## **Lecture 2: Memory System**

Main memory

Secondary memory

Memory hierarchy

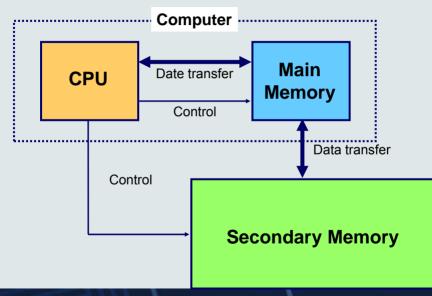
# **Many Different Technologies**



### Introduction

- The <u>main memory</u>, also called <u>primary memory</u>, is used to store the program and data which are <u>currently</u> manipulated by the CPU.
- The <u>secondary memory</u> provides the <u>long-term storage</u> of <u>large amounts</u> of data and program.
- Before the data and program in the secondary memory can be manipulated by the CPU, they must first be loaded into the main memory.

### **Internal and External Memories**



## **Main Memory Model**

A word (8, 16, 32, or 64 bits) **Memory** Control Unit address<sup>2</sup> a bit Read/write control **Address selection** MBR (in CPU) MAR (in CPU)

# Main Memory (MM)

- The MM can be viewed as a set of storage cells, each of which is used to store a word.
- Each cell is assigned a unique address and the addresses are numbered sequentially: 0, 1, 2,...



In this example(!)

1 word = 1 byte

1 word = 1 byte

This is a special case; very often a word consists of 2, 4, or 8 bytes.

# **Main Memory Capacity**

- A byte is traditionally used to encode a text character, and it is the smallest addressable unit of memory.
- Modern computers are usually byte-addressable, even if a word has 2, 4, or 8 bytes.
- The number of address bits determines the maximal size of the memory.
  - Ex.16 bits—64K, 24 bits—16M, 32 bits—4G.
- There are a read/write controller and an address selection mechanism, which are part of the memory control unit.

# **Memory Capacity Units**

Name	Abbreviation	Number of Bytes	No. of bytes	Approximation
Byte	В	1	1 byte	1 byte
Kilobyte	KB	1,024 Bytes	2 <sup>10</sup> bytes	10 <sup>3</sup> bytes
Megabyte	MB	1,024 Kilobytes (about 1 million)	2 <sup>20</sup> bytes	10 <sup>6</sup> bytes
Gigabyte	GB	1,024 Megabytes (about 1 billion)	2 <sup>30</sup> bytes	109 bytes
Terabyte	ТВ	1,024 Gigabytes (about 1 trillion)	2 <sup>40</sup> bytes	10 <sup>12</sup> bytes
Petabyte	PB	1,024 Terabytes (about 1 quadrillion)	2 <sup>50</sup> Bytes	10 <sup>15</sup> bytes

- Exabyte (EB) =  $2^{60}$  =  $10^{18}$ ;
- Zettabyte (ZB) =  $2^{70}$  =  $10^{21}$ ;
- Yottabyte (YB) =  $2^{80}$  =  $10^{24}$ .

# **Growing Data Volume**

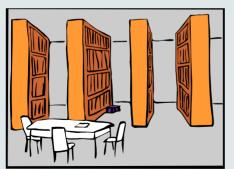
- In a life time, a programmer may generate code of the size:
  - 3 (lines/h) X 8 (hours/d) X 200 (days/y) X 40 (years)
     = 192 000 LOC ≈ 768 000 Bytes < 1 MB.</li>
- In a life time, a writer will at most fill memory of the size:
  - A book page ≈ 2000 characters (2 KB).
  - A 1000-page book ≈ 2 MB.
  - 500 books ≈ 1 GB.
- A digital photo  $\approx$  20 MB (a iPhone photo  $\approx$  3 MB).
- A minute of professional video ≈ 0.5 GB.
- A HD movie ≈ 50 GB after compression.

## **Memory Characteristics**

The most important characteristics of a memory:

- Speed as fast as possible;
- Size as large as possible;
- Cost reasonable price.

They are determined by the technology used for implementation.



#### Your personal library



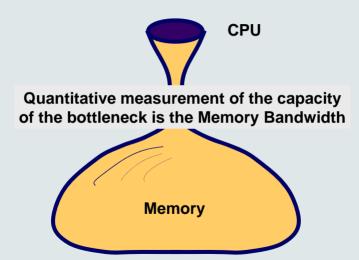
# **Memory Access**

- Reading or writing the content of a memory cell is called an access.
- The time needed to finish one reading or writing operation is called the access time.
- The access time plus any additional time required before a second access can start is called the memory cycle time.
- The memory cycle time is equal to the access time in most technologies.
- Memory access is the bottleneck, because the memory cycle time is much longer than the machine cycle time.
  - Inside a computer, operations are performed in lock-step, each lasting for a clock period.
  - Ex. a memory access can take 10 clock cycles to complete.

