Programming, Algorithms and Data Structures (210CT)

COURSEWORK REPORT

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INTRODUCTION:

The report is a record of all programming code written in accordance to the tasks objectives, alongside the comments for further description and clarification of functionalities as well as side notes entailing encountered drawbacks and applicable (or theoretical) solutions to them.

To view the live code from all tasks, see appendix.

CONTENTS:

1. Defining the shuffle function:

```
01  from random import *
02  input1 = [5,3,8,6,1,9,2,7]
03  print(input1)
04
05  def shuff(L):
06      '''Randomises order of elements in the taken array'''
07      for i in range(len(L)):
08          iRandom = randint(0,len(L)-1) # random index in the array
09          temp = L[i]          # swapping takes place
10          L[i] = L[iRandom]
11          L[iRandom] = temp
12      return L
13
14  print(shuff(input1))
```

2. Counting the trailing zeros in a factorial number:

```
01  num = int(input("num: "))
02
03  def trailingZero(t):
04      '''Returns number of trailing zeros of answer from a factorial int'''
05      def fact(n):
06          '''Returns factorial number of argument'''
07          if n == 1:
08              return n
09          else:
10              return n * fact(n-1)
11      count = 0
12      factAns = str(fact(t))[::-1] # factorial number, converted to string and sequentially reversed
13      for i in range(len(factAns)): # counts 0s: stops if no more at trailing end
14          if factAns[i] != "0":
15              break
16          else:
17              count+=1
18      return count
19
20  print(trailingZero(num))
```

3. Returning the highest square number less or equal to the function parameter:

```
01  from math import sqrt
02  num = int(input("number: "))
03
04  def highestSquareNumber(n):
05      ''' Returns closest perfect square number '''
06
07      psn = []                        ## generating perfect square numbers
08      for i in range(10000):
09          if sqrt(i) % 1 == 0:
10              psn.append(i)
11
12      for i in range(len(psn)):       ## returns closest perfect square number
13          if n <= psn[i]:
14              return psn[i]
15
16  print( highestSquareNumber(num) )
```

4. Pseudocode for questions 1 and 2:

Question 1:

Pseudocode	Time cost
`from random import *`	
`input1 = [5,3,8,6,1,9,2,7]`	1
`print(input1)`	1
`def shuff(L):`	
`  for i in range(len(L)):`	n
`  iRandom =`	n
`randint(0,len(L)-1)`	n
`  temp = L[i]`	n
`  L[i] = L[iRandom]`	n
`  L[iRandom] = temp`	1
`  return L`	
`print(shuff(input1))`	1
Runtime bound:	f(n) = 1 + 1 + n + n + n + n + n + 1 + 1 = 5n + 4
Complexity:	O(n)

Question 2:

Pseudocode	Time cost
`num = int(input("num: "))`	1
`def trailingZero(t):`	
`  def fact(n):`	
`  if n == 1:`	1
`  return n`	1
`  else:`	
`  return n * fact(n-1)`	n
`  count = 0`	1
`  factAns = str(fact(t))[::-1]`	1
`  for i in range(len(factAns)):`	n
`  if factAns[i] != "0":`	n
`  break`	1
`  else:`	
`  count+=1`	n
`  return count`	1
`print(trailingZero(num))`	1
Runtime bound:	f(n) = 2n + 2n? + 8 = 2n + 2mn + 8 = 2n + 2n + 8 = 4n + 8
Complexity:	O(n)

5. Matrix calculations:

Addition function

Pseudocode	Time cost
`ADD_MATRICES(matrix1, matrix2):` `    '''Returns a new matrix with each value equal to the sum of each index values of the given matrices'''`	
`    newMatrix <- new empty integer array [] of size length of matrix1 and matrix2`	1
`    FOR i <- 0 TO matrix1.length:`	
`        IF len(matrix1) != len(matrix2) or len(matrix1[i]) != len(matrix2[i]):`	
`            RETURN NONE`	n
`        ELSE:`	
`            FOR j <- 0 to matrix1[i].length:`	n
`                newMatrix[i][j] <- matrix1[i][j] + matrix2[i][j]`	1
`    RETURN newMatrix`	n^2
	n^2
Runtime bound:	$2n^2 + 2n + 2$
Complexity:	$O(n^2)$

[CONTINUED]

