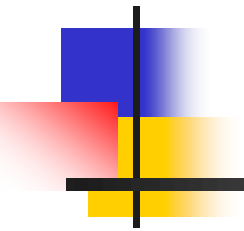
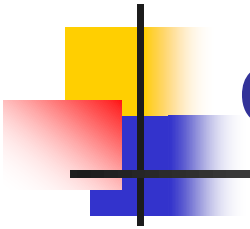


Embedded Software
(MicroC/OS-II kernel)

MicroC/OS-II kernel(scheduler)

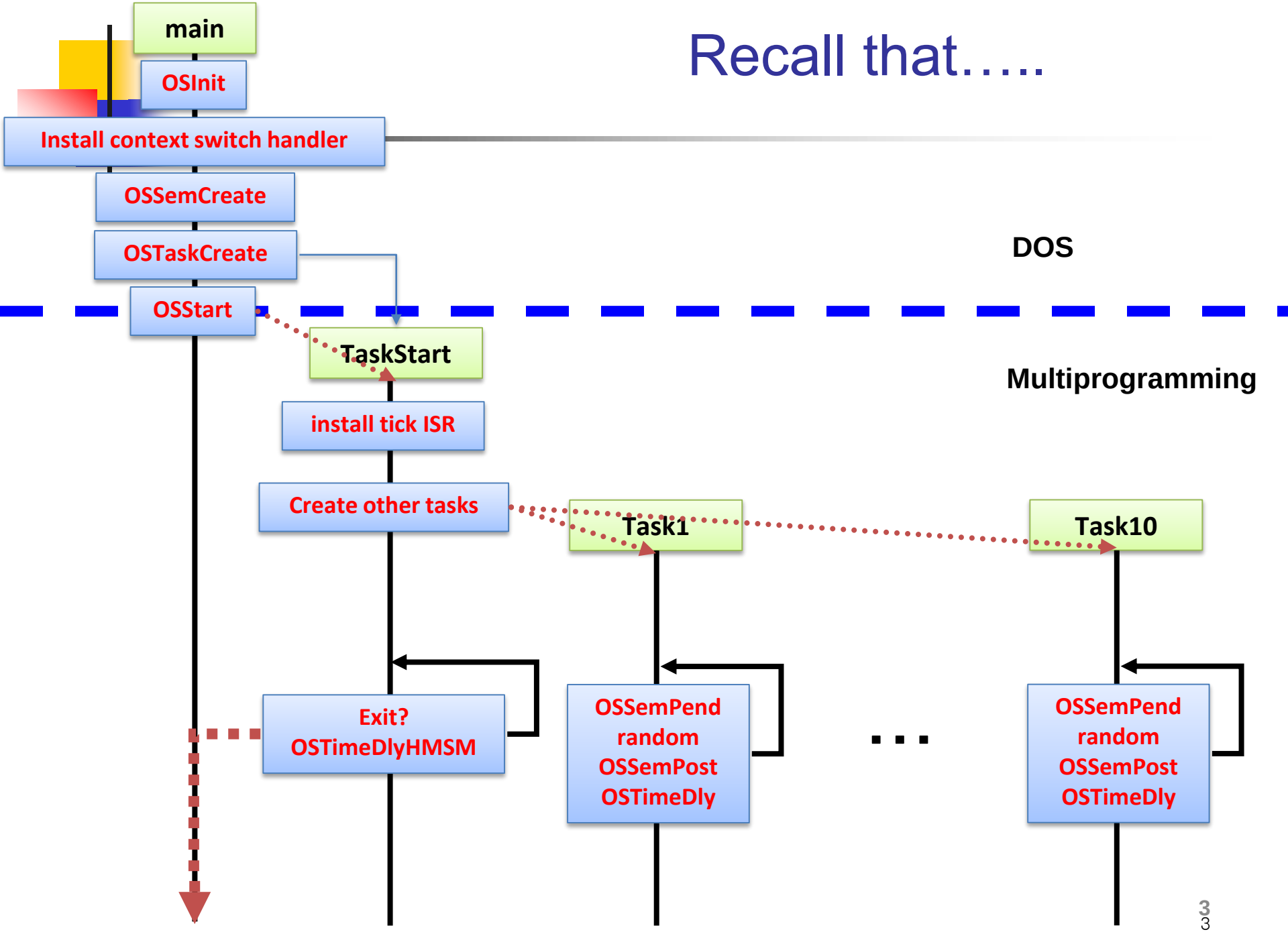
Fall, 2018

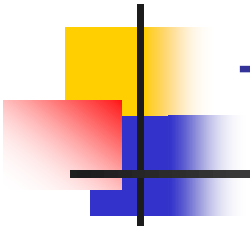


Outline

- Kernel
  - Task types
  - Task Scheduling
    - Overview
    - State
    - TCB (Task Control Block)
    - Ready list
    - Scheduling
    - ISR
  - Clock ticks

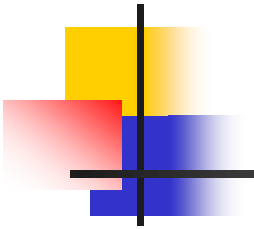
Recall that…..


main
OSInit
Install context switch handler
OSSemCreate
OSTaskCreate
OSStart
DOS
TaskStart
Multiprogramming
install tick ISR
Create other tasks
Task1
Task10
Exit?
OSTimeDlyHMSM
OSSemPend
random
OSSemPost
OSTimeDly
OSSemPend
random
OSSemPost
OSTimeDly



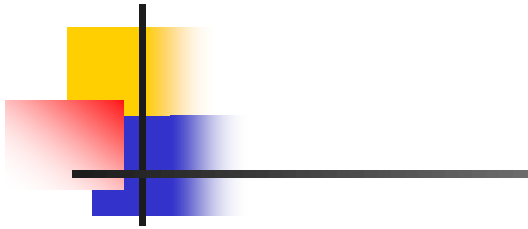
Task types

- Two types
  - Infinite loop
  - After work, task may delete itself



- Infinite loop

```
void MyTask(void *pdata)
{
   while(1)
   {
     Code…
     OSMboxPend();
     OSQPend();
     OSSemPend();
     OSTaskDel(OS_PRIO_SELF);
     OSTaskSuspend(OS_PRIO_SELF);
     OSTimeDly();
     OSTimeDlyHMSM();
     Code…
   }
}
```



