

# Container With Most Water

📅 2nd Review	@October 28, 2025
# Attempts	3
📅 Date Solved	@October 27, 2025
🕒 Difficulty	Medium
🕒 Status	Solved
☰ Topic/Pattern	Two Pointers

LINK →

## Problem

- You are given an integer array `height` where each element represents the height of a vertical line drawn on the x-axis.
- Any two lines together with the x-axis form a container.
- Return the **maximum amount of water** the container can store.

The container's capacity is determined by:

$$\text{area} = \min(\text{height}[i], \text{height}[j]) * (j - i)$$

## Example

Input	Output	Reason
<code>[1,7,2,5,4,7,3,6]</code>	<code>36</code>	Max area between lines at indices 1 and 7 → $\min(7,6) \times (7-1) = 6 \times 6 = 36$
<code>[2,2,2]</code>	<code>4</code>	Lines at 0 and 2 → $\min(2,2) \times (2-0) = 4$

## Approaches

## 1. Brute Force — Check Every Pair

```
def max_area(height):
    max_water = 0
    n = len(height)

    for i in range(n):
        for j in range(i + 1, n):
            # Width between bars
            width = j - i
            # Height is limited by the shorter line
            area = min(height[i], height[j]) * width
            # Update max area
            max_water = max(max_water, area)

    return max_water
```

- **Time:**  $O(n^2)$  — check all pairs of lines
- **Space:**  $O(1)$  — only uses constant extra space
- **Notes:** Simple but very slow for large inputs (TLE for big arrays)

## 2. Two-Pointer Approach (Optimal)

```
def max_area(height):
    i, j = 0, len(height) - 1
    max_water = 0

    while i < j:
        # Calculate current area
        width = j - i
        h = min(height[i], height[j])
        area = h * width
        max_water = max(max_water, area)
```

```

# Move pointer from the shorter line inward
if height[i] < height[j]:
    i += 1
else:
    j -= 1

return max_water

```

- **Time:**  $O(n)$  — each line is checked once
- **Space:**  $O(1)$  — no extra space used
- **Notes:**
  - Intuition: moving the smaller height inward might find a taller line and increase area.
  - Moving the taller line inward never helps because the limiting height doesn't increase.

## Summary

Approach	Time	Space	Notes
Brute Force	$O(n^2)$	$O(1)$	Checks every pair; very slow
Two-Pointer	$O(n)$	$O(1)$	Efficient and elegant

## Edge Cases

Input	Output	Reason
[1,1]	1	Only two lines $\rightarrow \min(1,1) \times (1) = 1$
[1,8,6,2,5,4,8,3,7]	49	Max area between lines 1 and 8 ( $8 \times 7 = 56 \rightarrow$ but min height limits it to $7 \times 7 = 49$ )
[5]	0	Not enough lines
[]	0	Empty input

## Mistakes

- Made a mistake while moving the pointer.
- Used long ass code instead of `min()`, `max()` function.

## Tip

- The **two-pointer pattern** is common in array problems involving distance and comparison.
- Always remember:
  - $\text{Area} = \min(\text{height}[i], \text{height}[j]) * (j - i)$
  - Move the pointer pointing to the **smaller height**.
- Great practice for mastering the **Greedy + Two-Pointer** technique.