Computational Decision Making for Regular People

01: Introduction

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Today's Outline



- 1. Mathematical Modeling
- 2. Optimization
- 3. What kinds of things can be optimized?
- 4. General form of an optimization problem
- 5. A note on optimization theory
- 6. How to formulate an optimization problem
 - 6.1 The objective function
 - 6.2 Decision variables
 - 6.3 Parameters
 - 6.4 Constraints
- 7. Basic Examples



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Mathematical Modeling



How we think about a problem

- Consider what a good outcome looks like
 - Consider different ramifications of different decisions
 - utilitarian ethics, virtue ethics, deontological ethics, ...
- ▶ Discrete decisions (yes/no, 1,2,3, etc.)
- Continuous decisions
- Constraints against undesirable or infeasible decisions

How a computer thinks about a problem

- Quantify the quality of a solution
 - Combine all ramifications into one quantity
 - utilitarian objective only*

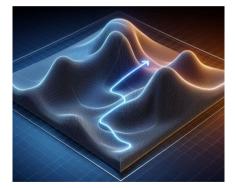
- Discrete decisions (yes/no, 1,2,3, etc.)
- Continuous decisions
- Constraints against undesirable or infeasible decisions

Optimization



the action of making the best or most effective use of a situation or resource.1

- Given a set of possible decisions, determine the best one(s)
- The way this determination is made depends on the nature of the set of decisions
 - If a human is executing this determination, the procedure will be unique to that person
 - If a computer algorithm is executing this determination, the procedure will be unique to that computer algorithm



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¹Oxford Languages

What Kinds of Things Can Be Optimized?

- ▶ What kinds of things can be optimized (using a computer)?
- If you can represent it using mathematical modeling, it (theoretically) can be optimized.
- Some problems are still too difficult for even modern computer algorithms to solve
 - Problems with lots and lots of variables
 - Problems with lots and lots of constraints
 - Problems that "don't behave well"
 - Lots of really good outcomes that lie very close to really bad outcomes
 - Lots of outcomes that are equally good
 - Good outcomes that are separated by bad outcomes
 - Really nuanced constraints
 - etc.

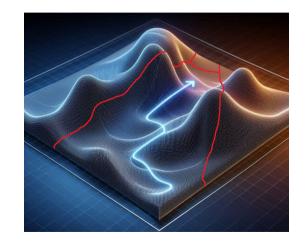
General Form of an Optimization Problem

$$\min f(\overline{X})$$

$$--- subject to (s.t.) ---$$

$$\overline{X} \in \mathbf{S}$$

- ► We must define:
 - \overline{X}
 - $ightharpoonup f(\overline{X})$
 - S



Optimization Theory



- ► This is a very active field of research
- It gets quite complicated and "mathy" very quickly
- ► I'll only emphasize two points:

If we can keep $f(\overline{X})$ and S as linear as possible, the computer algorithms are MUCH faster.

If we can't keep things linear, "Bilinear" is the next best thing.

 $\min f(\overline{X})$

 $s.t. \ \overline{X} \in \mathbf{S}$

- ▶ Linear: $\alpha X + \beta Y + \gamma Z$
- ▶ Bilinear: αXY or αX^2
- ► Non-linear:
 - ightharpoonup $\alpha\sqrt{X}$
 - $ightharpoonup \alpha \log(X)$
 - etc.

Formulating Optimization Problems



Key Elements:

In order how I would conceptualize them:

- 1. Objective Function $(f(\overline{X}))$
- 2. Decision Variables (\overline{X})
- 3. Parameters $(\overline{\alpha})$
- 4. Constraints (S)

In order how I would write / code them:

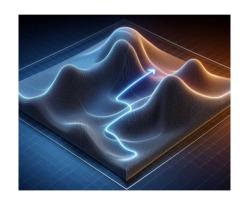
- 1. Parameters $(\overline{\alpha})$
- 2. Decision Variables (\overline{X})
- 3. Constraints (**S**)
- 4. Objective Function $(f(\overline{X}))$

Objective Function $(f(\overline{X}))$



- ▶ What am I trying to accomplish?
- What am I trying to minimize (or maximize)?
- How do different decisions change the outcome?

- Minimize cost
- Minimize risk
- Maximize probability of reaching a goal
- Maximize comfort



Decision Variables (\overline{X})

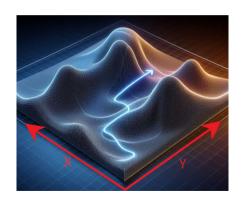


- What decisions can I make?
- What values change when I make different decisions?