- After work, task may delete itself

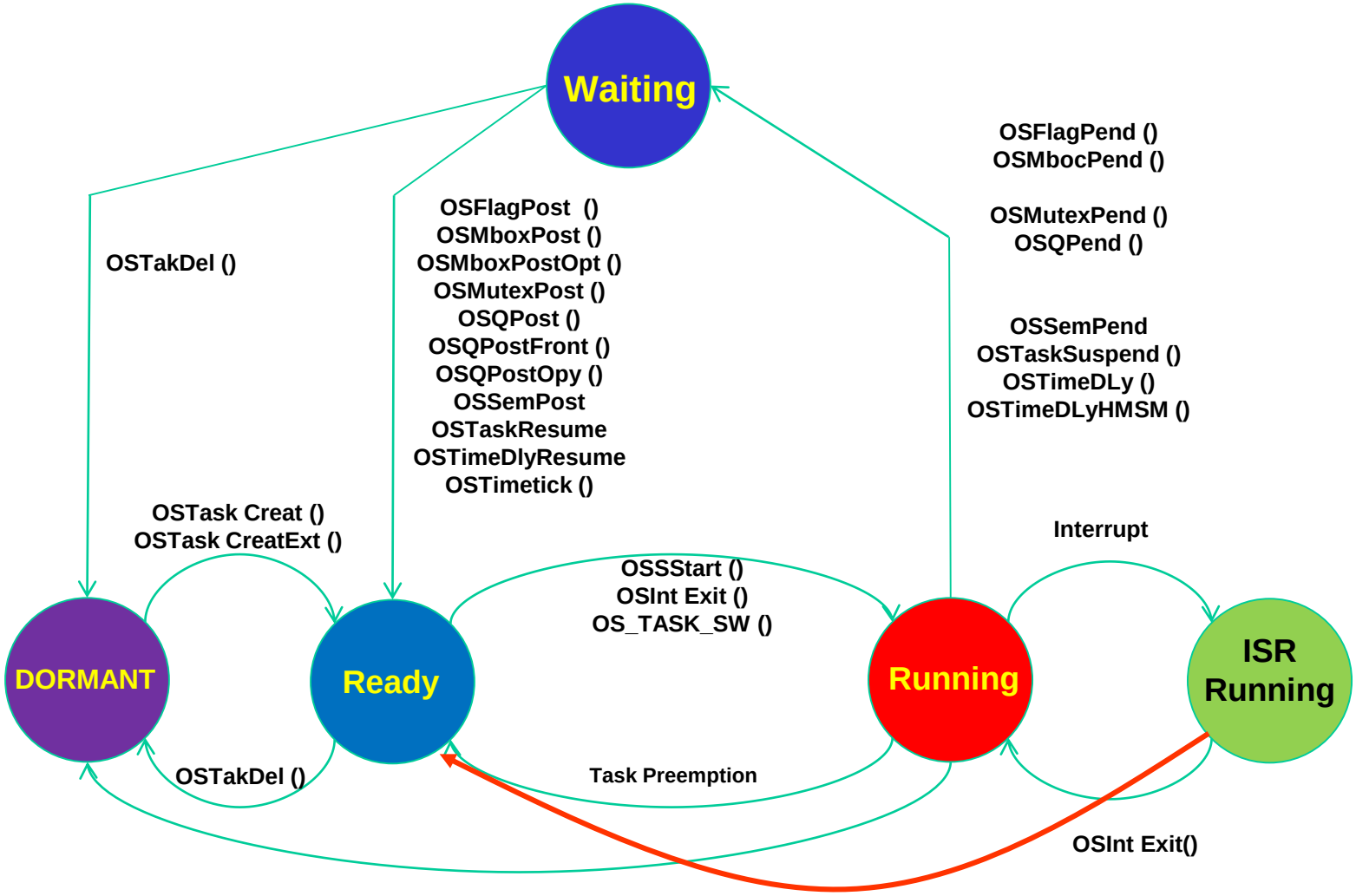
```
void MyTask(void *pdata)
{
   Code…
   OSTaskDel(OS_PRIO_SELF);
}
```

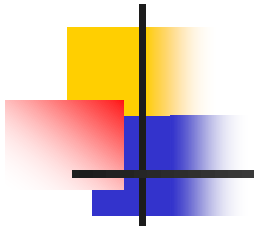
Task scheduling - overview

- Number of tasks
  - 64
- Static priority scheduling
- Priority
  - Unique identifier of tasks
  - 0 ~ 63
  - 0 : The highest priority,
  - 63: The lowest priority
  - Priority may be changed using OSTaskChangePrio() function
  - Assignment
    - 0, 1, 2, 3, …, OS_LOWEST_PRIO-3, OS_LOWEST_PRIO-2, OS_LOWEST_PRIO-1, OS_LOWEST_PRIO

Task scheduling – task state

- Task state transition



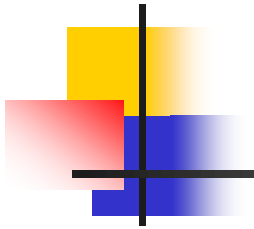


- Task state
  - DORMANT
    - Task is not registered to the OS
    - OSTaskCreate() or OSTaskCreateExt() changes the state of tasks to READY
  - READY
    - Task can be a candidate to be scheduled
  - RUNNING
    - CPU currently executes the task
  - Delayed
    - Task is temporarily stopped during some period
    - After expiration, its state is returned to READY
  - WAIT
    - Task is waiting for some event to occur
    - It cannot be scheduled
  - ISR
    - Interrupt routine is executed

Task scheduling –TCB

- Task control block (TCB)
  - Data structure that contains various information about task (one TCB per task)
  - When the task is scheduled, its information can be used for execution
  - Reside in main memory

- OS_TCB structure (* :OSTaskCreateExt())
  - OSTCBStkPtr
    - A pointer that indicates the top-of-stack
  - OSTCBExtPtr*
    - A pointer that indicates the extended TCB defined by a user
  - OSTCBStkBottom*
    - A pointer that indicates bottom-of-stack
  - OSTCBStkSize
    - Stack size



- OS_TCB structure (* :OSTaskCreateExt())
  - OSTCBNext, OSTCBPrev
    - Pointers used to maintain OS_TCB linked list
  - OSTCBOpt*
    - Optional flag used for task creation
      - *OS_TASK_OPT_STK_CHK*
        - *Investigate stack area when task is created*
      - *OS_TASK_OPT_STK_CLR*
        - *Make stack area with 0*
      - *OS_TASK_OPT_SAVE_FP*
        - *A flag indicating that task uses floating point operation*
  - OSTCBDly
    - Number of clock ticks until the timeout
  - OSTCBStat
    - Task state
  - Others
    - …

