Programming, Algorithms and Data Structures (210CT)

COURSEWORK REPORT

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INTRODUCTION:

The report is a record of all programming code written in accordance to the tasks objectives, alongside the comments for further description and clarification of functionalities as well as side notes entailing encountered drawbacks and applicable (or theoretical) solutions to them.

To view the live code from all tasks, see appendix.

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1. Defining the shuffle function:

```
01 from random import *
02 input1 = [5,3,8,6,1,9,2,7]
03 print(input1)
04
05 def shuff(L):
        '''Randomises order of elements in the taken array'''
06
        for i in range(len(L)):
07
            iRandom = randint(0, len(L)-1) # random index in the array
08
                               # swapping takes place
            temp = L[i]
09
            L[i] = L[iRandom]
10
            L[iRandom] = temp
11
12
        return L
13
14 print(shuff(input1))
```

2. Counting the trailing zeros in a factorial number:

```
num = int(input("num: "))
02
03
   def trailingZero(t):
        '''Returns number of trailing zeros of answer from a factorial int'''
04
05
        def fact(n):
            '''Returns factorial number of argument'''
06
07
            if n == 1:
08
                return n
09
            else:
                return n * fact(n-1)
10
11
        count = 0
12
        factAns = str(fact(t))[::-1] # factorial number, converted to string and sequentially reversed
        for i in range(len(factAns)): # counts 0s: stops if no more at trailing end
13
            if factAns[i] != "0":
14
                break
15
16
            else:
17
                count+=1
18
        return count
19
    print(trailingZero(num))
```

3. Returning the highest square number less or equal to the function parameter:

```
from math import sqrt
   num = int(input("number: "))
02
03
    def highestSquareNumber(n):
04
        ''' Returns closest perfect square number '''
05
06
                                        ## generating perfect square numbers
        psn = []
07
        for i in range(10000):
08
            if sqrt(i) % 1 == 0:
09
                psn.append(i)
10
11
        for i in range(len(psn)): ## returns closest perfect square number
12
            if n <= psn[i]:</pre>
13
                return psn[i]
14
15
    print( highestSquareNumber(num) )
16
```

4. Pseudocode for questions 1 and 2:

Question 1:

Pseudocode	Time cost
from random import *	
input1 = $[5,3,8,6,1,9,2,7]$	1
print(input1)	1
def shuff(L):	
for i in range(len(L)):	n
iRandom =	n
randint(0,len(L)-1)	n
temp = L[i]	n
L[i] = L[iRandom]	n
L[iRandom] = temp	1
return L	
<pre>print(shuff(input1))</pre>	1
Runtime bound:	f(n) = 1 + 1 + n + n + n + n
	+ n + 1 + 1
	= 5n + 4
Complexity:	O(n)
• •	

Question 2:

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Pseudocode	Time cost
<pre>num = int(input("num: "))</pre>	1
<pre>def trailingZero(t): def fact(n):</pre>	
if n == 1:	1
return n	1
else:	
return n * fact(n-1)	n
count = 0	1
factAns = str(fact(t))[::-1]	1
<pre>for i in range(len(factAns)): if factAns[i] != "0";</pre>	n
<pre>if factAns[i] != "0": break</pre>	n 1
else:	1
count+=1	n
return count	1
<pre>print(trailingZero(num))</pre>	1
Runtime bound:	f(n) = 2n + 2n? + 8 = 2n + 2mn + 8 = 2n + 2n + 8
	= 4n + 8
Complexity:	O(n)

5. Matrix calculations:

Addition function

```
Pseudocode
                                                                                                              Time cost
ADD MATRICES (matrix1, matrix2):
   -
'''Returns a new matrix with each value equal to the sum of each index values of the given
matrices'''
    newMatrix <- new empty integer array [] of size length of matrix1 and matrix2</pre>
                                                                                                              1
    FOR i <- 0 TO matrix1.length:
        IF len(matrix1) != len(matrix2) or len(matrix1[i]) != len(matrix2[i]):
            RETURN NONE
                                                                                                              n
        ELSE:
            FOR j <- 0 to matrix1[i].length:</pre>
                                                                                                              n
                newMatrix[i][j] <- matrix1[i][j] + matrix2[i][j]</pre>
                                                                                                              1
    RETURN newMatrix
                                                                                                              n^2
                                                                                                              n^2
                                                                                               Runtime bound: 2n^2 + 2n + 2
                                                                                                  Complexity: O(n^2)
```

