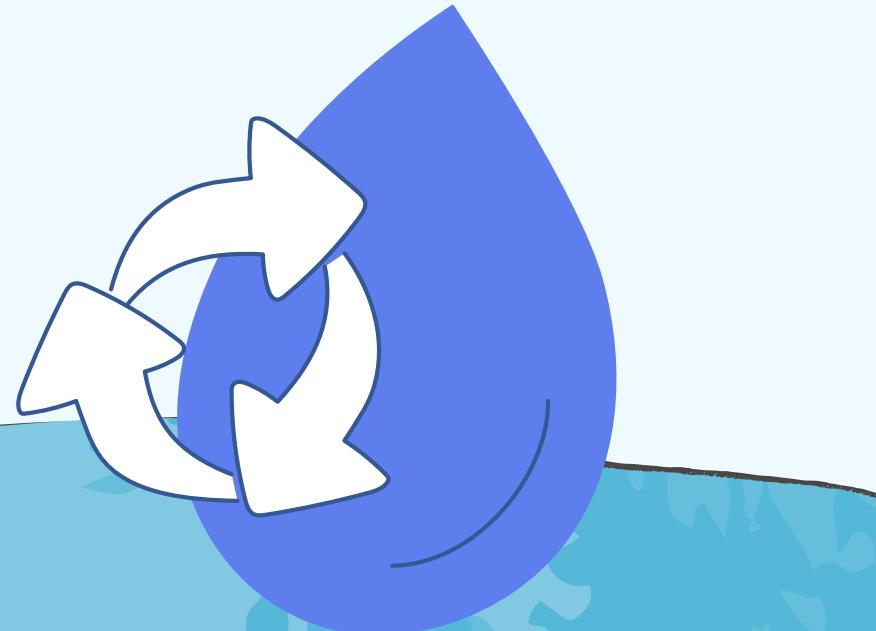
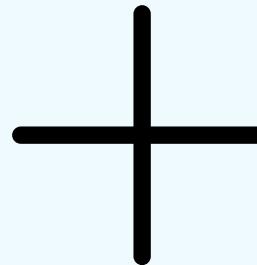
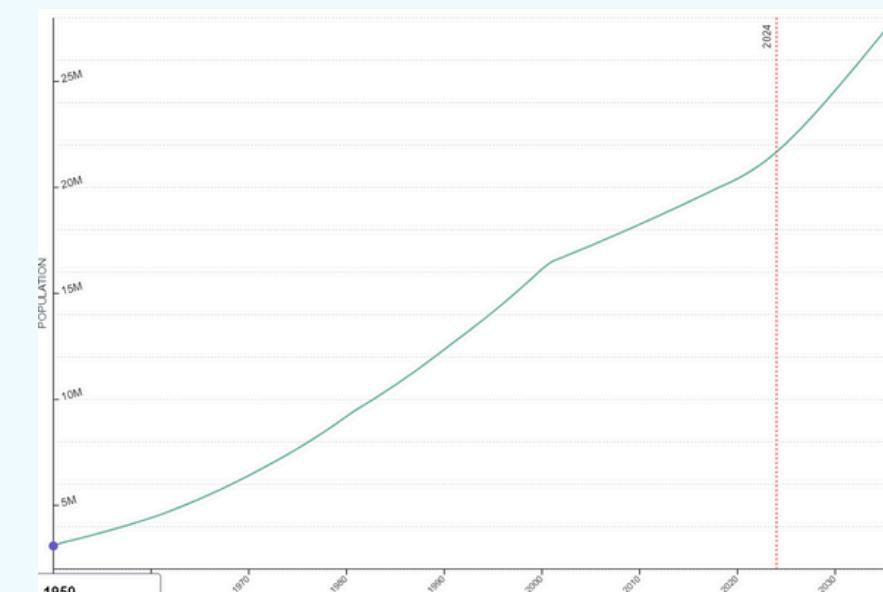


# ILP MODELS FOR WATER DISTRIBUTION SYSTEM OPTIMISATION

A PRESENTATION BY  
SPARSH AMARNANI



# Motivation



## Thane Water Crisis: City Faces Water Shortage as Main Pipeline Undergoes Emergency Repairs

By Nirmeeti Patole | Published: July 10, 2024 02:35 PM



The Thane Municipal Corporation (TMC) area is experiencing a perfect storm of water supply issues, as a newly detected ...

## Water cut announced for Mumbai

Richa Pinto / TNN / May 25, 2024, 16:40 IST



Mumbai faces water shortage with BMC announcing 5% and 10% water cuts. Citizens urged to conserve water until sufficient rainfall improves reservoir levels. Indian Meteorological Department forecasts timely monsoon arrival. Water cuts also apply to Thane, Bhiwandi-Nizampur Municipal Corporation. BMC provides list of water-

**Tanker economy revealed:  
Rajasthan's groundwater level takes  
a dip as water mafias emerge**  
Water shortages in Rajasthan need no introduction; but the extent of the crisis covers a lot more than what meets the eye

## Rapid Urbanisation

## Variable Water Supply

## Water problems



# WHAT IS WDS OPTIMISATION ?

## DECISION VARIABLE

Hydralic Components that we can choose like ,pipes, pumps and valves

## OBJECTIVE FUNTION

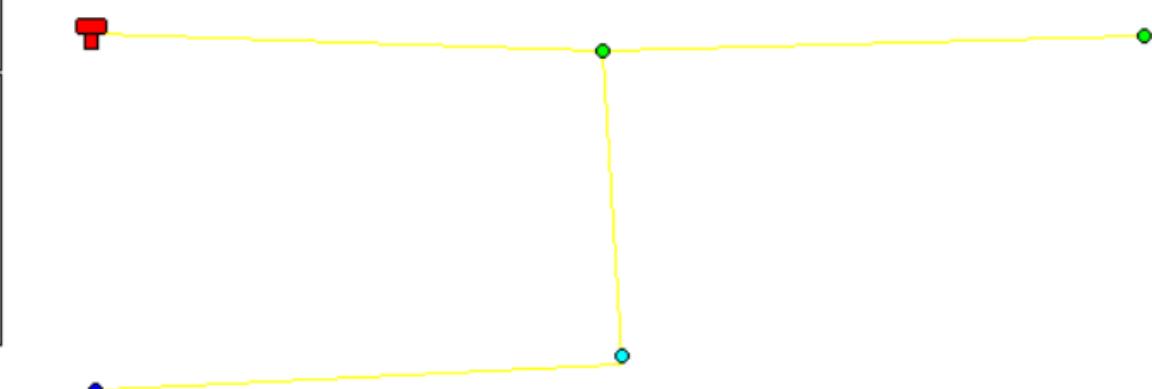
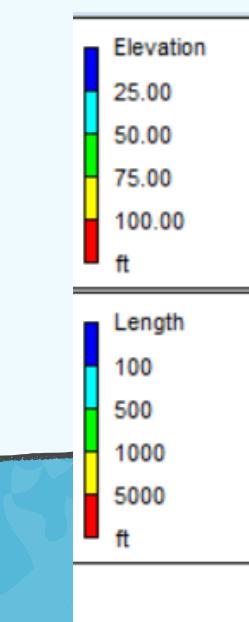
Variables that are dependent on our choices that we want to minimize or maximize, like cost, water quality or greenhouse gas emmisions.

## CONSTRAINTS

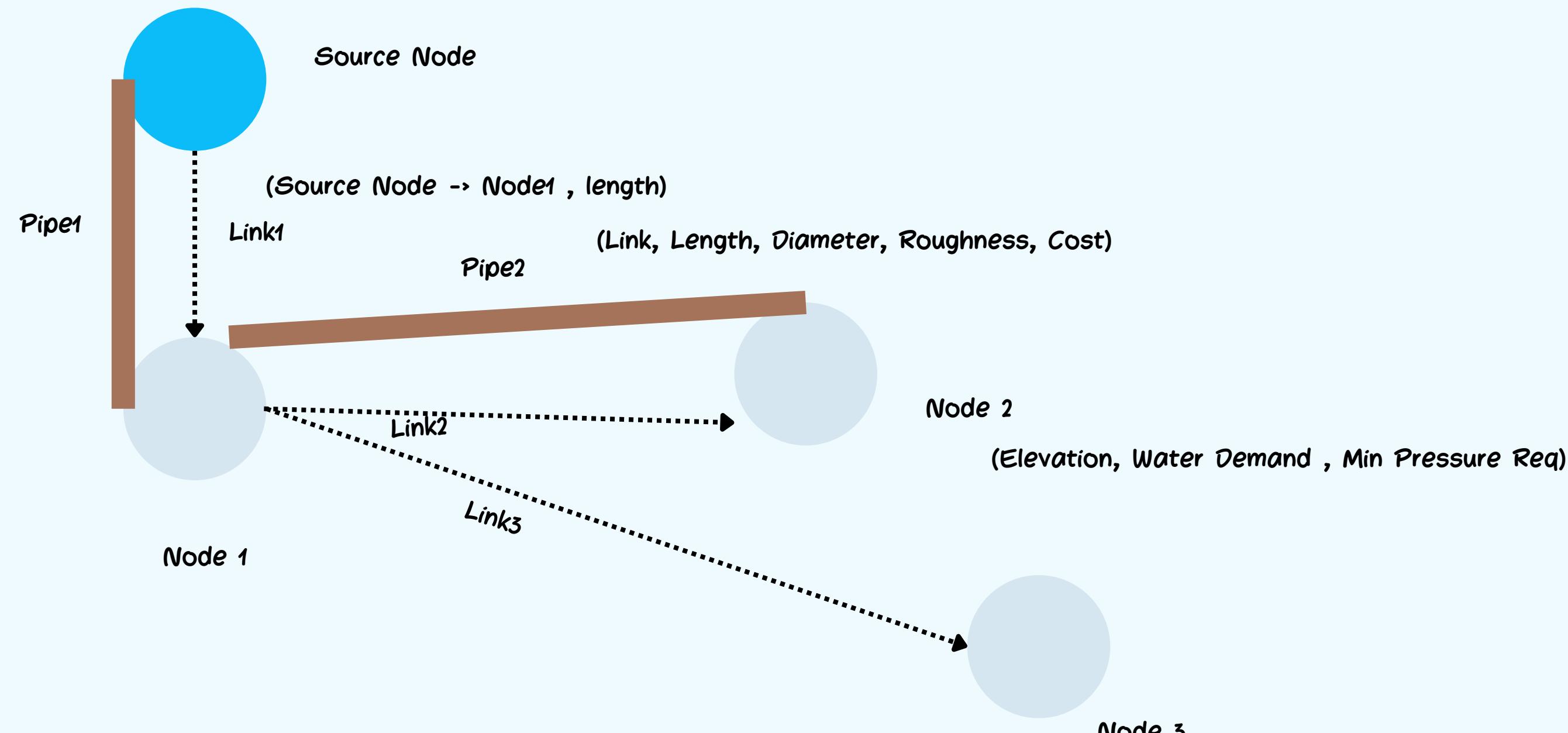
Hydraulic constraints are given by physical laws governing the fluid flow in a pipe network.

System Constraints like water demands and water quality requirement

THE GOAL IN DESIGNING A SOFTWARE WHICH CAN OPTIMISE THIS IS THAT WE CAN UPDATE OUR WATER NETWORKS WITH INCREASING WATER DEMANDS OR CHANGE IN WATER SUPPLY



# Inputs of the Optimisation Model



## Tank Table



Base Cost, Unit Cost, Upper Limit Capacity, Lower Limit Capacity

Tank 1



Tank 2



Tank 3

Other tank inputs : Maximum, Min Tank Height allowed,

Pump Cost, Pump Life Cycle, Minimum Pump Size, Efficiency

Pumps



Pump 1

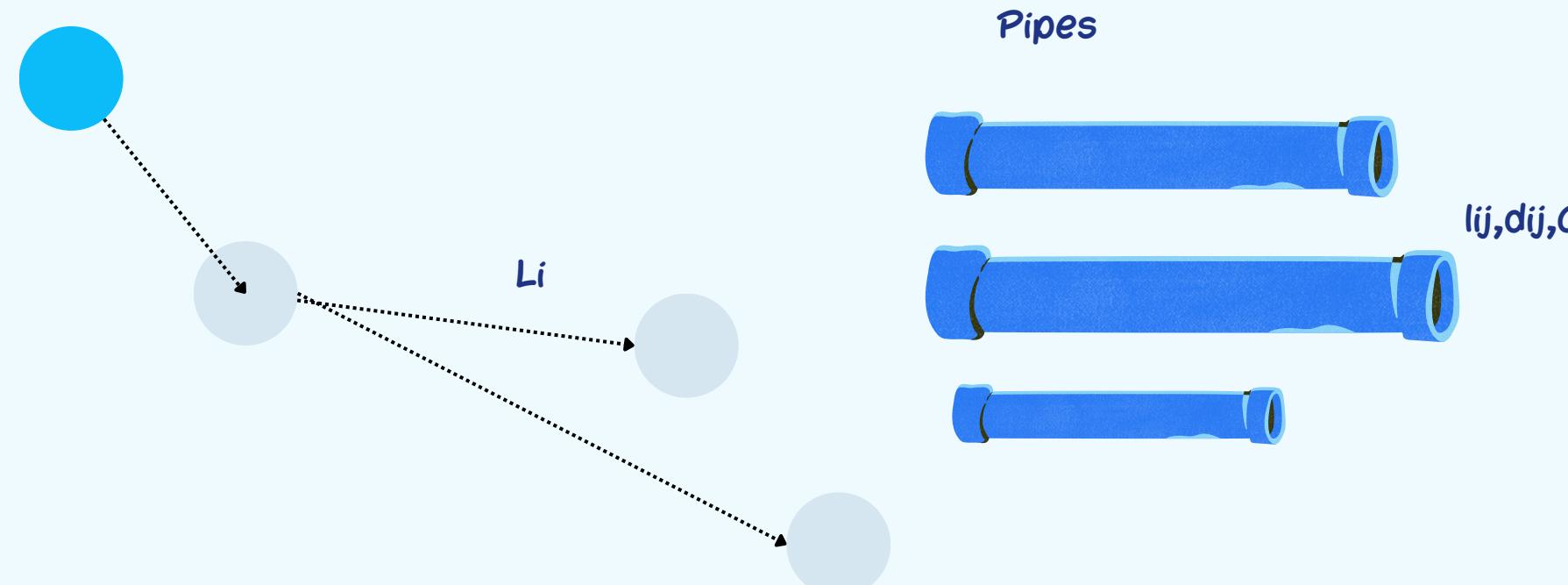


Pump 1

# A SYSTEM FOR OPTIMAL DESIGN OF PRESSURE CONSTRAINED BRANCHED PIPED WATER NETWORKS

N. Hoodaa, O. Damanib

# PIPE ONLY MODEL OPTIMISATION



## OBJECTIVE FUNCTION

minimize  $L : C(D_{i,j}) * L_{ij}$  sum  
over  $i$  and  $j$

## Length of Pipes

$$\begin{aligned} l_{ij} \text{ sum over } j &= L_i \\ l_{ij} &\geq 0 \end{aligned}$$