### **Performance Metrics**

- Clock frequency
  - Ex. 5GHz = a clock period is 0.2 nano-second long.
- Memory cycle time
  - Ex. A memory access can take 10 clock cycles = 2 nano-second.
- Machine cycle time = Instruction execution time
  - Ex. a computer can have a performance of 500 MIPS (Millions of Instructions Per Second).
  - Since different instructions need different time to execute, the average instruction execution time is often used.
- Very common, <u>FLOPS</u> (FLoating-point Operations Per Second) is used nowadays.
  - Ex. A 4-core PC can have a performance of 20 GigaFLOPS.

# **Memory Access Bottleneck**



# **Memory Bandwidth**

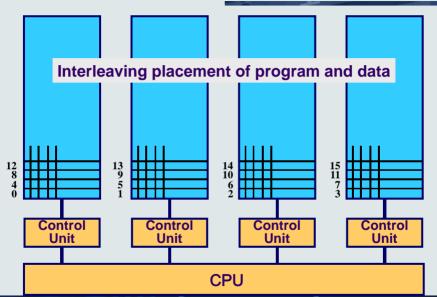
 Memory bandwidth denotes the amount of data that can be accessed from a memory per second:

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M-Bandwidth = \frac{1}{\text{memory cycle time}} • amount of data per access
```

Ex. MCT = 10 nano-second and 4 bytes (a 32-bit word) per access: M-Bandwidth = 400 megabytes per second.

- There are two basic techniques to increase the bandwidth of a given memory:
  - Reduce the memory cycle time
    - Expensive
    - · Memory size limitation
  - Divide the memory into several banks, each of which has its own control unit (using parallelism).

# **Memory Banks**



### **Semiconductor Memories**

- The most widely used technology to implement main memories is semiconductor memories.
  - Ex. CMOS technology: high density, low power consumption, and relatively cheap.
- They use flip-flops made of transistors as the basic unit to represent 0 and 1.
- They are random access memory (RAM).
  - Static RAM (SRAM): fast but expensive, no refreshing needed.
  - Dynamic RAM (DRAM): small and cheap, but need refreshing.
- The information stored in a semiconductor memory will be lost when electrical power is removed (volatile)!

# Read Only Memory (ROM)

- It is a <u>permanent memory</u> which can only be read but not written.
  - Hold instructions that start the computer when it is first switched on (BIOS).
  - A part of the main memory.
- Since ROM is usually very fast, they can also be used when fast reading of program/data is required.
  - Store library subroutines (e.g., for division operation).
  - Store dictionaries for spell checking.
- <u>Programmable ROM (PROM)</u> A ROM which can be programmed once, by either the vendor or a customer.

## **Read-Mostly Memory**

- <u>Erasable PROM (EPROM)</u> To erase the contents of the memory, the chip is exposed to extra-violet radiation.
  - It can take up to 20 minutes to do this.
- <u>Electrically Erasable PROM (EEPROM)</u> Electronic impulses (higher voltages) are used for write, but takes much longer time than the read operation.
- <u>Flash Memory</u> An entire memory or part of it can be erased electrically, based on floating-gate transistor technology.

# **Memory Classification**

### • Random-access memory (RAM):

- The time taken to read or write data does not depend on where the data is stored inside the memory.
- All main memories are of the RAM type.
- Ex. semiconductor memories.

### Sequential-access memory (SAM):

- If a data item is to be read, all data items before it must also be read.
- Ex. magnetic tape.

# **Memory Classification (Cont'd)**

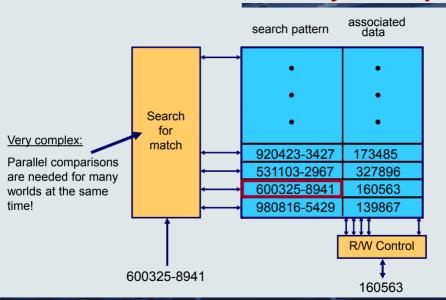
### Direct-access storage devices:

- Data can be obtained without having to read through masses of other data
- Ex hard disk

### Associative memory:

- A RAM where a word is retrieved based on a portion of its content rather than its address.
- Comparison of the given bits of a word with a specified pattern is made for each access, and this is performed for all words simultaneously.

