

Lost frames are sent several times wasting bandwidth

① Error correction & detection.

→ why errors?

↳ Ideal / errorless

↳ error-prone links

→ Types: → single (rare) (delay/noise is large)
↳ multiple (usual)
↳ burst (consecutive) Burst length (bits)

→ Redundancy or EDC bits

→ Detection vs correction

→ Forward error correction: receiver guesses the actual data using EDC bits (if small error)

→ Retransmission: sender retransmits (cost 1)

Parity check.

Data (D) | Parity (P)

↓
Even

↓
Odd

no. of 1's	P	Total 1's
odd	1	Even
even	0	Even

no. of 1's	P	total D+P
odd	0	odd
even	1	odd

$D+P = \text{Even ones}$

$D+P = \text{odd ones}$

PIC →

APC

0/0/2 → port number

1 → pic

flexible pic

→ Even parity detects odd number of bits errors

→ 2P parity add rowwise & columnwise both
it can detect & correct one bit errors

→ Hamming Code:

$$2^r \geq m + r + 1$$

$$m = 3 \text{ min } r = 3.$$

R.W. Hamming

detect and correct errors

Redundant bits are in between.

→ CRC (Cyclic)

→ Needs: 1) XOR

2) when n is multiplied by 2^k
then left shift by k

→ Sender and receiver agree on q (generator)
 $r+1$ bits. Sender appends r bits after
times left shifting data.

$$\text{data} \ll r = \text{codeword}$$

$$\text{codeword} \% q = 0 \text{ (no error)}$$

$$\text{Sender } (S) = \underbrace{(D \cdot 2^r)}_{\text{LS by } r} \text{ XOR } \underbrace{R}_{\text{appends } r \text{ bits.}}$$

$$d \ll r \% q = 0$$

So choose r accordingly.

$$\text{Sender } R = \text{remainder of } \frac{D \cdot 2^r}{q}$$

$$T = \text{remainder of } \frac{S}{q}$$

if $T=0$ no error

Else error

- ARP - plug & play - TTL (20 min - 1 hr). No admin role. Entry flushed if no traffic from the other host after TTL is expired.
- All routers / switches / hosts have own ARP table and logic to update & flush it. When system enters entry is made and removed after TTL. Time at which populated is there.

→ Process: (Same Subnet).

- No row for dest in sender's ARP table.
- Create ARP Query with dest IP and network adapter put broadcast mac address saying it's broadcast.
- So adapter broadcast ARP Query. When it reaches intended dest with matching IP.
- Broadcast: FF - (6) times

Different subnet

- Dest IP is dest from diff subnet
- MAC isn't broadcast its MAC of router's interface.

→ Router replace its MAC to dest MAC

Link layer switch

① When packet arrives from source switch does

MAC	Interface ID	Time.
m1	1	1:28
m2	3	1:30
m3	2	1:35

So if m2 send on 3 interface.
After TTL discard entry.

It is self-learning. It's play & play.

