



JOHNS HOPKINS  
WHITING SCHOOL  
*of* ENGINEERING

# Data Structures

Complexity – Part 1

# Complexity

---

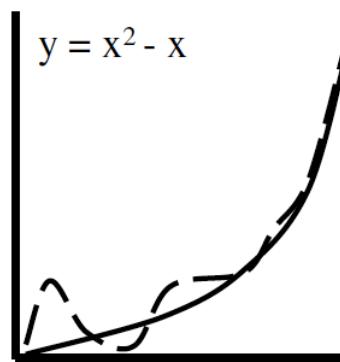
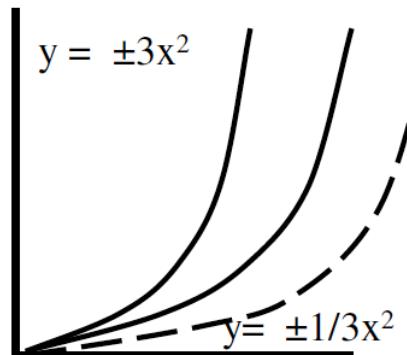
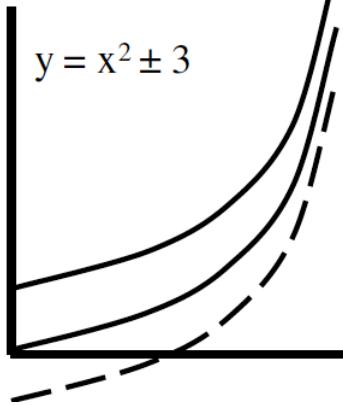
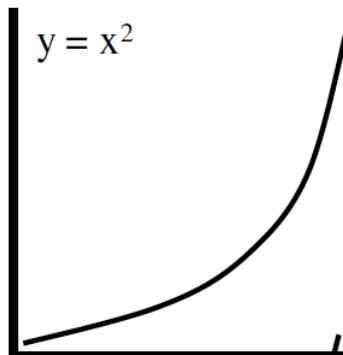
- People naturally make comparisons
  - e.g., Car A is better than Car B!
- How do we define “better”?
  - A. A has more horsepower than
  - B. B holds 7 people, but A only holds 2
- Each can be better than the other in a specific, useful context.

# Can We Do This With Programs?

---

- YES, we call this “Big O” Notation.
- Label code to allow comparisons What does *better* mean here?

# Review of Simple Planar Functions



# A Review of Simple, Planar Functions:

---

- What decides the shape of the curve?
- This is independent of:
  - the lower order terms
  - the coefficients used

# Deriving Work Done By Code (1)

---

Given a specific language,

a specific operating system and

a specific compiler:

Consider this assignment:

$$x = x + 1$$

How much does it take?

# Deriving Work Done By Code (2)

---

- Suppose we change the system?

- Suppose:

```
for (int i = 1; i<=n; i++)
```

```
x = x + 1;
```

- Suppose:

```
for (int i = 1; i<=n; i++)
```

```
for (int j = 1; j<=n; j++)
```

```
x = x + 1;
```

# Deriving Work Done By Code (3)

---

For any piece of code, generate a function to represent the work done:

For example:

$$f(n) = c_1n + c_2 + c_3n + c_4n^3 + c_5n^2 + c_6 + c_7n^2 + c_8$$

Simplifying:

$$f(n) = c_4n^3 + (c_5 + c_7)n^2 + (c_1 + c_3)n + (c_2 + c_6 + c_8)$$

# Deriving Work Done By Code (4)

---

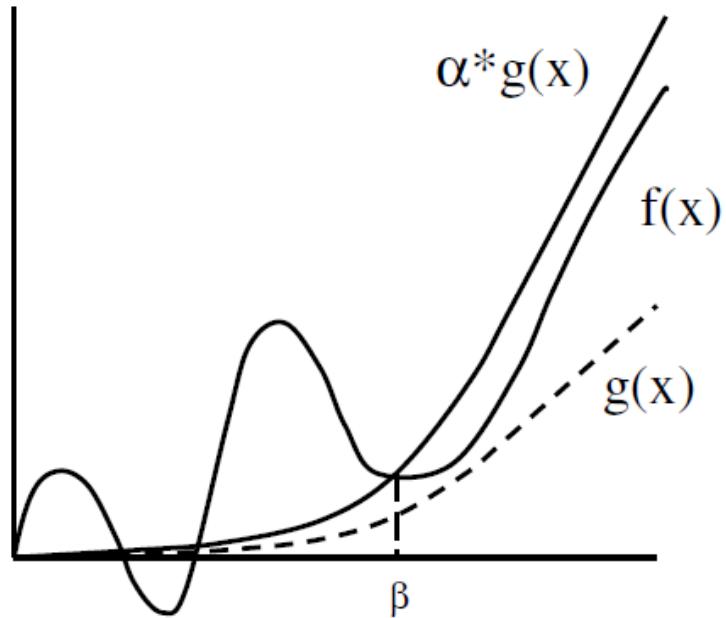
- **This is messy to graph.**
- If all we are interested in is the basic shape, we can simplify
- by using the dominant term.
- This gives us a label to use for the code whose work is represented by  $f(x)$ .