Important Note:

- Be open-minded here: often the best way to formulate a problem is to consider each individual part of a problem as it's own variable.
- Don't just consider the main decisions, consider the smaller decisions that contribute to (or even strictly define) the main decisions

- ► How much of a certain item to buy
- ► How much of a budget category to allot for that item

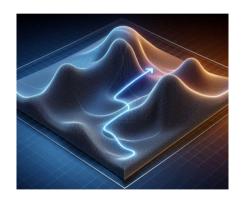


Parameters $(\overline{\alpha})$



- What parts of the problem are fixed or immovable?
- ▶ While the Objective function defines the general shape of the mountain, the parameters define the height of the peaks, the steepness of the slopes, etc.

- Cost of a certain item
- ► Total amount of time available
- ► The minimum probability of reaching a goal that we are willing to accept

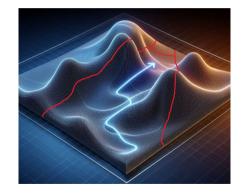


Constraints (S)



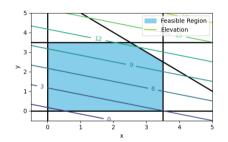
- ▶ What sets of decisions are incompatible?
- What is the nature of an individual decision? (Binary, Continuous, etc.)
- ► How do different decision variables relate to each other?

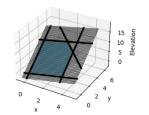
- We cannot spend more in a budget category than the total allotment for that category.
- The amount of money spent in a budget category is strictly equal to the sum of transactions that lies in that category.
- The calculated probability of reaching a goal must be greater than the minimum value we specified.



A Visual Example (Linear)







$$\max X + 3Y - 0.5$$

$$s.t. \quad X \ge 0$$

$$Y \ge 0$$

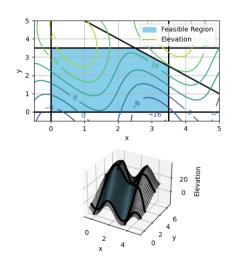
$$X \le 3.5$$

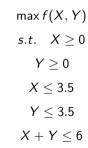
$$Y \le 3.5$$

$$X + Y \le 6$$

A Visual Example (Nonlinear)









Jack and Jill own an Etsy shop that sells t-shirts and wooden sculptures. Jack can either make 5 t-shirts per hour or he can make 2 sculptures per hour. Jill can either make 9 t-shirts per hour or she can make 3 sculptures per hour. Their most recent order indicates that they only need 25 t-shirts but would welcome as many sculptures as they can create. They each how 8 hours to work today. How should they spend their time today so that they can maximize the total number of t-shirts and sculptures they make?

Take a minute to think about the problem on your own:

- 1. Objective Function $(f(\overline{X}))$
- 2. Decision Variables (\overline{X})
- 3. Parameters $(\overline{\alpha})$
- 4. Constraints (**S**)



First conceptualize the problem:

- Objective: Maximize the total number of products produced
- Variables:
 - How many hours each person spends on producing each product
 - The amount of each product that gets made
- Parameters:
 - How much product can be made by each person per hour
 - How must time there is in a day
 - ► The number of t-shirts they need
 - ► The priority between each product

Constraints:

- Each person can only work between 0 and 8 hours in a day
- The amount of each product produced should be directly determined by how many hours each person spends making that product
- The number of t-shirts should be limited to 25



Let's write the problem using mathematical modeling:

- 1. Define Parameters:
 - How much product can be made by each person per hour
 - ho α Jack Sculptures = 2
 - $\alpha_{Jack,Shirts} = 5$
 - $\sim \alpha_{Jill,Sculptures} = 3$
 - $\alpha_{lill,Shirts} = 9$
 - $au^{DAY} = 8$: How much time there is to work in a day
 - $\kappa = 25$: The number of t-shirts they need

- 2. Define Decision Variables:
 - The number of hours each person should spend making each product in a day
 - H_{Jack}, Sculptures
 - ► H_{Jack}, Shirts
 - ► H_{Jill},Sculptures
 - ► H_{Jill},Shirts
 - The amount of each product that gets made
 - N_{Sculptures}
 - N_{Shirts}



- 3. Constraints:
 - Each person can only work between 0 and 8 hours in a day:

$$0 \le H_{Jack,Sculptures} + H_{Jack,Shirts} \le au^{DAY}$$

 $0 \le H_{Jill,Sculptures} + H_{Jill,Shirts} \le au^{DAY}$

► The amount of each product produced should be directly determined by how many hours each person spends making that product:

$$N_{Sculptures} = lpha_{Jack,Sculptures}H_{Jack,Sculptures} + lpha_{Jill,Sculptures}H_{Jack,Sculptures}$$
 $N_{Shirts} = lpha_{Jack,Shirts}H_{Jack,Shirts} + lpha_{Jill,Shirts}H_{Jack,Shirts}$

