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Computer Organization

Lab4 MIPS(3) - Switch, Loop, Function

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Topics

- **Branch**
- **Loop**
- **Function, Call, Return**

How does CPU execute the instructions

- Before executing an instruction, CPU **fetches** it from memory according to its address, then **analyze**, finally **execute**.
- Register **PC** stores the address of the instruction which is to be executed.

pc 0x00400000

Text Segment				
Bkpt	Address	Code	Basic	
<input type="checkbox"/>	0x00400000	0x00004820	add \$9,\$0,\$0	5: add \$t1,\$zero,\$zero
<input type="checkbox"/>	0x00400004	0x20080000	addi \$8,\$0,0x00000000	6: addi \$t0,\$zero,0
<input type="checkbox"/>	0x00400008	0x200f000a	addi \$15,\$0,0x0000000a	7: addi \$t7,\$zero,10
<input type="checkbox"/>	0x0040000c	0x21080001	addi \$8,\$8,0x00000001	9: addi \$t0,\$t0,1
<input type="checkbox"/>	0x00400010	0x01284820	add \$9,\$9,\$8	10: add \$t1,\$t1,\$t0
<input type="checkbox"/>	0x00400014	0x010f082a	slt \$1,\$8,\$15	11: bgt \$t7,\$t0,calcu
<input type="checkbox"/>	0x00400018	0x1420ffff	bne \$1,\$0,0xffffffff	
<input type="checkbox"/>	0x0040001c	0x3c011001	lui \$1,0x00001001	13: <5> la \$a0,pstr_M0
<input type="checkbox"/>	0x00400020	0x34240004	ori \$4,\$1,0x00000004	
<input type="checkbox"/>	0x00400024	0x24020004	addiu \$2,\$0,0x00000004	<6> li \$v0,4
<input type="checkbox"/>	0x00400028	0x0000000c	syscall	<7> syscall
<input type="checkbox"/>	0x0040002c	0x00092021	addu \$4,\$0,\$9	14: move \$a0,\$t1
<input type="checkbox"/>	0x00400030	0x24020001	addiu \$2,\$0,0x00000001	15: li \$v0,1
<input type="checkbox"/>	0x00400034	0x0000000c	syscall	16: syscall
<input type="checkbox"/>	0x00400038	0x2402000a	addiu \$2,\$0,0x0000000a	18: <12> li \$v0,10
<input type="checkbox"/>	0x0040003c	0x0000000c	syscall	<13> syscall

Label	Address ▲
Sum1to10.asm	
calcu	0x0040000c

Q1: What is the difference between the addresses of two adjacent instructions?

Q2: How does the value in \$PC change?

Branch, Jump Instruction

Conditional branch

- **beq** \$t0,\$t1,lable // branch to instruction addressed by the label if \$t1 and \$t2 are equal
- **bne** \$t0,\$t1,lable // branch to instruction addressed by the label if \$t1 and \$t2 are NOT equal
- **blt, ble, bltu, bleu, bgt, bge, bgtu, bgeu**

Unconditional jump

Jump (j)	Unconditionally jumps to a specified location. A symbolic address or a general register specifies the destination. The instruction j \$31 returns from the a jal call instruction.
Jump And Link (jal)	Unconditionally jumps to a specified location and puts the return address in a general register. A symbolic address or a general register specifies the target location. By default, the return address is placed in register \$31. If you specify a pair of registers, the first receives the return address and the second specifies the target. The instruction jal procname transfers to procname and saves the return address. For the two-register form of the instruction, the target register may not be the same as the return-address register. For the one-register form, the target may not be \$31.

Branch

Are the running results of two demos the same ?

