

Q1) What are the main advantages and disadvantages of the layered approach to system design.

### Advantages

Layering makes it easier to enhance the operating system. Implementation of a layer can be changed only without affecting the other layers.

\* It is very easy to perform debugging and system verification.

### Disadvantages

\* In this structure the application performance is degraded as compared to simple structure.

\* It requires careful planning for designing the layers as higher layers can't function unless the lower layers.

Q2) Differentiate b/w preemptive and nonpreemptive algorithm.

#### Preemptive

In this resources (CPU cycle) are allocated to a process for a limited time.

#### Non Preemptive

One resources (CPU cycle) are allocated to a process. It holds it till it completes its burst time or switches to waiting time.

Interrupt process can be interrupted in a process cannot be interrupted until it terminates itself or its time up.

b/w  
non preemptive it has overhead of scheduling & it does not have overheads of the processes.

#### Non Preemptive

Flexibility flexible

Eg: FCFS, SJF

eg: Round Robin, shortest remaining

time first

Q3) Some CPU provides more than two modes of operations. What are the use of the modes?

Although most S/W only distinguishes b/w user and kernel mode some CPU's have supported multiple modes which could be used to provide a security policy.

for eg, rather than cur. and kernel modes you could distinguish  
with b/w different types of cur mode

1. ~~FORK~~ # include <stdio.h>  
~~FORK~~ # include <unistd.h>

unshared mode instances

id (fork() off fork())

fork();

printf ("FORK");

return;

while the o/p of the program

2. ~~FORK~~  
~~FORK~~

5) what are the two models of ipc ? what are the strengths and  
weakness of these two approaches.

i) shared m/y message passing.

Shared m/y:- advantage of shared m/y model is that may communication is faster as compared to the msg passing model on the some machine. however, it may create problems such as synchronization and m/y protection that need be addressed.

message passing: it is easier to build II type, message passing is quite tolerate of higher communication latencies. It is slower than the shared m/y model because the connection set up time.

6) list and explain any five services provided by an operating system

#### Program Execution

→ reading program into main memory.  
→ all need of resources and allocations are identified,  
→ process synchronization at the concurrent programs.

→ process scheduling.

→ avoid deadlocks using process synchronization.

#### I/O operation

→ os. provides the access to required I/O devices when it is requires.

No handling operations means Read or write operation with any file or any I/O specific devices.

### File handling

- the OS gives the permission to the program for operation on files
- permission varies from read only or write only or read/write denied and OS provides an interface to the user to create and delete files.

### Error handling

Error can occur any time anywhere. and errors may occur in CPU in I/O devices or in the memory b/w. Some major activities of an OS with respect to error-handling are OS constantly checks the possible errors or takes appropriate action to ensure correct and constantly computing.

### Protection and Security

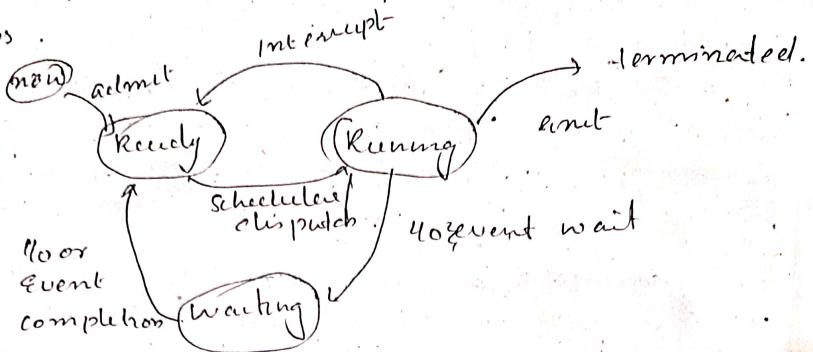
Considering a Computer system having multiple users and concurrent execution of multiple programs the various process must be protected from each other activities. Protection refers to an mechanism or a way to control the access of programs, processes or users to the resources define by a computer system. following are the major activities as required to protection.

B) Explain the actions taken by an OS to perform a context - switching b/w processes.

- \* save the context of the process that is currently running on the CPU: update the process control block and other important fields.
- \* move the process control block of the above process into the relevant queue such as the ready queue, I/O queue.
- \* Select a new process for execution.
- \* update the PCB of the selected process: this includes updating the process state to running.
- \* update the memory management data structures as required.
- \* Restore the context of the process that was previously

running when it's loaded again" on the processor. This is done by loading the previous value of the process control block and registers.

