Comprehensive Exam - Data Structure and Algorithm

Note: All questions are compulsory.

Total:-50 Marks

1. Prove or disprove the following:

(a)
$$n! = O(n^n)$$
 [1-marks]

(b) If
$$f(n) = O(g(n))$$
 and $g(n) = O(h(n))$
then, $h(n) = \Omega(f(n))$ [1-marks]

(c)
$$f(n) = \log_2 n$$
 is $O(n^{\alpha})$ for any $\alpha > 0$ [2-marks]

Note: All O's are (big-O) notation.

2. With the help of mathematical induction show that when n is an exact power of 2, the recurrence function:

$$T(n) = \begin{cases} 2, & \text{if } n = 2\\ 2T(n/2) + n, & \text{if } n = 2^i \text{ for } i > 0 \end{cases}$$

solution comes out to be $T(n) = n \log n$

[2-marks]

- 3. Given an array A with distinct element and sorted in decreasing order, a Quick-Sort is applied to it, show the running time of Quick-short is $\Theta(n^2)$. What will be running time of Quick-Sort when all elements of array A have same value. [2-marks]
- 4. If a queue Q is given, how many stacks are needed to replicate the queue Q, Implement it, and also give the running time of the implementation.

 [2-Marks]

- 5. Mr. B. C. Dull claims to have developed a new data structure for priority queues that supports the operations Insert, Maximum, and Extract-Max all in O(1) worst- case time. Prove that he is mistaken. [2-marks]
- 6. Give asymptotic upper and lower bounds for T(n) in each of the following recurrences. Assume that T(n) is constant for sufficiently small n. Make your bounds as tight as possible, and justify your answers.

(a)
$$T(n) = 4T(n/2) + n^2\sqrt{n}$$
 [2-marks]

(b)
$$T(n) = T(n/2) + T(n/4) + T(n/8) + n.$$
 [2-marks]

- 7. A Minimum Bottleneck Spanning Tree of an undirected graph G(V, E) is a spanning tree whose maximum weight edge is minimized. Explain how an MST is always a minimum bottleneck spanning tree, but the converse may not be true. [3-marks]
- 8. Given two BST's how do you check if they represent the same set of the elements? [3-marks]
- 9. Show a simple way to compute the \sqrt{n} smallest element among n elements in O(n) deterministic time, without using the Median of Medians algorithm. [4-marks]
- 10. Given a connected graph G, with n vertices and m edges, assume that the cost of all edges are distinct. A particular edge e of G is specified. Give an algorithm with running time O(m+n) to decide whether e is contained in a minimum spanning tree of G. [4-marks]
- 11. Bob algorithm enthusiastic wants to modify Strassen's algorithm to multiply $n \times n$ matrices in which n is not an exact power of 2? You as Bob friend derive an algorithm, and also show that the resulting algorithm runs in time $O(n^{\ln n})$. [4-marks]
- 12. Prove that, For any two nodes s and t in a directed graph G, their strong components are either identical or disjoint. [4-marks]
- 13. Show how to construct the min-heap in O(n) time. In other words, show how an arbitrary array can be transformed to make sure that it satisfies the min-heap property in O(n) time. [4-marks]

- 14. Consider an undirected graph G = (V, E) with a weight function W providing non-negative real valued weights, such that the weights of all the edges are different. Prove that, G has a unique Minimum Spanning Tree. [4-marks]
- 15. Topological short on a directed acyclic graph G=(V,E) can also be performed by repeatedly find a vertex of in-degree 0 (no incoming edges), output it, remove it and all of its outgoing edges from the graph. Explain how to put this concept into action so that it runs in time O(V+E). What happens if G has cycles in this algorithm?

[4-marks]