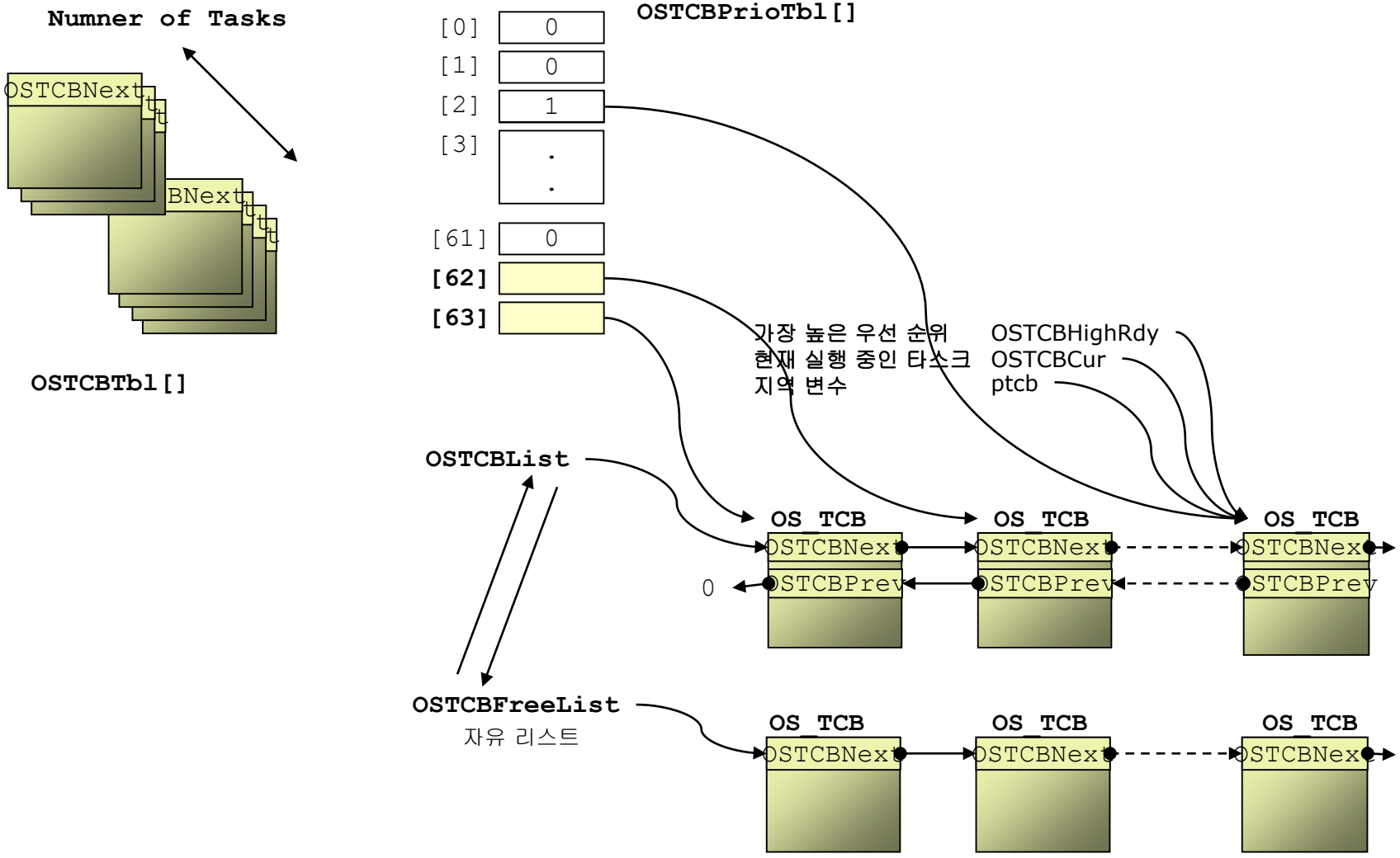
TCB특성

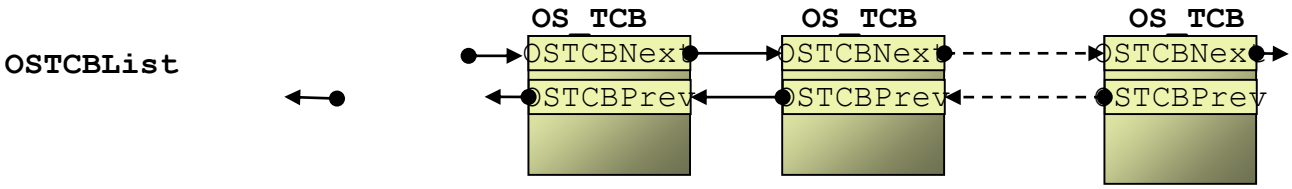
- 기타특징
  - All valid TCB's are ***doubly linked***.
  - Free TCB's are linked in a free list.

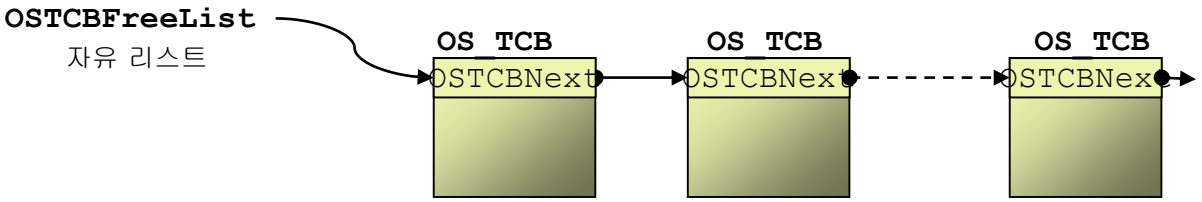
uCOS_II.h

.
.
.

```
460 OS_EXT   OS_TCB          *OSTCBCur;                    /* Pointer to currently running TCB          */
461 OS_EXT   OS_TCB          *OSTCBFreeList;               /* Pointer to list of free TCBs              */
462 OS_EXT   OS_TCB          *OSTCBHighRdy;                /* Pointer to highest priority TCB R-to-R    */
463 OS_EXT   OS_TCB          *OSTCBList;                   /* Pointer to doubly linked list of TCBs     */
464 OS_EXT   OS_TCB          *OSTCBPrioTbl[OS_LOWEST_PRIO + 1];/* Table of pointers to created TCBs      */
465 OS_EXT   OS_TCB           OSTCBTbl[OS_MAX_TASKS + OS_N_SYS_TASKS];   /* Table of TCBs               */
```


Numner of Tasks
OSTCBNext
BNext
OSTCBTbl[]
OSTCBPrioTbl[]
[0] 0
[1] 0
[2] 1
[3]
[61] 0
[62]
[63]
가장 높은 우선 순위
현재 실행 중인 타스크
지역 변수
OSTCBHighRdy
OSTCBCur
ptcb
OSTCBList
OS_TCB
OSTCBNext
OSTCBPrev
0
OSTCBFreeList
자유 리스트