*Modify them without changing the result by using **ble** or **blt** instead*

```
.include "macro_print_str.asm"
.data
.text
    print_string("please input your score (0~100):")
    li $v0,5
    syscall
    move $t0,$v0
case1:
    bge $t0,60,passLable
case2:
    j failLable

passLable:
    print_string("\nPASS (exceed or equal 60) ")
    j caseEnd
failLable:
    print_string("\nFailed(less than 60)")
    j caseEnd
caseEnd:
    end
```

```
.include "macro_print_str.asm"
.data
.text
    print_string("please input your score (0~100):")
    li $v0,5
    syscall
    move $t0,$v0
case1:
    bge $t0,60,passLable
    j case2
case2:
    j failLable

passLable:
    print_string("\nPASS (exceed or equal 60) ")
    j caseEnd
failLable:
    print_string("\nFailed(less than 60)")
    j caseEnd
caseEnd:
    end
```

Loop

Compare the operations of loop which calculates the sum from 1 to 10 in java and MIPS.

Java:

```
public class CalculateSum{  
    public static void main(String [] args){  
        int i = 0;  
        int sum = 0;  
        for(i=0;i<=10;i++){  
            sum = sum + i;  
            System.out.print("The sum from 1 to 10 : " + sum );  
        }  
    }  
}
```

MIPS:

```
.include "macro_print_str.asm"  
.data  
    tdata: .word 0  
.text  
    add $t1,$zero,$zero  
    addi $t0,$zero,0  
    addi $t7,$zero,10  
calcu:  
    addi $t0,$t0,1  
    add $t1,$t1,$t0  
    bgt $t7,$t0,calcu  
  
    print_string ("The sum from 1 to 10 : ")  
    move $a0,$t1  
    li $v0,1  
    syscall  
  
    end
```

Demo #1

The code in the next page is expected to get 10 integers from the input device, and print it as the following sample

*Will the code get desired result?
If not, what happened ?*

```
please input an array (no more than 10 integer): 1
2
3
4
5
6
7
8
9
0

the arrayx is:1 2 3 4 5 6 7 8 9 0
-- program is finished running --
```

Demo #1

#piece 1/3

.include "macro_print_str.asm"

.data

arrayx: .space 10
str: .asciiz "\nthe arrayx is:"

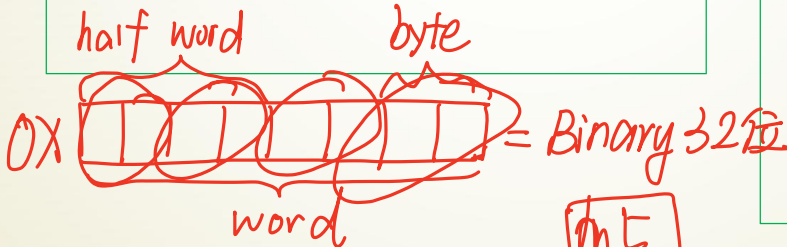
.text

main:

print_string("please input 10 integers: ")
add \$t0,\$zero,\$zero
addi \$t1,\$zero,10
la \$t2,arrayx

half word

byte



#piece 2/3

loop_r:

li \$v0,5
syscall
sw \$v0,(\$t2)
addi \$t0,\$t0,1
addi \$t2,\$t2,4
bne \$t0,\$t1,loop_r

la \$a0,str
li \$v0,4
syscall
addi \$t0,\$zero,0
la \$t2,arrayx

#piece 3/3

loop_w:

lw \$a0,(\$t2)
li \$v0,1
syscall
print_string(" ")
addi \$t2,\$t2,4
addi \$t0,\$t0,1
bne \$t0,\$t1,loop_w
end

The function of following code is to get 5 integers from input device, and find the min value and max value of them.