High-level language code (Python):

```
01 \quad m1 = [[40, 17, 82],
02
          [13, 30, 93],
03
          [2, 33, 6]]
04
05 \quad m2 = [[35, 89, 22],
06
          [16, 41, 4],
07
          [1, 55, 61]]
08
09
    def addM (matrix1, matrix2):
        '''Returns a new matrix with each value equal to the sum of each index values of the given matrices'''
10
        newMatrix = [[0 for i in range(len(matrix1[x]))] for x in range(len(matrix1))]
11
12
        # Os added to each index of inner list; inner list added to each index of main list
13
        # indicates empty spaces to new array, equal to size length of one of given matrices
14
15
        for i in range(len(matrix1)):
16
            if len(matrix1) != len(matrix2) or len(matrix1[i]) != len(matrix2[i]): # checks if given matrices
    are of the same size length for addition to take place
                return None
17
18
            else:
                for j in range(len(matrix1[i])):
19
                    newMatrix[i][j] = matrix1[i][j] + matrix2[i][j] # gives sum of values of both given
20
    matrices by index
21
        return newMatrix
22
23 print(addM(m1, m2))
```

Subtraction function:

```
Time cost
Pseudocode
SUBTRACT MATRICES (matrix1, matrix2):
    '''Returns a new matrix with each value equal to the subtraction of each index values of the
first matrix by that of the second matrix'''
    newMatrix <- new empty integer array [] of size length of matrix1 and matrix2</pre>
    FOR i <- 0 TO matrix1.length:
                                                                                                                1
        IF len(matrix1) != len(matrix2) or len(matrix1[i]) != len(matrix2[i]):
             RETURN NONE
                                                                                                                n
        ELSE:
             FOR j <- 0 to matrix1[i].length:</pre>
                                                                                                                n
                 newMatrix[i][j] <- matrix1[i][j] - matrix2[i][j]</pre>
                                                                                                                1
    RETURN newMatrix
                                                                                                                n^2
                                                                                                                n^2
                                                                                                 Runtime bound: 2n^2 + 2n + 2
                                                                                                    Complexity: O(n<sup>2</sup>)
```

High-level language code (Python):

```
01 \quad m1 = [[40, 17, 82],
          [13, 30, 93],
02
03
          [2, 33, 6]]
04
05 \quad m2 = [[35, 89, 22],
06
          [16, 41, 4],
          [1, 55, 61]]
07
08
09 def subM(matrix1, matrix2):
        '''Returns a new matrix with each value equal to the subtraction of each index values of the first
10
    matrix by that of the second matrix'''
        newMatrix = [[0 for i in range(len(matrix1[x]))] for x in range(len(matrix1))]
11
12
13
        for i in range(len(matrix1)):
            if len(matrix1) != len(matrix2) or len(matrix1[i]) != len(matrix2[i]):
14
15
                return None
16
            else:
                for j in range(len(matrix1[i])):
17
                    newMatrix[i][j] = matrix1[i][j] - matrix2[i][j] # subtracts each value in given matrices
18
19
        return newMatrix
20
21 print(subM(m1, m2))
```

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Multiplication function:

Pseudocode	Time cost
MULTIPLY MATRICES (matrix1, matrix2):	
'''Returns a new matrix with each value equal to the multiplication of given matrices'''	
IF matrix2 is a list type variable:	1
newMatrix <- new empty integer array [] equal to height (j) of matrix1 and width (i) of	
matrix2	1
FOR i <- 0 TO matrix1.length:	n
FOR j <- 0 to matrix1[i].length:	n ²
newMatrix[i][j] <- (matrix1[i][j] * matrix2[i][j]) + (matrix1[i][j+1] *	
matrix2[i+1][j])	n ²
ESLE IF matrix2 is a integer type variable:	1
newMatrix <- new empty integer array [] equal size length of matrix1	1
FOR i <- 0 TO matrix1.length:	n
FOR j <- 0 to matrix1[i].length:	n ²
newMatrix[i][j] <- matrix1[i][j] * 2	n^2
RETURN newMatrix	1
Runtime bound:	$4n^2 + 2n + 5$
Complexity:	O(n²)

```
01 \quad m1 = [7, 8],
02
          [0, 9],
03
          [3, 6]]
0.4
05 \quad m2 = 2
06
07 def multiM(matrix1, matrix2):
        '''Returns a new matrix with each value equal to the multiplication of given matrices'''
0.8
09
        if type (matrix2) == list: # multiplies values in indices from first and second matrix, if second
    argument is a matrix
10
            newMatrix = [[0 for i in range(len(matrix1))] for x in range(len(matrix1[0]))]
            # determines that the dimension product of new matrix will be height (i) of matrix1 and width (j)
11
    of matrix2
12
            for i in range(len(matrix1)-1):
13
                for j in range(len(matrix1[i])-1):
14
                    newMatrix[i][j] = (matrix1[i][j] * matrix2[i][j]) + (matrix1[i][j+1] * matrix2[i+1][j]) #
1.5
    multiplies each index value in given matrix by that of second matrix
        elif type (matrix2) == int: # multiplies values in indices from first matrix by number, if second
16
    argument is an integer
17
            newMatrix = [[0 for i in range(len(matrix1[x]))] for x in range(len(matrix1))]
18
            number = matrix2 # to identify second argument as number
19
            for i in range(len(matrix1)):
20
                for j in range(len(matrix1[i])):
21
                    newMatrix[i][j] = matrix1[i][j] * number
22
        return newMatrix
23
24 print(multiM(m1, m2))
```

6. Function reversing the words in a given sentence:

Pseudocode	Time cost
String word <- "This is awesome"	1
WORD REVERSE (word):	
word <- String [] word	1
reversedWord <- New String Array []	1
For i <- a.length-1 Downto 0 {	n
Append a[i] to reversedWord	n
}	
reversedWord <- joined String	1
reversedWord	1
RETURN reversedWord	
	1
PRINT WORD_REVERSE(word)	
Runtime bound:	f(n) = 2n + 6
Complexity:	O(n)

```
01 word = "This is awesome"
02
03 def wordReverse(w):
04    '''Returns the string argument reversed by
    word'''
05     w = w.split(' ')
06     reversedWord = []
07 ##     w = w[::-1]
08     for i in range(len(w)-1, -1, -1):
09         reversedWord.append(w[i])
10     reversedWord = ' '.join(reversedWord)
11    return reversedWord
12
13 print(wordReverse(word)
```

7. Checking if a given integer is a prime number:

Pseudocode	Time cost
number <- any given integer value	1
IS PRIME(num, test <- number)	
If $num \le 2$ OR test = 1 {	1
Return TRUE	1
} Else If num mod test = 0 {	1
<pre>Return IS_PRIME(num, test-1)</pre>	n
}	
Return FALSE	1
}	
PRINT IS_PRIME(number)	1
Runtime bound:	f(n) = n + 6
Complexity:	O(n)

```
01 number = int(input("number: "))
03 def isPrime(num, test = number-1):
        '''Determines whether an integer value is
   prime'''
05
        try:
06
            if num <= 2 or test == 1:
07
               return True
08
            elif num % test != 0:
09
               return isPrime(num, test-1)
10
            return False
11
        except RecursionError:
12
            return False
13
14 print ( "%d is a prime number: %s" % (number,
    isPrime(number)) )
```

