

3.5 Quadratic Programming

Since the work in Frank and Wolfe (1956), Quadratic Programming aims to effectively solve for the numerical solution of a board class of quadratic optimization problems. Conventionally, the corresponding convex optimization problem, usually coined as a Quadratic Program (QP), is formulated as follows:

$$\underset{\mathbf{x} \in \mathbb{R}^D}{\text{minimize}} \quad \frac{1}{2} \mathbf{x}^\top \mathbf{P} \mathbf{x} + \mathbf{q}^\top \mathbf{x} + r, \quad (3.5.1)$$

$$\text{subject to} \quad \text{i) } \mathbf{G} \mathbf{x} \leq \mathbf{h}; \quad (3.5.2)$$

$$\text{ii) } \mathbf{A} \mathbf{x} = \mathbf{b};$$

where $\mathbf{P} \in \mathbb{R}^{D \times D}$ is positive definite, $\mathbf{G} \in \mathbb{R}^{K \times D}$, and $\mathbf{A} \in \mathbb{R}^{M \times D}$, and the symbol \leq here means that every entry of vector $\mathbf{G} \mathbf{x}$ is less than or equal to the corresponding entry of the right-hand side vector $\mathbf{h} \in \mathbb{R}^K$. Moreover, $\mathbf{P} := (p_{ij})$ is usually assumed to be symmetric, since

$$\mathbf{x}^\top \mathbf{P} \mathbf{x} = (\mathbf{x}^\top \mathbf{P} \mathbf{x})^\top = \mathbf{x}^\top \mathbf{P}^\top \mathbf{x},$$

is a scalar. Hence, we could always replace the asymmetric matrix \mathbf{P} with $\tilde{\mathbf{P}} := (\mathbf{P} + \mathbf{P}^\top)/2$ as:

$$\mathbf{x}^\top \tilde{\mathbf{P}} \mathbf{x} := \frac{1}{2} \mathbf{x}^\top (\mathbf{P} + \mathbf{P}^\top) \mathbf{x} = \frac{1}{2} (\mathbf{x}^\top \mathbf{P} \mathbf{x} + \mathbf{x}^\top \mathbf{P}^\top \mathbf{x}) = \mathbf{x}^\top \mathbf{P} \mathbf{x}.$$

Particularly in a quadratic program, we minimize a convex quadratic function on a high dimensional polyhedron, *i.e.* the feasible set \mathcal{P} as specified by the constraints (3.5.2), as depicted in Figure 3.5.1.

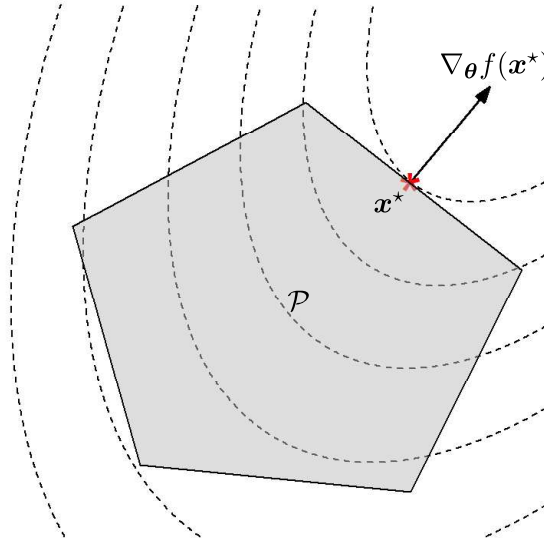


Figure 3.5.1: Geometric illustration of Quadratic Program

3.5.1 Inequality Constrained Lagrangian Multipliers

We first state the celebrated **von Neumann's Minimax Theorem** as below:

Theorem 3.5.1. (von Neumann's Minimax Theorem) [4]: Let $\mathcal{X} \subset \mathbb{R}^N$ and $\mathcal{Y} \subset \mathbb{R}^M$ be compact convex sets: If $f : \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$ is a continuous function that is concave-convex, *i.e.*

1. for each $y \in \mathcal{Y}$, $f(\cdot, y) : \mathcal{X} \rightarrow \mathbb{R}$ is concave in $x \in \mathcal{X}$; and
2. for each $x \in \mathcal{X}$, $f(x, \cdot) : \mathcal{Y} \rightarrow \mathbb{R}$ is convex in $y \in \mathcal{Y}$.