# Definition: Upper Bound

---

- Give two functions  $f(n)$  and  $g(n)$  and two real constants  $\alpha$  and  $\beta$ ;  
if  $\alpha^* g(n) \geq f(n)$ , for all  $n > \beta$  then  $g(n)$  is an upper bound for  $f(n)$   
 $f$  is said to be  $O(g(n))$

# Definition: Upper Bound (cont.)



# Upper Bounds

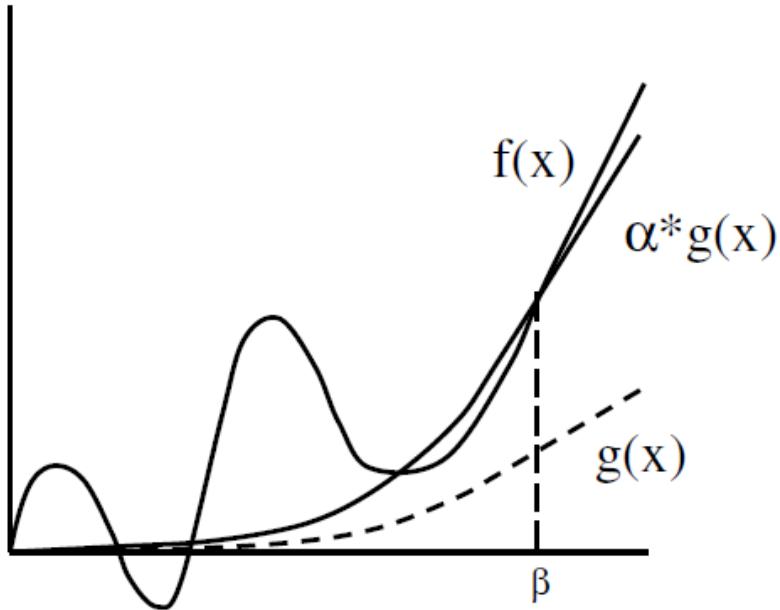
- In particular, if  $f(n)$  is a polynomial then  $g(n)$  is the dominant term.
- $g(n)$  is an estimate of how  $f(n)$  acts.
- We are guaranteed that  $f$  will do no worse than  $g$ . It might do better.

# Definition: Lower Bound

---

- Give two functions  $f(n)$  and  $g(n)$  and two real constants  $\alpha$  and  $\beta$ ;  
if  $\alpha * g(n) \leq f(n)$ , for all  $n > \beta$  then  $g(n)$  is a lower bound for  $f(n)$   
 $f$  is said to be  $\Omega(g(n))$

# Definition: Lower Bound (cont.)



# Lower Bounds

---

- In particular, if  $f(n)$  is a polynomial then  $g(n)$  is the dominant term.
- $g(n)$  is an estimate of how  $f(n)$  acts.
- We are guaranteed that  $f$  is no better than  $g$ . It might be worse.

# Both

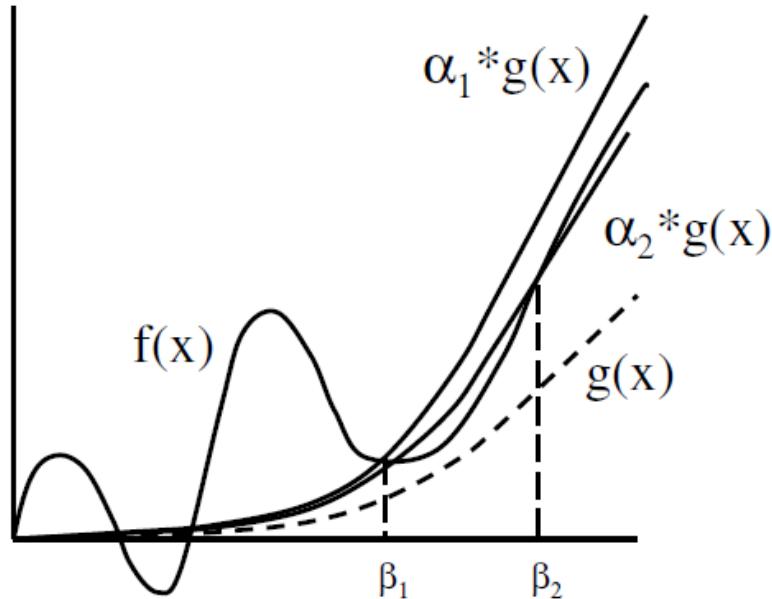
---

- If  $f(n)$  is  $O(g(n))$  and  $\Omega(g(n))$  then  $f$  is said to be  $\Theta(g(n))$   
 $g$  is an upper bound for  $f$  and  $g$  is a lower bound for  $f$ .

e.g.  $\alpha_2^*g(n) \leq f(n) \leq \alpha_1^*g(n)$

# Definition: Both Bounds

---





# JOHNS HOPKINS

## WHITING SCHOOL *of* ENGINEERING