8. Removing vowels from a given string:

Pseudocode	Time cost
word <- "beautiful"	1
DEMONIE MONIEL (1)	
REMOVE_VOWEL(wrd)	
wrd <- String [] wrd	1
For i <- 0 to length(wrd):	n
If w[i] = "a" OR	
w[i] = "e" OR	
w[i] = "i" OR	
w[i] = "o" OR	
w[i] = "u":	n
Remove wrd[i]	n
	1
wrd = String wrd	_
Return wrd	1
PRINT REMOVE_VOWEL(word)	1
Runtime bound:	f(n) = 3n + 5
Complexity:	O(n)
• • •	• •

```
01 word = "beautiful"
02
03 def removeVowel(wrd):
       wrd = list(wrd)
       for i in range(len(wrd)):
05
           if wrd[i] == "a" or wrd[i] == "e" or
   wrd[i] == "i" or wrd[i] == "o" or wrd[i] == "u":
              wrd[i] = ''
07
       wrd = ''.join(wrd)
08
       return wrd
09
10
11 print(removeVowel(word))
```

9. Adapting the binary search (searching by interval):

```
01  number1 = int(input("1st number: "))
02  number2 = int(input("2nd number: "))
   List = [4, 19, 23, 36, 40, 43, 61, 64, 78, 95]
04
05 def binarySearch(num1, num2, array):
06
        '''performs binary search to identify if there is
           a number in the List within a given interval'''
07
        mid = len(array)//2
0.8
09
        try:
            if num1 <= num2:
10
11
                if array[mid-1] >= num1 and array[mid-1] <= num2:</pre>
12
                    return True
13
                elif array[mid-1] > num2:
                                                                        ##if the value is larger than pivot (to
    the right of the array)
                    return binarySearch(num1, num2, array[:mid-1])
                                                                        ##calls itself with the array halved to
14
    the right side of the pivot
                elif array[mid-1] < num1:</pre>
                                                                        ##if the value is smaller than pivot
15
     (to the left of the array)
                    return binarySearch (num1, num2, array[mid:]) ##calls itself with the array halved to
16
    the left side of the pivot
17
                return False
18
            else:
19
                return "ERROR! Lower value > upper value"
        except IndexError or RecursionError:
20
                                                                        ##signifying the two common errors
    during runtime
21
            return False
22
    print("Is there an integer between %d and %d in the list? Answer: %s" % (number1, number2,
    binarySearch(number1, number2, List)))
```

Table 1:

Pseudocode	Time cost
number1 <- String input converted to Int	1
number2 <- String input converted to Int	1
List <- [4, 19, 23, 36, 40, 43, 61, 64, 78, 95]	1
<pre>binarySearch(num1, num2, array) {</pre>	
<pre>midpoint <- len(array)//2</pre>	1
<pre>If num1 <= num2 {</pre>	1
If $array[midpoint] \ge num2 AND array[midpoint] \le num2 {$	1
Return TRUE	1
}	
<pre>ElseIf midpoint > num2 {</pre>	1
Return binarySearch(num1, num2, right half of array)	n
}	
<pre>ElseIf midpoint < num2 {</pre>	1
Return binarySearch(num1, num2, left half of array)	n
Return FALSE	1
}	
Else {	
Return "Error! Lower value > upper value"	1
}	
}	
Runtime bound:	= 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + n + 1 + n + 1 + 1
	= 2n + 11
Complexity:	O(n)

Table2 – After considering how recursion affects runtime bound:

```
Pseudocode
                                                                       Time cost
number1 <- String input converted to Int</pre>
                                                                    1 1
number2 <- String input converted to Int</pre>
                                                                    1 1
List <- [4, 19, 23, 36, 40, 43, 61, 64, 78, 95]
                                                                    1 1
binarySearch(num1, num2, array) {
 midpoint <- len(array)//2
                                                                    n n
   If num1 <= num2 {</pre>
                                                                    n n
     If array[midpoint] ≥ num2 AND array[midpoint] ≤ num2 {
                                                                    n n
       Return TRUE
                                                                    1 1
     ElseIf midpoint > num2 {
                                                                    n n
       Return binarySearch(num1, num2, right half of array)
                                                                    n n
     ElseIf midpoint < num2 {</pre>
                                                                    n n
       Return binarySearch(num1, num2, left half of array)
                                                                    n n
       Return FALSE
                                                                    1 1
     Else {
       Return "Error! Lower value > upper value"
                                                                    1 1
                                                         = 7n + 6
                                                            Complexity: O(n)
```

Note 1:

If I specified the except statement as:

then the programme, when executed, would have first encountered an error that would be both an IndexError and RecursionError, by first identifying the IndexError and rendering it mute without the user's awareness and then proceeding to the identifying recursion error as it would also be the case.

