

Stack'simplicit

→ In memory we do this in recursion

Memory / computer

Explicit

Real Life Examples

1) Books →



2) Coins →



3) Stones →



Operations

1) PUSH O(1)

→ add

2) POP O(1)

→ delete

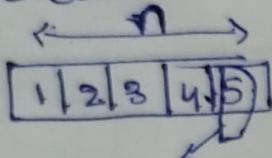
3) Peek O(1)

Note: Last in first out (LIFO)

Implementation

Arrays

Size → fixed

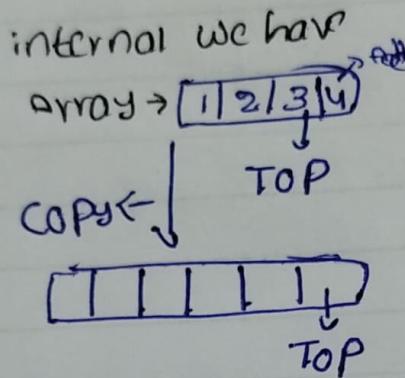


last index
(TOP)

Stack == full

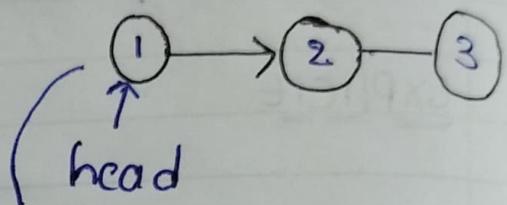
ArrayList

variable → size

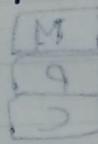


LinkedList

variable → size