▶ The number of shirts should be limited to 25:

$$N_{Shirts} = \kappa$$

- 4. Define Objective:
 - Maximize the total number of products produced

$$\max N_{Shirts} + N_{Sculptures}$$



FULL FORMULATION

$$\max N_{Shirts} + N_{Sculptures}$$

$$--- subject \ to \ (s.t.) \ ---$$

$$0 \leq H_{Jack,Sculptures} + H_{Jack,Shirts} \leq \tau^{DAY}$$

$$0 \leq H_{Jill,Sculptures} + H_{Jill,Shirts} \leq \tau^{DAY}$$

$$N_{Shirts} = \alpha_{Jack,Shirts} H_{Jack,Shirts} + \alpha_{Jill,Shirts} H_{Jack,Shirts}$$

$$N_{Sculptures} = \alpha_{Jack,Sculptures} H_{Jack,Sculptures}$$

$$+ \alpha_{Jill,Sculptures} H_{Jack,Sculptures}$$

$$N_{Shirts} = \kappa$$

- Notice how much repetition there is.
- How could I change this problem if I had 1000 people and 75 products?

Sets



- Often there are several variables, parameters, constraints, etc. that are repeated for a given set of values.
- We can write the general idea behind that variable, parameter, constraint, etc. by simply writing it once and indicating that it should be repeated for every element in a set.
- ► In math, it looks like this:

$$\forall e \in \mathbf{E}$$

"for all elements e in the set E"

$$0 \leq H_{Jack,Sculptures} + H_{Jack,Shirts} \leq \tau^{DAY}$$

$$0 \leq H_{Jill,Sculptures} + H_{Jill,Shirts} \leq \tau^{DAY}$$

$$\downarrow$$

$$0 \leq H_{p,Sculptures} + H_{p,Shirts} \leq \tau^{DAY} \quad \forall p \in \mathbf{P}$$

$$\downarrow$$

$$0 \leq \sum_{r \in \mathbf{R}} H_{p,r} \leq \tau^{DAY} \quad \forall p \in \mathbf{P}$$

 $p \in \mathbf{P}$: A set of all People

 $r \in \mathbf{R}$: A set of all Products

Sets



$$\begin{aligned} &\alpha_{\textit{Jack},\textit{Sculptures}} = 2 \\ &\alpha_{\textit{Jack},\textit{Shirts}} = 5 \\ &\alpha_{\textit{Jill},\textit{Sculptures}} = 3 \\ &\alpha_{\textit{Jill},\textit{Shirts}} = 9 \\ &\downarrow \\ &\alpha_{\textit{p,r}} \;\; \forall \textit{p} \in \mathbf{P}, \textit{r} \in \mathbf{R} \end{aligned}$$

$\alpha_{p,r}$	Jack	Jill
Shirts	5	9
Sculptures	2	3

$$\begin{split} N_{Shirts} &= \alpha_{Jack,Shirts} H_{Jack,Shirts} + \alpha_{Jill,Shirts} H_{Jack,Shirts} \\ N_{Sculptures} &= \alpha_{Jack,Sculptures} H_{Jack,Sculptures} \\ &+ \alpha_{Jill,Sculptures} H_{Jack,Sculptures} \\ &\downarrow \\ N_{r} &= \alpha_{Jack,r} H_{Jack,r} + \alpha_{Jill,r} H_{Jack,r} \quad \forall r \in \mathbf{R} \\ &\downarrow \\ N_{r} &= \sum_{p \in \mathbf{P}} \alpha_{p,r} H_{p,r} \quad \forall r \in \mathbf{R} \end{split}$$

Example: Etsy Shop Problem with Sets



$$\max N_{Shirts} + N_{Sculptures}$$

$$--- subject \ to \ (s.t.) \ ---$$

$$0 \leq H_{Jack,Sculptures} + H_{Jack,Shirts} \leq \tau^{DAY}$$

$$0 \leq H_{Jill,Sculptures} + H_{Jill,Shirts} \leq \tau^{DAY}$$

$$N_{Shirts} = \alpha_{Jack,Shirts}H_{Jack,Shirts} + \alpha_{Jill,Shirts}H_{Jack,Shirts}$$

$$N_{Sculptures} = \alpha_{Jack,Sculptures}H_{Jack,Sculptures}$$

$$+ \alpha_{Jill,Sculptures}H_{Jack,Sculptures}$$

$$N_{Shirts} = \kappa_{S}$$

$$\max \sum_{r \in \mathbf{R}} N_r$$
 $---$ subject to $(s.t.)$ $-- 0 \le \sum_{r \in \mathbf{R}} H_{p,r} \le \tau^{DAY} \quad \forall p \in \mathbf{P}$
 $N_r = \sum_{p \in \mathbf{P}} \alpha_{p,r} H_{p,r} \quad \forall r \in \mathbf{R}$
 $N_{Shirts} = \kappa$