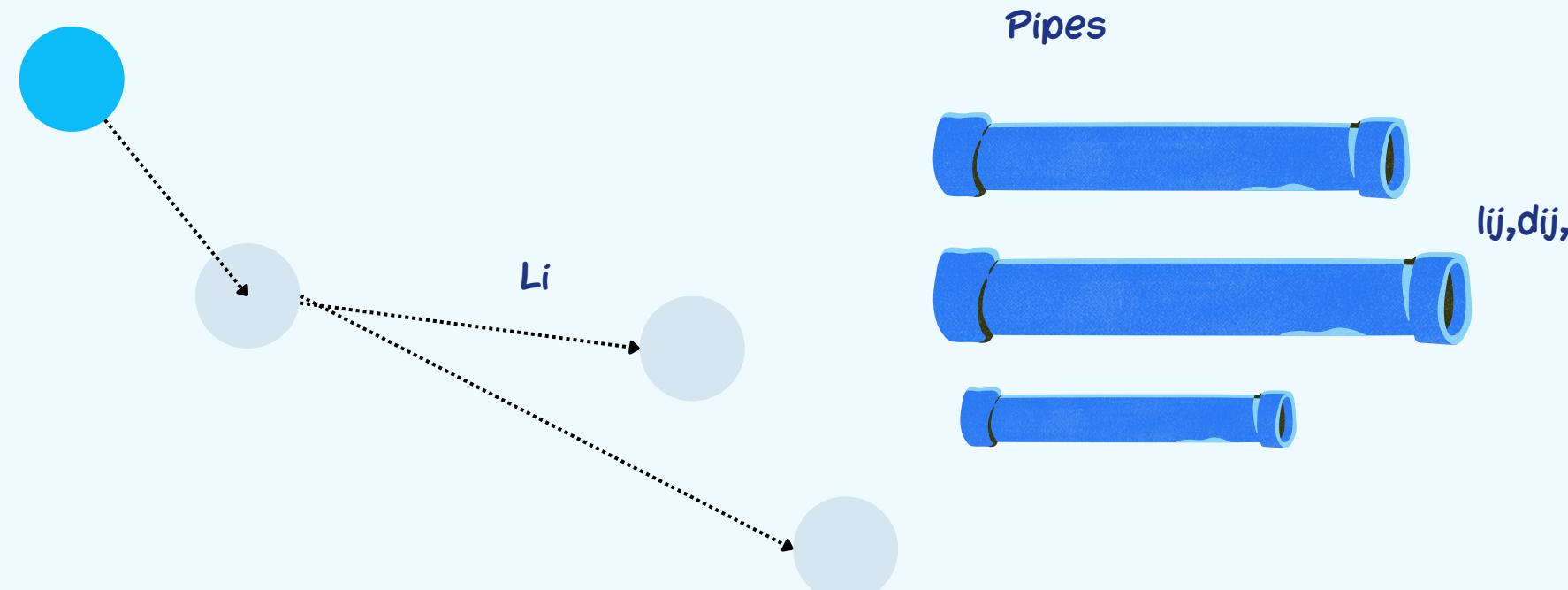
## CONSTRAINTS

Pressure requirement  
at each node

$$P_n \leq H_r - E_n - (H L_{ij} * L_i) \text{ sum over } j \text{ and } i \text{ which are parent links}$$

OUR GOAL IS TO  
PUT PIPES IN EACH  
LINK WHICH WILL  
MINIMIZE THE COST  
OF THE NETWORK

# HAZEN WILLIAMS EQUATION



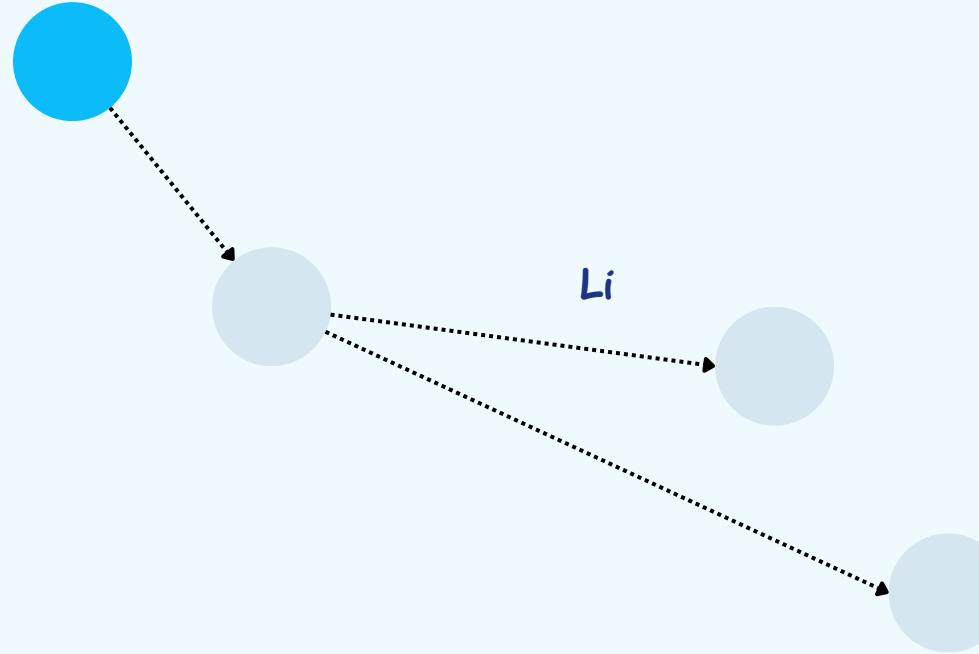
OUR GOAL IS TO  
PUT PIPES IN EACH  
LINK WHICH WILL  
MINIMIZE THE COST  
OF THE NETWORK

$$HL'_{ij} = \frac{10.68 * \frac{flow_i}{roughness_j}^{1.852}}{diameter_j^{4.87}}$$

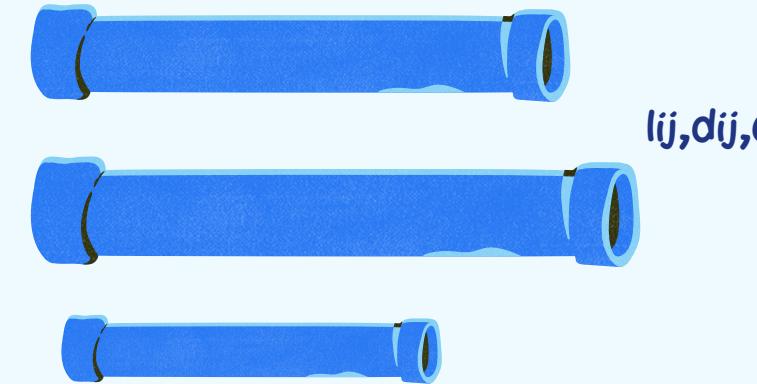
# A SERIES OF ILP MODELS FOR THE OPTIMIZATION OF WATER DISTRIBUTION NETWORKS

NIKHIL HOODA<sub>1</sub>,\*  
, ASHUTOSH MAHAJAN<sub>2</sub> and OM DAMANI

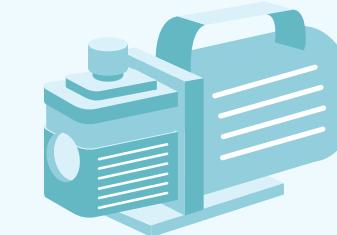
# PIPES AND PUMP MODEL OPTIMISATION



Pipes



Pump



What power pump should be applied at which link i

Efficiency of a pump, Max and minimum power allowed by a pump, Capital and Operational Cost of a Pump

## OBJECTIVE FUNCTION

minimize  $L, p$  : Pipe Cost +  
 $(C_p * p_i) + (EF * DF * SH * p_i)$  for all i

## Pressure on Link

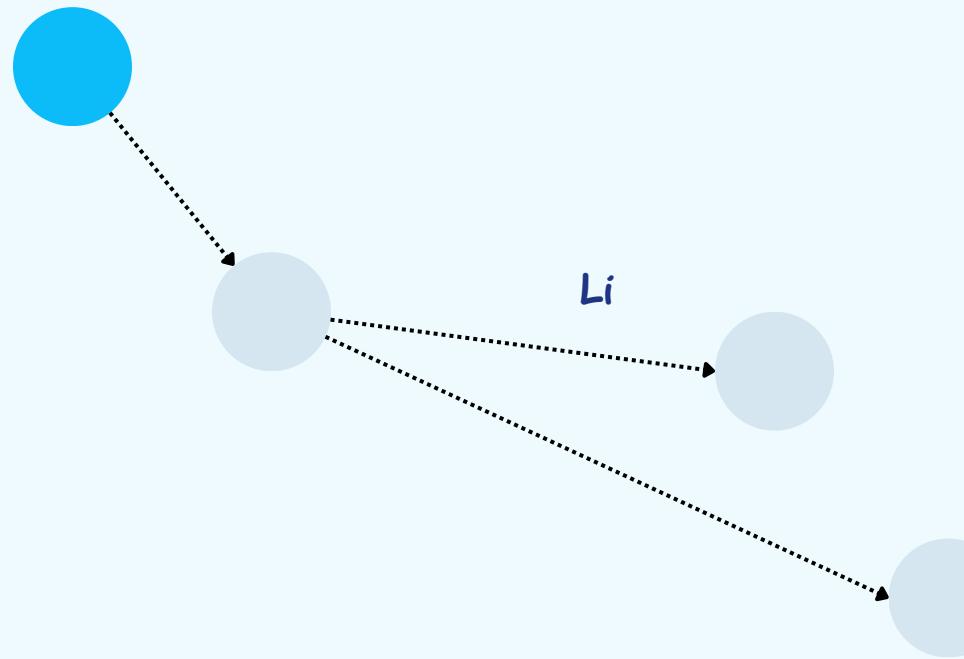
$$\phi_i = \frac{(eff * p_i)}{wtden} * g * F_{Li}$$

## CONSTRAINTS

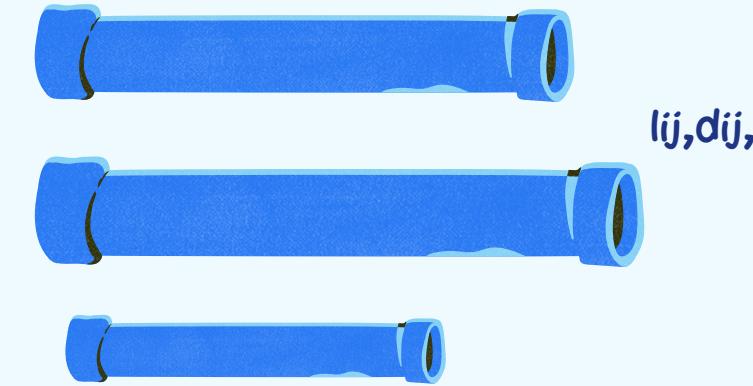
### Pressure requirement at each node

$$P_n \leq H_r - E_n - (H L_{ij} * L_i) + \phi_i \text{ sum over } j \text{ and } i \text{ which are parent links}$$

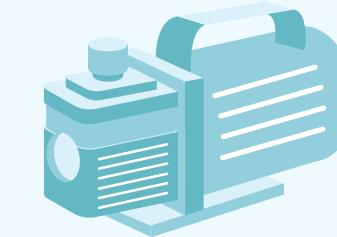
# PIPES AND PUMP MODEL OPTIMISATION



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What power pump should be applied at which link i