As a solution I substituted to above statement to:

The solution is mainly based on a test run of the programme and entering a negative number for the lower range bound. The attempt was carried out with the except statement written before the amendment, as shown in the first instance, and when running the programme, the stack report had returned an IndexError regardless of the written except statement. This led to the second instance as the solution, which indicates that the programme would the code would cause either error instead of both, thus finally displaying the expected result.

Note 2:

Initial analysis concluded that the time taken for some of the code is constant, yet when the first condition (base case) is not met and the programme executes the code in lines 14 or 16, the function calls itself thus triggering a recursion. Due to this, the programme re-executes the code before the call until the base case is met, which arguably therefore alters the time complexity of the said code to **O(log n)**. Nevertheless, after much consideration, the worst-case scenario is taken in account so the complexity of the programme algorithm would result to O(log n) regardless.

10. Extracting the sub-sequence of maximum length (ascending order):

```
array = [71, 41, 15, 68, 49, 9, 26, 46, 15, 53, 96, 23, 54, 17, 11, 5]
01
02
       print(array, end="\n\n")
03
04
05
       def subSeq(list1):
06
07
           '''Finds and returns ordered sub-sequence within a given list'''
0.8
09
           newList = [] # new list for ordered sub-sequence
10
11
           while len(list1) >= 2:
12
               try:
13
                   lowest = list1[0]
                                           # lowest value is first value in given list by default (temp)
14
                   if lowest > list1[1]:
15
                       lowest = list1[1]
                                           # value in second index of the list is now the lowest
                       list1 = list1[1:] # list sequence now starts from point of new lowest value
16
17
                   else:
18
                       for i in range(1, len(list1)):
                                                       # checks for smaller integers than lowest value in given
19
                           if lowest > list1[i]:
       list
20
                               lowest = list1[i]
                                                       # value in index i is now the lowest value
21
                               newList.append(lowest)
                                                         # adds new lowest value to new list
22
                               list1 = list1[i+1:]
                                                       # list now starts from index adjacent to that of lowest
       value
                               for j in range(len(newList[:i])): # code in for-loop deletes every value before
23
       last value of new list (lowest) if a smaller value is found
24
                                   if newList[-1] < newList[j]:</pre>
25
                                       del newList[:i]
26
                                       break
27
                               break
28
                       if len(list1) == 2:
29
                           if newList[-1] < list1[-2] and newList[-1] < list1[-1]:
30
                               # checks if last value in new list is smaller than last two elements in given list;
31
                               # if otherwise, following statements are disregards them + returns new list
```

```
32
                                 if list1[0] < list1[1]:</pre>
                                                                  # orders remaining values if size length of list is
       two
                                     newList.append(list1[0])
33
                                     newList.append(list1[1])
34
35
                                 else:
                                     newList.append(list1[1])
36
37
                                 break
               except IndexError:
38
39
                    break
40
41
           return newList
42
       print(subSeq(array))
43
```

Note:

I could have adjusted the program to run through both the main and new list to check if there is a smaller value in the main list than the last value added to the new list, but bigger than the second to last value and onwards (in the new list). This would have made the function find a more closely consecutive subsequence. This could not have been fully accomplished due to great difficulty and time-consuming attempts.

12. Rendering IN_ORDER function as iterative from recursive:

```
def in order(tree):
35
        currentNode = tree
36
        stack = []
37
38
        while currentNode:
39
40
            if currentNode.left != None:
                stack.append(currentNode.value)
41
                currentNode = currentNode.left
42
43
            print( currentNode.value )
44
45
46
            if currentNode.right != None:
                print(stack[-1])
47
                stack.pop()
48
                currentNode = currentNode.right
49
50
51
            ##if currentNode.right != None:
52
53
            ##
                  stack.pop()
54
            if stack == []:
55
                print("break")
56
                break
57
```

Appendix:

- 1. Binary Search tree template http://pastebin.com/LXdWF0KW
- 2. Git repository https://github.com/Nathan-Zenga/210CT-Coursework-tasks