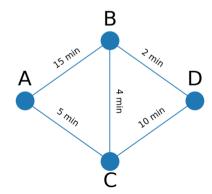
This one can handle any number of people or products.



Given the roads indicated in the graph below, what's the fastest way to get from point A to point D?

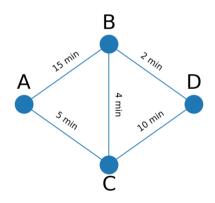
Take a minute to think about the problem on your own:

- 1. Objective Function $(f(\overline{X}))$
- 2. Decision Variables (\overline{X})
- 3. Parameters $(\overline{\alpha})$
- 4. Constraints (S)





First conceptualize the problem:



- Objective: Minimize the time spent traveling
- Variables:
 - Whether or not to travel down each stretch of road
 - Whether or not to visit each point
- Parameters:
 - How long each stretch of road is
 - Which roads connect to which points
- Constraints:
 - If I travel into a point, I must travel out of it.
 - ▶ I have to visit point A and point D



Then write the problem using mathematical modeling:

1. Define Sets:

- $p \in P = \{A, B, C, D\}$: A set of all points
- ▶ $p \in \mathbf{P}^{NON-TERM} = \{B, C\}$: A set of all points other than the terminal (starting and ending) points
- ▶ $p \in \mathbf{P}^{TERM} = \{A, D\}$: A set of all the terminal points
- ▶ $r \in \mathbf{R} = \{AB, AC, BC, BD, CD\}$: A set of all roads
- $r \in \mathbf{R}_p$: A set of all roads that touch point p

р	R_p
Α	$\{AB,AC\}$
В	$\{AB, BC, BD\}$
С	$\{AC, BC, CD\}$
D	$\{BD,CD\}$

2. Define Parameters:

 $ightharpoonup \underline{\delta_r}$: The time it takes to travel road r

r	$\delta_r(minutes)$
AB	15
AC	5
ВС	4
BD	2
CD	10

3. Define Decision Variables:

- ▶ $X_r \ \forall r \in \mathbf{R}$: Whether or not to travel down road r (Binary Variable)
- ▶ Y_p $\forall p \in \mathbf{P}$: Whether or not to visit point p (Binary Variable)



4. Define Constraints:

If I travel into a point, I must travel out of it

$$\sum_{r \in \mathbf{R}_p} X_r = 2Y_p \quad \forall p \in \mathbf{P}^{NON-TERM}$$

$$\sum_{r \in \mathbf{R}_p} X_r = Y_p \quad orall p \in \mathbf{P}^{TERM}$$

► I have to visit point A and point D

$$Y_p = 1 \quad \forall p \in \mathbf{P}^{TERM}$$

- 5. Define Objective:
 - Minimize the amount of time spent traveling

$$\min_{X_r, Y_p} \sum_{r \in \mathbf{R}} \delta_r X_r$$

$$--- \text{ FULL FORMULATION } ---$$

$$\min_{X_r, Y_p} \sum_{r \in \mathbf{R}} \delta_r X_r$$

$$s.t. \quad \sum_{r \in \mathbf{R}_p} X_r = 2Y_p \quad \forall p \in \mathbf{P}^{NON-TERM}$$

$$\sum_{r \in \mathbf{R}_p} X_r = Y_p \quad \forall p \in \mathbf{P}^{TERM}$$

$$Y_p = 1 \quad \forall p \in \mathbf{P}^{TERM}$$

 $X_r, Y_n \in \{0, 1\}$

^{*} Notice how the objective and all constraints are linear

Next Class



▶ Visit the course website:



QR Code to Course Website

These slides are posted in the Lecture 01 folder.

- Python / Optimization coding crash course
- Go the the Lecture 02 folder on the course website
 - Follow the instructions on "slides.pdf" in the Lecture 02 folder.
 - The more you can do on your own before class the more time we'll have to answer questions and do example next class.
- ▶ Please bring your laptops to class!