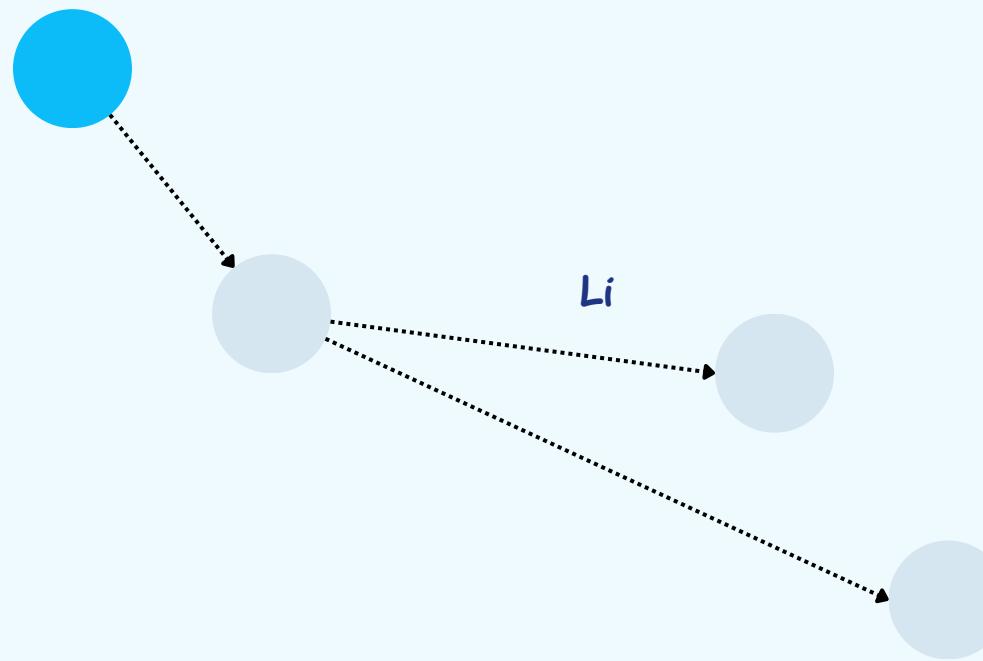
Efficiency of a pump, Max and minimum power allowed by a pump, Capital and Operational Cost of a Pump

## CONSTRAINTS CONTINUED

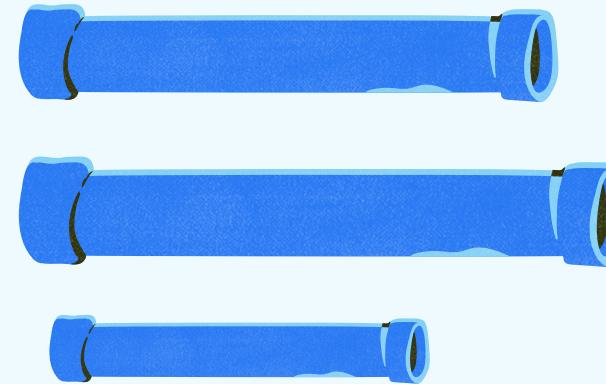
Power Constraints

$$PP_{min} \leq p_i \leq PP_{max}$$

# FINAL MODEL INPUTS

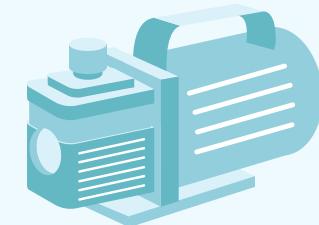


Pipes



$l_{ij}, d_{ij}, C$

Pump



Efficiency of a pump, Max and minimum power allowed by a pump, Capital and Operational Cost of a Pump

Tank Cost Table

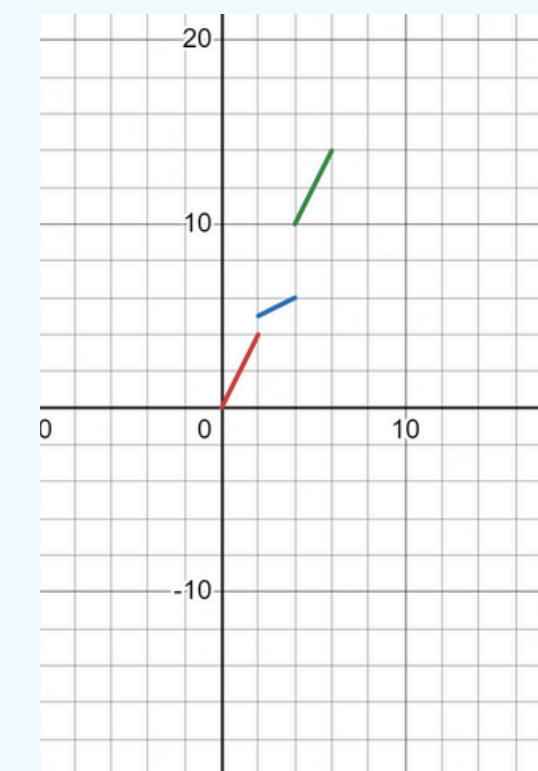
Base Cost, Unit Cost, Upper Limit Capacity, Lower Limit Capacity

Tank 1

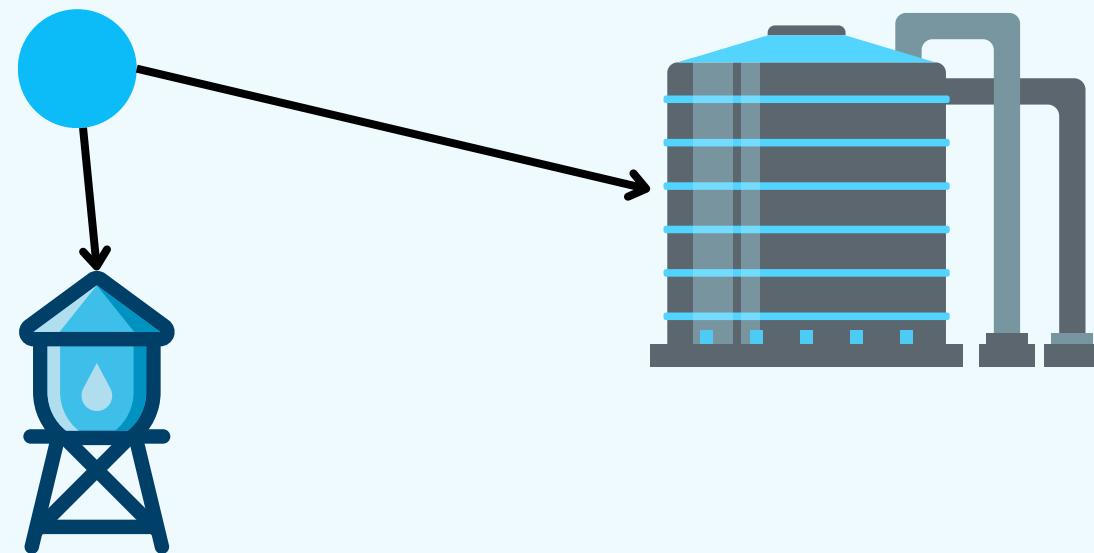
Tank 2

Tank 3

Other tank inputs : Maximum, Min Tank Height allowed,

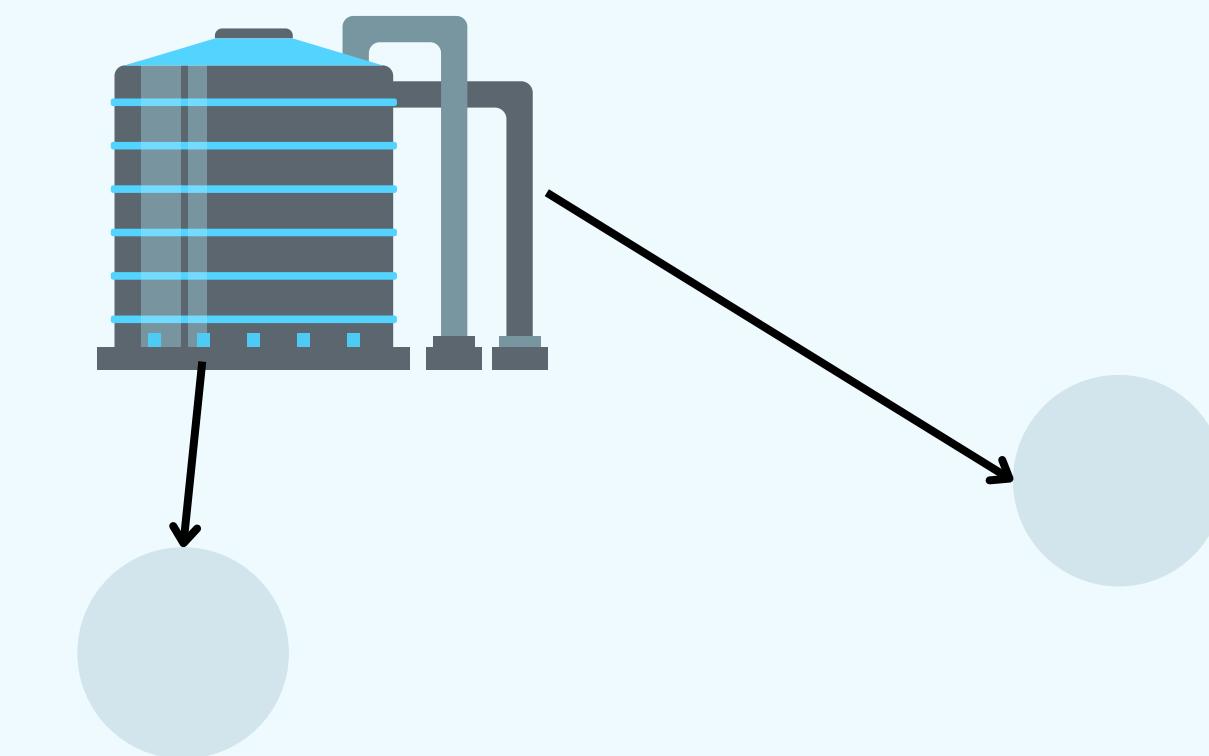


# PRIMARY



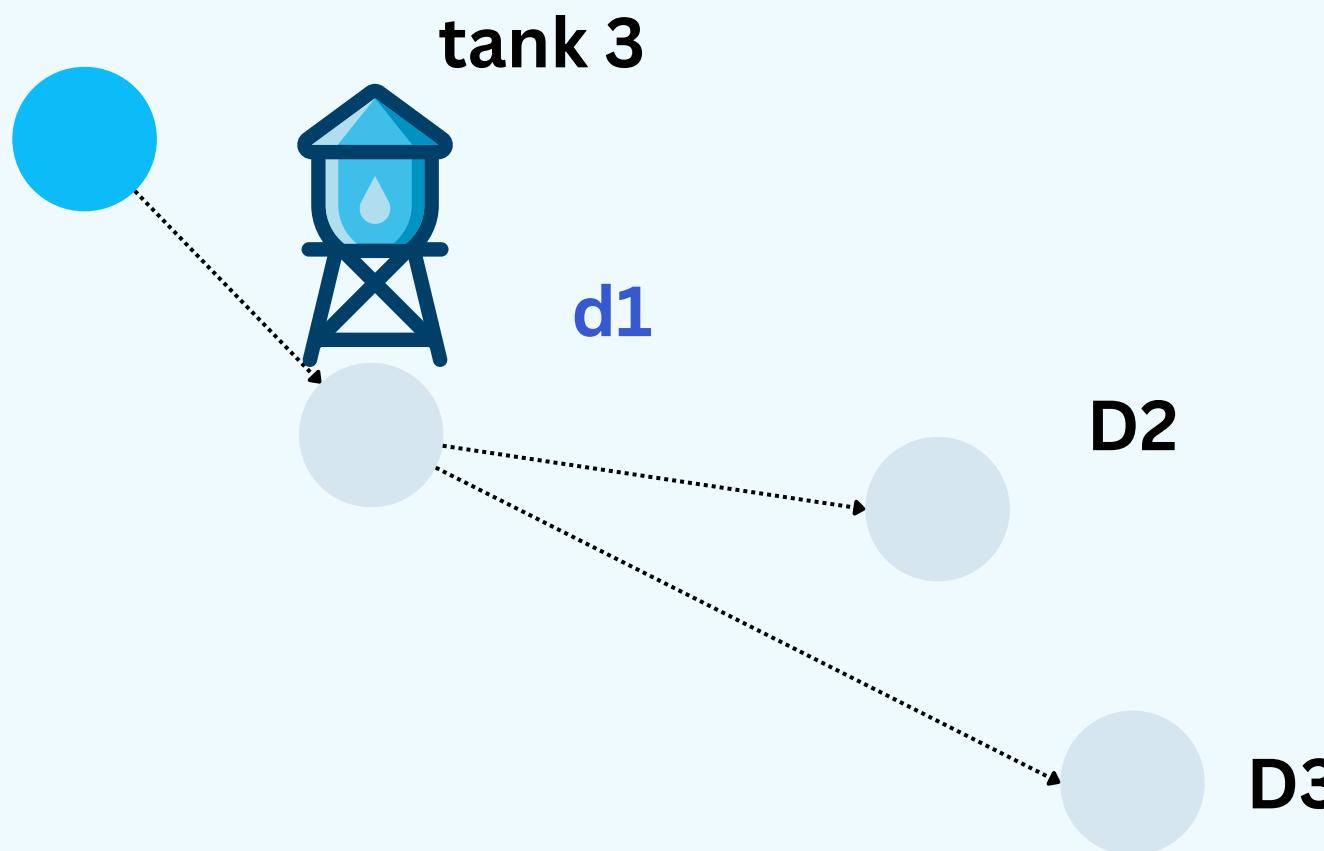
The primary network is the main system of pipes and infrastructure that transports water from the source (such as a water treatment plant or reservoir) to various storage facilities

The secondary network is the subsystem that distributes water from the primary network (or storage tanks) to individual demand nodes, such as households, businesses, and public facilities.