## **Associative Memory Example**



# **Lecture 2: Memory System**

Main memory

Secondary memory

Memory hierarchy

## **Magnetic Tape**

- Magnetic tape is made up from a layer of plastic that is coated with iron oxide.
  - The oxide can be magnetized in two different directions to represent 0/1.
- Its operation uses a similar principle as in the case of a tape recorder.
- Main features:
  - Sequential access (access time about 1-5 s).
  - Relatively high capacity of storage (ca. 80 MB per tape).
  - Inexpensive.
- It is often used for backup or archive purpose.

### **Diskette**

- Data are recorded on the surface of a <u>floppy disk</u> made of polyester coated with magnetic material.
- A special diskette drive must be used to access data stored in the floppy disk.
  - It works much like a record turntable of a gramophone.
- Main features:
  - Direct-access memory
  - Cheap
  - Portable, convenient to use
- Main standards:
  - 5 1/4-inch. Capacity ≈ 360 KB/disk
  - 3 1/2-inch. Capacity ≈ 1.44 MB/disk (about 700 pages of A4 text)



### **Disk Performance**

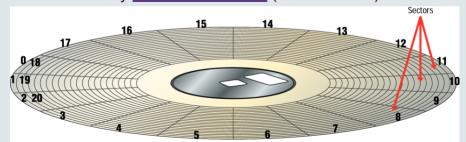
#### Access time:

- Seek time the time required to spin the disk to a constant rotation speed and to position the read/write head at the right track.
- Rotational delay the time required for the read/write head to position at the beginning of the sectors where data are stored.
- <u>Read/write time</u> the time required to read/write a basic unit of data. Often, the data transfer rate is given instead:

$$DTR = \frac{1}{\text{read/write time}}$$

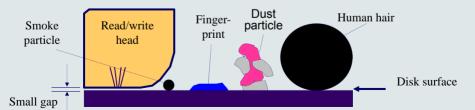
# **Accessing Secondary Memories**

- A secondary memory is usually divided into large blocks (of the size of a few kilobytes, for example).
- Each block has a unique address and can be individually addressed.
- Data are moved between the secondary memory and the main memory one block at a time (not one word!).



### **Hard Disk**

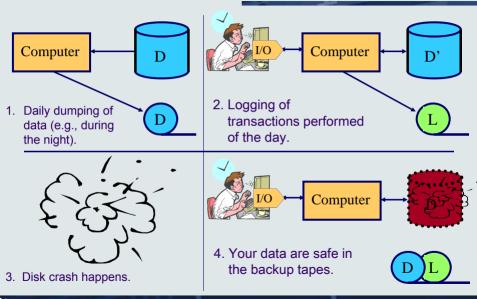
- Data are recorded on the surface of a hard disk made of metal coated with magnetic material.
- A hard disk spins constantly to reduce seek time.
- The disk spins also at a very high speed (up to 10,000 rpm) to reduce the rotational delay and read/write time.
- The disks and the drive are usually built together and encased in a air tight container.
  - It protects the disks from pollutants, such as smoke particle and dust.



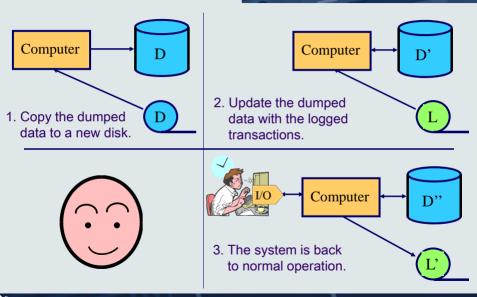
## Hard Disk (Cont'd)

- Several disks are usually stacked on a common drive shaft with each disk having its own read/write head.
- Main features:
  - Direct access (not random access!).
  - Fast access:
    - seek time ≈ 8 ms (vs. 100 ms for floppy)
    - rotational delay ≈ 3 ms (vs. 100 ms for floppy)
    - data transfer rate ≈ 1 Gbits/s (0,5 Mbits/s f. floopy)
  - Huge storage capacity (ca. 500 GB for a compact unit)
- The huge amount of data stored in hard disks must be backed up regularly.

# **Backup Procedure**

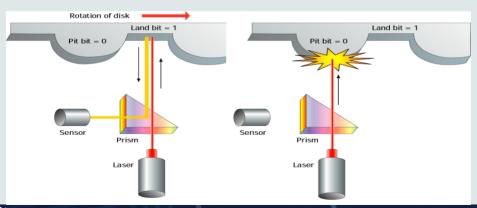


## **Recovery from Backup**



## **Optical Memory**

- An optical disk's surface is imprinted with microscopic holes which record digital information.
- When a low-powered laser beam shines on the surface, the intensity of the reflected light changes, representing 0 or 1.



## **Optical Memory Devices**

#### CD-ROM (Compact Disk ROM):

- Large capacity: > 650 MB/disk (ca. 460 diskettes).
- Inexpensive replication, cheap for mass production.
- Removable.
- Long access time (could be half a second).
- Read-only.