② filtering & forwarding

↓
discarding
↓
not sending

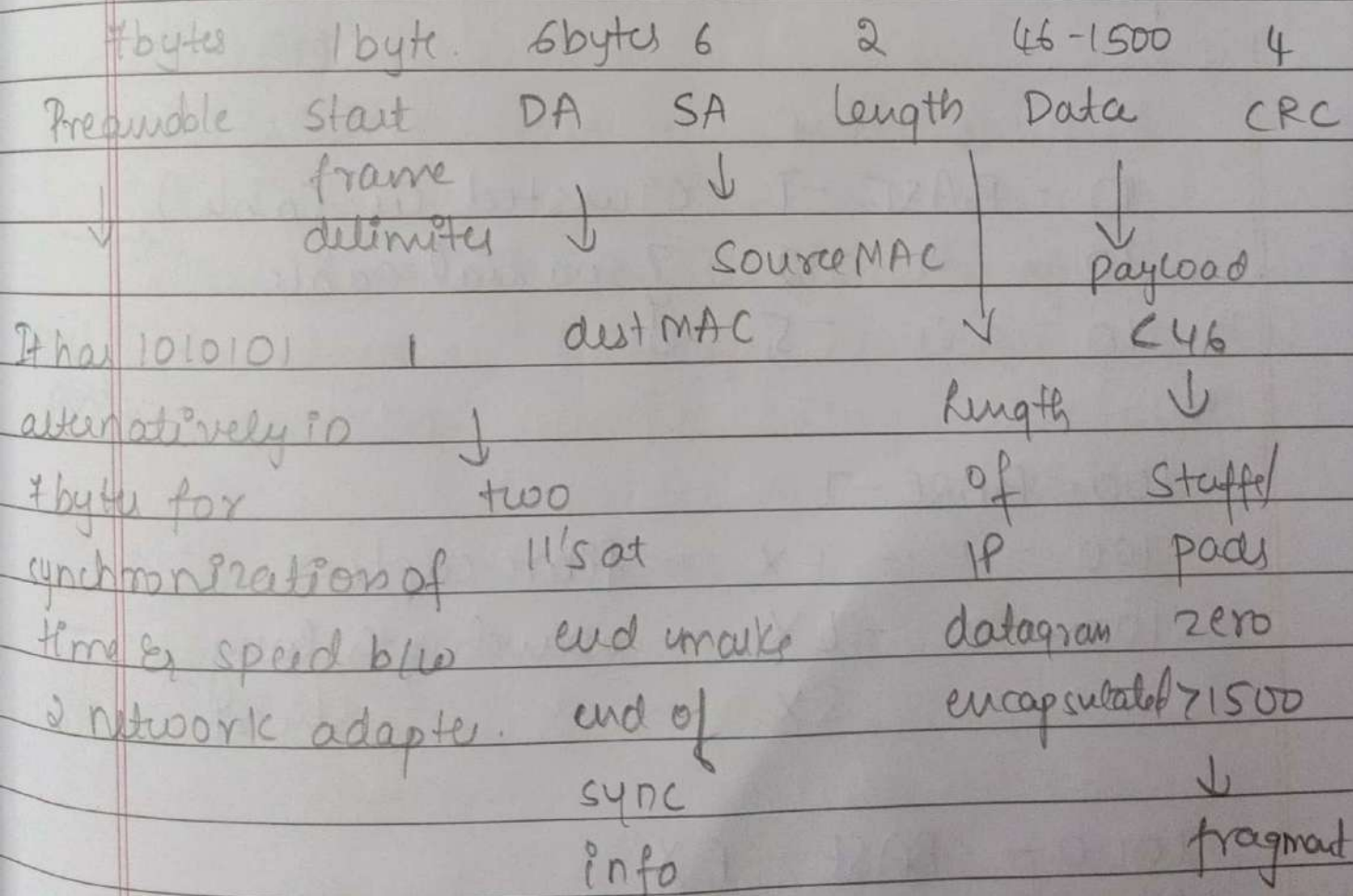
→ Send to all nodes in network if mac isn't known so that it gets discarded if no matching MAC

→ If matching MAC check table transfer on interface

→ If packet has dest with interface link same as sender discarded

Ethernet

- It helps in switched LAN's or mac address forwarding.
- IP header is encapsulated in Ethernet frame
- frame format



Repeaters

→ A microwave link is used as medium for 802.3 Ethernet

Ethernet technology

number	-	base	-	distance / type of cable
↓		↓		↓
Mbps of data can be trans- ferred		baseband network		max distance signal can reach without repeaters

10 - BASE - T (Twisted Cu Cable)
10 - " - 2 } coaxial cable
10 - " - 5 }

100 - BASE - T

100 - " - FX - fibre cable
100 - " - LX - wavelength long
SX - " short

1000 - BASE - FX
1 Gbps

10 G - BASE - FX

- Traffic Isolation when more switches for few hosts in subnet
- Only some ports of switch works for LAN's and traffic isolation is for ports
- So when host have virtual LAN we can shrink ports no for than LAN make it free, or give to another one.
- But how it replaces switch managing LAN?
- Efficient use of switches
- last port of one switch to another switch to increase multiple VLANs — VLAN Trunking.

New Ethernet frame

SA VLAN TAG length

4 bytes



identified to which VLAN packet should go

MPLS (multi protocol label switching)

- ① Combines both circuit & packet switching networks
- ② MP - Independent of lower layer protocol (OSI)
- ③ LS - based on packets i/p labels it determines o/p label & sends just like in virtual circuit
- ④ They come with MPLS / link switched routers
- ⑤ Extremely fast forwarding
- ⑥ MPLS table

i/p	o/p
Replaces	

④ Border protocol (L2 & L3)

- ⑤ Take IP packet from L3 then add label transmit to L2

How is it better?

- Packet arrives (IP) ^{at router} → does longest prefix match for dest ip in router's table
- find o/p interface & forwards
- so it takes long times

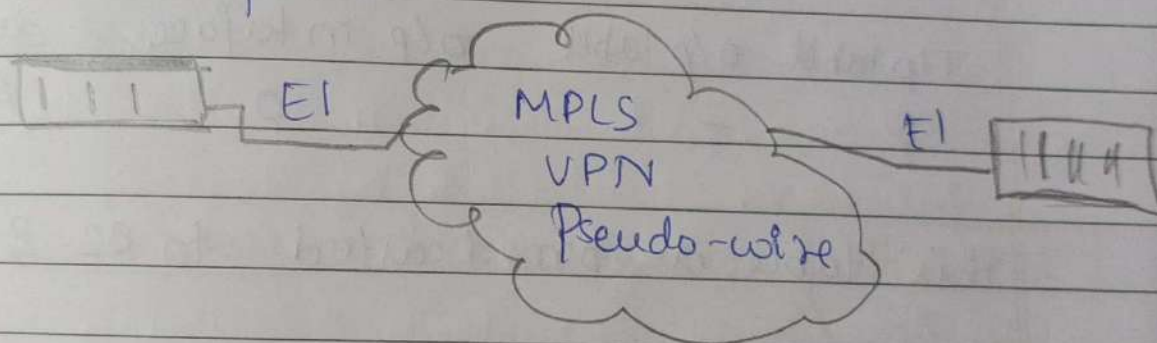
② Virtual / Circuit switched network
with VCI - Virtual circuit Identifier

