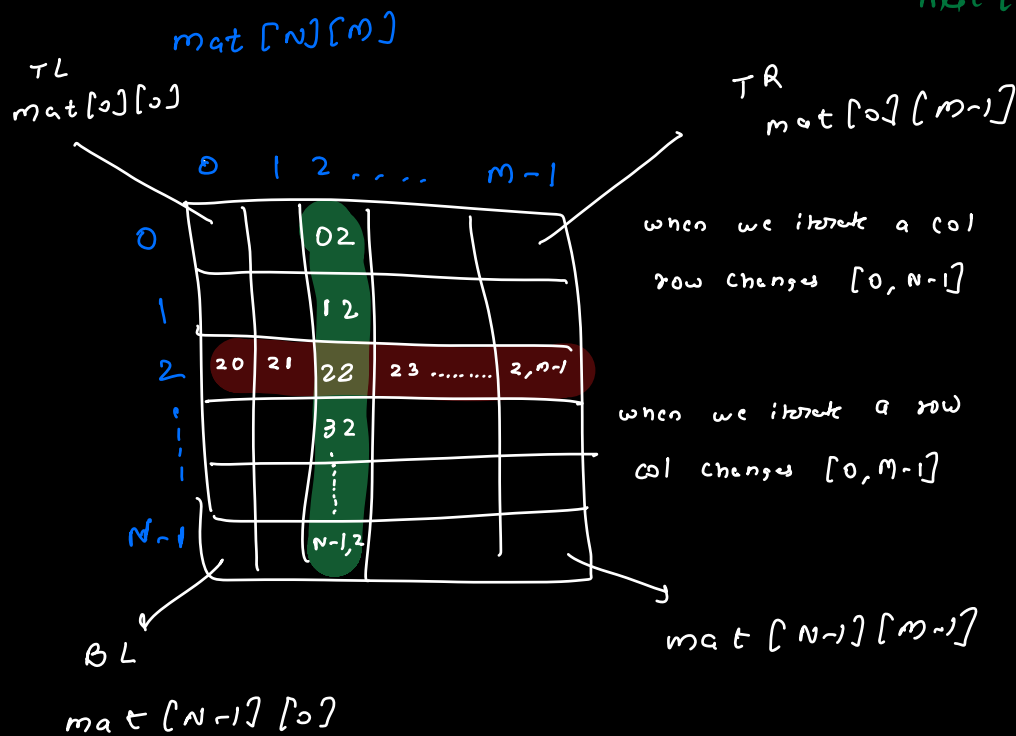
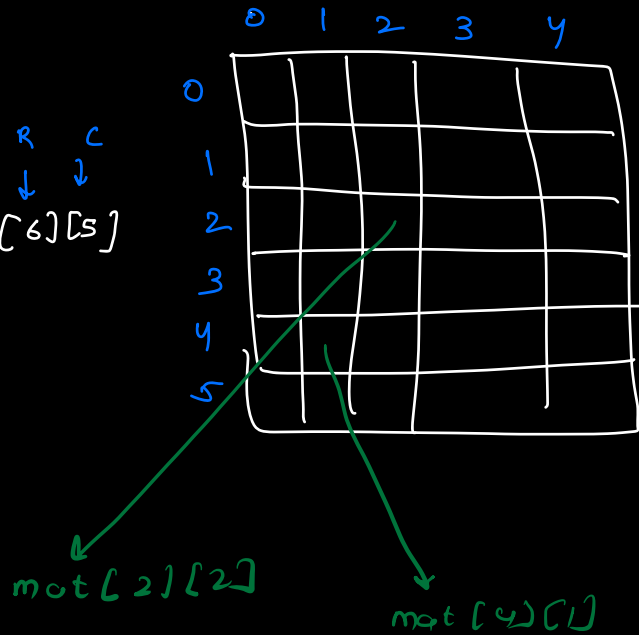


```
int arr[]
int mat[][] = new int[6][5]
```

rows columns



Q. Given $mat[N][M]$, print row-wise sum.

mat[3][4]

	0	1	2	3	Output
0	3	8	9	2	22
1	1	2	3	6	12
2	4	10	11	17	42

```
for(i=0; i < N; i++)
```

```

{
    // ith row
    sum = 0
    for(j=0; j < m; j++)
    {
        sum = sum + mat[i][j]
    }
    print(sum)
}

```

TC: $O(N \times m)$

SC: $O(1)$

Q. Given $mat[N][m]$, find max col-sum.

mat[3][4]

	0	1	2	3
0	3	8	9	2
1	1	2	3	6
2	4	10	11	17
	8	20	23	25

ans = 25

maxi = INT_MIN

for (j = 0; j < M; j++)

// jth column

sum = 0

for (i = 0; i < N; i++)

{ sum = sum + mat[i][j]

if (sum > maxi) maxi = sum

maxi = Math.max(maxi, sum)

Tc: $O(N \times M)$

Sc: $O(1)$

OR

return maxi

void fun (int mat[][7])