# SECONDARY

# OBJECTIVE FUNCTION WITH TANK



Head on link if tank is its source

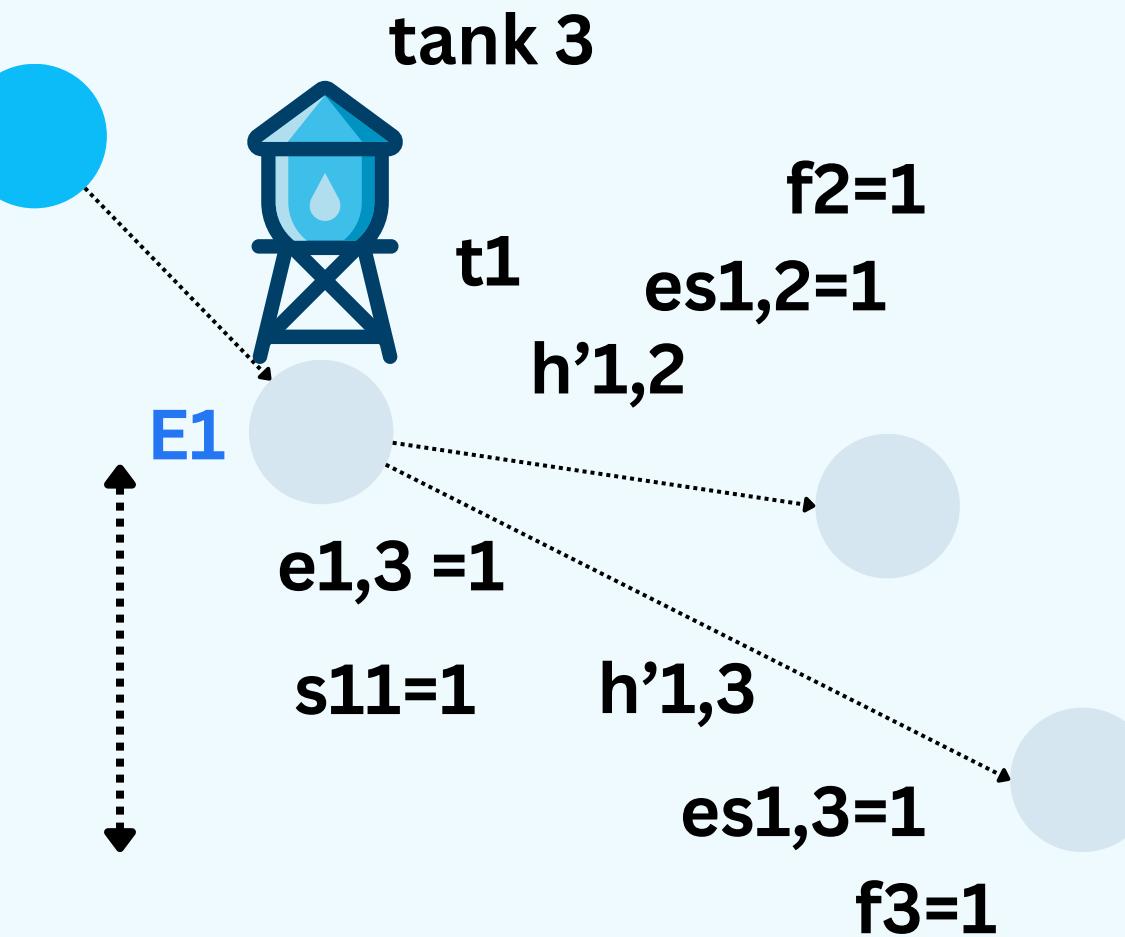
$$d_1 = D_2 + D_3$$

$z_{nk} = d_n * e_{nk}$  (linearisation of  
the )

min  $\sum l_i p_i$  Cost of pipes + Cost of Pumps +  
 $e_{nk} * (\beta_k + \text{UnitCk} * (d_n - L_{Ok}))$  over all  
nodes (n) of the network and rows of  
the tank table (k)

# PRESSURE CONSTRAINT WITH TANK

Head on link if tank is its source



$$h'^{ni} = (t_n + E_n) * es_n + h_n * (1 - es_n)$$

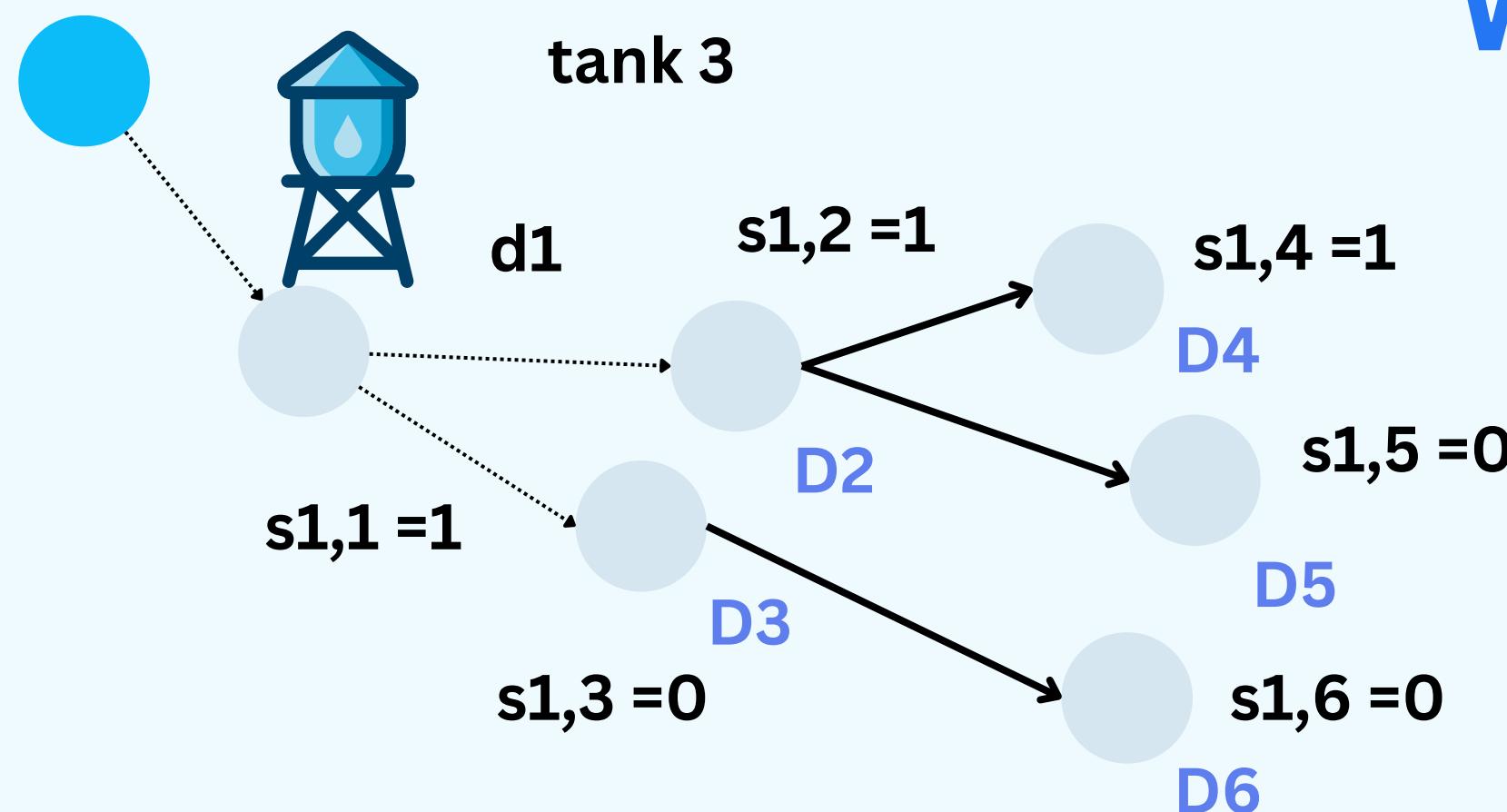
$$h'^{1,2} = t_1 + E_1$$

$$es_{ni} = s_{nn} * (1 - f_i)$$

$$p_n \leq h'^{mi} - h^{li}$$

If the starting node is a source node  
then the head loss equation changes

# CONSTRAINTS ON TANK AND THEIR DESCENDANTS



Where m is a descendant of n

$$smm \leq snn$$

$$snm \leq snn$$

sum over all Ancestors of m  $snm = 1$

$$dn = snm * DEM$$

$snm \leq 1 - soo$  for all o in path of n,m

Nodes can serve their descendants if they serve their own

# OTHER CONSTRAINTS

Only 1 tank per row:

sum over k enk =1

Capacity-demand Constraint:

$$dn \geq LOk * enk$$

$$UPk + DE(1-enk) \geq dn$$

# EFFECT ON OTHER CONSTRAINTS

Splitting headloss calculation:

$hli = \text{sum over all pipe segments}$

$$(HLPij * lpij + HLSij * [lij - lpij])$$

Primary and Secondary Flows:

$$FLPi / FLSi = SH / PH$$

Primary and Secondary Pumps:

$$pi = ppi + psi$$

## CONCLUSION

The paper goes on to create tighter relaxed feasible set of the model established along with giving proofs as to why are they tighter

The first ILP model took 45 mins to solve a 150 nodes network while this one took 5sec

Future research opportunities with creating a pump scheduling algorithm and designing a model for looped networks

# OPTIMAL SCHEDULING OF WATER DISTRIBUTION SYSTEMS

By Manish K. Singh Student Member, IEEE  
and Vassilis Kekatos, Senior Member, IEEE

# OPTIMAL WATER FLOW PROBLEM

## DECISION VARIABLE

Water Flow in pumps and Pressure given out by pumps for a given time period

## OBJECTIVE FUNCTION

Minimize operational cost of pumps, which in our case mainly include electricity cost

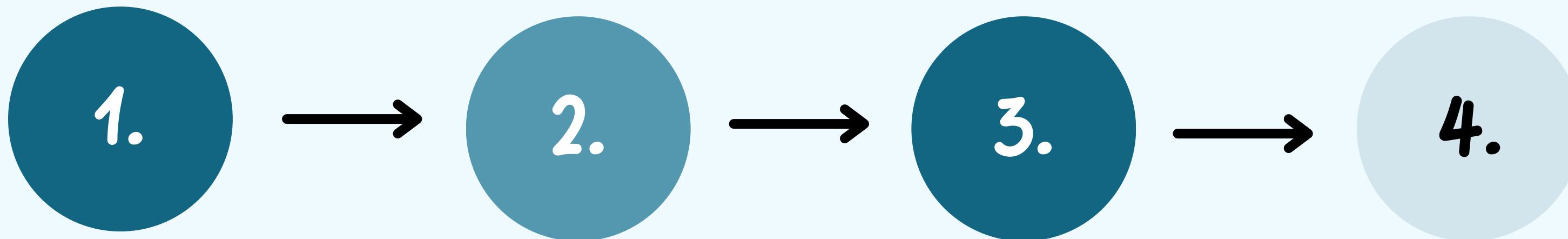
## CONSTRAINTS

Hydraulic constraints are given by physical laws governing the fluid flow in a pipe network.

System Constraints like water demands and water quality requirement

THE GOAL IN DESIGNING SOFTWARE WHICH CAN COME UP WITH AN OPTIMAL SCHEDULE FOR A GIVEN WATER DISTRIBUTION NETWORK , CONSISTING OF PIPES ,TANKS AND PUMPS