$O(1) \rightarrow$ change → add
pop

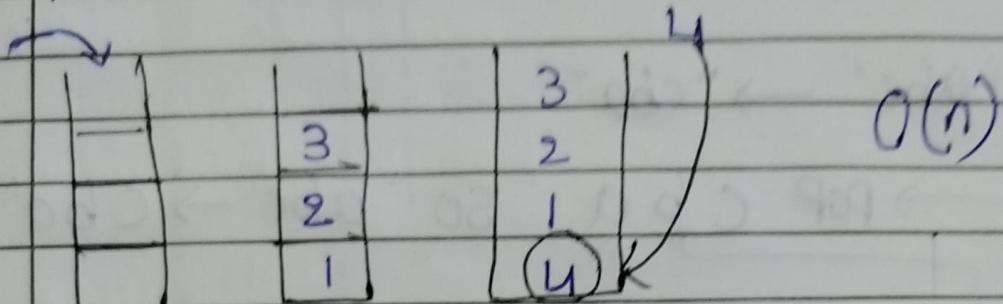
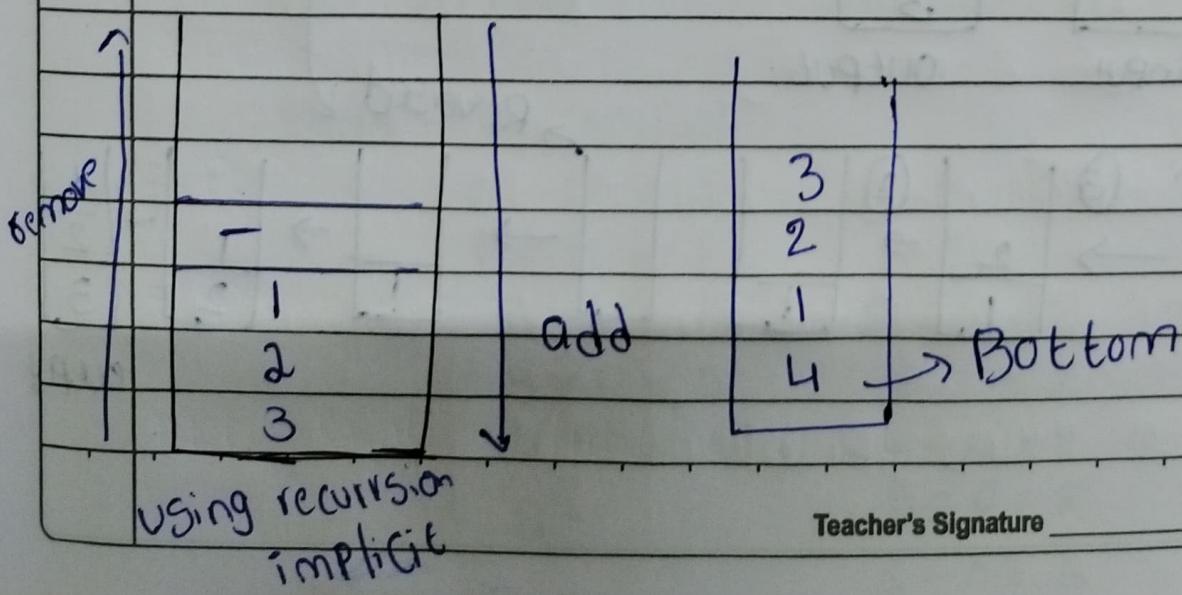
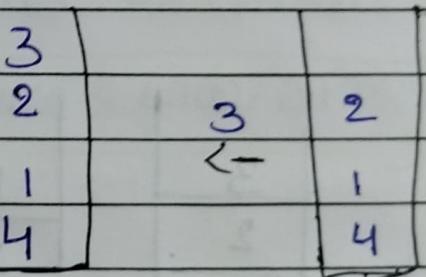
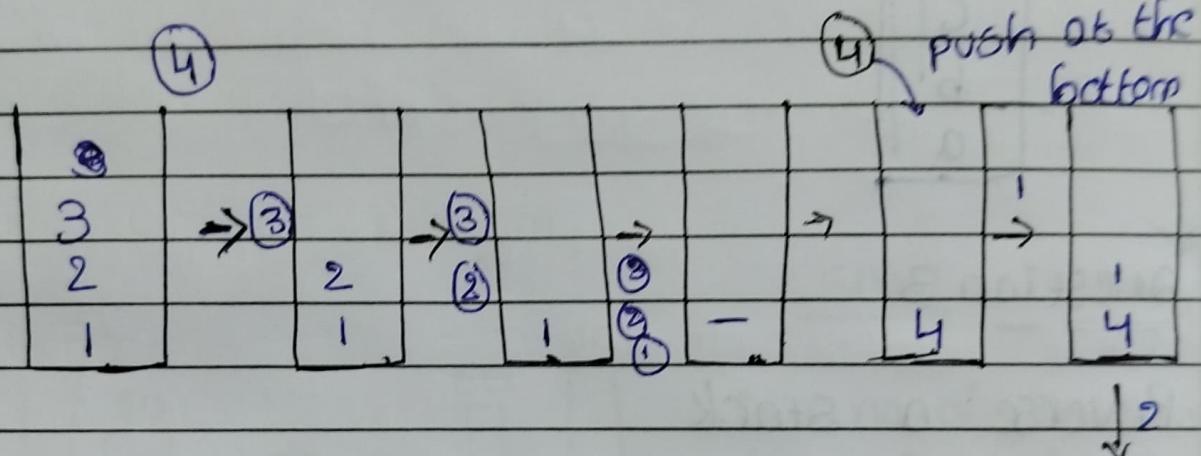


Stack - using ArrayList

Expt. No.:

Question 1

push at the bottom of the stack

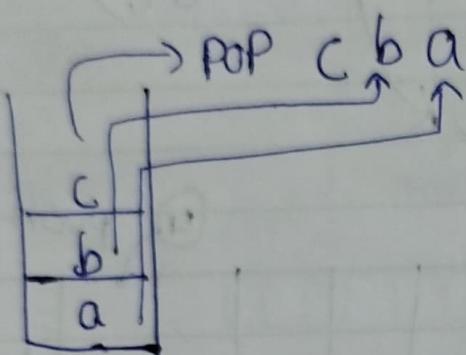
 $O(n)$ Using recursion
implicit

Teacher's Signature _____

Question 2

Reverse a string using a stack

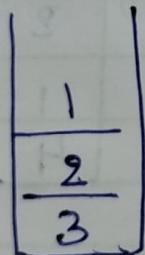
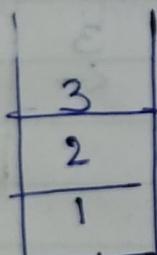
$$\text{"abc"} \rightarrow \text{"cba"}$$



