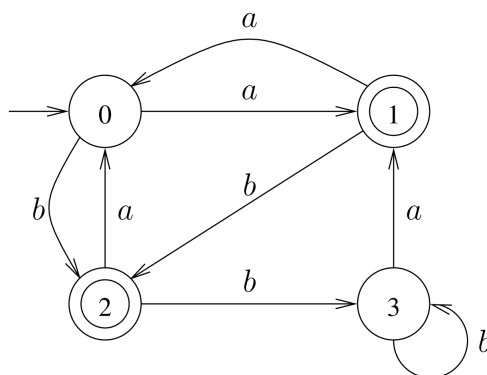


## CSCE 355 Homework 2

**Problem 1** Consider the following DFA:



- (a) For each of the strings below, say which state the DFA is in after reading the string, and say whether or not the DFA accepts the string.

*aaa    bb    bbb    abab    bbbbbbbaaa     $\epsilon$     aabbbababbaaabbababbbb*

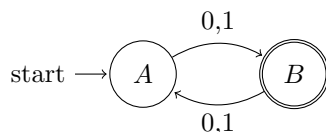
- (b) Give two different strings of length 4 that each make the DFA go from state 0 to state 1.

- (a)
- *aaa*: the DFA ends in state 1, where the DFA accepts the string.
  - *bb*: the DFA ends in state 3, where the DFA rejects.
  - *bbb*: the DFA ends in state 3, where the DFA rejects.
  - *abab*: the DFA ends in state 2, where the DFA accepts.
  - *bbbbbbbaaa*: the DFA ends in state 1, so DFA accepts.
  - $\epsilon$ : the DFA ends in state 0, so the DFA rejects.
  - *aabbbababbaaabbababbbb*: the DFA ends in state 3, so the DFA rejects.

- (b) The strings *abaa* and *abba* both end in state 1 and are length 4.

**Problem 2** Draw a DFA with alphabet  $\{0, 1\}$  that accepts a binary string  $x$  iff  $x$  has odd length, i.e., iff  $|x|$  is odd.

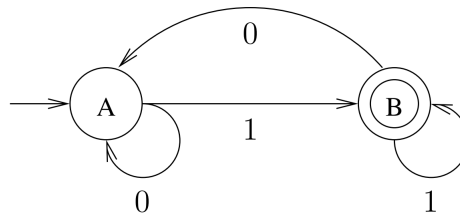
The following DFA works:



**Problem 3** Let  $A$  be the DFA given by the following tabular form:

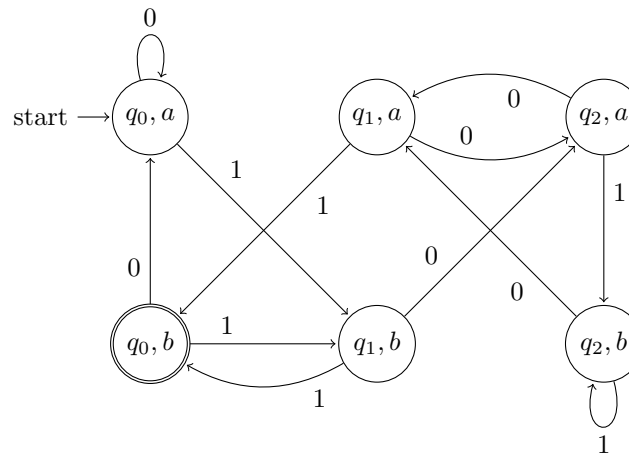
	0	1
$\rightarrow *q_0$	$q_0$	$q_1$
$q_1$	$q_2$	$q_0$
$q_2$	$q_1$	$q_2$

( $A$  accepts a binary string iff it represents a multiple of 3.) Recall the DFA described in class (here we'll call it  $B$ ) that accepts a binary string iff the string ends with 1:



