



## Today's agenda

↳ Sorting basics (recap)

↳ Selection sort

↳ Merge sort

↳ merge 2 sorted arr[ ]

↳ merge 2 sorted arr[ ] in single arr

↳ merge sort.

↳ 1 Problem

↳ you didn't come this far only to come this far.



AlgoPrep



Sorting: arranging data in inc/dec order based on Parameters.

ex: 1 7 2 9 24 : inc → on the basis of factors.  
factors: 1 2 2 3 8

### //Bubble Sort

↳ sort the arr in asc. order but we can swap adjacent elements only.

T.C:  $O(n^2)$

S.C:  $O(1)$



## ↳ Selection Sort

$\text{arr}[8]: \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 2 & 8 & 4 & -1 & 7 & 10 & 5 & 6 \end{matrix}$

iter 0:

min ele  
2-1

min idm  
0 3

$\rightarrow \text{swap}(0, 3)$

$\text{arr}[8]: \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 2 & 8 & 4 & 7 & 10 & 5 & 6 & -1 \end{matrix}$

iter 1:

min ele  
8 4 2

min idm  
1 2 3

$\rightarrow \text{swap}(1, 3)$

$\text{arr}[8]: \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ -1 & 8 & 4 & 2 & 10 & 5 & 6 & 2 \end{matrix}$

iter 2:

min ele  
4

min idm  
2

$\rightarrow \text{swap}(2, 2)$

$\text{arr}[8]: \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ -1 & 2 & 4 & 8 & 7 & 10 & 5 & 6 \end{matrix}$

iter N-2



## NPpseudo code

```
void SelectionSort (int arr[N]) {
```

```
    for (int i=0; i<N-1; i++) {  
        int minEle = arr[i], minIdn = i;
```

```
        for (int j=i+1; j<N; j++) {  
            if (arr[j] < minEle) {  
                minEle = arr[j];  
                minIdn = j;
```

```
                int temp = arr[i];  
                arr[i] = arr[minIdn];  
                arr[minIdn] = temp;
```

T.C:  $O(N^2)$

S.C:  $O(1)$



Tracing

```
for (int i=0; i<N-1; i++) {  
    int minele = arr[i], minidn = i;  
    for (int j=i+1; j<N; j++) {  
        if (arr[j] < minele) {  
            minele = arr[j];  
            minidn = j;  
        }  
    }  
}
```

arr[8]: 0 1 2 3 4 5 6 7  
-1 8 4 ~~12~~ 2 7 10 5 6

minele = 1 minidn = 3

temp = 2

3  
int temp = arr[i];  
arr[i] = arr[minidn];  
arr[minidn] = temp;



AlgoPrep



Q) Merge two sorted array

↳ Given 2 sorted arrays  $a[n]$ ,  $b[m]$  create  $c[n+m]$  which contains overall sorted data.

Ex:  $a[4] = \{ 7 \ 10 \ 11 \ 14 \}$   
 $b[3] = \{ 3 \ 8 \ 9 \}$

$c[4+3] = \{ 3 \ 7 \ 8 \ 9 \ 10 \ 11 \ 14 \}$

1/idea1

↳ Create a new array of  $(n+m)$  length, copy the data of  $a[]$  and  $b[]$  and sort the new array.

T.C:  $O((n+m) * \log(n+m))$

1/idea2

$a[4] = \{ \cancel{7} \ \cancel{10} \ \cancel{11} \ \cancel{14} \}$

$b[3] = \{ \cancel{3} \ \cancel{8} \ \cancel{9} \}$

$c[4+3] = \{ 3 \ 7 \ 8 \ 9 \ 10 \ 11 \ 14 \}$



## // Pseudo Code

```
int [] merge (int a[n], int b[m]) {
```

```
    int [] c = new int [n+m];
```

```
    int p1 = 0;
```

```
    int p2 = 0;
```

```
    int p3 = 0;
```

```
    while (p1 < n && p2 < m) {
```

```
        if (a[p1] < b[p2]) {
```

```
            c[p3] = a[p1]; p3++, p1++;
```

```
}
```

```
        else {
```

```
            c[p3] = b[p2]; p3++, p2++;
```

```
}
```

```
    while (p1 < n) { c[p3] = a[p1]; p3++, p1++; }
```

```
    while (p2 < m) { c[p3] = b[p2]; p3++, p2++; }
```

```
    return c;
```

```
}
```

