Chapter 2 - Types and Functions

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1 Haskell's type system

Types provide *abstraction* - we can know something is a string without having to understand the underlying implementation of strings.

Haskell's type system is:

- Strong here meaning that you can only write expressions that obey Haskell's typing rules, and that values are not automatically cast to their correct type if they are incorrect.
- Static The compiler knows the types in the program at compile time, so if there are any errors, the code won't run.
- using Type inference declaring the types of expressions is optional, theu can also be deduced by the compiler.

2 Other things to note about the type system

Function application has higher precedence than using operators, so compare 2 3 == LT is the same as (compare 2 3) == LT

Function application is left associative, so a b c d is the same as ((a b) c) d. Type signatures are right associative, so $Int \rightarrow [a] \rightarrow [a]$ is the same as $Int \rightarrow ([a] \rightarrow [a])$.

To specify polymorphic types, we use a type variable, which must begin with a lowercase letter. For example, we can write functions on all lists as $[a] \to a$ or $[a] \to a$, etc. (because list is a polymorphic type).

Parametric polymorphism is where the type variable is substituted with the actual type on evaluation. A parameterised type in Haskell is similar to a type variable in Java generics. C++ templates also bear a resemblance to parametric polymorphism.

The empty tuple is the unit type, (). It has one value, () and is a bit like void in other languages.

The record that we use to track an unevaluated expression is referred to as a thunk.

3 Examples of types

```
ex1_1 :: Bool

ex1_1 = False

ex1_2 :: ([String], Char)

ex1_2 = (["foo", "bar"], 'a')

ex1_3 :: [(Bool, [[Char]])]

ex1_3 = [(True, []), (False, [['a']])]
```

4 Functions used in this chapter

4.1 take

take is a function in the Prelude. take n, applied to a list xs, returns the prefix of xs of length n, or xs itself if n ξ length xs:

```
 \begin{array}{l} take' :: Int \rightarrow [\,a\,] \rightarrow [\,a\,] \\ take' \ 0 \ \_ = [\,] \\ take' \ n \ [\,] \ = [\,] \\ take' \ n \ (x : xs) \\ \mid \ n < 0 = [\,] \\ \mid \ n \leqslant length \ (x : xs) = x : take' \ (n-1) \ xs \\ \mid \ otherwise = [\,] \end{array}
```

4.2 drop

drop is a function in the Prelude. drop n xs returns the suffix of xs after the first n elements, or [] if n $\tilde{\iota}$ length xs:

```
\begin{array}{l} \operatorname{drop'}::\operatorname{Int}\to [\,a\,]\to [\,a\,]\\ \operatorname{drop'}0\,\operatorname{xs}=\operatorname{xs} \end{array}
```

```
drop' \ n \ [] = []
drop' \ n \ xs
| \ n < 0 = xs
| \ n \leqslant length \ xs = drop' \ (n-1) \ (tail \ xs)
| \ otherwise = []
```

4.3 (

fst and snd) fst takes the first element of a pair

$$fst' :: (a, b) \rightarrow a$$

 $fst' (a, _) = a$

snd takes the second element of the pair

$$snd' :: (a, b) \rightarrow b$$

 $snd' (_, b) = b$

4.4 null

null indictes if a list is empty:

$$null' :: [a] \rightarrow Bool$$

 $null' [] = True$
 $null' _ = False$

4.5 last

last will extract the last element of a list, which must be finite and non-empty:

```
\begin{array}{l} last' :: [a] \rightarrow a \\ last' \ [] = error \text{ "empty list"} \\ last' \ [x] = x \\ last' \ (x:xs) = last' \ xs \\ \\ lastButOne :: [a] \rightarrow a \\ lastButOne \ (x:y:[]) = x \\ lastButOne \ (x:xs) = lastButOne \ xs \\ lastButOne \ \_ = error \text{ "Not enough elements"} \end{array}
```

if the list is empty, it throws an exception.