# HOW THE PAPER IS DIVIDED

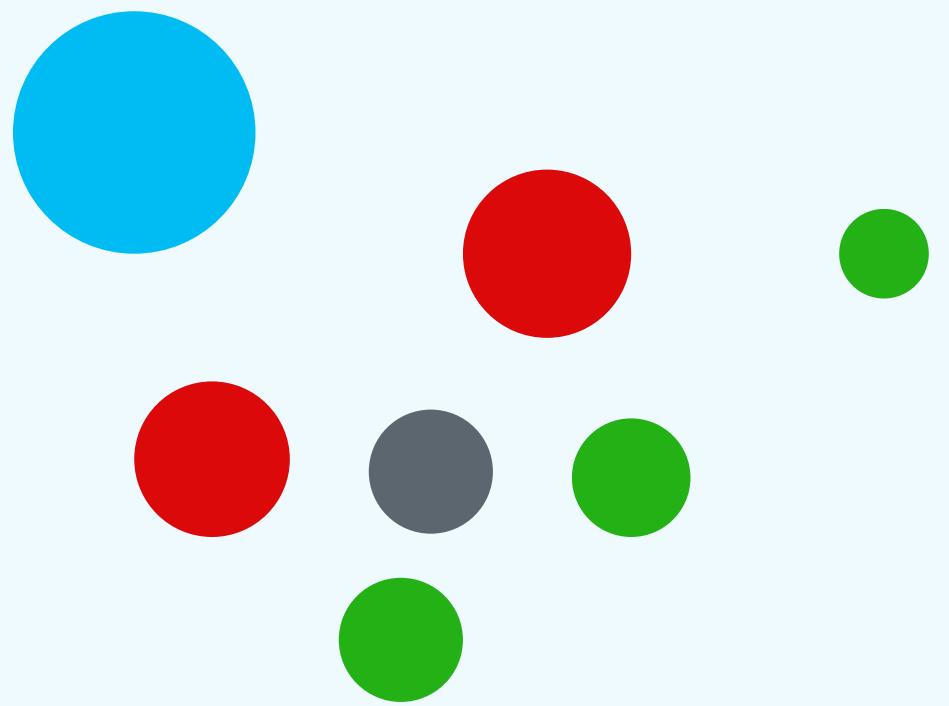


WDS Modelling Optimal Water  
Flow

Convex  
Relaxation

Penalty Term  
Tests and  
Results

## NODES (M)



**Flow (d) for each node (m) at time (t)**

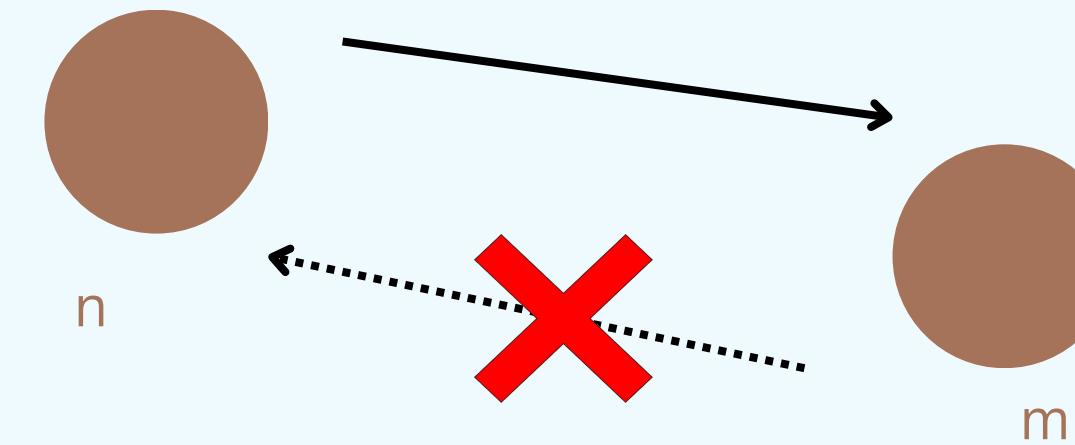
$dm \geq 0$

$dm$  belongs to R

$dm < 0$

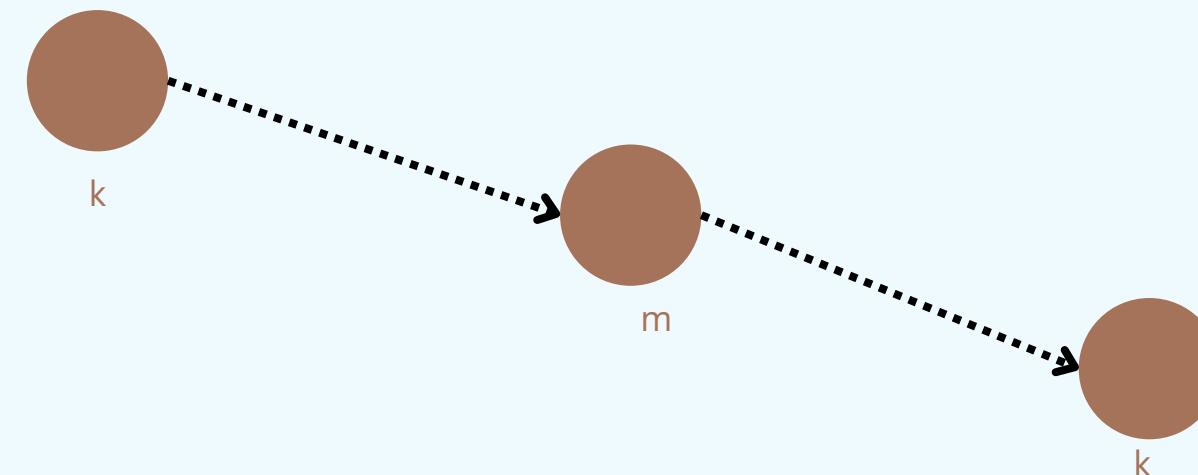
$dm = 0$

## PIPES (P)



if  $(n,m)$  in P then  $(m,n)$  not in

**Flow between nodes (dkmt)**



$dmt = \text{sum}(dkmt) \text{ (for all incoming pipes)} - \text{sum}(dkmt) \text{ (for all outgoing pipes)}$

# PRESSURE IN OUR NETWORK (PIPES)

Pressure at each node is referred to as pressure head, a quantity that is linearly correlated with pressure

For water to flow each node requires to be above a certain pressure level

$$h_m^t \geq h_{min}$$

When water flows through a pipe, it loses some pressure due to frictional losses proportional to **square(dnm)**

We express this using **Darcy-Wesbach** equation

$$h_m^t - h_n^t = c_{mn} \text{sign}(d_{mn}^t) (d_{mn}^t)^2$$

$$c_{mn} := \frac{\ell_{mn} f_{mn}}{4\pi^2 r_{mn}^5 \tilde{g}}$$

$$-M(1 - x_{mn}^t) \leq d_{mn}^t \leq Mx_{mn}^t \quad (4a)$$

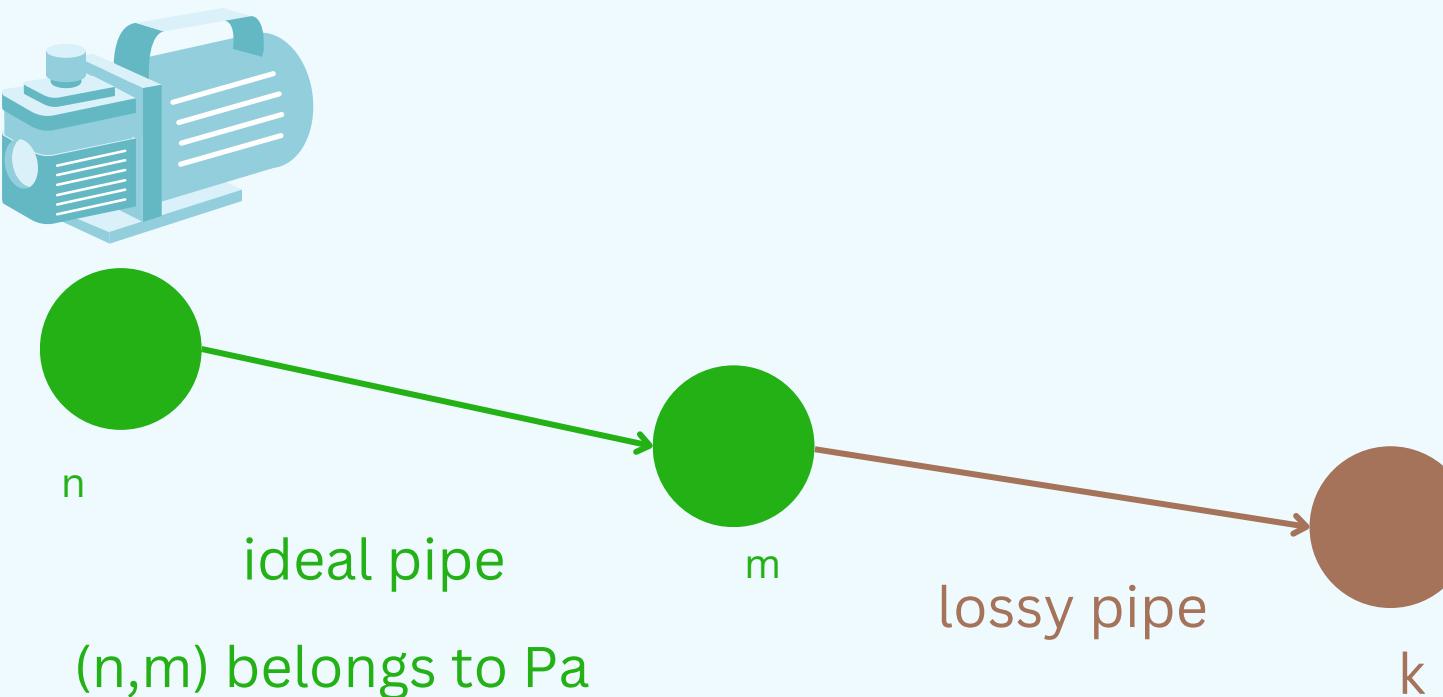
$$-M(1 - x_{mn}^t) \leq h_m^t - h_n^t - c_{mn}(d_{mn}^t)^2 \leq M(1 - x_{mn}^t) \quad (4b)$$

$$-Mx_{mn}^t \leq h_m^t - h_n^t + c_{mn}(d_{mn}^t)^2 \leq Mx_{mn}^t \quad (4c)$$

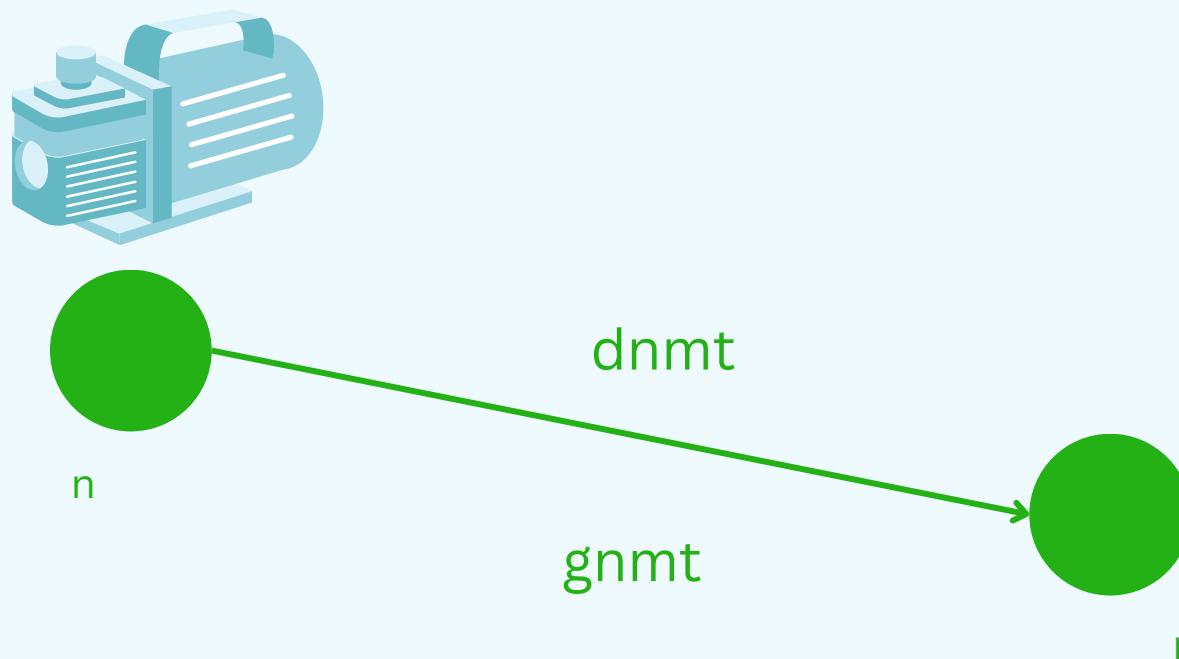
$$x_{mn}^t \in \{0, 1\} \quad (4d)$$

Linearize to remove **sign**

# PRESSURE IN OUR NETWORK (PUMPS)



**Pumps over a pipe are modelled as  
two pipe**



**Relation between pressure gain (gnmt) and  
water flow (dnmt)**

**if the pump is operational**

$$xmnt=1$$

$$dnmmmin \leq dnmt \leq dnmmmax$$

$$hnt - hmt = hnmt$$

**if the pump is not operational**

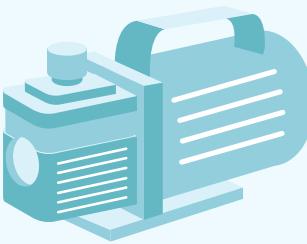
$$xmnt=0$$

$$dnmt \in R$$

$$hnt = hmt$$

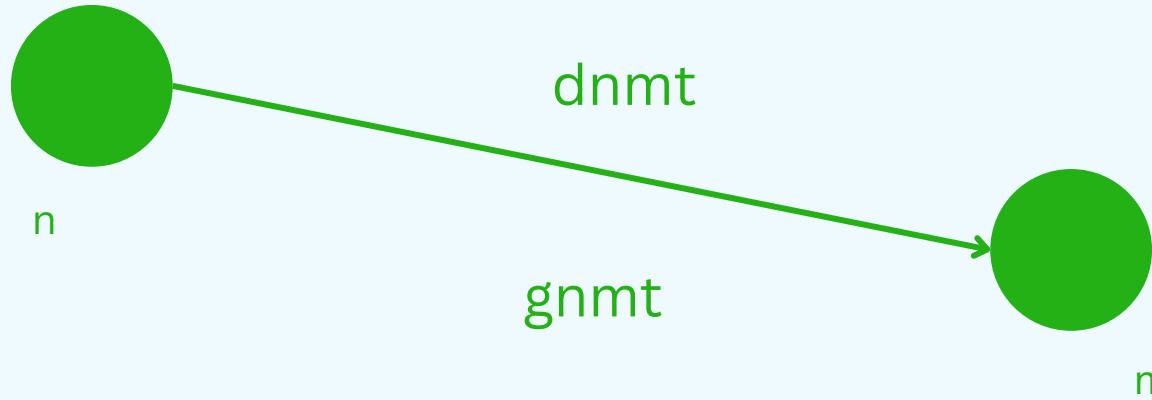
**Pressure gain (gnmt) is considered constant for a tank**

