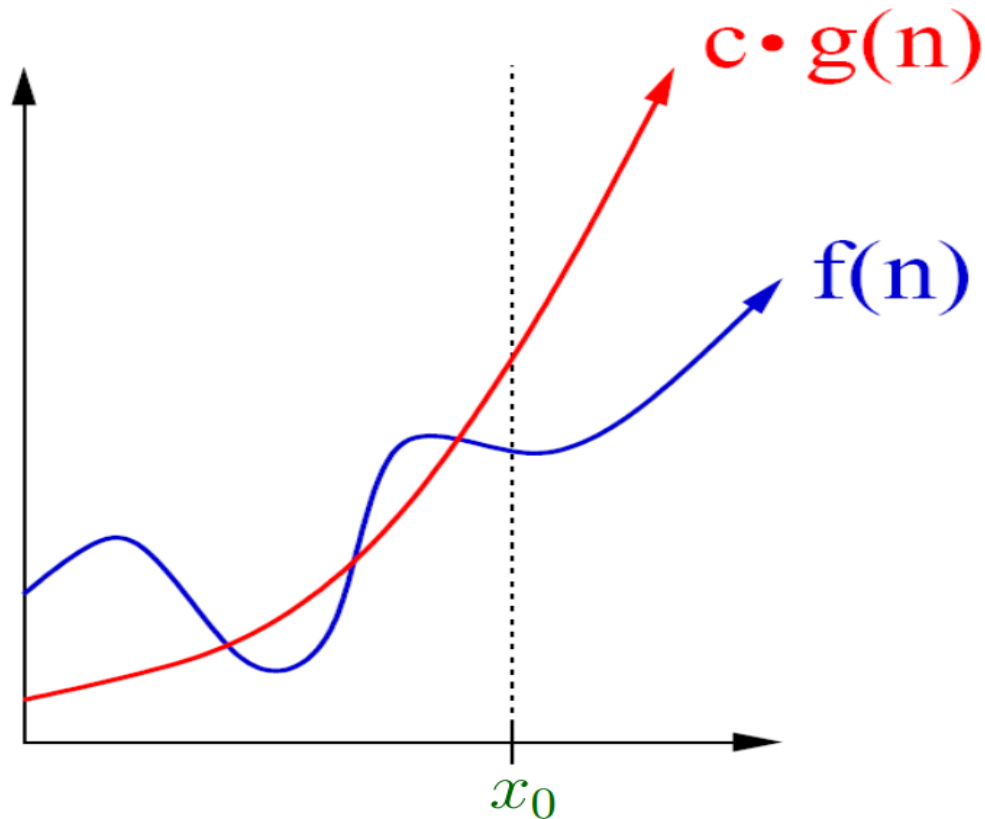
bijection function (one-to-one correspondence)?

- counting the number of such functions?



# Big- $O$ Notation

- Let  $f$  and  $g$  be functions from the set of integers or the set of real numbers to the set of real numbers. We say that  $f(n) = O(g(n))$  (reads:  $f(n)$  is  $O$  of  $g(n)$ ), if there exist some positive constants  $C$  and  $k$  such that  $|f(n)| \leq C|g(n)|$ , whenever  $n > k$ .





# Number Theory

- Divisibility



# Number Theory

- Divisibility

Congruence relation



# Number Theory

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Congruence relation

Primes



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GCD and Euclidean Algorithm



# Number Theory

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Congruence relation

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GCD and Euclidean Algorithm

Modular Inverse



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Modular Inverse

When does an inverse of  $a$  modulo  $m$  exist?

How to find inverses?



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Back substitution





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GCD and Euclidean Algorithm

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Chinese Remainder Theorem

Back substitution

$$\begin{aligned}x &\equiv 2 \pmod{3} \\x &\equiv 3 \pmod{5} \\x &\equiv 2 \pmod{7}\end{aligned}$$


# Cryptography

- Fermat's Little Theorem



# Cryptography

- Fermat's Little Theorem

## Euler's Theorem

Primitive roots, multiplicative order



# Cryptography

- Fermat's Little Theorem

Euler's Theorem

Primitive roots, multiplicative order

RSA cryptosystem

DLP, Diffie-Hellman protocol



# Mathematical Induction

- A *typical* proof by mathematical induction, showing that a statement  $P(n)$  is true for all integers  $n \geq b$  consists of three steps:



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2. We then,  $\forall n > b$ , show either

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or

$$(**) \quad P(b) \wedge P(b+1) \wedge \cdots \wedge P(n-1) \rightarrow P(n)$$



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3. We conclude on the basis of the principle of **mathematical induction** that  $P(n)$  is true for all  $n \geq b$ .



# Recurrence

- Iterating a recurrence



# Recurrence

- Iterating a recurrence  
bottom up or top down



# Recurrence

- Iterating a recurrence

bottom up or top down

prove by induction, complexity, ...



# Counting

- The sum rule and product rule



# Counting

- The **sum rule** and **product rule**  
The **Inclusion-Exclusion Principle**



# Counting

- The sum rule and product rule
- The Inclusion-Exclusion Principle
- The Pigeonhole Principle



# Counting

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The Inclusion-Exclusion Principle

The Pigeonhole Principle

**Theorem** If  $N$  is a positive integer and  $k$  is an integer with  $1 \leq k \leq n$ , then there are

$$P(n, k) = n(n-1)(n-2) \cdots (n-k+1)$$

$k$ -element permutations with  $n$  distinct elements.





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Pascal's Triangle, Identity

The Binomial Theorem, Trinomial



# Binary Relations

- Properties of relations



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Representing relations



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Equivalence relation

**Definition** A relation  $R$  on a set  $A$  is called an *equivalence relation* if it is reflexive, symmetric, and transitive.



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Partial ordering

**Definition** A relation  $R$  on a set  $A$  is called a *partial ordering* if it is reflexive, antisymmetric, and transitive.



# Graphs & Trees

- Basic concepts



# Graphs & Trees

## ■ Basic concepts

connected graph, simple graph, isomorphism, chromatic number, Euler circuit, Hamilton circuit, shortest path, bipartite graph, complete graph, special graphs ( $K_n$ ,  $K_{m,n}$ ,  $C_n$ ,  $W_n$ ), m-ary tree, tree traversal, spanning tree ...



# Next Lecture

- Good Luck!