#### CD-Recordable (CD-R):

- Write-once read-many (WORM).
- A laser beam of modest intensity is used to imprint holes.
- Good for archival storage by providing a permanent record of large volumes of data.

# **Optical Memory Devices (Cont'd)**

#### CD-Rewritable (CD-RW):

- Based on different reflectivities in two different phase states.
- Erasable (500,000 to 1,000,000 erase cycles possible).
- Getting cheaper.

#### Digital Versatile Disk (DVD):

- Huge capacity: 17 GB per disk (ca. 12,000 diskettes).
- Full length movie on a single disk (with MPEG compression).
- Read-only, but DVD Recordable and DVD Rewritable are coming.

#### High Definition Optical Disks (HD-DVD):

- Much higher capacity than DVD: 50 GB per disk.
- Shorter wavelength laser (blue-violet range, called Blue-ray DVD).



### **USB Flash Drive**

- A small, portable flash memory card.
- It plugs into a computer's USB port.
  - USB (Universal Serial Bus): an industry standard that defines the cables, connectors and communications protocols for computers and electronic devices.
  - Off-line memory, disconnected or removable memory.
- It functions as a portable hard drive.
- Large capacity is possible: e.g., 256 GB.
- Convenient to store and transfer data.

## **Lecture 2: Memory System**

Main memory

Secondary memory

• Memory hierarchy

### **Motivation**

#### What do we need?

 A memory to store very large programs/data and to work at a speed comparable to that of the CPU.

#### • The reality is:

- The larger a memory, the slower it will be;
- The faster a memory, the greater the cost per bit.

#### A solution:

- To build a composite memory system which combines a small and fast memory with a large and slow memory, and behaves, most of the time, like a large and fast memory.
- This two-level principle can be extended to a hierarchy of many levels.

# **Memory Hierarchy**

**CPU** 

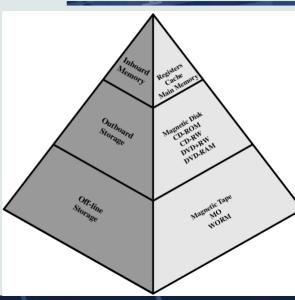
Registers

Cache

**Main Memory** 

**Secondary Memory** of direct access type

**Secondary Memory** of archive type



Technical data are always out of date! **Memory Hierarchy** Access time

example: **CPU** 0.25 ns

1 ns

8 ns

1 ms

(4KB)

100 ms

(100KB)

Registers

Cache

**Main Memory** 

**Secondary Memory** of direct access type

Secondary Memory of archive type

Capacity example:

1 KB

**4 MB** 

**16 GB** 

**8 TB** 

As one goes down the hierarchy, the following occurs:

- Increasing capacity.
- Increasing access time.
- Decreasing cost/bit.
- Decreasing frequency of access by the CPU.

(100 MB/tape)

Zebo Peng, IDA, LiTH

## **Locality of Reference**

- Programs access a small portion of their address space at any short period of time.
- **Temporal locality:** If an item is accessed, it will tend to be accessed again soon.
- Spatial locality: If an item is accessed, items whose addresses are close by will tend to be accessed soon.
- This access pattern is an intrinsic features of the von Neumann architecture:
  - Sequential instruction storage and execution.
  - Loops and iterations (e.g., subroutine calls).
  - Sequential data storage (e.g., array).

## **Summary**

- A memory system has to store very large programs and a huge amount of data and still provide fast access.
- No single type of memory can provide all such need for a computer system.
- Therefore, several different storage mechanisms are organized in a layer hierarchy.
  - The main memory stores the program and data which are currently manipulated.
  - The secondary memory provides the long-term storage of large amounts of data and program.
- The layer structure works very well due to the <u>locality of</u> <u>reference</u> principle.