There are 4 pieces of code, write your code based on them, Can it find the real min and max?

```
#piece ?/4
.include "macro_print_str.asm"
.data
    min: .word 0
    max: .word 0
.text
    lw $t0,min
    lw $t1,max
    li $t7,5
    li $t6,0
    print_string("please input 5 integer:")
loop:
    li $v0,5
    syscall
    bgt $v0,$t1,get_max
    j get_min
```

```
#piece ?/4
get_max:
    move $t1,$v0
    j get_min
get_min:
    bgt $v0,$t0,judge_times
    move $t0,$v0
    j judge_times
```

```
#piece ?/4
judge_times:
    addi $t6,$t6,1
    bgt $t7,$t6,loop
```

```
#piece ?/4
    print_string("min : ")
    move $a0,$t0
    li $v0,1
    syscall
    print_string("max : ")
    move $a0,$t1
    li $v0,1
    syscall
end
```

Function

110 10

- ▶ **jal function_label**
 - ▶ #Unconditionally jump to the instruction at function_label. Save the address of the next instruction in register \$ra
 - ▶ Used in caller while calling the function
- ▶ **jr \$ra**
 - ▶ #Unconditionally jump to the instruction whose address is in register ra
 - ▶ Used in callee while returning to the caller
- ▶ **lw / sw with \$sp**
 - ▶ #protects register data by using stack in memory

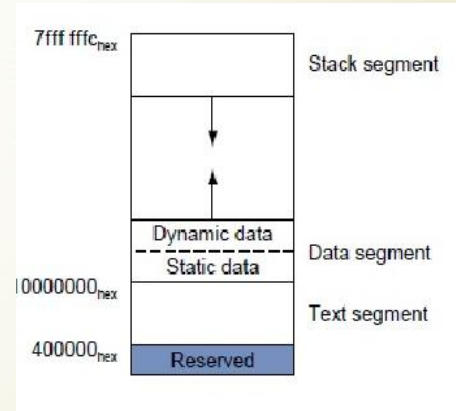
Stack Segment

stack segment: The portion of memory used by a program to hold procedure call frames.

The program **stack segment**, resides **at the top of the virtual address space** (starting at address **7fffffff_{hex}**).

Like dynamic data, the maximum size of a program's stack is not known in advance.

As the program **pushes values on the stack**, the operating system **expands** the stack segment **down**, **toward the data segment**.



Demo #2

read string

8

\$a0 = address of input buffer
\$a1 = maximum number of
characters to read

```
.data                                     #piece 1/3
tdata: .space 6
str1: .asciiz "the original string is: "
str2: .asciiz "\nthe last two character of the string is: "

.text
la $a0,tdata
addi $a1,$zero,6
addi $v0,$zero,8
syscall
```

```
print_string: #piece 3/3
    addi $sp,$sp,-8
    sw $a0,4($sp)
    sw $v0,0($sp)

    addi $v0,$zero,4
    syscall

    lw $v0,0($sp)
    lw $a0,4($sp)
    addi $sp,$sp,8
    jr $ra
```

```
la $a0,str1 #piece 2/3
```

```
jal print_string
```

```
la $a0,tdata
jal print_string
```

```
la $a0,str2
jal print_string
```

```
la $a0,tdata+3
jal print_string
```

```
addi $v0,$zero,10
syscall
```

Q1. IS it ok to remove the push and pop processing of **\$a0** on the stack in "print_string" ?

Q2. IS it ok to remove the push and pop processing of **\$v0** on the stack in "print_string" ?

Demo #2

What's the value of \$ra while jumping and linking to the print_string (at line 12,15,18,21)?

```

print_string:
    addi $sp,$sp,-8
    sw $a0,4($sp)
    sw $v0,0($sp)

    addi $v0,$zero,4
    syscall

    lw $v0,0($sp)
    lw $a0,4($sp)
    addi $sp,$sp,8

    jr $ra
  