# PRESSURE IN OUR NETWORK (PUMPS)



**Linearize our relationship**

d'tmn is an auxilary variable



**Relation between pressure gain (gnmt) and water flow (dnmt)**

**if the pump is operational**

$$xmnt=1$$

**if the pump is not operational**

$$xmnt=0$$

$$dnmmmin \leq dnmt \leq dnmmmax$$

$$hnt - hmt = hnmt$$

$$dnmt \text{ belongs to } R$$

$$hnt = hmt$$

**Pressure gain (gnmt) is considered constant for a tank**

$$h_m^t - h_n^t = -g_{mn}x_{mn}^t \quad (6a)$$

$$-M(1 - x_{mn}^t) \leq d_{mn}^t - \tilde{d}_{mn}^t \leq M(1 - x_{mn}^t) \quad (6b)$$

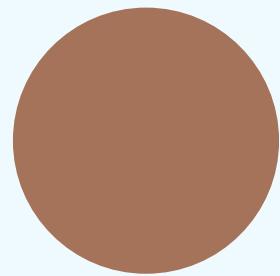
$$\underline{d}_{mn}x_{mn}^t \leq \tilde{d}_{mn}^t \leq \bar{d}_{mn}x_{mn}^t \quad (6c)$$

$$x_{mn}^t \in \{0, 1\}. \quad (6d)$$

# PRESSURE IN OUR NETWORK (RESEVIORS)



$h'm$



$m$

$hmt$

$h'm$  does not change over time

$hmt$  is the pressure drawn to the node from the reservoir

When the reservoir is connected

$\text{Alpha}(mt)=1$

$0 \leq hmt \leq h'm$

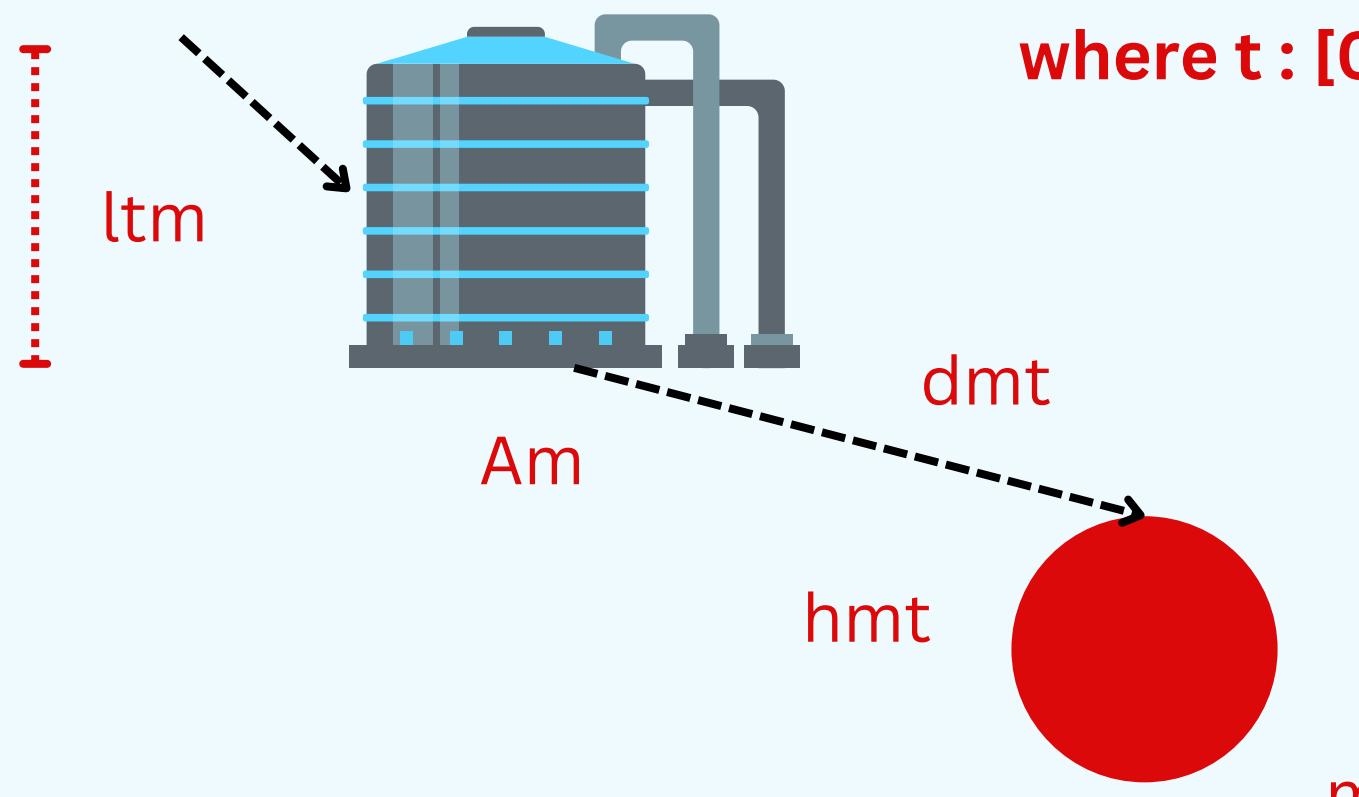
Linearize this relationship

$$0 \leq d_m^t \leq M\alpha_m^t \quad (7a)$$

$$h_m^t \leq \bar{h}_m + M(1 - \alpha_m^t) \quad (7b)$$

$$\alpha_m^t \in \{0, 1\} \quad (7c)$$

# PRESSURE IN OUR NETWORK (TANKS)



where  $t : [0.....T]$

Water Level after a given time interval

$$l_{tm} = l_{t-1} - m \cdot d_{mt} * \Delta t / A_m$$

Tank should be between a min and max level

$$l_{mmin} \leq l_{tm} \leq l_{mmax}$$

Net water exchange is kept at zero

$$l_0m = l_{Tm}$$

When tanks is being filled  
 $\beta_{amt}=1$  and  $\alpha_{amt}=1$

$$h_{tm} \leq l_{mmax}$$

When tank is not being filled  
 $\beta_{amt}=0$  and  $\alpha_{amt}=1$

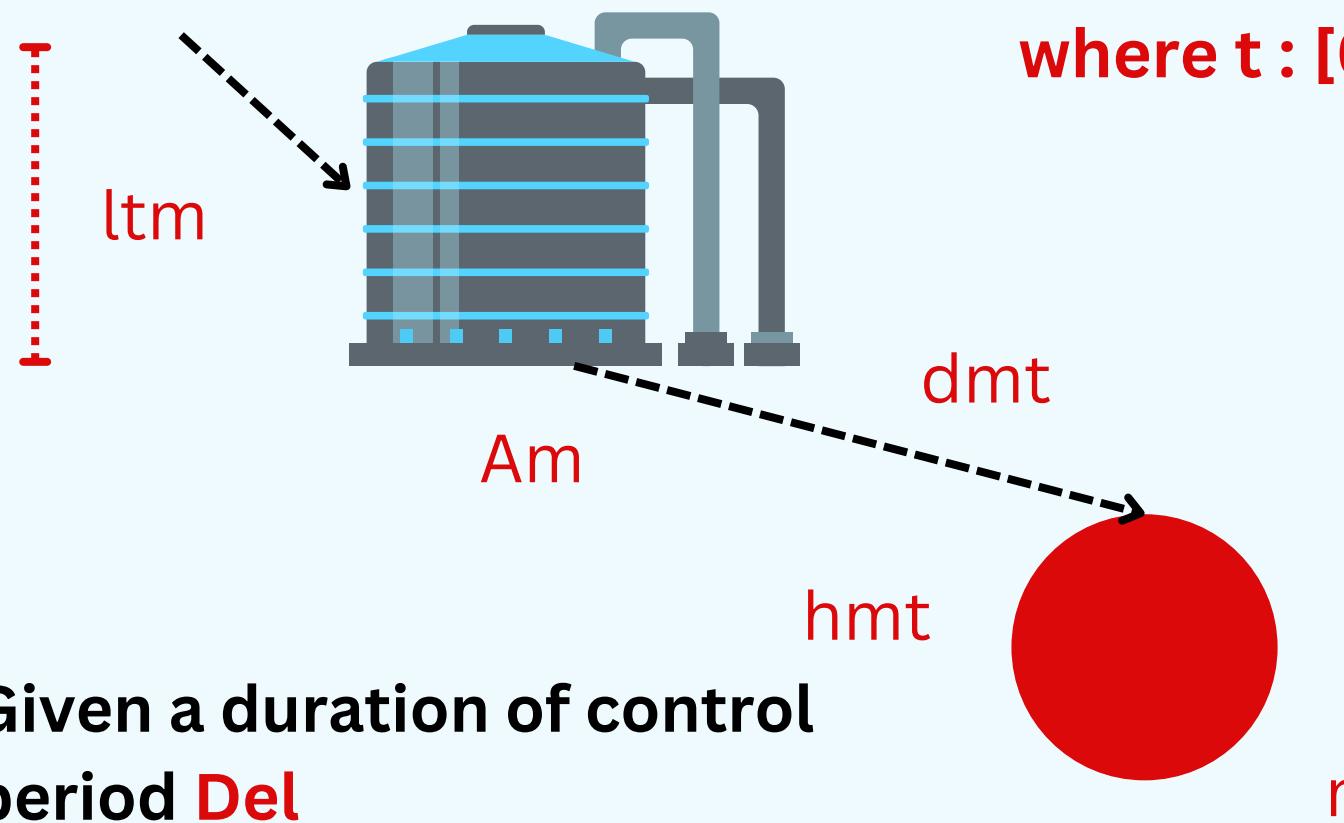
$$h_{tm} \leq l_{tm}$$

When tank is not connected  
 $\alpha_{amt}=0$

**$h_{tm}$  is not related to  $l_{tm}$**

Given a duration of control period  $\Delta t$

# PRESSURE IN OUR NETWORK (TANKS)



where  $t : [0.....T]$

When tanks is being filled  
 $\text{Betamt}=1$  and  $\text{Alphamt}=1$

When tank is not being filled  
 $\text{Betamt}=0$  and  $\text{Alphamt}=1$

When tank is not connected  
 $\text{Alphamt}=0$

$htm \leq lmmax$

$htm \leq lmt$

**htm is not related to lmt**

Given a duration of control period **Del**

Linearizing our tank level and pressure relationship

$$-M(1 - \alpha_m^t) \leq \tilde{h}_m^t - h_m^t \leq M(1 - \alpha_m^t) \quad (11a)$$

$$-M\alpha_m^t \leq d_m^t \leq M\alpha_m^t \quad (11b)$$

$$-M\beta_m^t \leq d_m^t \leq M(1 - \beta_m^t) \quad (11c)$$

$$\bar{\ell}_m - M(1 - \beta_m^t) \leq \tilde{h}_m^t \leq \bar{\ell}_m + M\beta_m^t \quad (11d)$$

$$\alpha_m^t, \beta_m^t \in \{0, 1\}. \quad (11e)$$

## HOW WE MODEL CONTROL PERIODS

- We discretize our continuous system into time chunks  $t=0,1,2,\dots,T$
- **0** represents the start of our system, while **T** represents the end
- All of these chunks are of a constant duration **Δt**

## HOW CAN WE CALCULATE ELECTRICITY CONSUMPTION

- The overall energy consumed by a pump in **Δt** is proportional to the **pressure (gnmt)** and the **flow-induced (d'nmt)** at a pipe

- Energy Consumption Coefficient for **Δt**

$$c_{mn} := \frac{\delta \rho \tilde{g} g_{mn}}{\eta_{mn}}, \quad \forall (m, n) \in \mathcal{P}_a.$$

- **$E_{mn}(t)=c_{mn}*d'_{mn}(t)$**

PROBLEM FORMATION

## COST AND OBJECTIVE FUNCTION

- Our electricity cost is  $\pi_{mn}(t)$  for pump located at  $(m,n)$
- We assume this cost to be invariant for all time for simplicity

Objective Function

$$f(\tilde{\mathbf{d}}) := \sum_{t=1}^T \sum_{(m,n) \in \mathcal{P}_a} c_{mn} \pi_t \tilde{d}_{mn}^t$$

## DECISION VARIABLE AND CONSTRAINTS

- Lemma1 = The variable  $h_m^t, d_m^t$  and  $d_{mn}^t$  are sufficient to represent any feasible point in our constraints
- Other variables are really there to avoid discontinuous equations and multiplication of binary and continuous variables

over  $\{h_m^t\}_{m \in \mathcal{M}}, \{d_m^t\}_{m \in \mathcal{M}_b \cup \mathcal{M}_r}, \{d_{mn}^t\}_{(m,n) \in \mathcal{P}},$   
 $\{\tilde{h}_m^t\}_{m \in \mathcal{M}_b}, \{\ell_m^t\}_{m \in \mathcal{M}_b}, \{\tilde{d}_{mn}^t\}_{(m,n) \in \mathcal{P}_a},$   
 $\{x_{mn}^t\}_{(m,n) \in \mathcal{P}}, \{\alpha_m^t\}_{m \in \mathcal{M}_r \cup \mathcal{M}_b}, \{\beta_m^t\}_{m \in \mathcal{M}_b}, \quad \forall t$

- Our Constraints were the ones we discussed before

# PROBLEM FORMATION

# MAKING OUR RELAXATION CONVEX

- The equations 4b and 4c make our relaxation non convex

$$-M(1 - x_{mn}^t) \leq d_{mn}^t \leq Mx_{mn}^t \quad (4a)$$

$$-M(1 - x_{mn}^t) \leq h_m^t - h_n^t - c_{mn}(d_{mn}^t)^2 \leq M(1 - x_{mn}^t) \quad (4b)$$

$$-Mx_{mn}^t \leq h_m^t - h_n^t + c_{mn}(d_{mn}^t)^2 \leq Mx_{mn}^t \quad (4c)$$

$$x_{mn}^t \in \{0, 1\} \quad (4d)$$

P1

- We get rid of the righside inequality for 4b and the left side inequality for 4c

- This makes our relaxation convex (MISOCP), but since P2 is a relaxation of P1, they only share a minimizer when 13b and 13c's equalities are satisfied

$$-M(1 - x_{mn}^t) \leq \dot{d}_{mn}^t \leq Mx_{mn}^t \quad (13a)$$

$$-M(1 - x_{mn}^t) \leq h_m^t - h_n^t - c_{mn}(d_{mn}^t)^2 \quad (13b)$$

$$h_m^t - h_n^t + c_{mn}(d_{mn}^t)^2 \leq Mx_{mn}^t. \quad (13c)$$

P2

## CONCLUSION

We can solve this model with off the shelves MISOCP solvers, but the paper creates a penalty term for our objective to help minimizers

The numerical tests show how this formulation is implemented in MATLAB and Gurobi and simulated in EPANET.