High-level language code (Python):

```
01  m1 = [[40, 17, 82],
02        [13, 30, 93],
03        [2, 33, 6]]
04
05  m2 = [[35, 89, 22],
06        [16, 41, 4],
07        [1, 55, 61]]
08
09  def addM(matrix1, matrix2):
10      '''Returns a new matrix with each value equal to the sum of each index values of the given matrices'''
11      newMatrix = [[0 for i in range(len(matrix1[x]))] for x in range(len(matrix1))]
12      # 0s added to each index of inner list; inner list added to each index of main list
13      # indicates empty spaces to new array, equal to size length of one of given matrices
14
15      for i in range(len(matrix1)):
16          if len(matrix1) != len(matrix2) or len(matrix1[i]) != len(matrix2[i]): # checks if given matrices
    are of the same size length for addition to take place
17              return None
18          else:
19              for j in range(len(matrix1[i])):
20                  newMatrix[i][j] = matrix1[i][j] + matrix2[i][j] # gives sum of values of both given
    matrices by index
21      return newMatrix
22
23  print(addM(m1, m2))
```

[CONTINUED]

Subtraction function:

Pseudocode	Time cost
`SUBTRACT_MATRICES(matrix1, matrix2):`	
`    '''Returns a new matrix with each value equal to the subtraction of each index values of the first matrix by that of the second matrix'''`	
`    newMatrix <- new empty integer array [] of size length of matrix1 and matrix2`	
`    FOR i <- 0 TO matrix1.length:`	1
`        IF len(matrix1) != len(matrix2) or len(matrix1[i]) != len(matrix2[i]):`	
`            RETURN NONE`	n
`        ELSE:`	
`            FOR j <- 0 to matrix1[i].length:`	n
`                newMatrix[i][j] <- matrix1[i][j] - matrix2[i][j]`	1
`    RETURN newMatrix`	
	n^2
	n^2
	1
Runtime bound:	$2n^2 + 2n + 2$
Complexity:	$O(n^2)$

[CONTINUED]

High-level language code (Python):

```
01  m1 = [[40, 17, 82],
02        [13, 30, 93],
03        [2, 33, 6]]
04
05  m2 = [[35, 89, 22],
06        [16, 41, 4],
07        [1, 55, 61]]
08
09  def subM(matrix1, matrix2):
10      '''Returns a new matrix with each value equal to the subtraction of each index values of the first
    matrix by that of the second matrix'''
11      newMatrix = [[0 for i in range(len(matrix1[x]))] for x in range(len(matrix1))]
12
13      for i in range(len(matrix1)):
14          if len(matrix1) != len(matrix2) or len(matrix1[i]) != len(matrix2[i]):
15              return None
16          else:
17              for j in range(len(matrix1[i])):
18                  newMatrix[i][j] = matrix1[i][j] - matrix2[i][j] # subtracts each value in given matrices
19      return newMatrix
20
21  print(subM(m1, m2))
```

[CONTINUED]

Multiplication function:

Pseudocode	Time cost
`MULTIPLY_MATRICES(matrix1, matrix2):`	
`    '''Returns a new matrix with each value equal to the multiplication of given matrices'''`	
`      newMatrix <- new empty integer array [] equal to height (j) of matrix1 and width (i) of matrix2`	
	1
`    FOR i <- 0 TO matrix1.length:`	
`            FOR j <- 0 to matrix1[i].length:`	n
`            newMatrix[i][j] <- (matrix1[i][j] * matrix2[i][j]) + (matrix1[i][j+1] * matrix2[i+1][j])`	n^2
`    RETURN newMatrix`	n^2
	1
Runtime bound:	$2n^2 + n + 2$
Complexity:	$O(n^2)$

[CONTINUED]

High-level language code (Python):

```
01  m1 = [[7, 8],
02        [0, 9],
03        [3, 6]]
04
05  m2 = [[5, 9, 2],
06        [6, 1, 4]]
07
08  def multiM(matrix1, matrix2):
09      '''Returns a new matrix with each value equal to the multiplication of given matrices'''
10      newMatrix = [[0 for i in range(len(matrix1))] for x in range(len(matrix1[0]))]
11      # determines that the dimension product of new matrix will be height (j) of matrix1 and width (i) of
    matrix2
12
13      for i in range(len(matrix1)-1):
14          for j in range(len(matrix1[i])-1):
15              newMatrix[i][j] = (matrix1[i][j] * matrix2[i][j]) + (matrix1[i][j+1] * matrix2[i+1][j]) #
    multiplies values in given matrices by index
16      return newMatrix
17
18  print(multiM(m1, m2))
```

6. Function reversing the words in a given sentence:

Pseudocode	Time cost
`String word <- "This is awesome"`	1
`WORD_REVERSE(word):`	
`    word <- String [] word`	1
`    reversedWord <- New String Array []`	1
`    For i <- a.length-1 Downto 0 {`	n
`        Append a[i] to reversedWord`	n
`    }`	
`    reversedWord <- joined String reversedWord`	1
`    RETURN reversedWord`	1
`PRINT WORD_REVERSE(word)`	1
Runtime bound:	f(n) = 2n + 6
Complexity:	O(n)

