

## Hashing, Hash table, Hashing Techniques

### Why Hashing:

- ▶ Internet has grown to millions of users and terabytes of data every day.
- ▶ It is impossible to find anything in the internet, unless we develop a new data structure to store and access the data very quickly.
- ▶ The amount of time required to look up an element in an array or linked list is either  $O(\log n)$  or  $O(n)$  based on the list is sorted or not.
- ▶ New data structure called Hashing used to store and retrieve any entry with constant time  $O(1)$ .
- ▶ This technique is irrelevant to size of the list and order.
- ▶ To increase the search efficiency the items to be stored in such a way as to make it easy to find them later.

### Hash Table

- ▶ Hash table is a data structure to store and retrieve data very fast.
- ▶ Hash table consist of key and its value.
- ▶ Each location in the hash table is called cell or bucket.
- ▶ Hash table is implemented using array.
- ▶ An element is accessed very fast if we know the key or its index.

**Example:** To store the Student record in hash table, Student rollno is used as a key

Following are the basic primary operations of a hash table.

Search – Searches an element in a hash table.

Insert – inserts an element in a hash table.

Delete – Deletes an element from a hash table

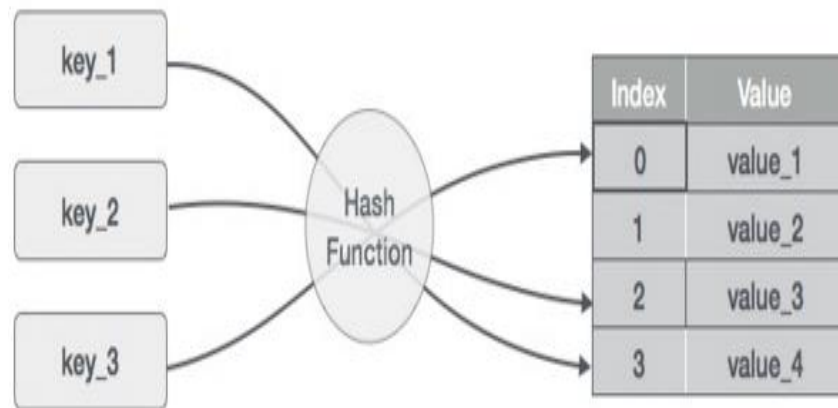
### Hashing

- ▶ Hashing is a technique used to generate key where an element is to be inserted or to be located from.

**Note:** It minimizes the number of comparisons while performing the search.

### Hash function:

- To map the key value into its corresponding index in hash table hash function is used.
- A hash function  $h$  transforms a key into an index in a hash table  $T[0 \dots m-1]$   
Where  $m$  is size of hash table.



- Use hash function to compute the index in the hash table for the given key value.
- Hash function returns integer value which give the index value in the hash table.

### Characteristics of Hash Function

The characteristics / properties of a good hash function are

- It is efficiently computable.
- It minimizes the number of collisions.
- It distributes the keys uniformly over the table.

### Types of Hash Functions :

There are various types of hash functions available such as

1. Division method
2. Mid square method

**1 Division Method:** The hash function depends upon the remainder of division. Typically the divisor is table length (M).

$$h(\text{key}) = \text{record} \% M$$

For eg; If the record 54, 72, 89, 37 is placed in the hash table and if the table size is 10 then

$$54\%10 = 4$$

$$72\%10 = 2$$

$$89\%10 = 9$$

$$37\%10 = 7$$

0	
1	
2	72
3	
4	54
5	
6	
7	37
8	
9	89

Consider the following elements to be placed in the hash table of size 10,

37,90,45,22,17,49,55

$$H_1(37) = 37\%10 = 7$$

$$H_1(90) = 90\%10 = 0$$

$$H_1(45) = 45\%10 = 5$$

$$H_1(22) = 22\%10 = 2$$

$$H_1(17) = 17\%10 = 7$$

$$H_1(49) = 49\%10 = 9$$

$$H_1(55) = 55\%10 = 5$$

0	90
1	
2	22
3	
4	
5	45
6	
7	37
8	
9	49

In above example 17 and 55 are hashed to same location this condition is called **collision**.

**Note:** 1) If the hash function returns same hash key for more than one record, then this situation is called 'collision' and need to resolve it.

## 2 Mid Square:

In the mid square method, the key is squared and the middle or mid part of the result is used as the index. If the key is a string, it has to be preprocessed to produce a number. Consider that if we want to place a record 3111 then

$$3111^2 = 9678321$$

for the hash table of size 1000

$$H(3111) = 783 \text{ (the middle 3 digits)}$$

## Collision resolution techniques

There are several collision resolution techniques

### 1) Separate chaining: an array of linked list representation

All keys that map to the same hash value are kept in a list (or “bucket”).

- ▶ Hash table is implemented as array of linked list.
- ▶ All the records which are mapped to same hash address are lined together to form a linked list.

**Example:** Load the keys **23, 13, 21, 14, 7, 8, and 15** , in this order, in a hash table of size **7** using separate chaining with the hash

function:  **$h(\text{key}) = \text{key} \% 7$**

$$h(23) = 23 \% 7 = 2$$

$$h(13) = 13 \% 7 = 6$$

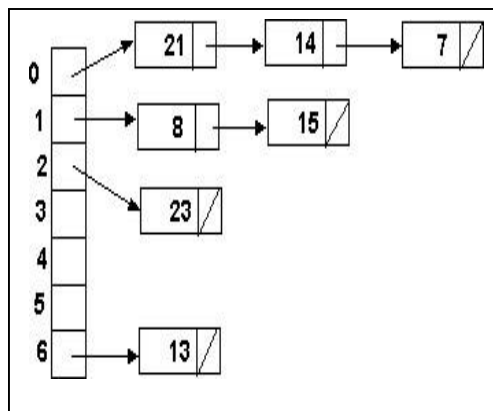
$$h(21) = 21 \% 7 = 0$$

$$h(14) = 14 \% 7 = 0 \quad \text{collision}$$

$$h(7) = 7 \% 7 = 0 \quad \text{collision}$$

$$h(8) = 8 \% 7 = 1$$

$$h(15) = 15 \% 7 = 1 \quad \text{collision}$$



### 2) Open addressing: array based implementation

- Linear probing (linear search)
- Quadratic probing (nonlinear search)
- Double hashing (uses two hash functions)