$$\text{so } \text{"abc"} \rightarrow \text{"cba"}$$

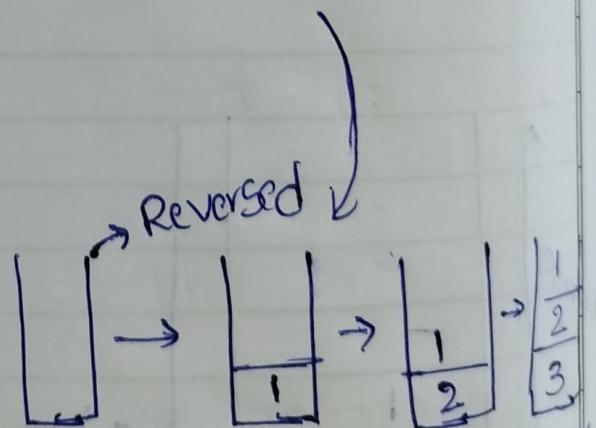
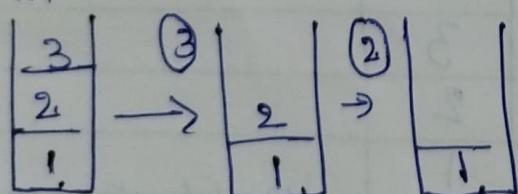
Question 3

Reverse a stack



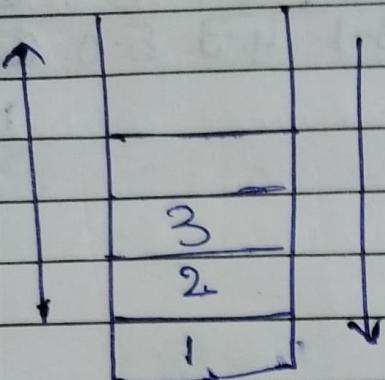
Push(at top)

in



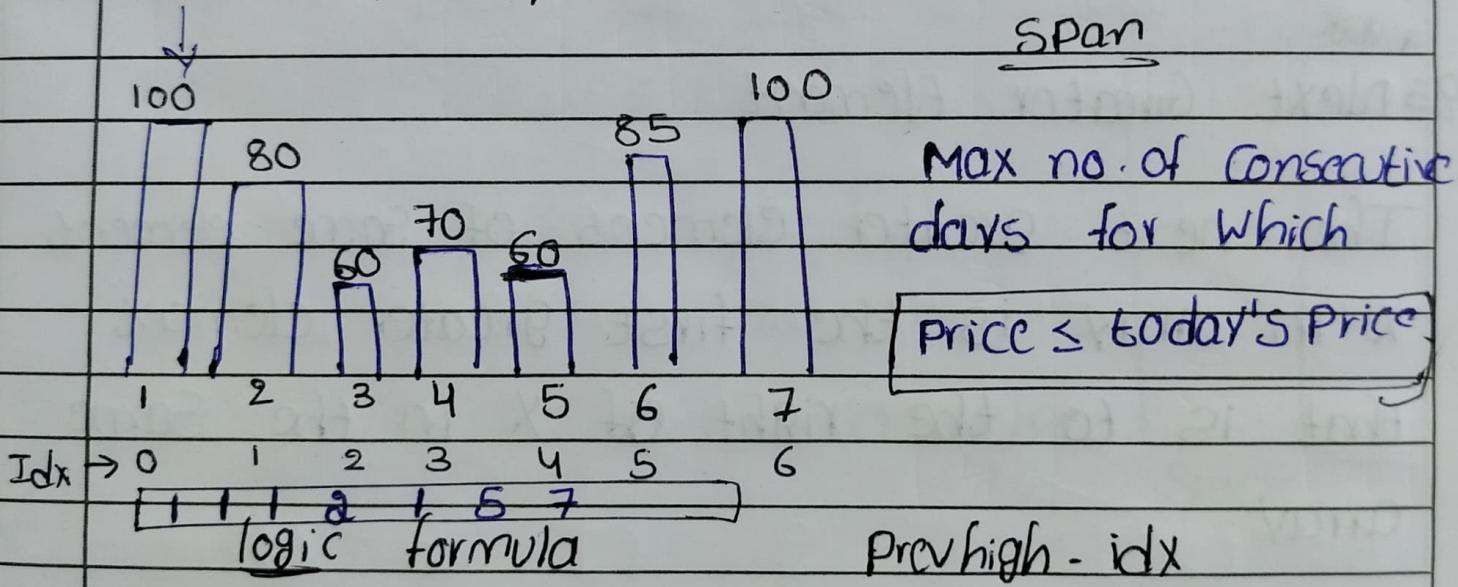
Expt No.:

Recursion



QUESTION 4

Stock Span problem



$$\text{Span} = i - \text{prevhigh}$$

$$0 - 0 = 0 \quad 0 - 0 = 0$$

~~$$1 - 0 = 1 \quad 1 - 0 = 1$$~~

~~$$2 - 1 = 1 \quad 2 - 1 = 1$$~~

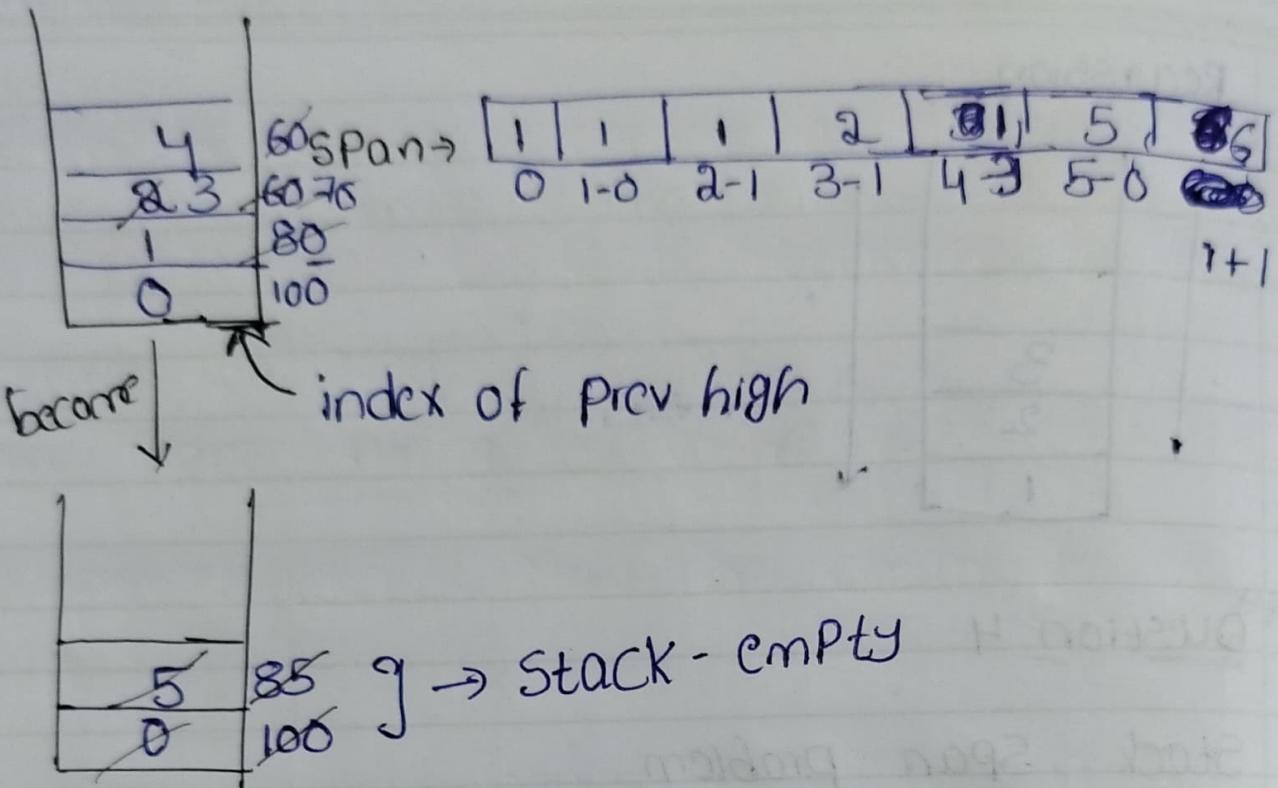
~~$$3 - 1 = 2 \quad 3 - 1 = 2$$~~

~~$$4 - 1 = 3 \quad 4 - 3 = 1$$~~

~~$$5 - 0 = 5$$~~

Teacher's Signature

$$6 - 0 = 6$$



25 Next Greater Element

The next greater element of some element x in array is the first greater element that is to the right of x in the same array.

$$\text{arr} = [6, 8, 0, 1, 3]$$

$$\text{next Greater} = [8, -1, 1, 3, -1]$$

Logic is very important we can use this in so many problems.