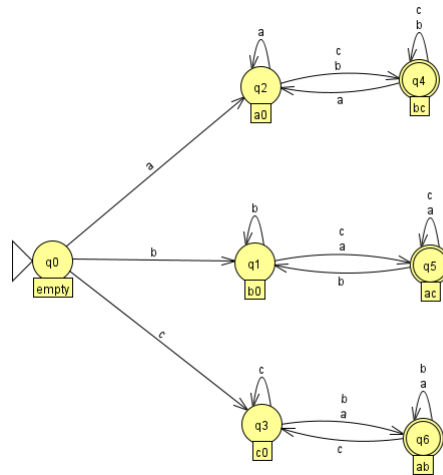
Recall the product construction from class. Draw the diagram for the product of  $A$  and  $B$  so the resulting DFA recognizes the language  $L(A) \cap L(B)$ .

Here is the product of DFAs  $A$  and  $B$ , that will accept a binary string iff it is a multiple of 3 and ends in a 1:



**Problem 4** Describe a DFA  $B$  that accepts a string over the alphabet  $\{a, b, c\}$  iff its first and last symbols are different.

The following DFA works:



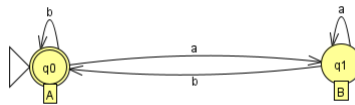
**Problem 5** Consider the following two languages over the alphabet  $\{a, b\}$ :

$$L_1 = \{w \mid w \text{ is either the empty string or ends with } b\},$$

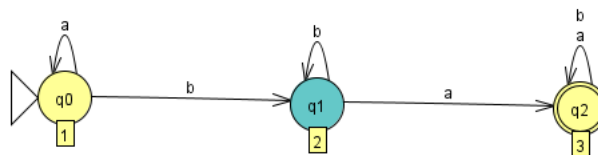
$$L_2 = \{w \mid \text{there is a } b \text{ followed by an } a \text{ somewhere in } w\}.$$

- (a) Draw a 2-state DFA recognizing  $L_1$  and a 3-state DFA recognizing  $L_2$ .
- (b) Using your answer and the product construction, draw a DFA recognizing  $L_1 \cap L_2$ . Do *not* perform any optimizations (e.g., removing unreachable states or transitions, or merging indistinguishable states).

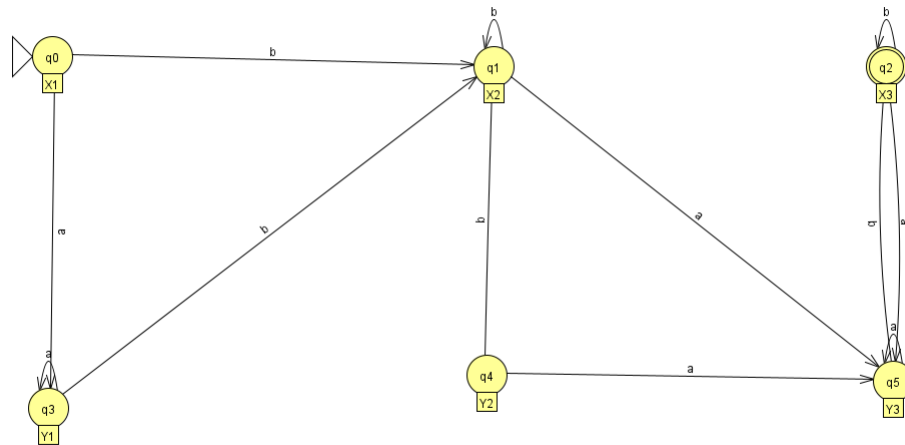
(a) This DFA recognizes  $L_1$ :