High-level language code (Python):

```
01  word = "This is awesome"
02
03  def wordReverse(w):
04      '''Returns the string argument reversed by
    word'''
05      w = w.split(' ')
06      reversedWord = []
07  ##    w = w[::-1]
08      for i in range(len(w)-1, -1, -1):
09          reversedWord.append(w[i])
10      reversedWord = ' '.join(reversedWord)
11      return reversedWord
12
13  print(wordReverse(word)
```

7. Checking if a given integer is a prime number:

Pseudocode	Time cost
number <- any given integer value	1
IS_PRIME(num, test <- number)	
If num ≤ 2 OR test = 1 {	1
Return TRUE	1
} Else If num mod test = 0 {	1
Return IS_PRIME(num, test-1)	n
}	
Return FALSE	1
}	
PRINT IS_PRIME(number)	1
Runtime bound:	f(n) = n + 6
Complexity:	O(n)

High-level language code (Python):

```
01  number = int(input("number: "))
02
03  def isPrime(num, test = number-1):
04      '''Determines whether an integer value is
    prime'''
05      try:
06          if num <= 2 or test == 1:
07              return True
08          elif num % test != 0:
09              return isPrime(num, test-1)
10          return False
11      except RecursionError:
12          return False
13
14  print( "%d is a prime number: %s" % (number,
    isPrime(number)) )
```

8. Removing vowels from a given string:

Pseudocode	Time cost
word <- "beautiful"	1
REMOVE_VOWEL(wrd)	
wrd <- String [] wrd	1
For i <- 0 to length(wrd):	n
If w[i] = "a" OR	
w[i] = "e" OR	
w[i] = "i" OR	
w[i] = "o" OR	
w[i] = "u":	n
Remove wrd[i]	n
wrd = String wrd	1
Return wrd	1
PRINT REMOVE_VOWEL(word)	1
Runtime bound:	f(n) = 3n + 5
Complexity:	O(n)

High-level language code (Python):

```
01  word = "beautiful"
02
03  def removeVowel(wrd):
04      wrd = list(wrd)
05      for i in range(len(wrd)):
06          if wrd[i] == "a" or wrd[i] == "e" or
    wrd[i] == "i" or wrd[i] == "o" or wrd[i] == "u":
07              wrd[i] = ''
08      wrd = ''.join(wrd)
09      return wrd
10
11  print(removeVowel(word))
```

9. Adapting the binary search (searching by interval):

```
01  number1 = int(input("1st number: "))
02  number2 = int(input("2nd number: "))
03  List = [4, 19, 23, 36, 40, 43, 61, 64, 78, 95]
04
05  def binarySearch(num1, num2, array):
06      '''performs binary search to identify if there is
07         a number in the List within a given interval'''
08      mid = len(array)//2
09      try:
10          if num1 <= num2:
11              if array[mid-1] >= num1 and array[mid-1] <= num2:
12                  return True
13              elif array[mid-1] > num2:                              ##if the value is larger than pivot (to
    the right of the array)
14                  return binarySearch(num1, num2, array[:mid-1])     ##calls itself with the array halved to
    the right side of the pivot
15              elif array[mid-1] < num1:                              ##if the value is smaller than pivot
    (to the left of the array)
16                  return binarySearch(num1, num2, array[mid:])       ##calls itself with the array halved to
    the left side of the pivot
17              return False
18          else:
19              return "ERROR! Lower value > upper value"
20      except IndexError or RecursionError:                           ##signifying the two common errors
    during runtime
21          return False
22
23  print("Is there an integer between %d and %d in the list? Answer: %s" % (number1, number2,
    binarySearch(number1, number2, List)))
```

[CONTINUED]

Table 1:

Pseudocode	Time cost
number1 <- String input converted to Int	1
number2 <- String input converted to Int	1
List <- [4, 19, 23, 36, 40, 43, 61, 64, 78, 95]	1
binarySearch(num1, num2, array) {	
midpoint <- len(array)//2	1
If num1 <= num2 {	1
If array[midpoint] ≥ num2 AND array[midpoint] ≤ num2 {	1
Return TRUE	1
}	
ElseIf midpoint > num2 {	1
Return binarySearch(num1, num2, right half of array)	n
}	
ElseIf midpoint < num2 {	1
Return binarySearch(num1, num2, left half of array)	n
Return FALSE	1
}	
Else {	
Return “Error! Lower value > upper value”	1
}	
}	
Runtime bound:	= 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + n + 1 + n + 1 + 1 = 2n + 11
Complexity:	O(n)

[CONTINUED]

Table2 – After considering how recursion affects runtime bound:

Pseudocode		Time cost
number1 <- String input converted to Int	1	1
number2 <- String input converted to Int	1	1
List <- [4, 19, 23, 36, 40, 43, 61, 64, 78, 95]	1	1
binarySearch(num1, num2, array) {		
midpoint <- len(array)//2	n	n
If num1 <= num2 {	n	n
If array[midpoint] ≥ num2 AND array[midpoint] ≤ num2 {	n	n
Return TRUE	1	1
}		
ElseIf midpoint > num2 {	n	n
Return binarySearch(num1, num2, right half of array)	n	n
}		
ElseIf midpoint < num2 {	n	n
Return binarySearch(num1, num2, left half of array)	n	n
Return FALSE	1	1
}		
Else {		
Return “Error! Lower value > upper value”	1	1
}		
}		
Runtime bound:		f(n) = 1 + 1 + 1 + n + n + n + 1 + n + n + n + n + 1 + 1 = 7n + 6
Complexity:		O(n)

[CONTINUED]

Note 1:

If I specified the except statement as:

```
20        except IndexError and RecursionError:
```

then the programme, when executed, would have first encountered an error that would be both an IndexError and RecursionError, by first identifying the IndexError and rendering it mute without the user's awareness and then proceeding to the identifying recursion error as it would also be the case.

As a solution I substituted to above statement to:

```
20        except IndexError or RecursionError:
```

The solution is mainly based on a test run of the programme and entering a negative number for the lower range bound. The attempt was carried out with the except statement written before the amendment, as shown in the first instance, and when running the programme, the stack report had returned an IndexError regardless of the written except statement. This led to the second instance as the solution, which indicates that the programme would the code would cause either error instead of both, thus finally displaying the expected result.

Note 2:

Initial analysis concluded that the time taken for some of the code is constant, yet when the first condition (base case) is not met and the programme executes the code in lines 14 or 16, the function calls itself thus triggering a recursion. Due to this, the programme re-executes the code before the call until the base case is met, which arguably therefore alters the time complexity of the said code to **O(log n)**. Nevertheless, after much consideration, the worst-case scenario is taken in account so the complexity of the programme algorithm would result to O(log n) regardless.

10. Extracting the sub-sequence of maximum length (ascending order):

```
01      array = [71, 41, 15, 68, 49, 9, 26, 46, 15, 53, 96, 23, 54, 17, 11, 5]
02
03      print(array, end="\n\n")
04
05      def subSeq(list1):
06
07          '''Finds and returns ordered sub-sequence within a given list'''
08
09          newList = []  # new list for ordered sub-sequence
10
11          while len(list1) >= 2:
12              try:
13                  lowest = list1[0]       # lowest value is first value in given list by default (temp)
14                  if lowest > list1[1]:
15                      lowest = list1[1]   # value in second index of the list is now the lowest
16                      list1 = list1[1:]   # list sequence now starts from point of new lowest value
17                  else:
18                      for i in range(1, len(list1)):
19                          if lowest > list1[i]:       # checks for smaller integers than lowest value in given
        list
20                              lowest = list1[i]       # value in index i is now the lowest value
21                              newList.append(lowest)    # adds new lowest value to new list
22                              list1 = list1[i+1:]     # list now starts from index adjacent to that of lowest
        value
23                              for j in range(len(newList[:i])): # code in for-loop deletes every value before
        last value of new list (lowest) if a smaller value is found
24                                  if newList[-1] < newList[j]:
25                                      del newList[:i]
26                                      break
27                              break
28                  if len(list1) == 2:
29                      if newList[-1] < list1[-2] and newList[-1] < list1[-1]:
30                          # checks if last value in new list is smaller than last two elements in given list;
31                          # if otherwise, following statements are disregards them + returns new list
```

```
32                          if list1[0] < list1[1]:         # orders remaining values if size length of list is
    two
33                              newList.append(list1[0])
34                              newList.append(list1[1])
35                          else:
36                              newList.append(list1[1])
37                          break
38              except IndexError:
39                  break
40
41          return newList
42
43      print(subSeq(array))
```

Note:

I could have adjusted the program to run through both the main and new list to check if there is a smaller value in the main list than the last value added to the new list, but bigger than the second to last value and onwards (in the new list). This would have made the function find a more closely consecutive sub-sequence. This could not have been fully accomplished due to great difficulty and time-consuming attempts.

12. Rendering IN_ORDER function as iterative from recursive:

```
34  def in_order(tree):
35      currentNode = tree
36      stack = []
37
38      while currentNode:
39
40          if currentNode.left != None:
41              stack.append(currentNode.value)
42              currentNode = currentNode.left
43
44          print( currentNode.value )
45
46          if currentNode.right != None:
47              print(stack[-1])
48              stack.pop()
49              currentNode = currentNode.right
50
51
52          ##if currentNode.right != None:
53          ##    stack.pop()
54
55          if stack == []:
56              print("break")
57              break
```

Appendix:

1. Binary Search tree template - http://pastebin.com/LXdWF0KW
2. Git repository – https://github.com/Nathan-Zenga/210CT-Coursework-tasks