Brute force

Nested loop

for (int i=0; i<arr.length; i++) { //outer loop
 for (int j=i+1; j<arr.length; j++)) { //inner loop

if (@arr[i] < arr[j]) {

1st time

next Greater El

T.C $\rightarrow O(n^2)$

Optimize Time complexity using Stack

Approach arr = [6, 8, 0, 1, 3]

nextGreater[] = [] 0 1 2 3 4

* ① while (!S.isEmpty()) && stack[top] < arr[i])
 IMP S.pop()

Stack

② if stack.isEmpty()
 nextGreater = -1

Note! Code K:
 three index P

else

nextGreater = S.peek()

③ S.push(c0)

Valid Parentheses

Given a string s containing just the characters '`(`', '`)`', '`{`', '`}`', '`[`' & '`]`', determine if the input string is valid.

An input string is valid if:

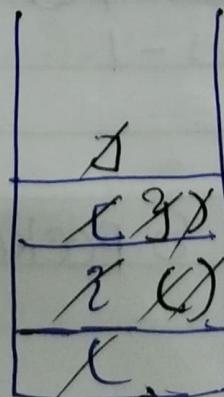
1. open brackets must be closed by the same type of brackets.
2. open brackets must be closed in the correct order.
3. Every close bracket has a corresponding open bracket of the same type.

Ex:-

$s = "()"$, $s = "([)]"$, $s = "({})"$, $s = "((()))"$

Approach

$(\{[)]\}()$)



→ EMPTY!

1) opening bracket.
S.push()

2) closing bracket.

↳ stack → top

Pair ✓ True

Not Found False

Stack.isEmpty() ✓ True

Code

```
for(int i=0; i<str.length(); i++) {  
    ith → ch
```

Opening → S.push

Closing

↳ ~~if~~ Pair ↔ S.pop()
S.pop()
↳ return false (INVALID)

3

S → Empty() → True

X → false

Duplicate Parentheses

Given a balanced expression, find if it contains duplicate parentheses or not. A set of parentheses are duplicate if the same subexpression is surrounded by multiple parentheses
Return a true if it contains duplicates else return false.

Ex: $((a + (b)) + (c + d)) \rightarrow \text{true}$

Ex: $((((a) + (b)) + c + d)) \rightarrow \text{true}$

Ex: $((a+b) + (c+d)) \rightarrow \text{false}$

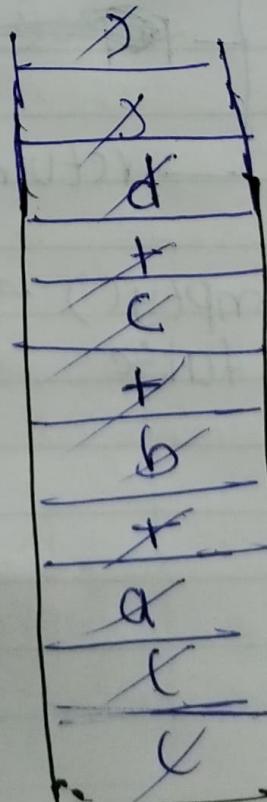
Ex: $((a+b)) + c) \rightarrow \text{true}$

Approach

Ex: $((a+b) + (c+d))$

false

Count = $\emptyset \times 2^3$



opening ✓
closing ✓

stack

Expt. No.:

① opening operator operand (a, b, c)

S.push()