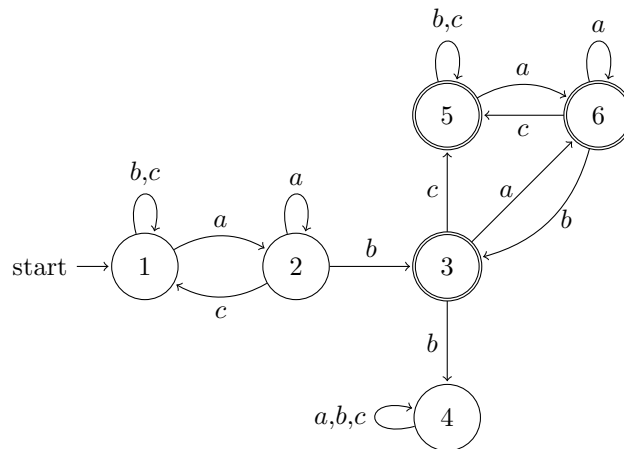
And this DFA recognizes  $L_2$ :



(b) We use the product construction to draw the DFA recognizing the intersection:

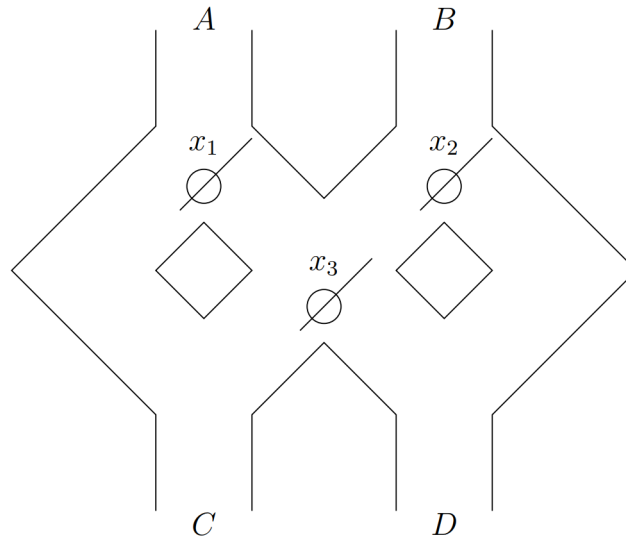


**Problem 6** Give the transition diagram for a DFA over the alphabet  $\Sigma = \{a, b, c\}$  that accepts a string  $w$  iff  $w$  contains  $ab$  as a substring but does not contain  $abb$  as a substring. What is the least number of states you need?



The least number of states you need is 6 because you need 3 accepting states for when you've seen  $ab$  and are still checking to see if you see any  $abbs$ , you need two states for finding the first  $ab$ , and you need a *trash* state for when you've seen  $abb$  in the string.

**Problem 7** This exercise is adapted from Exercise 2.2.1 on pages 52–53, which is formulated somewhat vaguely. Consider the marble-rolling toy (redrawn from Figure 2.8):



A marble is dropped at  $A$  or  $B$ . Levers  $x_1$ ,  $x_2$ , and  $x_3$  cause the marble to fall either to the left or to the right. whenever a marble encounters a lever, it causes the lever to reverse after the marble passes, so the next marble will take the opposite branch.

Model this toy as a finite automaton. An input to the automaton is a string over the alphabet  $\{A, B\}$ , which represents a sequence of marbles being dropped into the toy. The toy is initially in the configuration above before any marbles are dropped (so that the first ball will exit at  $C$  regardless of where it is dropped). Say that a sequence of marble drops is *accepted* exactly in the case that if one additional marble were to be dropped in, it would go out through  $D$  regardless of where it was dropped.

Each of the three switches can either point left or right, so we will have  $2^3 = 8$  states. We call the states  $d_1d_2d_3$ , where  $d_i$  is  $R$  if  $x_i$  is pointing right or  $L$  if it is pointing left. Then, the tabular form is:

	$A$	$B$
$\rightarrow RRR$	$LRR$	$RLL$
$RRL$	$LRL$	$RLR$
$RLR$	$LLR$	$RRR$
$RLL$	$LLL$	$RRL$
$LRR$	$RRL$	$LLL$
$LLR$	$RLL$	$LRR$
$*LRL$	$RRR$	$LLR$
$*LLL$	$RLR$	$LRL$