OSTaskCreate

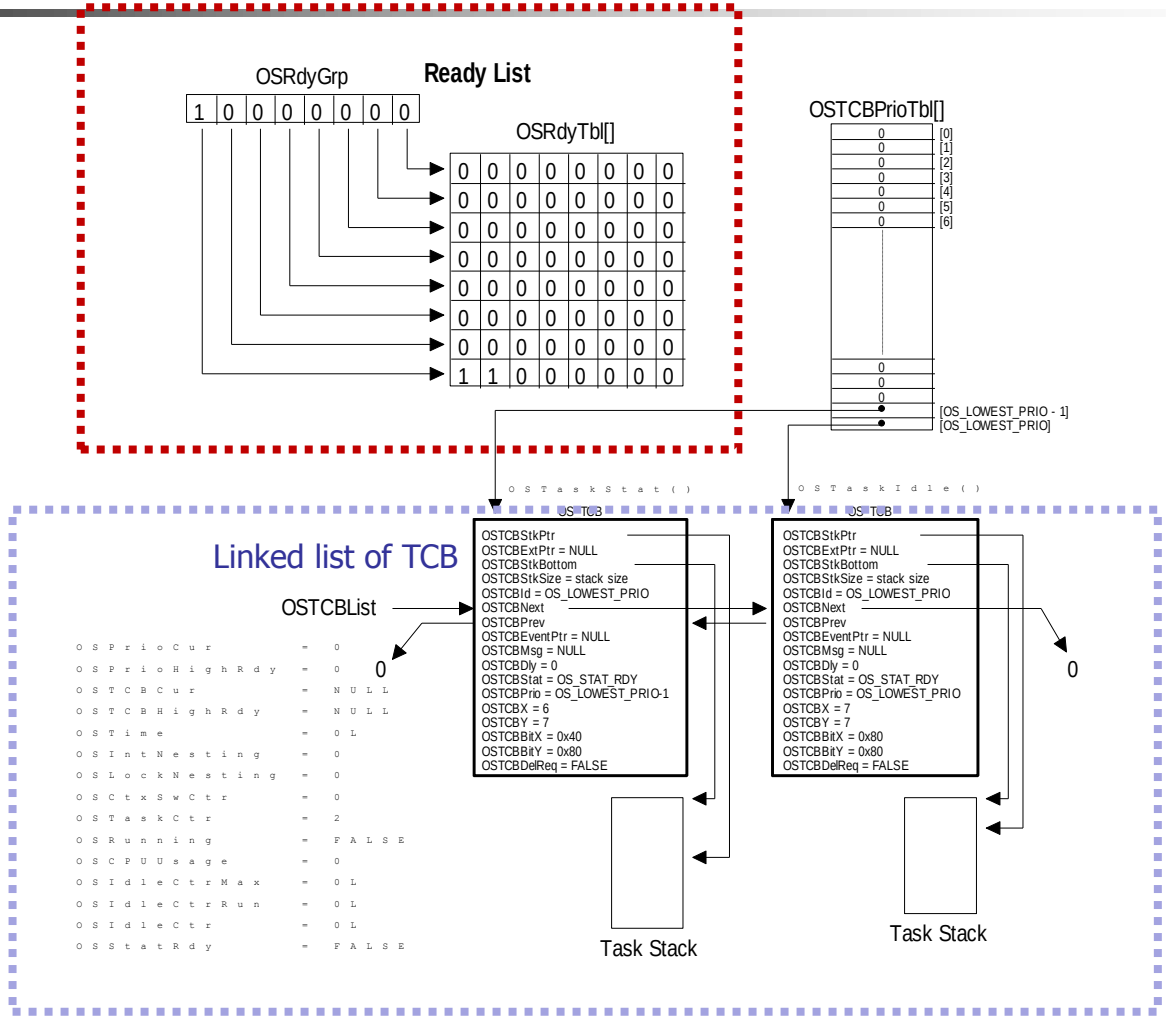
Ready List

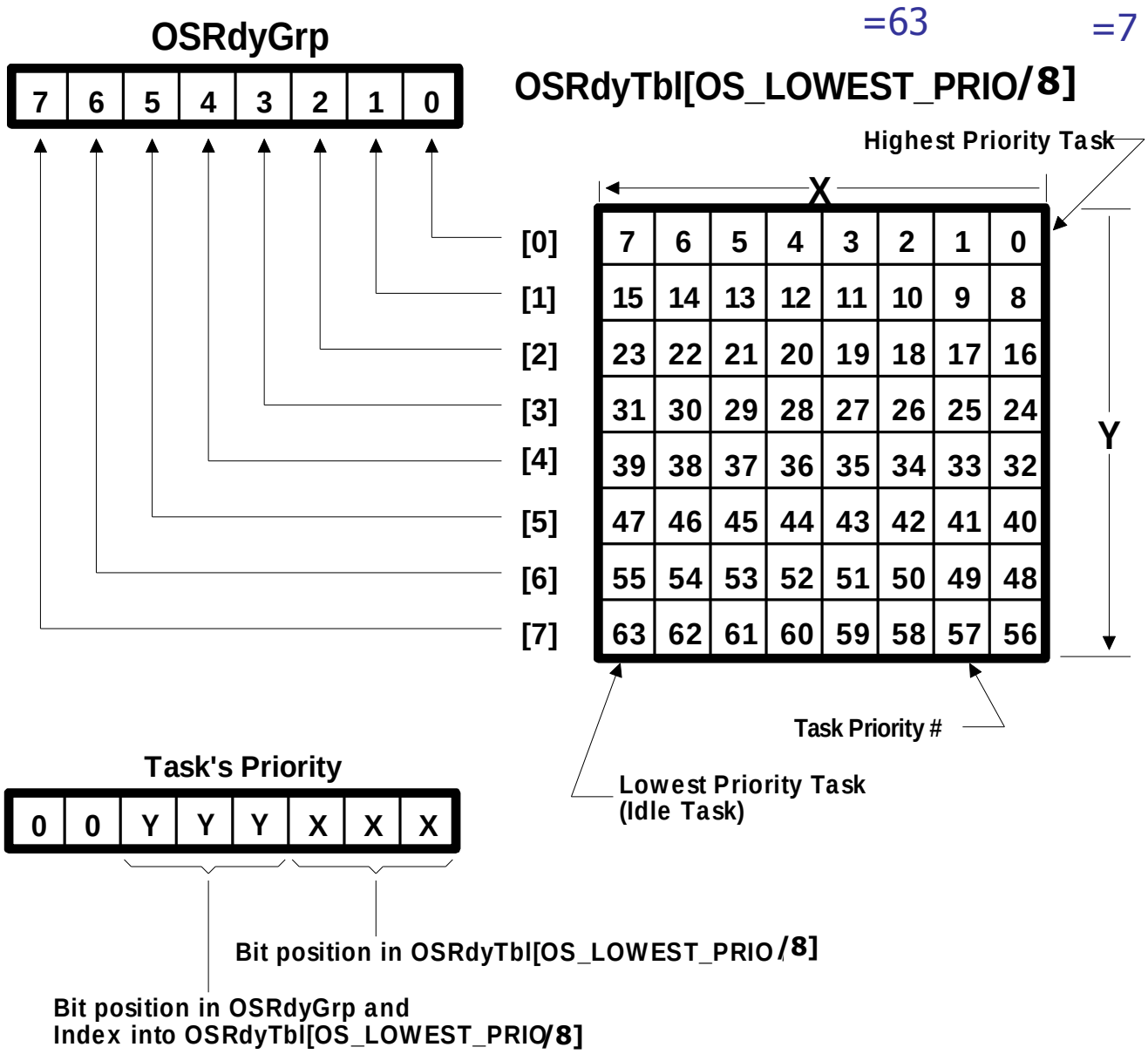
- Ready list is a special bitmap to reflect which task is currently in the ready state.
  - Each task is identified by its unique priority in the bitmap.
- A primary design consideration of the ready list is how to efficiently locate the highest-priority ready task.
  - The designer could trade some ROM space for an improved performance.
- If a linear list is adopted, it takes O(n) to locate the highest-priority ready task.
  - It takes O(log n) if a heap is adopted.
  - Under the design of ready list of µC/OS-II, it takes only O(1).
    - Note that the space consumption is much more than other approaches, and it also depends on the bus width.

Task scheduling–ready list

- Why Bitmap?
  - To find a prior task in a constant time
- Ready list
  - Data structure
    - OSRdyGrp
      - Bit indicating group
    - OSRdyTbl[]
      - 8 task in the same group
    - OSMapTbl[]
      - Array used for bit mask
    - OSUnMapTbl[]
      - Array used to search the highest priority in ready list
  - Recall that there are total 64 tasks max !!