Future research directions suggest extensions of the formulation to variable speed pumps and electricity prices



# HOW IS OUR MODEL IMPROVED

$\min x \mathbf{C} \cdot \mathbf{x}$   
 $A\mathbf{x} \leq \mathbf{B}$   
 $\mathbf{x} \geq 0$   
 $\mathbf{X}$  is Integer

$\min x \mathbf{C} \cdot \mathbf{x}$   
 $A\mathbf{x} \leq \mathbf{B}$   
 $\mathbf{x} \geq 0$   
 $\mathbf{X}$  is Rational

SOLVING A LINEAR  
INTEGER PROGRAM  
INVOLVES LINEAR  
RELATION

ILP'

New  
Relaxation'

ILP

Old  
Relaxation

IF NEW RELAXATION IS A STRICT  
SUBSET OF THE OLD RELAXATION  
THEN IT'S A SMALLER SOLUTION  
SPACE FOR THE SAME PROBLEM

# HEADLOSS IMPROVEMENT

## Old model

$$lijp = lij * fi$$

↓  
Linearize

$$lijp \geq 0$$

$$lijp \leq Li * fi$$

$$lij - Li * (1 - fi) \leq lij p_j$$

$$lijp \leq lij$$

## New model

We introduce  $lijs$  which represents length of link for the secondary network

$$lijp + lijs = lij$$

sum over a link  $lijp = Li * fi$

sum over a link  $lijs = Li * (1 - fi)$

# HEADLOSS IMPROVEMENT STRICT SUBSET PROVE

Consider  
the point :

Link Length = L



Primary pipe segment  
 $=[L/2, L/2]$ ,  $f = 0.5$



Secondary pipe  
segment  $=[0, 0]$

Link/Pipe Length =  $[L/2, L/2]$

Old model  
Satisfied

$$lijp = lij * fi$$



$$L/2 = L * 0.5$$

$$lijp + lijs = lij$$

$$L/2 = L/2$$

$$\text{sum over a link } ljp = L * fi$$

$$L/2 = L/2$$

$$\text{sum over a link } lijs = Li * (1 - fi)$$

$$0 = L/2$$

New model Unsatisfied

# TANK COST IMPROVEMENT

## Old model

$$znk = enk * dn \quad \text{eq0}$$

↓  
Linearize

$$znk \geq 0$$

$$znk \leq DE * enk$$

$$znk - DE * (1 - enk) \leq dn$$

$$dn \leq znk$$

znk is the water demand of  
tank k at node n

DE is total wated demand of the  
network

## New model

$$LOk * enk \leq znk \quad \text{eq1}$$

$$znk \leq UPk * enk \quad \text{eq2}$$

$$\text{eq3} \quad \text{sum for all row for } n \text{ enk} = 1$$

$$\text{eq4} \quad \text{sum for all row for } n \text{ znk} = dn$$

lets say for node n we have 4 tanks

applying eq1&2 to enk

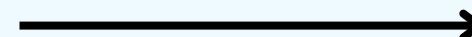
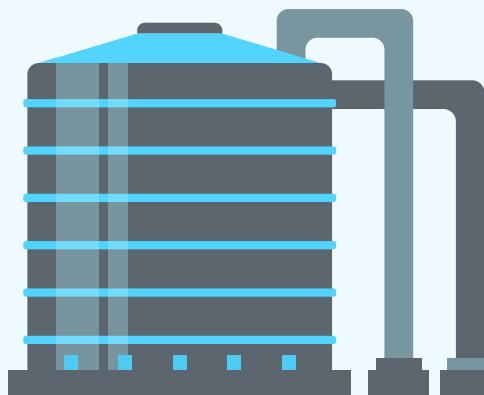
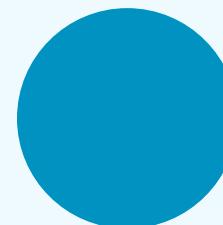
applying eq4 to znk

enk : 0 1 0 0

znk : 0 zn2 0 0

zn2=dn

# TANK COST IMPROVEMENT : STRICT SUBSET PROOF



$$\begin{aligned} z_{1k} &: [d, d] \\ e_{1k} &: [1/2, 1/2] \\ d_1 &= d \end{aligned}$$

$$\begin{aligned} L_{0k} &: [0, d] \\ U_{P_k} &: [d, 2d] \\ DE &= 2d \end{aligned}$$



**Old model is satisfied**

$$z_{nk} \geq 0$$

$$z_{nk} \leq DE * e_{nk}$$

$$z_{nk} - DE * (1 - e_{nk}) \leq d_n$$

$$d_n \leq z_{nk}$$

**New model is violated**

eq1  $L_{0k} * e_{nk} \leq z_{nk}$

eq2  $z_{nk} \leq U_{P_k} * e_{nk}$

eq3 sum for all row for n  $e_{nk} = 1$

eq4 sum for all row for n  $z_{nk} = d_n$

# TANK COST IMPROVEMENT : BEST RELAXATION PROOF

$z_{1k} : [1,1]$   
 $e_{1k} : [1/2, 1/2]$   
 $d_1 = 1$

P1

$t = 1/2$   
 $n' = 1$   
 $k' = 1$

$z_{1k} : [0,2]$   
 $e_{1k} : [0,1]$   
 $d_1 = 0$

P2

$LO_k : [2,0]$   
 $UP_k : [1,0]$   
 $DE = 2$

P3

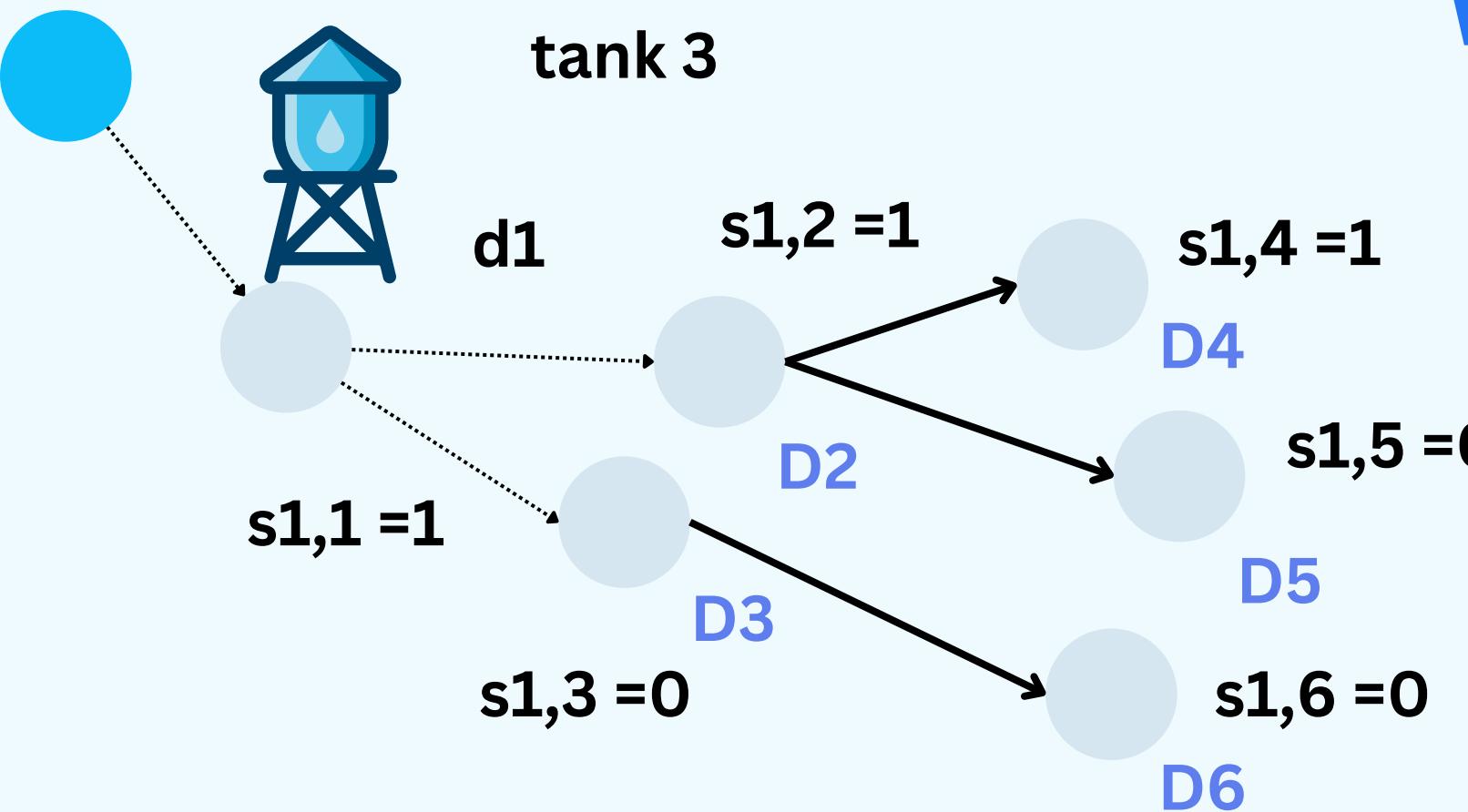
New model is Satisfied

$$P_1 = P_2/2 + P_3/2$$

- eq1       $LO_k * e_{nk} \leq z_{nk}$   
eq2       $z_{nk} \leq UP_k * e_{nk}$   
eq3      sum for all row for n  $e_{nk} = 1$   
eq4      sum for all row for n  $z_{nk} = d_n$

Since  $P$  can be represented as a linear combination of two other points belonging to  $R_2$ ,  $P$  cannot be a corner point of  $R_2$ .