T.C: O(n+m)

S.C: O(n+m) O(1)



## Tracing

```
int[] merge (int a[n], int b[m]) {
    int c = new int [n+m];
    int p1=0;
    int p2=0;
    int p3=0;
```

```
while (p1 < n && p2 < m) {
    if (a[p1] < b[p2]) {
        c[p3] = a[p1];
        p3++;
        p1++;
    } else {
        c[p3] = b[p2];
        p3++;
        p2++;
    }
}
```

```
while (p3 < n) { c[p3] = a[p1]; p3++; p1++; }
```

```
while (p3 < m) { c[p3] = b[p2]; p3++; p2++; }
```

```
} return c;
```

$a[4] = \{ \textcircled{1}, \textcircled{2}, \textcircled{3}, \textcircled{4} \}$

$b[3] = \{ \textcircled{2}, \textcircled{3}, \textcircled{4} \}$

$c[7] = \{ \textcircled{1}, \textcircled{2}, \textcircled{3}, \textcircled{4}, \textcircled{5}, \textcircled{6}, \textcircled{7} \}$

$p_3$

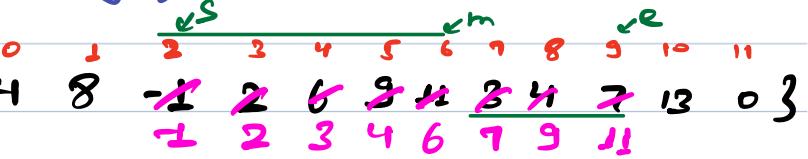
$p_2$

$p_1$



Q) merge 2 sorted subarray

Given arr[n] elements and 3 indices s, m and e  
and subarray [s, m] is sorted, Subarray [m+1, e] is  
sorted. Sort subarray from [s, e]. {s, e} is continuous.

Ex: arr[12]: { 4 8 -1 2 6 3 12 3 4 7 13 0 }  


s m e

2 6 9

temp[g-2+1] = temp[8]: { 0 2 2 3 4 5 6 7 }



AlgoPrep



//  
//

// Pseudo code

int[] merge (int a[s], int s, int m, int e) {

int[] temp = new int [e-s+1];  
int p1 = s;  
int p2 = m+1;  
int p3 = 0;

T.C: O(N)

S.C: O(N)

while ( $p1 \leq m \text{ and } p2 \leq e$ ) {  
if ( $a[p1] < a[p2]$ ) {  
temp [ $p3$ ] =  $a[p1]$ ;  $p3++$ ,  $p1++$ ;  
}  
else {  
temp [ $p3$ ] =  $a[p2]$ ;  $p3++$ ,  $p2++$ ;  
}  
}

while ( $p1 \leq m$ ) { temp [ $p3$ ] =  $a[p1]$ ;  $p1++$ ,  $p3++$ , }  
}

while ( $p2 \leq e$ ) { temp [ $p3$ ] =  $a[p2]$ ;  $p2++$ ,  $p3++$ , }

// Copy temp [0 - 7] back to a [s - e] → todo

return a;

}



```
int[] temp = new int[e-s+1];
```

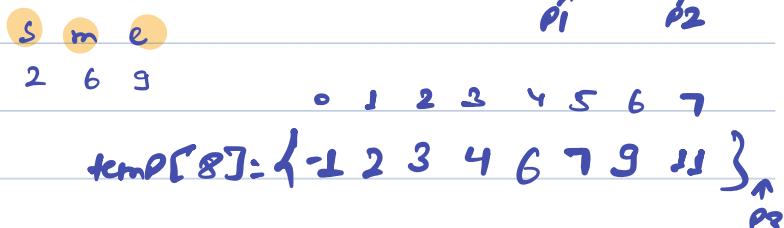
```
int p1 = s;
```

```
int p2 = m+1;
```

```
int p3 = 0;
```

```
while (p1 <= m && p2 <= e) {  
    if (a[p1] < a[p2]) {  
        temp[p3] = a[p1]; p3++, p1++;  
    }  
    else {  
        temp[p3] = a[p2]; p3++, p2++;  
    }  
}
```