```

Text Segment				
Bkpt	Address	Code	Basic	Source
<input type="checkbox"/>	0x0040001c	0x0c100013	jal 0x0040004c	12: jal print_string
<input type="checkbox"/>	0x00400020	0x3c011001	lui \$1, 0x00001001	14: la \$a0, tdata
<input type="checkbox"/>	0x00400024	0x34240000	ori \$4, \$1, 0x00000000	
<input type="checkbox"/>	0x00400028	0x0c100013	jal 0x0040004c	15: jal print_string
<input type="checkbox"/>	0x0040002c	0x3c011001	lui \$1, 0x00001001	17: la \$a0, str2
<input type="checkbox"/>	0x00400030	0x3424001e	ori \$4, \$1, 0x0000001e	
<input type="checkbox"/>	0x00400034	0x0c100013	jal 0x0040004c	18: jal print_string
<input type="checkbox"/>	0x00400038	0x3c011001	lui \$1, 0x00001001	20: la \$a0, tdata+3
<input type="checkbox"/>	0x0040003c	0x34240003	ori \$4, \$1, 0x00000003	
<input type="checkbox"/>	0x00400040	0x0c100013	jal 0x0040004c	21: jal print_string
<input type="checkbox"/>	0x00400044	0x2002000a	addi \$2, \$0, 0x0000000a	23: addi \$v0, \$zero, 10
<input type="checkbox"/>	0x00400048	0x0000000c	syscall	24: syscall

pay attention to the value of \$pc

Recursion

fact is a function to calculate the factorial.

C

```
int fact(int n) {
    if(n<1)
        return 1;
    else
        return (n*fact(n-1));
}
```

Q1. While calculate *fact(6)*, how many times does push and pop processing on stack happen?

Q2. How does the value of *\$a0* change when calculate *fact(6)*?

MIPS

fact:

```
addi $sp,$sp,-8    #adjust stack for 2 items
sw   $ra, 4($sp)   #save the return address
sw   $a0, 0($sp)   #save the argument n
```

```
slti  $t0,$a0,1    #test for n<1
beq   $t0,$zero,L1 #if n>=1,go to L1
```

```
addi $v0,$zero,1   #return 1
addi $sp,$sp,8     #pop 2 items off stack
jr    $ra          #return to caller
```

```
L1: addi $a0,$a0,-1 #n>=1; argument gets(n-1)
     jal  fact      #call fact with(n-1)
```

```
lw   $a0,0($sp)    #return from jal: restore argument n
lw   $ra,4($sp)    #restore the return address
addi $sp,$sp,8     #adjust stack pointer to pop 2 items
```

```
mul  $v0,$a0,$v0   #return n*fact(n-1)
```

```
jr   $ra          #return to the caller
```

Practice

1. Print out a 9*9 multiplication table.
 1. define a function to print $a*b = c$, the value of "a" is from parameter \$a0,the value of "b" is from parameter \$a1.
 2. less syscall is better
2. Get a positive integer from input, calculate the sum from 1 to this value by using recursion, output the result
3. Get a positive integer from input, output an integer in reverse order using loop and recursion seperately.
4. Answer the questiones on page 5,13 and14.

Tips

caller-saved register A register saved by the routine being called.

callee-saved register A register saved by the routine making a procedure call.

- ✓ Registers **\$a0~\$a3** are used to **pass the first four arguments to routines** (remaining arguments are passed on the stack).
- ✓ Registers **\$v0~\$v1** are used to **return values from functions**.
- ✓ Registers **\$t0~ \$t9** are **caller-saved registers** that are used to hold temporary quantities that need not be preserved across calls.
- ✓ Registers **\$s0~\$s7** are **callee-saved registers** that hold long-lived values that should be preserved across calls.
- ✓ Register **\$sp (29)** is the **stack pointer**, which points to the last location on the stack.
- ✓ Register **\$fp (30)** is the frame pointer.
- ✓ The jal instruction writes register **\$ra (31)**, the return address from a procedure call.

Tips : macro_print_str.asm

```
.macro print_string(%str)
    .data
    pstr: .asciiz %str
    .text
    la $a0,pstr
    li $v0,4
    syscall
.end_macro

.macro end
    li $v0,10
    syscall
.end_macro
```

get help from Mars about definition and usage of macro.