while(S.pop() == '(') {

S.pop

count++

y

count < 1 →

② closing

Count = 0 || items

try to find opening pair

(count < 1) → Duplicate → True

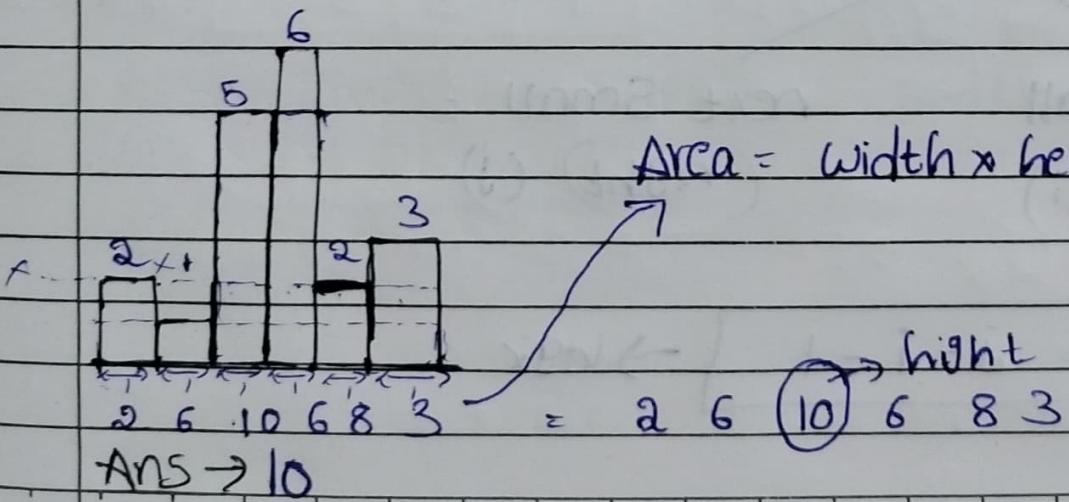
3 < 1 → False Duplicate not exist

S.pop() ← '('

IMP 40-50 AdvancedMax Area in Histogram

Given an array of integers heights representing the histogram's bar height. Where the width of each bar is 1, return the area of the largest rectangle in the histogram

height = [2, 1, 5, 6, 2, 3]



Approach

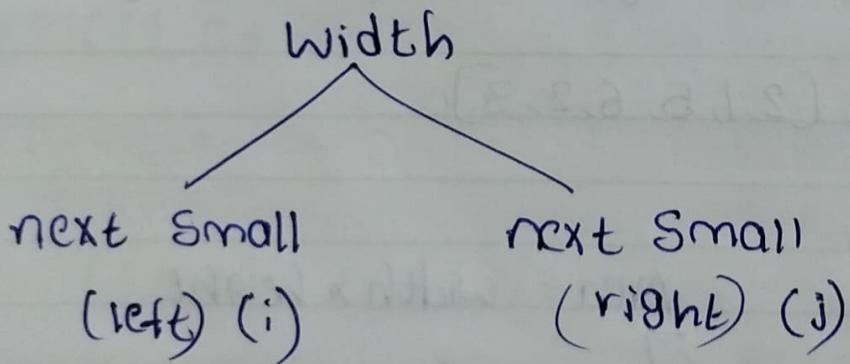
$$\text{Area} = \text{height} \times \text{width}$$
$$\downarrow \qquad \qquad \downarrow$$
$$\text{height}(i) \qquad \text{left, right}$$

Note:

For example $h = [2, 1, 5, 6, 2, 3]$

We take ~~the~~ 5 → In this left side small (remove eliminate) right side (remove the small).

index (area = height × width)



$$\boxed{\text{Width} = j - i - 1} \rightarrow \text{logic}$$

Expt. No.:

\rightarrow For ex: $h[\text{index}] = 6 - i^{\text{th}}$

$\rightarrow \text{0 } \text{0 } \text{0 } \text{0 } \text{0 } \text{4} \quad 4-3 = 1 \rightarrow \text{width}$

$$\begin{aligned}\text{For ex: } h[i] &= j-i-1 \\ &= 6 - (-1) - 1 \\ &= 6 + 1 - 1 \\ &= 6\end{aligned}$$

 $h[5]$

$$4-1-1 = 4-2 = 2 \quad h-[2, 1, 5, 6, 2, 3]$$

 $(-1 \ 1 \ 2 \ 1 \ 4) \rightarrow \text{index}$ $\text{nsL} \rightarrow [-1 \ 1 \ 1 \ 5 \ 1 \ 2] \quad \text{nsL} \rightarrow \text{next smaller left}$ $\text{nsr} \rightarrow [1 \ 6 \ 2 \ 2 \ 6 \ 6] \quad \text{nsr} \rightarrow \text{next smaller right}$ $(1 \ n \ 4 \ 4 \ n \ n) \rightarrow \text{index}$

Note: nsr \rightarrow Where a next smaller element is $\textcircled{1}$ next smaller element not their but you didn't take -1 instead of $6 \rightarrow n$

nsL \rightarrow start with "3" next smaller element is $\textcircled{2}$ next smaller $\textcircled{3}$ next smaller element is $\textcircled{5}$

$$\begin{aligned}\text{Area} &= \text{height} \times \text{width} \quad \text{width} = (j-i-1) \\ &= 2 \times (1+1-1) \\ &= 2 \times 1 = 2\end{aligned}$$