Ex:  $a[12] = \{ 4, 8, 1, 2, 6, 9, 11, 3, 4, 7, 13, 0 \}$



```
while (p1 <= m) { temp[p3] = a[p1]; p1++, p3++ }
```

```
while (p2 <= e) { temp[p3] = a[p2]; p2++, p3++ }
```

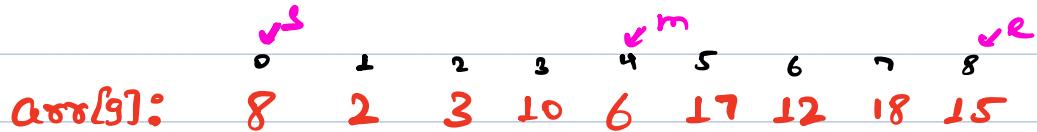


# AlgoPrep



## Q) merge sort

↳ Given  $\text{arr}[n]$  sort them.



## Pseudo code

```
void mergesort (int arr[n], int s, int e){  
    if (s == e) { return; }  
    int m =  $\frac{s+e}{2}$ ;
```

T.C:  $O(n \log n)$

S.C:  $O(n)$

```
        mergesort (arr, s, m);  
        mergesort (arr, m+1, e);
```

```
        merge (arr, s, m, e);
```

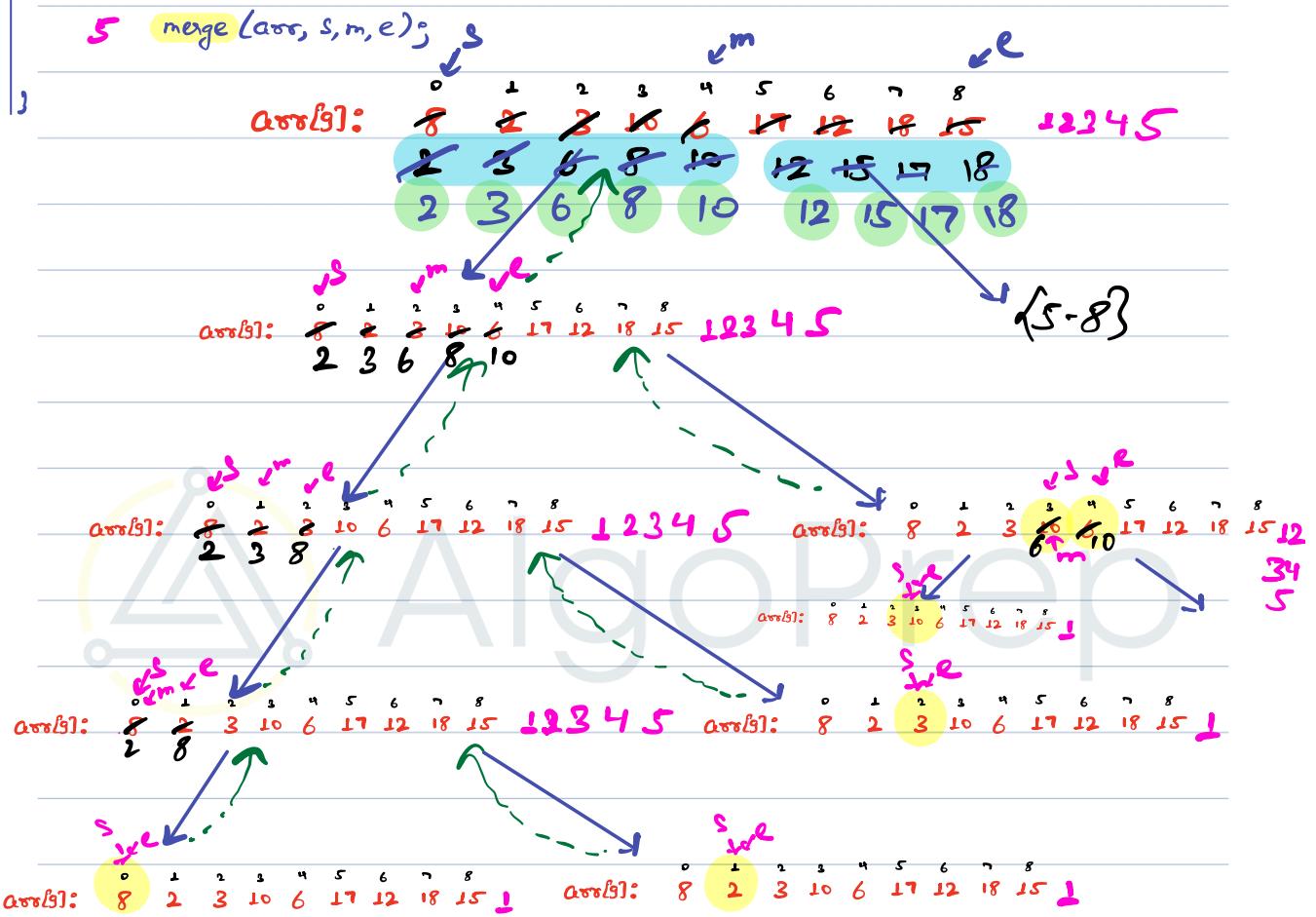
}

```

void mergesort (int arr[n], int s, int e) {
    1 if (s == e) { return; }

    2 int m =  $\frac{s+e}{2}$ ;
    3 mergesort (arr, s, m);
    4 mergesort (arr, m+1, e);
}