Then, we have

$$\max_{x \in \mathcal{X}} \min_{y \in \mathcal{Y}} f(x, y) = \min_{y \in \mathcal{Y}} \max_{x \in \mathcal{X}} f(x, y). \quad (3.5.3)$$

Remark 3.5.1. Certainly, the existence of a saddle point implies the validity of (3.5.3), as a saddle point (x_0, y_0) of f is such that

$$\min_{y \in \mathcal{Y}} \max_{x \in \mathcal{X}} f(x, y) = \min_{y \in \mathcal{Y}} f(x_0, y) = f(x_0, y_0) = \max_{x \in \mathcal{X}} f(x, y_0) = \max_{x \in \mathcal{X}} \min_{y \in \mathcal{Y}} f(x, y), \quad (3.5.4)$$

where the leftmost expression follows the red path in Figure 3.5.2; the rightmost term follows the blue path in the same figure. They both arrive at the same point (x_0, y_0) . Nevertheless, in general, even if the saddle point of f does not exist, the Theorem remains valid.

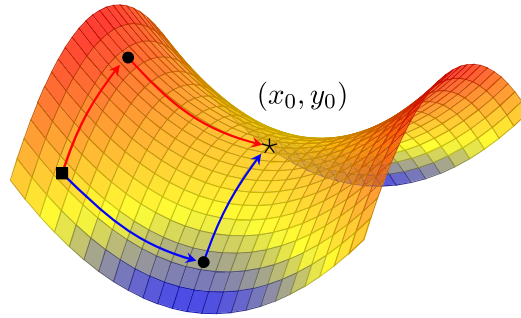


Figure 3.5.2: The two paths of the von Neumann's Minimax Theorem over a function f .

With this theorem in mind, we then consider a more general minimization problem:

$$\text{minimize}_{\mathbf{x} \in \mathbb{R}^D} \quad f(\mathbf{x}), \quad (3.5.5)$$

$$\text{subject to} \quad g_k(\mathbf{x}) \leq 0, \quad k = 1, \dots, K; \quad (3.5.6)$$

where $f, g_k : \mathbb{R}^D \rightarrow \mathbb{R}$, for $k = 1, \dots, K$, and we denote $\mathbf{g}(\mathbf{x}) := (g_1(\mathbf{x}), \dots, g_K(\mathbf{x}))^\top$.

By Lagrangian multiplier approach, the lagrangian with a multiplier $\boldsymbol{\lambda}$ is given by:

$$\mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}) := f(\mathbf{x}) + \boldsymbol{\lambda}^\top \mathbf{g}(\mathbf{x}). \quad (3.5.7)$$

Define $\mathcal{X} := \{\mathbf{x} \in \mathbb{R}^D : \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\}$. $\mathbf{g}(\mathbf{x}) \leq \mathbf{0}$, for every feasible $\mathbf{x} \in \mathcal{X}$, the lagrangian $\mathcal{L}(\mathbf{x}, \boldsymbol{\lambda})$ is a decreasing function in λ_k 's. Almost by definition, it is easy to see the equality that:

$$\min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x}) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\boldsymbol{\lambda} \geq \mathbf{0}} \left(f(\mathbf{x}) + \boldsymbol{\lambda}^\top \mathbf{g}(\mathbf{x}) \right) = \min_{\mathbf{x} \in \mathcal{X}} \max_{\boldsymbol{\lambda} \geq \mathbf{0}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}), \quad (3.5.8)$$

which is usually called the **primal problem**. We first note that $\boldsymbol{\lambda}$ lies in the unbounded domain, hence not a compact convex set, we therefore cannot directly apply Theorem 3.5.1 to (3.5.8) immediately. Under some mild regularity condition on f , we can regard $\boldsymbol{\lambda}_{\mathbf{x}} : \mathcal{X} \rightarrow [0, \infty)^K$ as a continuous function in $\mathbf{x} \in \mathcal{X}$ such that

the image $\lambda_x(\mathcal{X})$ can then be a closed and bounded, and so compact, set, let say now $\sup_{x \in \mathcal{X}} |\lambda_x| \leq M < \infty$, with this in mind, we can therefore apply Theorem 3.5.1 such that:

$$\min_{x \in \mathcal{X}} f(x) = \max_{0 \leq \lambda \leq M} \min_{x \in \mathcal{X}} \mathcal{L}(x, \lambda).$$

