

### 1. R-16.22

In our circular array implementation of a queue, can you compute the value of the `currentSize` from the values of the `head` and `tail` fields? Why or why not?

No. Because there are several edge cases which we can not use `tail minus head` to get size. For example, if `head = 0`, `tail = 0`, in this case, queue size can be full or 0, it is non-deterministic.

### 2. R-16.24

Suppose you are stranded on a desert island on which stacks are plentiful, but you need a queue. How can you implement a queue using two stacks? What is the big-Oh running time of the queue operations?

For two stacks:

```
//initialize size
size = 0;
//initialize two empty stack
Stack<T> A = new Stack<>();
Stack<T> B = new Stack<>();
//Implement my own add()
void add(T e) {
    while(!B.isEmpty()) {
        A.push(B.pop());
    }
    A.push(e);
    size++;
}
//Implement my own remove()
T remove() throws Exception {
    while(!A.isEmpty()) {
        B.push(A.pop());
    }
    //get return value
    T ret = B.pop();
    size--;
    return ret;
}
//Implement my own size()
int size() {
    return size;
}
//Implement my own isEmpty()
boolean isEmpty() {
    return size==0;
}
```

For running time:

`add()`: amortized  $O(1)$ , `remove()`: amortized  $O(1)$ , `size()`:  $O(1)$ , `isEmpty()`:  $O(1)$

### 3. R-16.25

Suppose you are stranded on a desert island on which queues are plentiful, but you need a stack. How can you implement a stack using two queues? What is the big-Oh running time of the stack operations?

```
//initialize size
size = 0;
//initialize two empty queue
Queue<T> A = new Queue<>();
Queue<T> B = new Queue<>();
//implement my own push()
void push(T e) {
    if(!A.isEmpty()) {
        A.add(e);
    } else {
        B.add(e);
    }
}
//implement my own pop()
T pop () throws Exception{
    int temp = size - 1;
    if (!A.empty()) {
        while (temp>0) {
            B.add(A.remove());
            temp --;
        }
        return A.remove();
    } else if (!B.empty()) {
        while (temp>0) {
            A.add(B.remove());
            temp --;
        }
        return B.remove();
    } else {
        throw new EmptyStackException();
    }
}

//Implement my own size()
int size() {
    return size;
}
```

```
//Implement my own isEmpty()
boolean isEmpty() {
    return size==0;
}
```

For running time:

push():  $O(1)$ , pop():  $O(n)$ , size():  $O(1)$ , isEmpty():  $O(1)$

4. Suppose you have a stack  $S$  containing  $n$  elements and a queue  $Q$  that is initially empty. Describe how you can use  $Q$  to scan  $S$  to see if it contains a certain element  $x$ , with the additional constraint that your algorithm must return the elements back to  $S$  in their original order. You may not use an array or linked list—only  $S$  and  $Q$  and a constant number of reference variables. What is the running time of your algorithm? (5 points)

$O(n)$ .

```
boolean contains(T x) {
    boolean ret = false;
    //reverse the order to insert elements into Q and check elements
    while(!S.isEmpty()) {
        temp = S.pop();
        if (temp .equals(x)) ret = true;
        Q.add(temp);
    }
    //to insert elements into S
    while(!Q.isEmpty()){
        S.push(Q.remove())
    }
    //reverse the order
    while(!S.isEmpty()){
        Q.add(S.pop())
    }
    //to insert elements into S
    while(!Q.isEmpty()){
        S.push(Q.remove())
    }
    return temp;
}
```