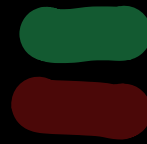
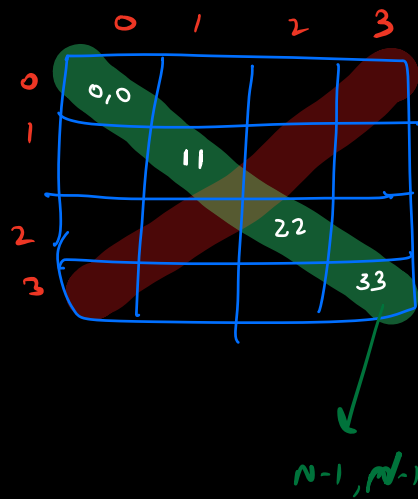
int N = mat.length

int M = mat[0].length

Q. Given a mat[N][N].

print both diagonal (TL → BR, TR → BL)

mat[4][4]



only one
is enough

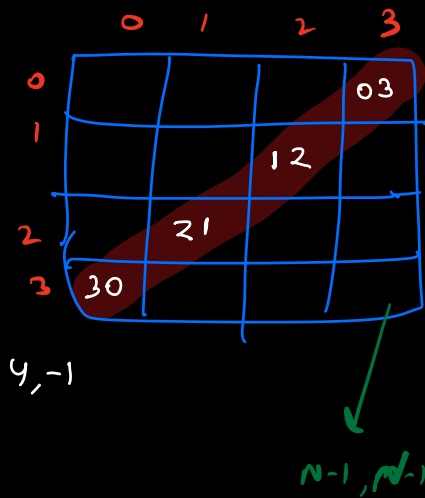
$i=0, j=0$

while ($i < N$ & $j < N$)

{
 print (mat[i][j])
 i++, j++
}

TC: $O(N)$

SC: $O(1)$



only one is
sufficient

$i=0, j=N-1$

while ($i < N$ & $j \geq 0$)

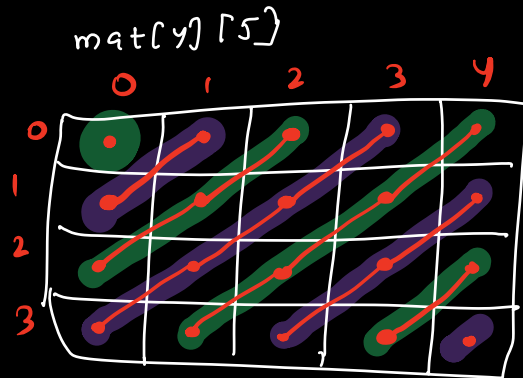
{
 print (mat[i][j])
 i++, j--
}

TC: $O(N)$

SC: $O(1)$

Total TC: $O(N)$
SC: $O(1)$

Q. Given $\text{mat}[N, M]$, print all the diagonals going
 for one diagonal $\rightarrow (R-L)$ (T-B)



Observation \rightarrow Starting points lie on first row or last column.

// 0^{th} row starting diagonals.

for ($j=0$; $j < m$; $j++$)

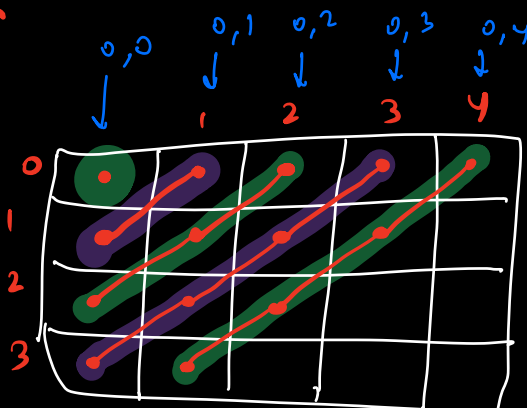
// print the diagonal starting at 0^{th} row, j^{th} col

$x=0, y=j$

while ($x < N$ & $y \geq 0$)

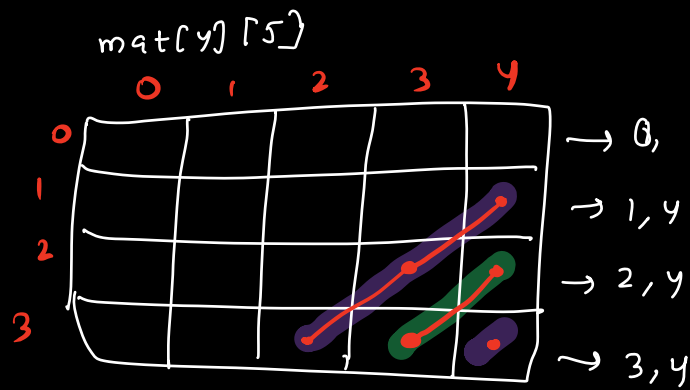
print ($\text{mat}[x][y]$)

$x++, y--$



```
for ( i = 1 ; i < N ; i++)
```

```
{  
    // print diagonal starting at  $i^{\text{th}}$ ,  $m-1$   
  
    x = i, y = m-1  
    while ( x < N & y ≥ 0 )  
    {  
        print ( mat[x][y] )  
        x++, y--  
    }  
}
```