Scheduler가 관리하는 ready list




OSRdyGrp
7 6 5 4 3 2 1 0
=63
=7
OSRdyTbl[OS_LOWEST_PRIO/8]
Highest Priority Task
X
Y
[0] 7 6 5 4 3 2 1 0
[1] 15 14 13 12 11 10 9 8
[2] 23 22 21 20 19 18 17 16
[3] 31 30 29 28 27 26 25 24
[4] 39 38 37 36 35 34 33 32
[5] 47 46 45 44 43 42 41 40
[6] 55 54 53 52 51 50 49 48
[7] 63 62 61 60 59 58 57 56
Task Priority #
Lowest Priority Task
(Idle Task)
Task's Priority
0 0 Y Y Y X X X
Bit position in OSRdyTbl[OS_LOWEST_PRIO/8]
Bit position in OSRdyGrp and
Index into OSRdyTbl[OS_LOWEST_PRIO/8]

Quiz

OSRdyGrp

?	?	?	?	?	?	?	?

OSRdyTbl

	1	1	1				
		1	1				
1	1						

OSMapTbl

Index	Bit mask (Binary)
0	00000001
1	00000010
2	00000100
3	00001000
4	00010000
5	00100000
6	01000000
7	10000000

Bit 0 in **OSRdyGrp** is 1 when any bit in **OSRdyTbl[0]** is 1.
Bit 1 in **OSRdyGrp** is 1 when any bit in **OSRdyTbl[1]** is 1.
Bit 2 in **OSRdyGrp** is 1 when any bit in **OSRdyTbl[2]** is 1.
Bit 3 in **OSRdyGrp** is 1 when any bit in **OSRdyTbl[3]** is 1.
Bit 4 in **OSRdyGrp** is 1 when any bit in **OSRdyTbl[4]** is 1.
Bit 5 in **OSRdyGrp** is 1 when any bit in **OSRdyTbl[5]** is 1.
Bit 6 in **OSRdyGrp** is 1 when any bit in **OSRdyTbl[6]** is 1.
Bit 7 in **OSRdyGrp** is 1 when any bit in **OSRdyTbl[7]** is 1.

- Make a task ready to run: (set proper bit in rdygrp&table)

<Task insertion into ready list>
1. OSRdyGrp bit -> 1
2. Associated bit in OSRdyTbl[] -> 1

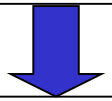
<Task insertion>

```
OSRdyGrp |= OSMapTbl[prio >> 3];
OSRdyTbl[prio >> 3] |= OSMapTbl[prio & 0x07];
```

- Remove a task from the ready list:

<Task deletion from ready list>

1. Associated bit in OSRdyTbl[] -> 0
2. If there is no 1 in OSRdyTbl[], then OSRdyGrp bit -> 0



<Task deletion from ready list>

```
if ((OSRdyTbl[prio >> 3] &= ~OSMapTbl[prio & 0x07]) == 0)
        OSRdyGrp &= ~OSMapTbl[prio >> 3];
```

```
char x,y,mask;

  x = prio & 0x07;
  y = prio >> 3;
  mask = ~(OSMapTbl[x]);            // a mask for bit clearing
  if((OSRdyTbl[y] &= mask) == 0) // clear the task's bit
  {                                 // the group bit should be cleared too
     mask = ~(OSMapTbl[y]);         // another bit mask…
     OSRdyGrp &= mask;                                  // clear the group bit
  }
```

To Find the highest priority task

- Finding the highest-prior ity task ready to run:

```
y    = OSUnMapTbl[OSRdyGrp];
x    = OSUnMapTbl[OSRdyTbl[y]];
prio = (y << 3) + x;
```

This matrix is used to locate the fi rst LSB which is '1', by given a va lue.

For example, if 00110010 is given, then '1' is returned.

```
INT8U const OSUnMapTbl[] = {
   0, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   6, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   7, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   6, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,
   4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0
};
```

Link

EX

LSB 기준 최초 1인 비트 번호를 표시

OSRdyGrp contains 0x68 = **01101000**

```
INT8U  const  OSUnMapTbl[] = {
    0, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x00 to 0x0F */
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x10 to 0x1F */
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x20 to 0x2F */
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x30 to 0x3F */
    6, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x40 to 0x4F */
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x50 to 0x5F */
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x60 to 0x6F */
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x70 to 0x7F */
    7, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x80 to 0x8F */
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0x90 to 0x9F */
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0xA0 to 0xAF */
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0xB0 to 0xBF */
    6, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0xC0 to 0xCF */
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0xD0 to 0xDF */
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,       /* 0xE0 to 0xEF */
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0        /* 0xF0 to 0xFF */
};
```

OSRdyTbl[3] contains 0xE4 = **11100100**

```
 3 = OSUnMapTbl[  0x68 ];
 2 = OSUnMapTbl[  0xE4 ];
26 = ( 3 << 3) +  2;
```

<Searching the highest priority >

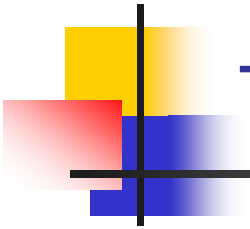
<example> OSRdyGrp: 01101000, 0x68 OSRdyTbl[3]: 11100100, 0xE4

3 << 3 + 2 = 26

Link

<Searching the highest priority >

```
y = OSUnMapTbl[OSRdyGrp];
x = OSUnMapTbl[OSRdyTbl[y]];
prio = (y << 3) + x;
```



Task scheduling–scheduler

- Task scheduling
  - Task 수준의 Scheduling은 OS_Sched() 함수에 의해 수행
  - ISR 수준의 Scheduling은 OSIntExit() 함수에 의해 수행

- OS_Sched()

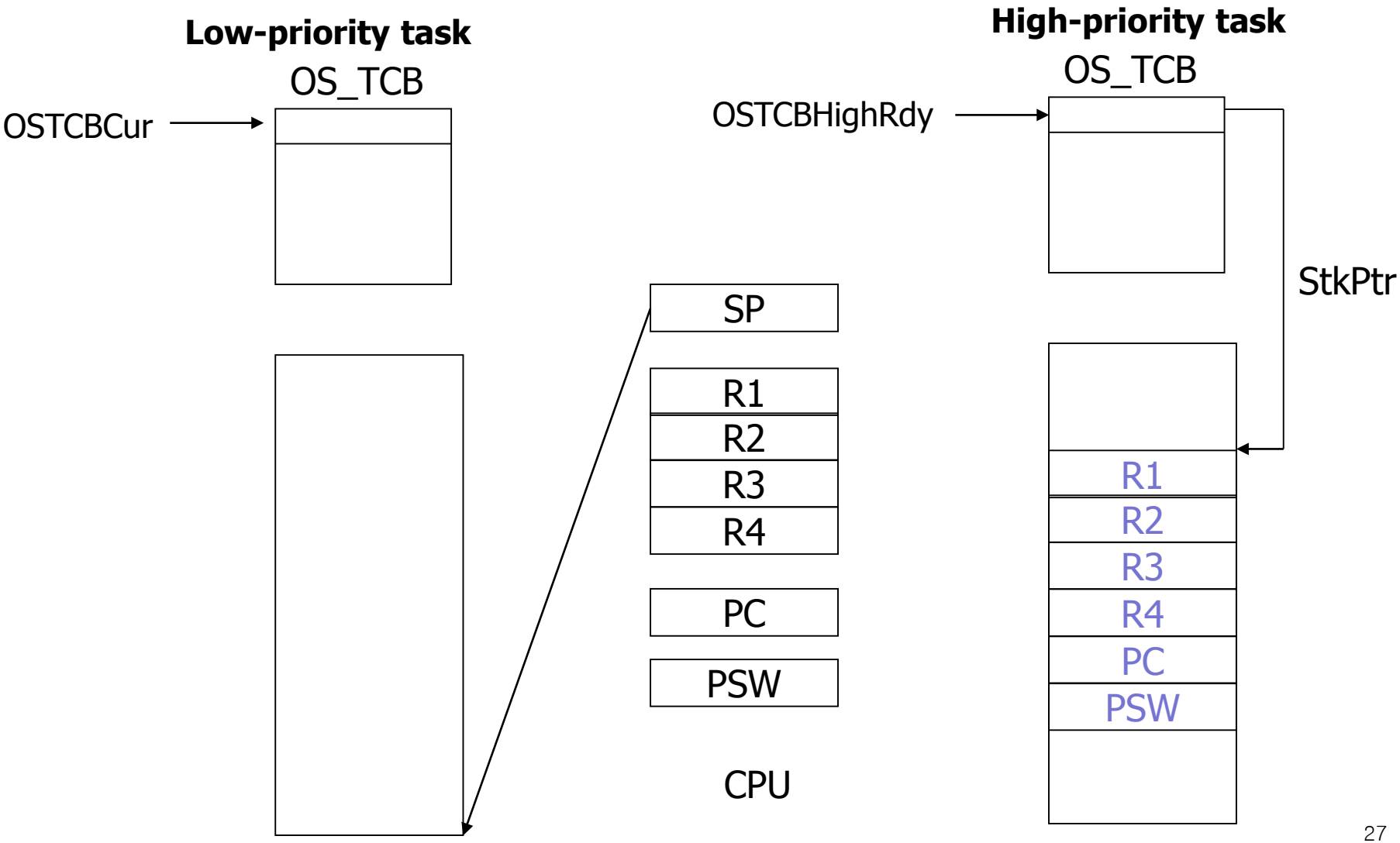
다른 task가 lock을 걸어 강제로 schedule을 막은 상태가 아님