```



At each level merge combined =  $O(n)$

$N$

$\downarrow$

$\frac{N}{2}$

$\downarrow$

$\frac{N}{4}$   
⋮  
1

How many such levels?  $\Rightarrow \log N$

overall T.C:  $O(N \log N)$ .



## Q) inversion count

Given  $\text{arr}[n]$ , find the inversion count.

Two elements  $a[i]$  and  $a[j]$  form an inversion if  $a[i] > a[j]$  and  $i < j$ .

o 1 2 3

$$\text{Ex: } \text{arr}[] = \{ 8 \ 4 \ 2 \ 1 \} \rightarrow \text{ans} = 6$$

(8,4) (8,2) (8,1) (4,2) (4,1) (2,1)

$$\text{Ex: } \text{arr}[] = \{ 1 \ 20 \ 6 \ 4 \ 5 \} \rightarrow \text{ans} = 5$$

Ideas

Check all Possible Pairs and Count the Inverted Pairs.

T.C:  $O(N^2)$

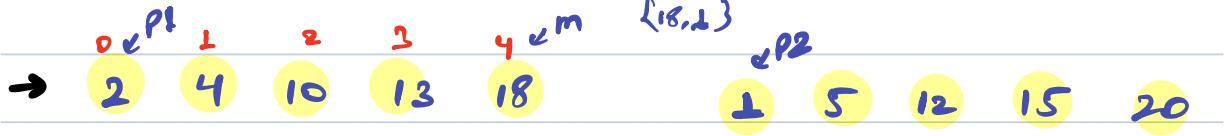
S.C:  $O(1)$

Count = 0 + 5 + 3 + 2

$\{2, 1\}$	$\{10, 5\}$	$\{12, 12\}$
$\{4, 13\}$	$\{13, 5\}$	$\{18, 12\}$
$\{10, 13\}$	$\{18, 5\}$	
$\{13, 1\}$	$\{18, 13\}$	
$\{18, 1\}$		