Furthermore, if one can find a function $\ell(\lambda)$ such that $\min_{x \in \mathcal{X}} \mathcal{L}(x, \lambda) \geq \ell(\lambda)$ and $\ell(\lambda)$ is coercive in λ , i.e. $\ell(\lambda) \uparrow \infty$ as $\lambda \uparrow \infty$, then we finally deduce that

$$\min_{x \in \mathcal{X}} f(x) = \max_{\lambda \geq 0} \min_{x \in \mathcal{X}} \mathcal{L}(x, \lambda).$$

Remark 3.5.2. However, if the regularity conditions are not met or in the absence of suitable concavity and convexity, we still have the **max–min inequality** [5]:

$$\max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} f(x, y) \leq \min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} f(x, y), \quad (3.5.9)$$

where $f : \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$; indeed, to see this, for all $x \in \mathcal{X}$ and $y \in \mathcal{Y}$, we clearly have $\min_{x' \in \mathcal{X}} f(x', y) \leq f(x, y)$.

Taking maximum with respect to y on both sides yields:

$$\max_{y \in \mathcal{Y}} \min_{x' \in \mathcal{X}} f(x', y) \leq \max_{y \in \mathcal{Y}} f(x, y), \quad \forall x \in \mathcal{X}.$$

Since the inequality above holds for any $x \in \mathcal{X}$, (3.5.8) can be obtained by considering only the minimum of the right-hand side. Finally, with the **max–min inequality** in (3.5.9), we can re-write the primal problem of (3.5.8) into a **dual problem**:

$$\min_{x \in \mathcal{X}} \max_{\lambda \geq 0} \mathcal{L}(x, \lambda) \geq \max_{\lambda \geq 0} \min_{x \in \mathcal{X}} \mathcal{L}(x, \lambda) =: \max_{\lambda \geq 0} \mathcal{D}(\lambda). \quad (3.5.10)$$

This result warrants the duality nature of primal and dual problems, such that the solution to the dual (maximization) problem provides a lower bound to the solution of the primal (minimization) problem, and the difference between the optimal values of primal and dual is called the duality gap; see the illustration in Figure 3.5.3, which can be referred to the Fenchel-Rockafellar duality (see Theorem 1.9 in Villani (2021)) of convex optimization problems, and this duality gap can be removed under a convex constraints as specified in the Minimax Theorem 3.5.1.

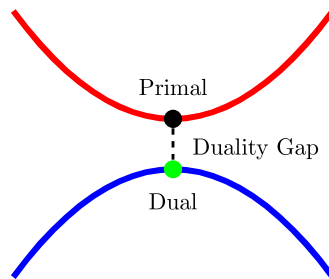


Figure 3.5.3: Duality Gap

3.5.2 Lagrange Multiplier in Quadratic Programming

Next, we shall reformulate the problem in (3.5.1) with constraint in (3.5.2) by using the Lagrangian multiplier approach, and the Lagrangian \mathcal{L} with $\lambda_1 \in \mathbb{R}^K, \lambda_2 \in \mathbb{R}^M$ is given by:

$$\begin{aligned} \mathcal{L}(x, \lambda_1, \lambda_2) &= \frac{1}{2} x^\top P x + q^\top x + r + \lambda_1^\top (Gx - h) + \lambda_2^\top (Ax - b) \\ &= \frac{1}{2} x^\top P x + (c + G^\top \lambda_1 + A^\top \lambda_2)^\top x - \lambda_1^\top h - \lambda_2^\top b + r, \end{aligned} \quad (3.5.11)$$