# TANK CONFIGURATION IMPROVEMENT : INITIAL MODEL



Where m is a descendant of n

$$smm \leq snn$$

$$snm \leq snn$$

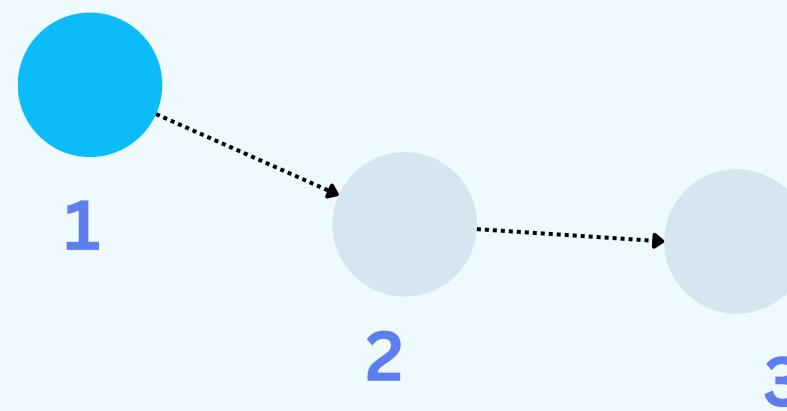
sum over all Ancestors of m  $snm = 1$

$$dn = snm * DEM$$

$snm \leq 1 - soo$  for all o in path of n,m

Nodes can serve their descendants if they serve their own

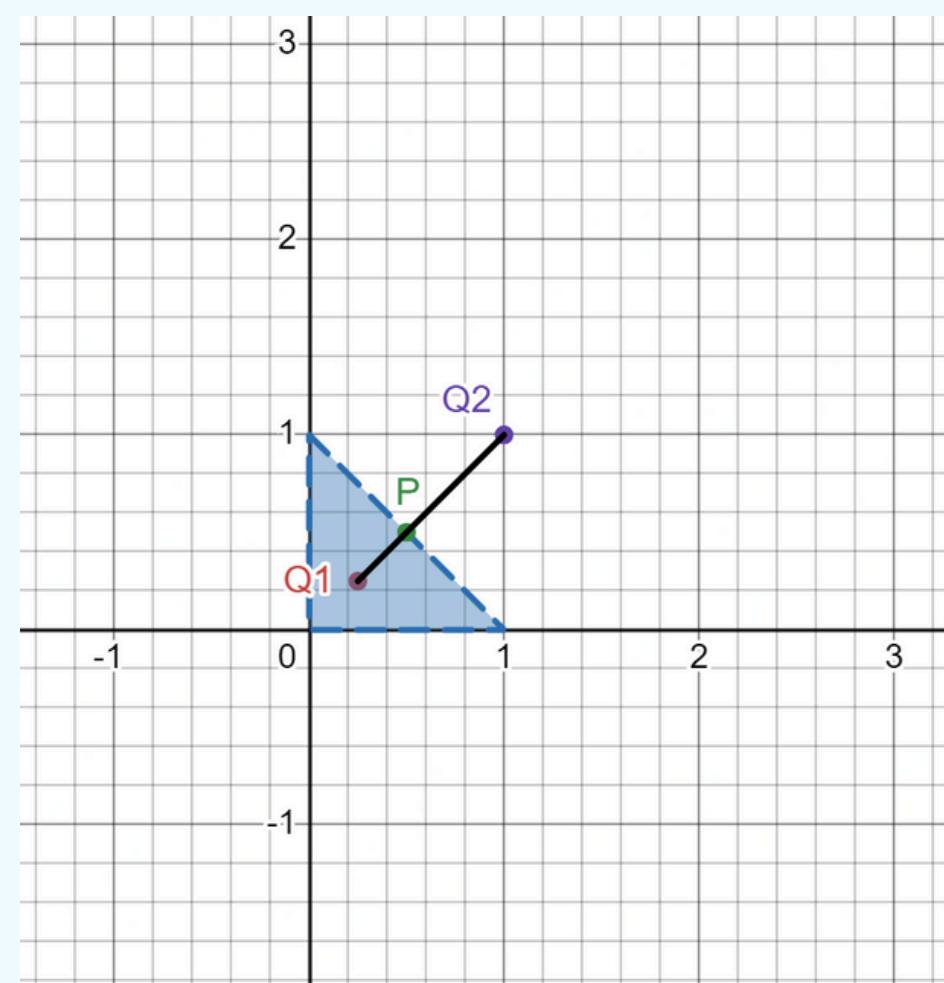
# TANK CONFIGURATION IMPROVEMENT : TIGHTNESS PROOF



$$P : \{1, 0.5, 0, 0.5, 0.5, 0.5\}$$

$$Q_1 : \{s_{11}, s_{12}, s_{13}, s_{22}, s_{23}, s_{33}\}$$

$$Q_2 : \{s_{11}, s_{12}, s_{13}, s_{22}, s_{23}, s_{33}\}$$



$$P = Q_1 + t(Q_2 - Q_1)$$

$P, Q_1, Q_2$  belong in  $R^2$

To prove

$P$

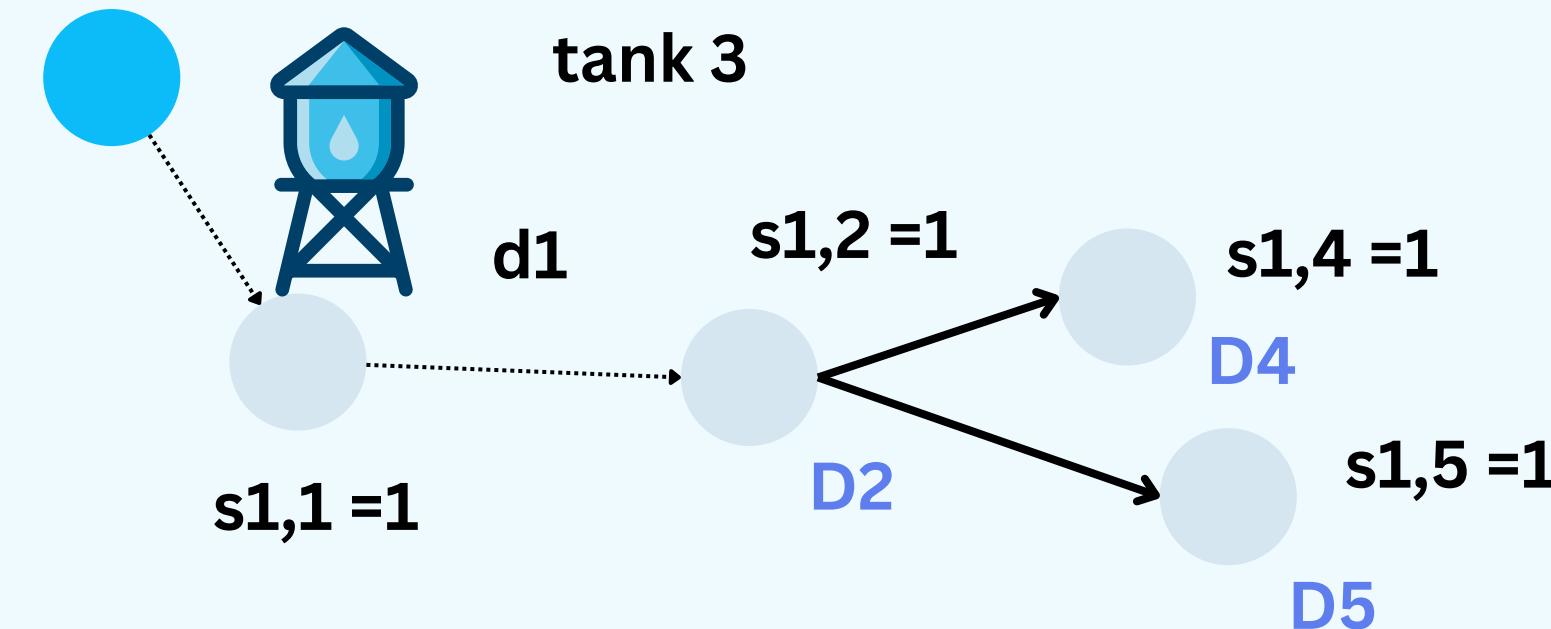
cannot be described as a linear combination of two distinct points that belong to  $R$ .



The paper's proof shows that this condition is only possible when  $Q_1 = Q_2 = P$

Thus  $P$  is a corner point with fractional value

# TANK CONFIGURATION IMPROVEMENT : NEW MODEL



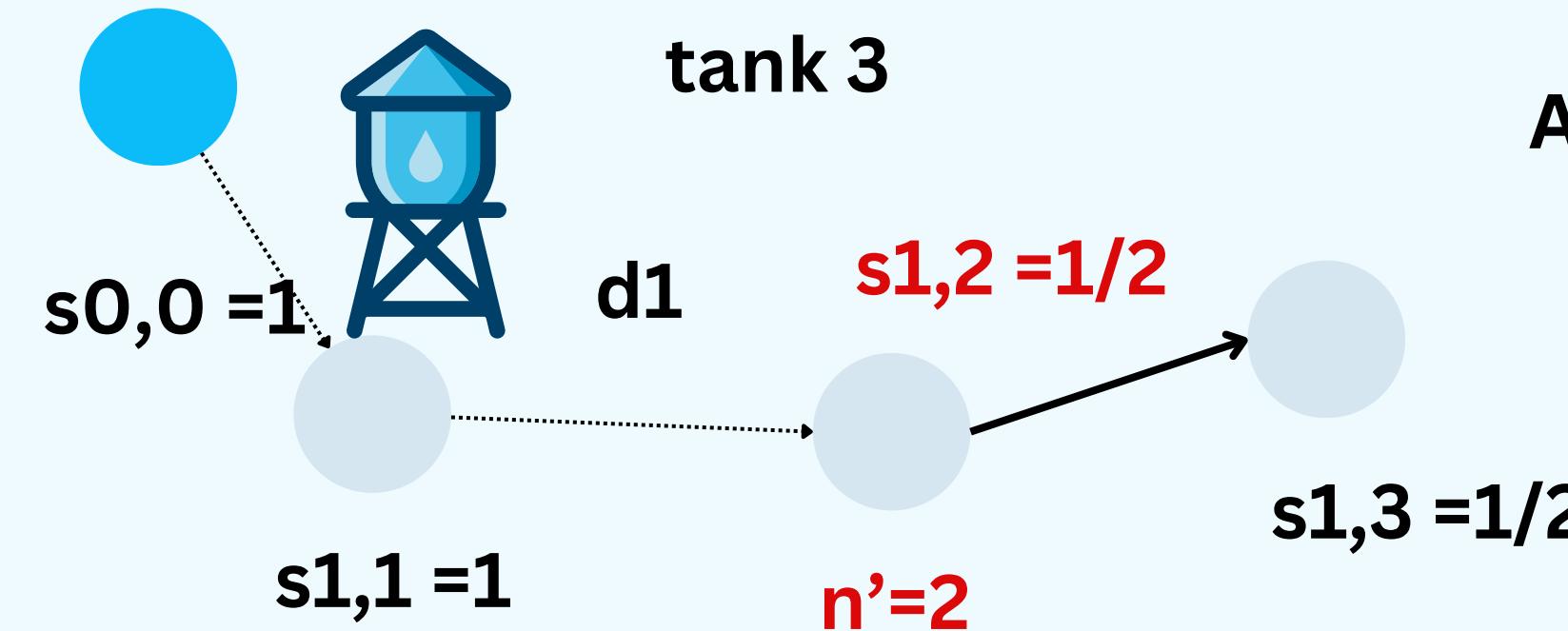
## New Constraints

sum over n  $s_{nm} = 1$  for all ancestors of m

$s_{nm} = s_{nk}$  where m is the child node and k is the descendant node

$$0 \leq s_{nm} \leq 1$$

# TANK CONFIGURATION IMPROVEMENT : NEW MODEL TIGHTNESS



Assume this allocation of tanks is P

## Claim 1

$s_{0,0} = 1$ , root node can never be fractional or zero

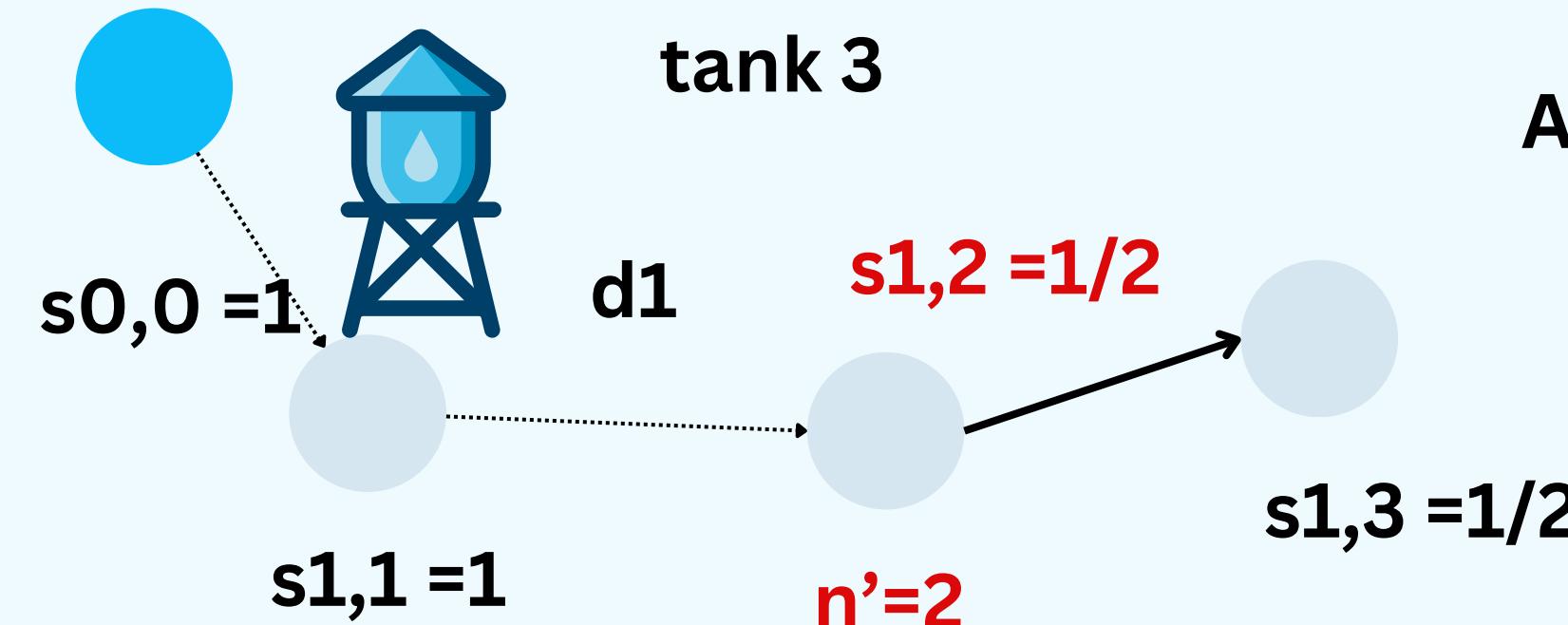
## Claim 2

$s_{0,2} = 0$ , descendants of secondary networks do not depend on nodes above their source node

## Claim 3

- Q1 is a valid allocation of our tank
- $s_{nm} = s_{nm}$  of P for all m not in 2's path
- $s_{nm} = s_{nm} / 1 - s_{1,2}$  for all n which are ancestors of 2
- $s_{2m} = 0$  for all m which are descendants of 2

# TANK CONFIGURATION IMPROVEMENT : NEW MODEL TIGHTNESS



Assume this allocation of tanks is  $P$

## Claim 4

- $Q_2$  is a valid allocation of our tank
- $s_{nm} = s_{nm}$  of  $P$  for all  $m$  not in  $2$ 's path
- $s_{nm} = s_{nm} / s_{1,2}$  for all  $n$  which are ancestors of  $2$
- $s_{2m} = 0$  for all  $m$  which are descendants of  $2$

## Claim 5

$$P = s_{1,2} * Q_1 + (1 - s_{1,2}) * Q_2$$