Ip port VCI O/p port VCI
Replace

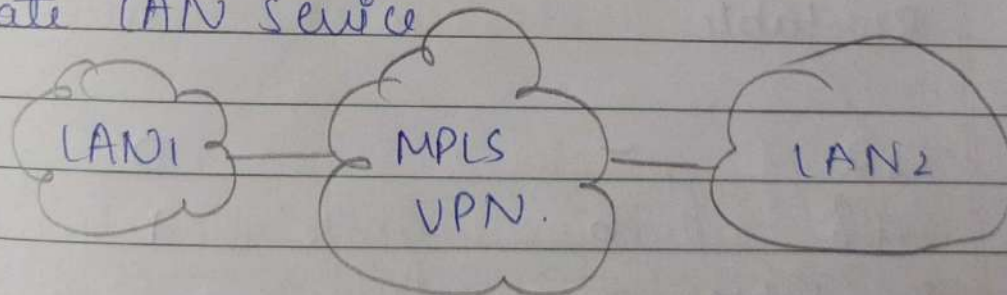
Router + Switch \Rightarrow Label Switch router.

Why?

\rightarrow Point to point



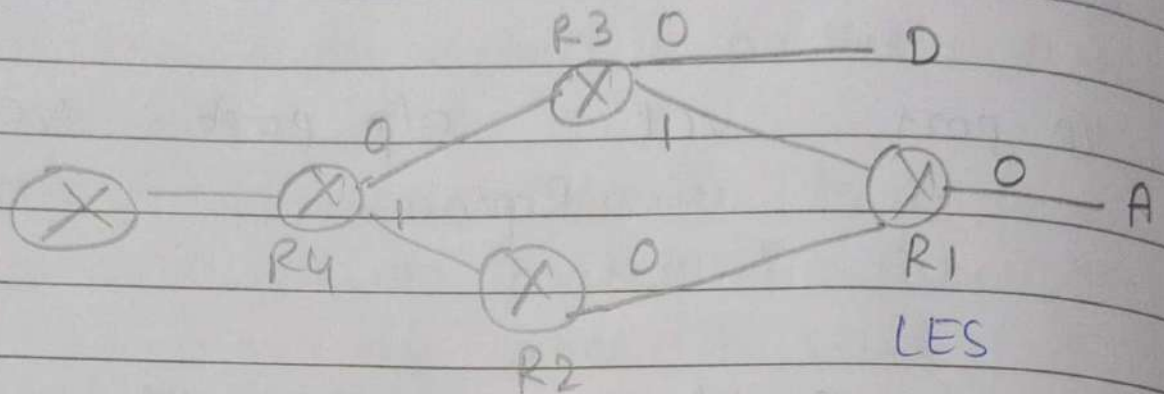
\rightarrow Private LAN service



- \rightarrow LER - label edge router (at border)
- \rightarrow LSR - label switch router (within network)
- \rightarrow It depends on IP MPLS IP label & not IP address IP label according to dest IP

LER, LSR has MPLS table.

All router have MPLS table.



Router R1

l/p label	o/p label	o/p interface	dest
6	-	0	A

This table is broadcasted to R2 & R3

R3 table

12	-	0	D
18	6	1	A

R2 table

15	6	0	A
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R2 & R3 broadcast to R4

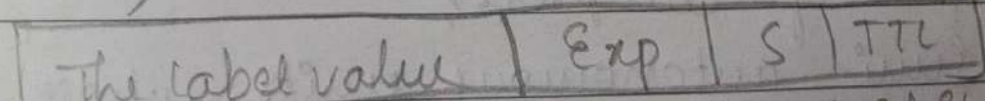
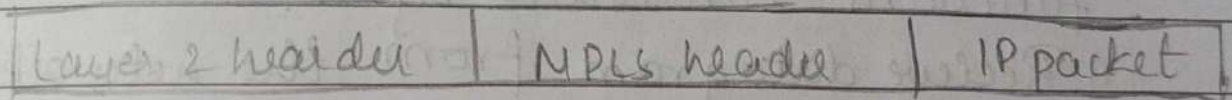
R4 is incoming edge router
So no i/p label.

R4's table

i/p L	o/p L	o/p i	dest
12	0	0	D
18	0	0	A
15	1	0	A R2

So quick forwarding

MPLS header is added b/w L2 & L3



← 20 bits → ← 3 bits → ← 8 bits →

4 byte long end of stack

28 Exp → experimental bit