Semaphores

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- Semaphores
- Producer/Consumer
- Implementing Semaphores
- Bounded Semaphores
- Using Java Semaphores to Implement a Bounded Buffer

Introduction

1: Semaphores

1.1: Disadvantages of synchronisation by communication

- complex algorithms
- time wasting busy waits
- primitive (atomic) actions are too weak

We need an atomic read/write action, and a mechanism to allow processes to "sleep" — semaphores.

Semaphores

1.2: Semaphores

Introduced by Edsgar Dÿkstra.

- P() Passeren (Pass)
 [A better name might be prolagen, from proberen (try) and verlagen (reduce)].
- V() Vrÿgeven (Free)

English mnemonic — Poll and Vote

A process that calls P() when s=0 can only complete the call when $s\neq 0$, after a V() operation — they must "sleep".

The order in which processes are woken is unspecified, but it must be *fair*.

Binary semaphores

1.3: Binary and split binary semaphores

Binary semaphore:

Can only take values ${f 0}$ or ${f 1}$

$$0 \leq \text{s.value}() \leq 1$$

Split binary semaphore:

Total value is 0 or 1

$$0 \le s_1.value() + s_2.value() \le 1$$

Properties can be *programmer's* responsibility — but sometimes only binary (split) semaphores provided.

Producer/consumer problem

2: Producer/Consumer

2.1: Introduction

Models many "buffering" problems

- Producer produces data, snacks, sproglets, . . .
- Consumer consumes data, snacks, sproglets, . . .
- Buffer
 - synchronises exchange
 - "smooths out" speed differences

Producer/consumer problem

- Initial assumptions
 - infinite buffer
 - simultaneous access for producer/consumer
- Later
 - access to buffer in critical section
 - finite buffer bounded buffer problem

Producer and Consumer

2.2: The producer and consumer

```
producer: buffer.put(item)
public void producer() {
    while (true) {
        // build item to add
        buffer.put(item);
    }
}
consumer: buffer.get()
public void consumer() {
    while (true) {
        T item = buffer.get();
        // do something with item
    }
}
```

Basic algorithm

2.3: Basic algorithm

infinite buffer, simultaneous access

Here noOfElements.value() = no. of elements in buffer.

Producer/consumer problem

2.4: Forbidding simultaneous access

putItem and getItem in critical section

Initialise noOfElements = 0, criticalSection = 1.

Bounded buffer

2.5: Bounded buffer

```
producer: buffer.put(item)
                             consumer: buffer.get()
                            public void get() {
 public void put() {
                               noOfElements.P():
    noOfSpaces.P();
                                criticalSection.P();
    criticalSection.P();
                               T item = getItem();
    putItem(item);
                                criticalSection.V();
    criticalSection.V():
                               noOfSpaces.V();
    noOfElements.V();
                               return item;
```

3: Implementing Semaphores 3.1: class Semaphore: Fields

```
// Value of the semaphore
private int value = 0;
```

3.2: class Semaphore: Constructors

```
// Initialise value to default value of zero
protected Semaphore() {
   value = 0;
}

// Initialise to the specified value
protected Semaphore(int initial) {
   value = initial;
}
```

```
3.3: class Semaphore: Methods
// Poll method
public synchronized void poll()
  throws InterruptedException {
   value--;
   if (value < 0) {
      wait()
// Vote method
public synchronized void vote() {
   value++;
   if (value <= 0) {
      notify();
```

Synchronized Methods

An Aside: synchronized Methods

A **synchronized** method is one that may only be executed by at most one process at any one time.

The wait() and notify() methods can be used to further refine processes' behaviour in synchronized methods:

- A process calling wait() will go to sleep
- The sleeping process will sleep until another process calls notify().

4: Bounded Semaphores

Three possibilities:

4.1: Absorbing Semaphores

Once the semaphore reaches its limit it stops incrementing its value.

4.2: Crashing Semaphores

Once the semaphore reaches its limit any attempt to increase its value by a V() causes it to crash.

4.3: Blocking Semaphores

Once the semaphore reaches its limit any attempt to increase its value by a V() causes it to block.

```
Need to import java.util.concurrent.Semaphore;
5.1: class Buffer(T): Fields

private Semaphore noOfSpaces, noOfElements, criticalSection
private T[] buffer;
private int putIndex = 0, getIndex = 0;
```

5: Using Java Semaphores to Implement a Bounded Buffer

```
5.2: class Buffer(T): Constructor

public Buffer(int size) {
   buffer = new T[size]; // can't actually do this in Java
   noOfSpaces = new Semaphore(size,true); // is fair
   noOfElements = new Semaphore(0);
   criticalSection = new Semaphore(1);
}
```

5.3: class Buffer $\langle T \rangle$: The put method public void put(T item) { try { noOfSpaces.acquire(); criticalSection.acquire(); addItem(item); criticalSection.release(); noOfElements.release(); catch (InterruptedException ie) { throw new BufferError("Data item not added\n" + ie.getMessage());

5.4: class Buffer $\langle T \rangle$: The get method public T get() { try { noOfElements.acquire(); criticalSection.acquire(); T item = getItem(); criticalSection.release(); noOfSpaces.release(); return item; catch (InterruptedException ie) { throw new BufferError("Data item not fetched\n" +

ie.getMessage());

5.5: class Buffer(T): Update methods

```
private void addItem(T item) {
   buffer[putIndex] = item;
  putIndex = (putIndex + 1) % buffer.length;
private T getItem() {
   T item = buffer[getIndex];
   getIndex = (getIndex + 1) % buffer.length;
  return item;
```

End of semaphores lecture