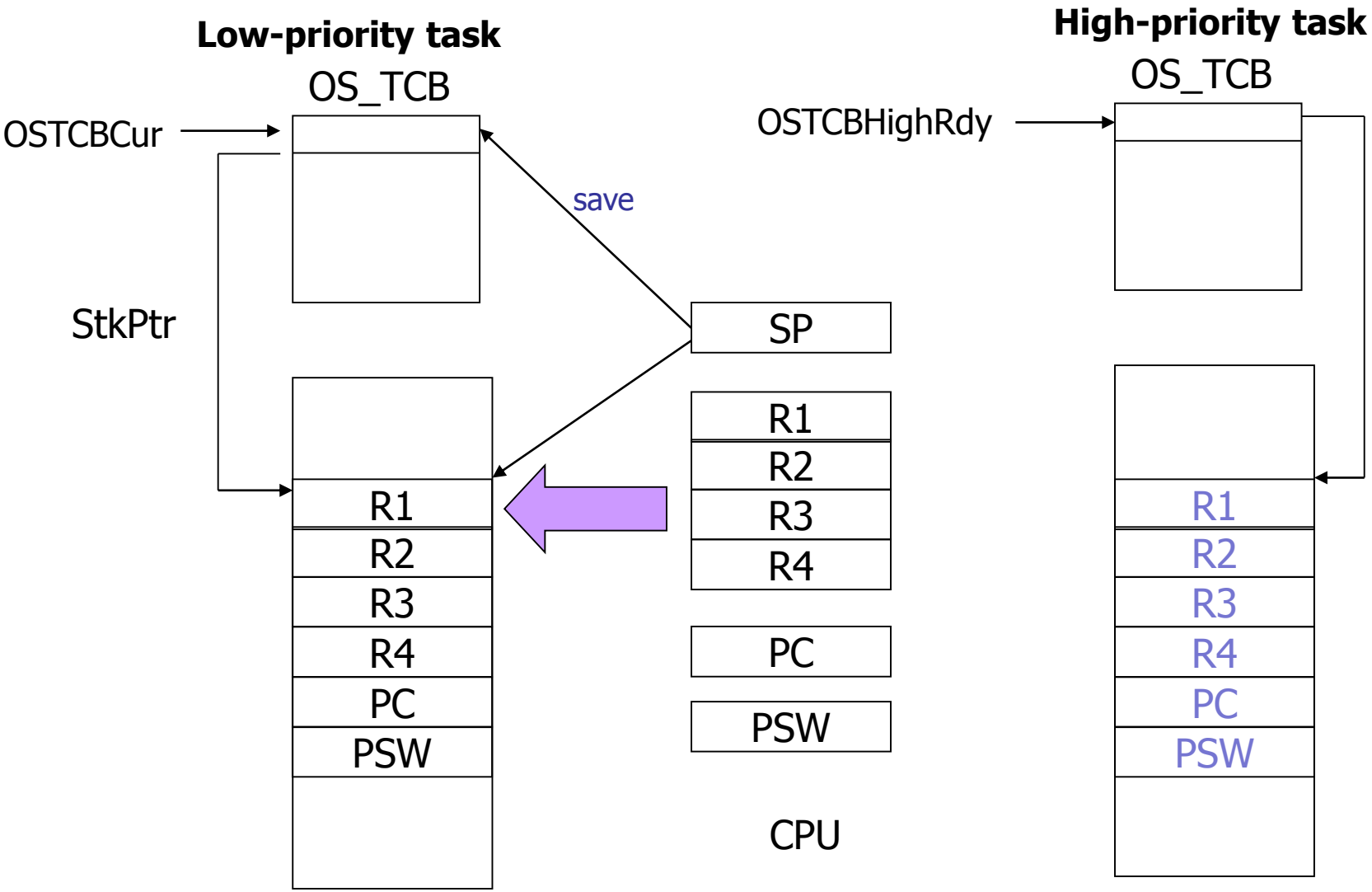
Interrupt가 호출되지 않았음

```
void OSSched (void)
{
   INT8U y;
   OS_ENTER_CRITICAL();
   if ((OSLockNesting | OSIntNesting) == 0) {
      y = OSUnMapTbl[OSRdyGrp];
      OSPrioHighRdy = (INT8U)((y << 3) + OSUnMapTbl[OSRdyTbl[y]]);
      if (OSPrioHighRdy != OSPrioCur) {
         OSTCBHighRdy = OSTCBPrioTbl[OSPrioHighRdy];
         OSCtxSwCtr++;
         OS_TASK_SW();
      }
   }
   OS_EXIT_CRITICAL();
}
```