TC: $O(N \times M)$

SC: $O(1)$

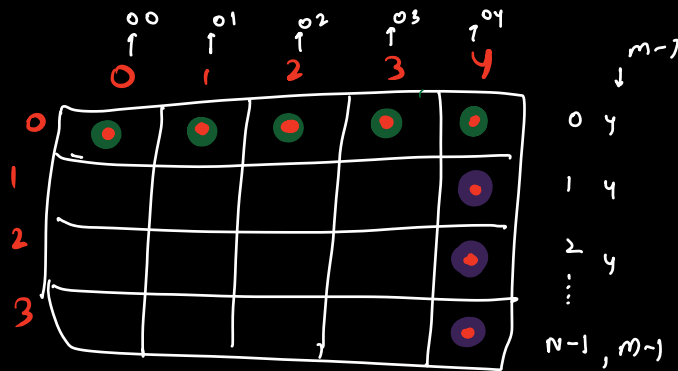
Rewriting above solution

```
void printdia (int i, int j)
```

```

{
    x = i, y = j
    while (x < N && y >= 0)
        print (mat[x][y])
        x++, y--
}

```



```

for (j = 0; j < m; j++)
{
    printdia (0, j)
}

```



TC: $O(N \times m)$
SC: $O(1)$

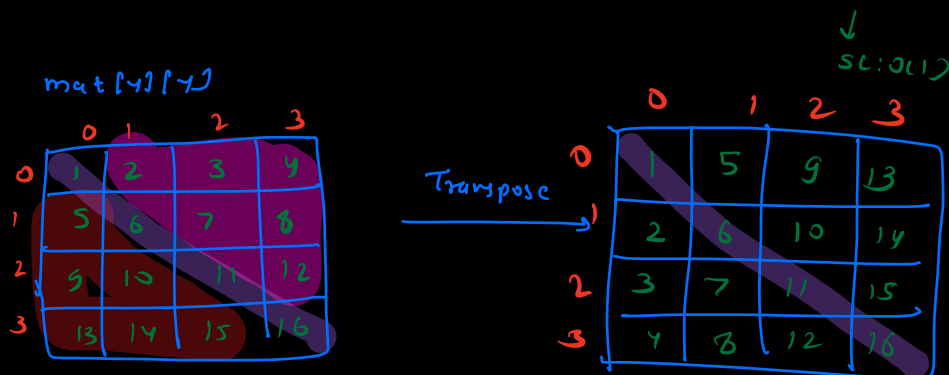
```

for (i = 1; i < N; i++)
{
    printdia (i, m-1)
}

```



Q. Given a $mat[N][N]$, find the transpose.



$mat[i][j] \longleftrightarrow mat[j][i]$

3,0 \rightarrow 0,3

3,1 \rightarrow 1,3

$mat[i][j] \leftrightarrow mat[j][i]$

for (i=0; i<n; i++)

{
 for (j=0; j<i; j++)
 {
 swap (mat[i][j], mat[j][i])
 }
}

4

	0	1	2	3
0	0	1	2	3
1	1	5	6	7
2	2	6	10	11
3	3	7	11	15
...				
N-1				

1 + 2 + 3 + ... N-1

$$\frac{(N)(N-1)}{2} = O(N^2)$$

Tc: $O(N^2)$

Sc: $O(1)$

void swap (int a, int b)

{
 c = a
 a = b
 b = c
}

Q. Given a square mat[N][N]. Rotate it
by 90° clockwise.

↓
SOLUTION

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	16

13	9	5	1
14	10	6	2
15	11	7	3
16	12	8	4

Transpose

1	5	9	13
2	6	10	14
3	7	11	15
4	8	12	16

Reverse
each row

13	9	5	1
14	10	6	2
15	11	7	3
16	12	8	4

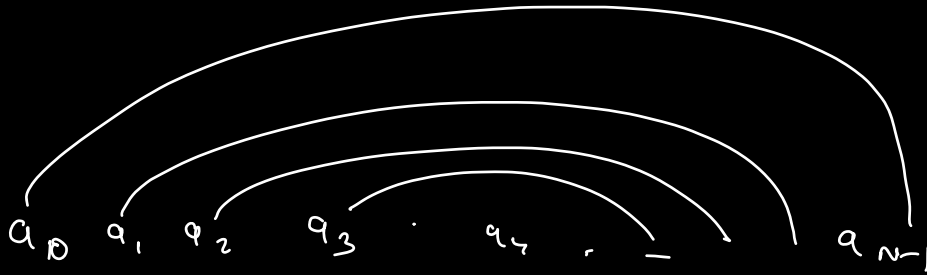
Rotate 90° = Transpose + Reverse each row
TC: $O(N^2)$ TC: $O(N^2)$

SC: $O(1)$

SC: $O(1)$

= TC: $O(N^2)$

SC: $O(1)$



$j=0, i=n-1$

while ($j < i$)

{ swap ($arr[i], arr[j]$)
 $i++, j--$

$arr[2]$

$arr[n-1]$

5
4

4
5