TA) with help of diagram, Explain different states of a process.



As a process executes, it changes states. Each process may be in one of the following states.

- ① New: the process is being executed.
- ② Ready: the instruction are being executed.
- ③ Running: the process is waiting for some event to occur.
- ④ Blocked: the process waiting to be assigned to a process.
- ⑤ Terminated: the process has finished execution.

TB) Explains the information held by the PCB associated with each process.

process state: the state may be New, Ready, Running, waiting, blocked, and so on.

PC: the counter indicates the address of the next instruction to be executed for this process.

CPU Registers: the Registers vary in no. and type, depending on the computer architecture. They include accumulators, index Registers, stack pointers, and general-purpose registers, plus any condition code information. Along with the program Counter, this state information must be saved when an interrupt occurs, to allow the process to be continued correctly afterwards.

CPU scheduling information: this information includes a process priority pointer to scheduling queues, and many others.

## Scheduling Parameters

Memory management information: this information may include such items as the values of the base and limit registers and the page tables, or the segment tables, depending on the memory system used by the operating system.

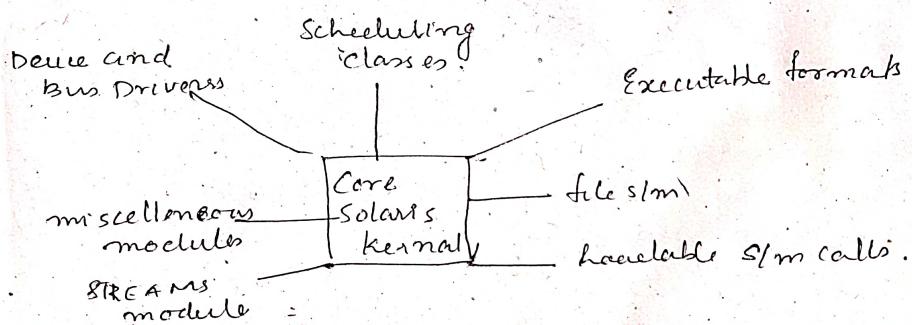
Q What are the disadvantages of the micro-kernel approach to S/m design? How do the user programs and S/m service interact via micro-kernel architecture?

Disadvantages of microkernel structure

\* increased level of inter module communication degrades S/m performance.

Q3) Explain kernel modules approach to S/m design with advantages.

- \* It is considered as the best approach for an OS.
- \* It involves designing a modular kernel.
- \* The kernel has only a set of core component and other services are added as dynamically loadable modules to the kernel either during runtime or boot time.
- \* It resembles layered structure due to the fact that each kernel has defined and protected interfaces but it is more flexible than the layered structure as a module can call any other module.
- \* For example Solaris OS is organized as shown.



#### Advantages of modular structure

- \* Advantages of layered structure but with more flexible.
- \* Advantages of microkernel approach - without message passing overhead.

#### Disadvantages of modular structure

- \* not as clean a design as the layered approach.
- \* not as small a kernel as a microkernel.
- \* But achieves best of both worlds as far as possible.

10 A) Explain any three activities of an operating s/m with regard to memory management.

1. Keeping track of which portion of mly are currently being used and who is using them.
2. Determining which process and data to move in and out of mly.
3. Allocatn and de-allocatn of mly blocks as needed by the program in mean mly (Caching collection).

10 B) Explain any five activities of an operating s/m with regard to process management.

- 1) the creation and delet of both user and s/m processes.
- 2) the suspension and resumption of processes.
- 3) The provision of mechanisms for process communication.
- 4) the provision of mechanism for process synchronization.
- 5) the provision of mechanism for dead lock handling

8 A) consider the FCFS, SJF and RR (Quantum = 10 milisecond) scheduling algorithms for the set of given processes, which algorithm would give the minimum average waiting time.

process      Burst time (ms.)

P<sub>0</sub>            10

P<sub>1</sub>            29

P<sub>2</sub>            3

P<sub>3</sub>            7

P<sub>4</sub>            12

FCFS

P <sub>0</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>3</sub>	P <sub>4</sub>
10	39	42	49	61

SJF

P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>	P <sub>0</sub>	P <sub>1</sub>	
3.	10	32	48	61	

RR

P <sub>0</sub>	P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>	P <sub>1</sub>	P <sub>1</sub>	P <sub>4</sub>
10	20	33	30	40	50	59	61

BB