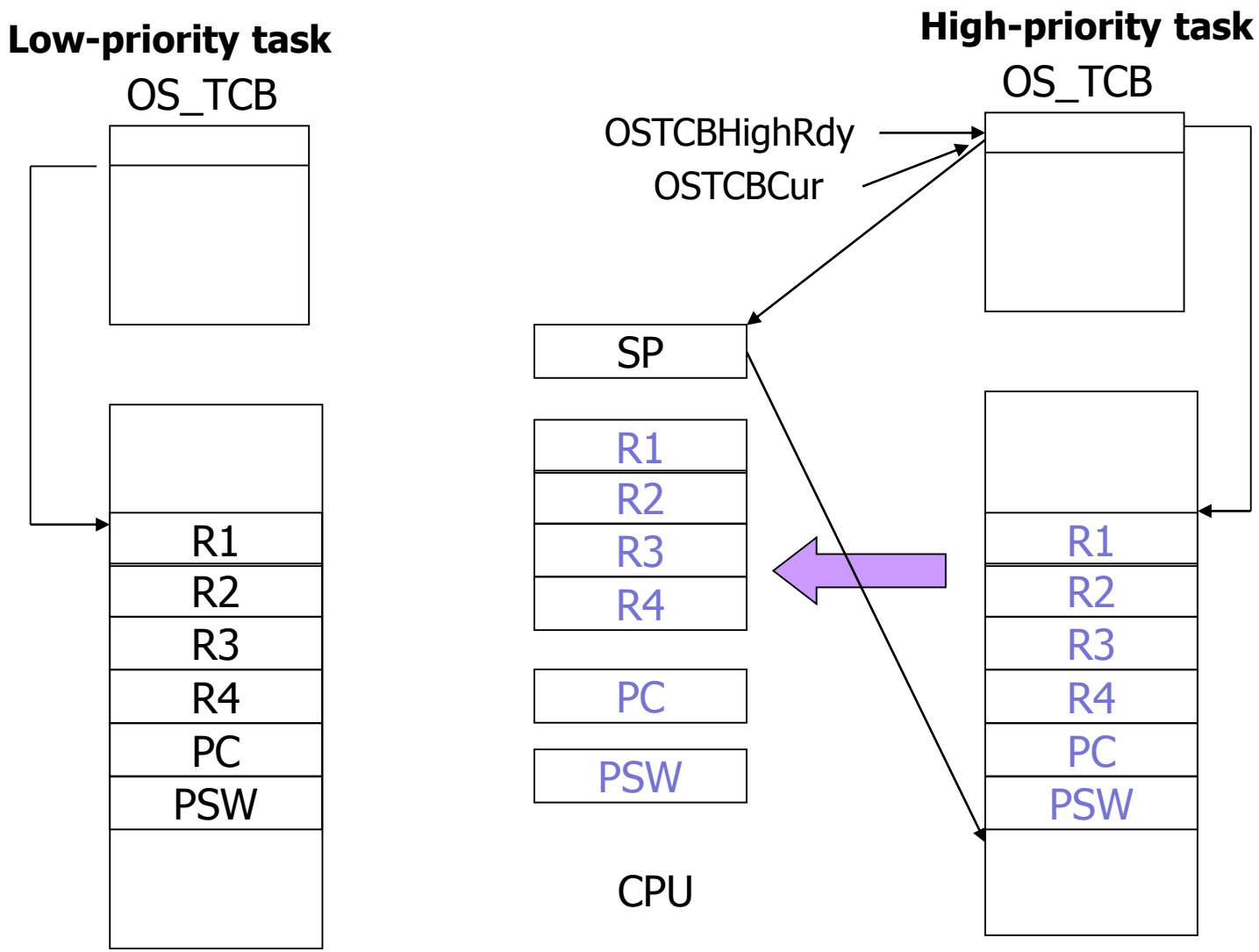
현재task가 priority가 제일 높으면 switching 없음

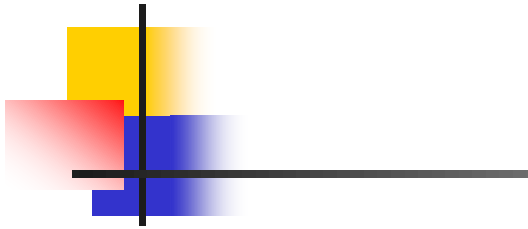
Before OS_TASK_SW()




OS_TASK_SW()
Low-priority task
OS_TCB
OSTCBCur
StkPtr
save
SP
R1
R2
R3
R4
PC
PSW
CPU
High-priority task
OS_TCB
OSTCBHighRdy
R1
R2
R3
R4
PC
PSW

OS_TASK_SW()


Low-priority task
OS_TCB
R1
R2
R3
R4
PC
PSW
High-priority task
OS_TCB
OSTCBHighRdy
OSTCBCur
SP
R1
R2
R3
R4
PC
PSW
CPU
R1
R2
R3
R4
PC
PSW



Task 수준에서의 context switch

```
void OSCtxSW (void)
{
   Values of R1, R2, R3, R4  and so on are pushed to stack;
   OSTCBCur->OSTCBStkPtr= SP;
   OSTCBCur                    = OSTCBHighRdy;
   SP                          = OSCBHighRdy->OSTCBStkPtr;
   Values of R4, R3, R2, R1 and so on are popped;
}
```

Locking and Unlocking the Scheduler

Locking scheduler

```
void OSSchedLock (void)
{
   if (OSRunning == TRUE) {
     OS_ENTER_CRITICAL();
     OSLockNesting++;
     OS_EXIT_CRITICAL();
   }
}
```

Unlocking scheduler

```
void OSSchedUnlock (void)
{
   if (OSRunning == TRUE) {
      OS_ENTER_CRITICAL();
      if (OSLockNesting > 0) {
         OSLockNesting--;
         if ((OSLockNesting | OSIntNesting) == 0) {
            OS_EXIT_CRITICAL();
            OSSched();
         } else {
            OS_EXIT_CRITICAL();
         }
      } else {
         OS_EXIT_CRITICAL();
      }
   }
}
```

Task scheduling -ISR

- ISR
  - µC/OS-II requires an ISR being written in assembly if your compiler does not support in-line assembly!

```
MyISR:
   Save all register values to TCB
   Call OSIntEnter function

   Execute ISR program

   Call OSIntExit()
   Restore all register values to the CPU
```

```
void OSIntEnter (void)
{
   OS_ENTER_CRITICAL();
   OSIntNesting++;
   OS_EXIT_CRITICAL();
}
```

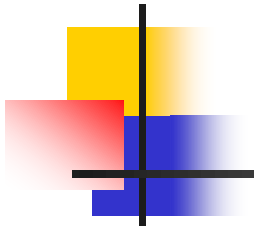
ISR 수준의 Scheduling은 OSIntExit() 함수에 의해 수행

```
void OSIntExit (void)
{
   OS_ENTER_CRITICAL();
   if ((--OSIntNesting | OSLockNesting) == 0) {
      OSIntExitY    = OSUnMapTbl[OSRdyGrp];
      OSPrioHighRdy = (INT8U)((OSIntExitY << 3) +
                      OSUnMapTbl[OSRdyTbl[OSIntExitY]]);
      if (OSPrioHighRdy != OSPrioCur) {
         OSTCBHighRdy  = OSTCBPrioTbl[OSPrioHighRdy];
         OSCtxSwCtr++;
         OSIntCtxSw();
      }
   }
   OS_EXIT_CRITICAL();
}
```

