· Dijskstras Algorithm :-	(L)
-) This is other kind of shortest path problem. -) This is other kind of shortest path problem. -) there there exists source vertex from where the parties are calculates its shortest path from all v	
-) This algorithm is given by E.W. DIS FIRM a	nd is Gulad
Didskstrås Algandhum. The steps to be followed in this dijskstrås Algande som the steps to be followed in this dijskstrås Algande som the steps to be followed in this dijskstrås Algande som the steps to be followed in this dijskstrås Algande som the steps to be followed in this dijskstrås Algande som the steps to be followed in this dijskstrås Algande som the steps to be followed in this dijskstrås Algande som the steps to be followed in this dijskstrås Algande som the steps to be followed in this dijskstrås Algande som the steps to be followed in the steps to be steps to be followed in the steps to be steps t	
J. Britially Mark all distances from the source vertex/hode as is except the source vertex/hode as is mark the distance of source vertex/hode as is mark the distance of source vertex/hode as is. Then cheek all the Possible ordering edges of a source vertex as,	yex,ie,
calculate their	
(curent light of source vertex to next vertex) + (1.	v. (source verta)
< ib R	serious Listance
il-, let V, be some vertex & vz, vz are is at	tracked nodes.
[length (V1 to V2) + V, distance (Frev.)] < length (i).	
where i' is the Prevy distance of the	at node
which was calculated already	<i>'</i> ,
3. Then must the node v, . Also it above condi	tion is true
then LHS Part is falken as distance @ RHS Part is	taken.
y, Continue the same method till it reaches all nodes of graph & J	ind Shortest distance a path Jum Jinal table

(4.A)

Step! Enitially set all the vertices length to is except the first node in graph (take it as 0).

	0,130	1 0 00
Status	node vertices	consth/distance
>>	VI	0
	VL	∞ .
	٧ _٤	20
	V4	∞
	VS	\propto

Step2: check all onlyoing edges of node V, (Status intrated node)

staty node	ListanCl
→ V)	1+0<0 ('' V ₁ -V ₂ it is I and prev. V ₁ value is 6+0<0 (and prev value is 5).
٧L	1+0<0 (. VI-V2 want value is so).
Vs	6+0< 2
V4.	
V.S	status node distance
	VIO
	→ V2 1

 $\begin{array}{c|cccc}
 & V_1 & O \\
 & V_2 & 1 \\
 & V_3 & 6 \\
 & V_4 & 20 \\
 & V_5 & \infty
\end{array}$

Step3: Now V, node is marked. Then deek all outgoing edges of node v_2 (ie, the next least).

Saby	node	distance	Study	nodl	distance
				′ V)	O
\(\rightarrow \)	Vz Vz	$\begin{pmatrix} 1 \\ 2+1 < 6 \end{pmatrix}$	(t) =) \(\sqrt{1}	V2 V3	\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\
	V4	1 2+1<6 3+1<\infty 5+1<\infty		\sqrt{\gamma}	4
	VS 1	5+1<0		VS	6

Stef 4: Now VIIVz one marked. From check all outgring edges of node v3 (ie, next loost).

Status	node	distance	Staty	node	distance.
	VI VI VS	0 1 3	=) 7	V) V2	0. 1. 3. 4. 5. 6.

SEP 5:- Now V11/2, vs are malked. Then decek all outgoing edges of node v4 (ie, next least).

3	, ,					
Status	node	distance		Status	node	distance
	VI	0			V	0
	1/2)	-)		V).	1
	1	3			7	3
	V4 VE	2+4 & S (F)			VS	4
	, v3	645(4)			1.	

so finally the Shortest path for verlex's from source verlex 1 is of length 's' ie, V1-V2-V3-V5.