## Idea 2



if ( $\text{arr}[\rho_1] > \text{arr}[\rho_2]$ ) {

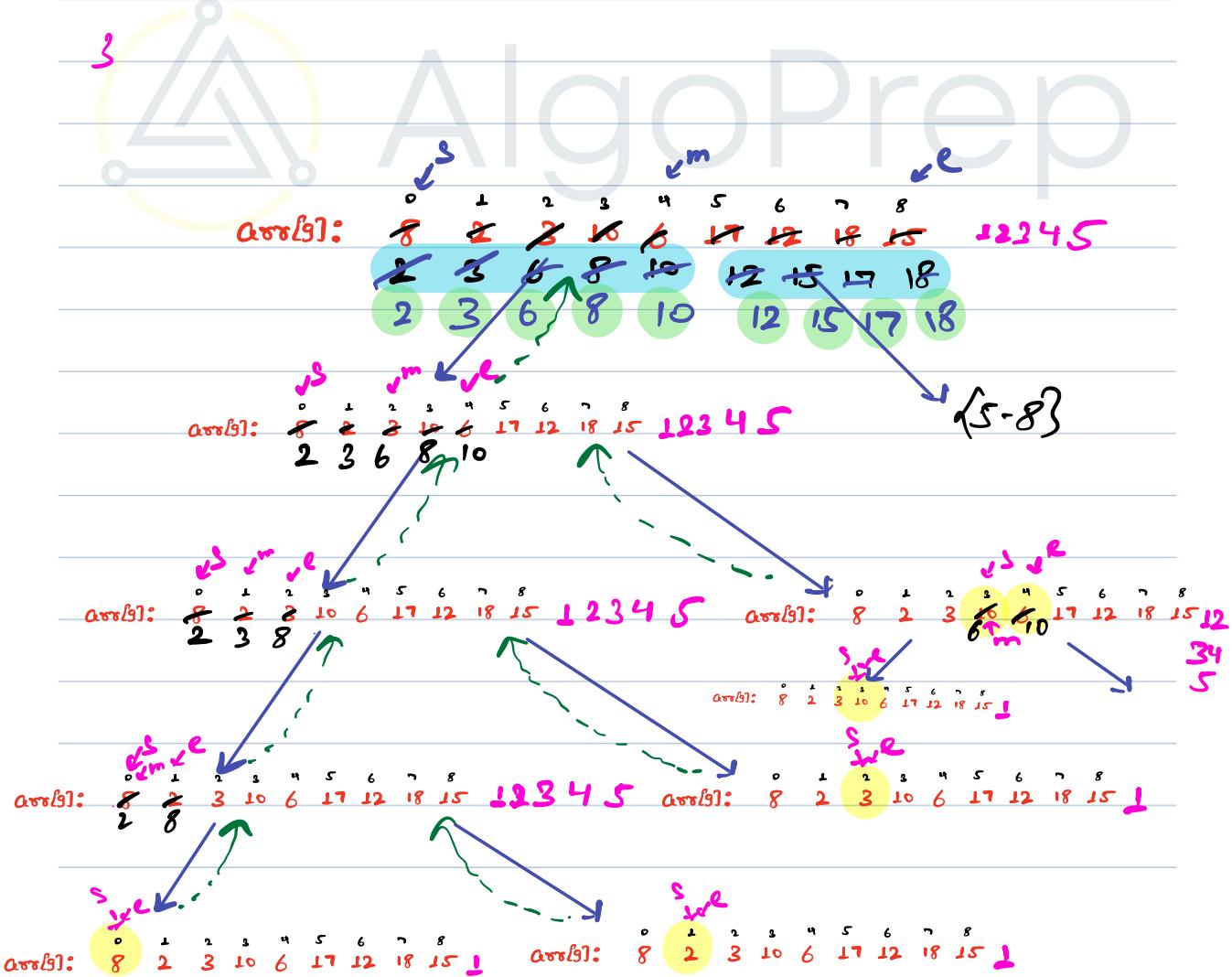
## update count

$$\text{Ansatz: } (m - p_1 + l),$$

3

else { if  $\cos(\theta_1) \leq \cos(\theta_2)$

2





```
void mergesort (int arr[n], int l, int e){  
    if (l == e) { return; }
```

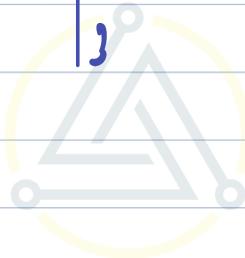
T.C:  $O(n \log n)$

S.C:  $O(n)$

$$\text{int } m = \frac{s+e}{2};$$

```
    mergesort (arr, s, m);  
    mergesort (arr, m+1, e);
```

merge (arr, s, m, e);



AlgoPrep



int ans = 0;

int[] merge(int a[], int s, int m, int e){

int[] temp = new int[e-s+1];

int p1 = s;

int p2 = m+1;

int p3 = 0;

T.C: O(N)

S.C: O(N)

while ( $p1 \leq m \text{ and } p2 \leq e$ ) {

if ( $a[p1] \leq a[p2]$ ) {

temp[p3] = a[p1]; p3++, p1++;

}

else { if  $a[p1] > a[p2]$

ans = ans + (m-p1+1);

temp[p3] = a[p2]; p3++, p2++;

}

}

while ( $p1 \leq m$ ) { temp[p3] = a[p1]; p1++, p3++; }

while ( $p2 \leq e$ ) { temp[p3] = a[p2]; p2++, p3++; }

/copy temp[0 - 7] back to a[s - e] → todo

return a;

}

$$ans = 0 + 1 + 2 + 1 + 1 \rightarrow \{8, 6\}$$