If no scheduler locking and no interrupt nesting

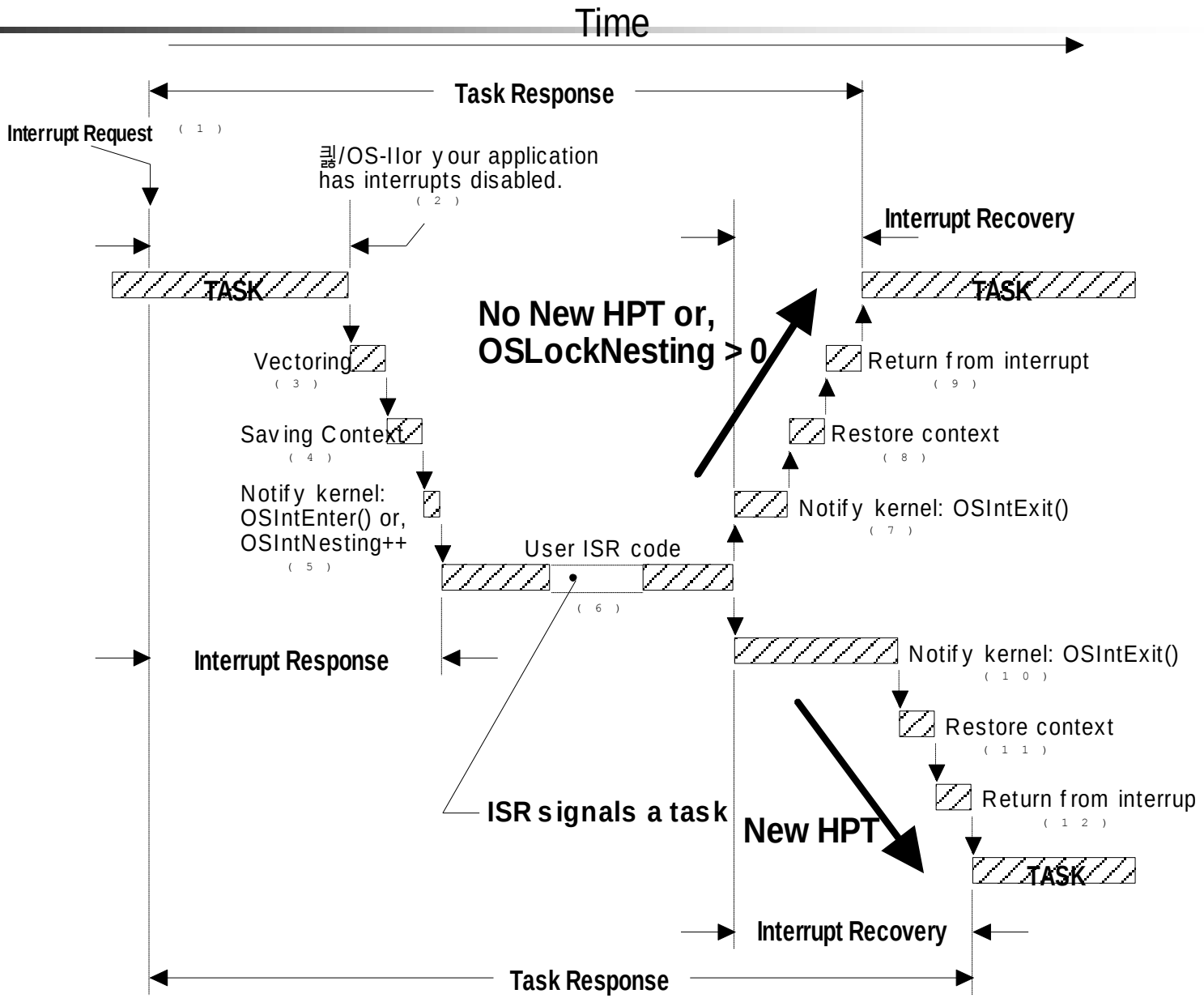
If there is another high-priority task ready

A context switch is executed. Not OSCtxSW()



- OSIntCtxSw()
  - OSIntExit()에서 context switch 수행
  - OSCtxSw()와 유사
    - 차이점: ISR에서 task를 수행시키므로, 이전 task의 문맥 저장 (stroing all CPU regiseters onto stack) 이 없음

```
void  OSIntCtxSw (void)
{
    Call OSTaskSwHook();
    OSTCBCur   = OSTCBHighRdy;
    OSPriorCur = OSPrioHighRdy;
    SP         = OSTCBHighRdy->OSTCBStkPtr;
    POP R4, R3, R2 and R1 from the task's stack;
    Execute a return from interrupt instruction;
}
```


Time
Task Response
Interrupt Request
(1)
킬/OS-IIor your application
has interrupts disabled.
(2)
Interrupt Recovery
TASK
TASK
No New HPT or,
OSLockNesting > 0
Vectoring
(3)
Return from interrupt
(9)
Saving Context
(4)
Restore context
(8)
Notify kernel:
OSIntEnter() or,
OSIntNesting++
(5)
Notify kernel: OSIntExit()
(7)
User ISR code
(6)
Interrupt Response
Notify kernel: OSIntExit()
(1 0)
Restore context
(1 1)
Return from interrup
(1 2)
ISR signals a task
New HPT
TASK
Interrupt Recovery
Task Response

Interrupts under µC/OS-II

```
void OSIntExit (void)
{
    OS_ENTER_CRITICAL();
    if ((--OSIntNesting | OSLockNesting) == 0) {
        OSIntExitY    = OSUnMapTbl[OSRdyGrp];
        OSPrioHighRdy = (INT8U)((OSIntExitY << 3) +
                        OSUnMapTbl[OSRdyTbl[OSIntExitY]]);
        if (OSPrioHighRdy != OSPrioCur) {
            OSTCBHighRdy  = OSTCBPrioTbl[OSPrioHighRdy];
            OSCtxSwCtr++;
            OSIntCtxSw();
        }
    }
    OS_EXIT_CRITICAL();
}
```

If no scheduler locking and no interrupt nesting

If there is another high-priority task ready

A context switch is executed.

Note that OSIntCtxSw() is called, instead of OS_TASK_SW(), because the ISR already saves the CPU registers onto the stack.

```
void OSIntEnter (void)
{
    OS_ENTER_CRITICAL();
    OSIntNesting++;
    OS_EXIT_CRITICAL();
}
```

Clock ticks

- Time source in OS
- OSTickISR() (in OS_CPU_A.ASM)
  - uC/OS-II에서 타임 아웃 기능과 시간지연 기능 등을 위해 Clock Tick을 사용
  - Multitasking을 시작한 후, 즉 반드시 OSStart()를 호출한 뒤에 task내에서 Clock Tick Interrupt를 활성화 해야 함 (ex. in the TaskStart())
  - Tick ISR이 OSTimeTick()을 호출함

```
void OSTickISR(void) {
    Save processor registers;
    Call OSIntEnter()

    Call OSTimeTick();

    Call OSIntExit();
    Restore processor registers;
    Execute a return from interrupt instruction;
}
```

Tick ISR

```
void OSTimeTick (void)
{
    while (ptcb->OSTCBPrio != OS_IDLE_PRIO) {
      OS_ENTER_CRITICAL();
      if (ptcb->OSTCBDly != 0) {
         if (--ptcb->OSTCBDly == 0) {
            if (!(ptcb->OSTCBStat & OS_STAT_SUSPEND)) {
               OSRdyGrp               |= ptcb->OSTCBBitY;
               OSRdyTbl[ptcb->OSTCBY] |= ptcb->OSTCBBitX;
            } else {
               ptcb->OSTCBDly = 1;
            }
         }
      }
      ptcb = ptcb->OSTCBNext;
      OS_EXIT_CRITICAL();
   }
}
```

Decrement delay-counter if needed

If the delay-counter reaches zero, make the task ready. Or, the task remains waiting.

```
OS_ENTER_CRITICAL();
 OSTime++;
 OS_EXIT_CRITICAL();
```