## Assignment no-6



## Title - The Dictionary ADT

Problem Statement - Implement all the functions of a dictionary using hashing Data: set of (key, value) pairs, keys are mapped to values, keys must be comparable, keys must be unique standard operations: Insert (key, value), find (key), Delete (key)

Objective - To understand the implementation of all the functions of a dictionary and standard operations on Dictionary.

Outcome - At the end of this assignment students will be able to perform standard operations on Dictionary ADT

Theory.

The Dictionary ADT: A dictionary is an ordered or unordered list of key-element pairs; where keys are used to locate elements in the list

Dictionary is a data structure, which is generally an association of unique keys with some values. One may bind a value to a key, delete a key (and naturally an associated value) and look up for a value by the key.

Values are not required to be unique.

A Dictionary can be implemented in various ways:
using a list, binary search tree, hash table etc

Hashing: Hashing is a technique to convert a range of tey values into a range of indexes of an array

by using hash functions. The values are then stored in a

data structure called hash table The idea of hashing
is to distribute entries uniformly across an array. Each
element is assigned a key by using that key you
can access the element in o(1) times.

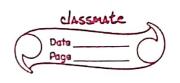
Hash function. A hash function is any function that can be used to map a data set of an arbitary set to a data set of a fixed size, which falls into the hash table. The values returned by a hash values hash codes, hash sums or simply hashes.

Hash Table - A hash table is a data structure that is used to store key/value poirs. It uses a hash function to compute an index into an array in which an element will be inserted or searched.

Collision Handling - Since a hash function gets us a small number for a big key, there is possibility that two keys, result in same value. The situation where a newly inserted key maps to an already occupied slot in hash table is called callision.

Following are the ways to handle collision

O Chaining - The ideo is to make each cell of hash table point to a linked list of records that have same hash function value chaining is simple, but requires additional memory outside the table.



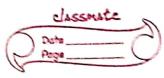
O open Addressing - In open addressing all elements are stored in the hash table itself Each table entry contains either a record or NULL. When searching for an element, we one by one examine table slots until the desired element is found or it is dear that the element is not in the table.

Implementation of Hash Table

Consider a dictionary where keys are integers in the range [O, N-1]. Then an array of size N can be used to represent the dictionary Each entry in this array is thought of as a "bucket" An element 'e' with key 'k' is inserted in A[k]. Bucket entries associated with keys not present in dictionary contains a special No such KEY object If the dictionary confains elements with the same key then two or more different elements may be mapped to the same bucket. of A. In this case, we say that a collision between these elements has occurred one easy way to deal with collision is to allow a sequence of elements with the same key, k, to be stored in ACE? Assuming that an arbitary element with key k satisfy queries find Item (k) and removeitem (k). these operation are now performed in o(1) time, while insertitem (k,e) needs only to find where an the existing list o(k) to insert the new item, e The drawback of this is that the size of the bucket set from which key are drawn, which way be huge.

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Algorithm.
HashNode class Declaration
 class Hash Node ?
 public:
  int key;
   int value;
   HashNode * next;
  HashNode (int key, int value) ?
     this ->key = key;
this -> value = value;
  this - next = NULL;
Insertions
void insert (int key, int value) ?
int k = key-hash (key);

if (arr (k)!=1) ?
   int t = traverse(k);
    orr[t] ->key=key;
    arr [t] -> meaning = value; 3
    orr [k] -> key = key;
    gir [k] -> meaning = value;
```



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neletion.
   void remove (int key) E
      int k = key_hash(key);
if (arr[k] \rightarrow key = = key)?
          arr[k]->key = -1;
           arr[k] -> meaning = " ";
       else f
         int t = traverse(k):
          if (arr[t]→Ley==Ley) {
         orr(t) → Ley = -1;
         arr[t] → meaning = " "; }
           else ?
         cout < " key not found";
Sear ch
       void search (int key) ?
        int k = key hash(key);
if (arr(k) → key = =key)
         cout <= "tey found";
         else E
           int t = traverse(k);
         if (arr (t] -> Ley = = Ley)

cout << "key found";
           else
           cout <= "key not found";
```

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	Testcases			ncas ora proteomicularicin fasticinis i substituti succe
	Test case	Expected	Out come	Result
0	Insert - 25,26,35,	0 - 30	some as	6955
	30,36,40,41,49,39,			
	33	2 - 41		
	Hash function:	3 - 3 g	5 5 45	1 1
	key 7.10	4-33	eta tuli	þ
	3	5 - 25		
		6 - 26	3 31 -	
		7 - 35	1 12 12-	
		8-36	ř. –	
	. A tree	9-49	And I was	
2	Insert - 21,24,30,	0 - 30	same as	pass
	09,04,14,28,18		expected	
	15,12		-	a e e f. e
		3 - 12	trees her	
		4 - 24	N 2 2 mil	
		5 - 4	and the same of	
		6 - 14	-3	
		7 - 15	E Sas	
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	Conclusion - Afk	d inclease	y completing	this assign
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