

Algorithms and Data Structures



COMP261 **Tutorial 4**

Yi Mei

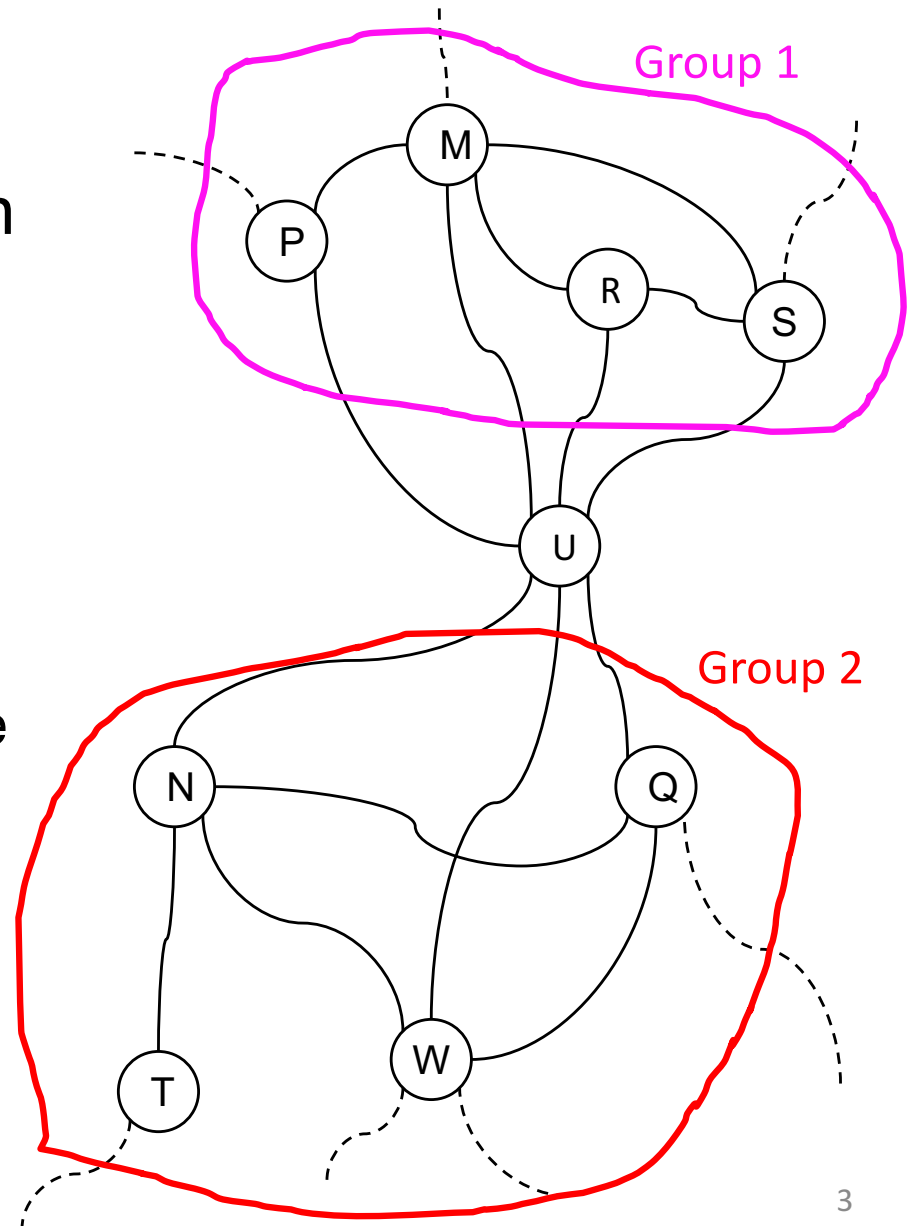
yi.mei@ecs.vuw.ac.nz

Outline

- Finding all articulation points
 - Idea
 - Implementation: recursive and iterative

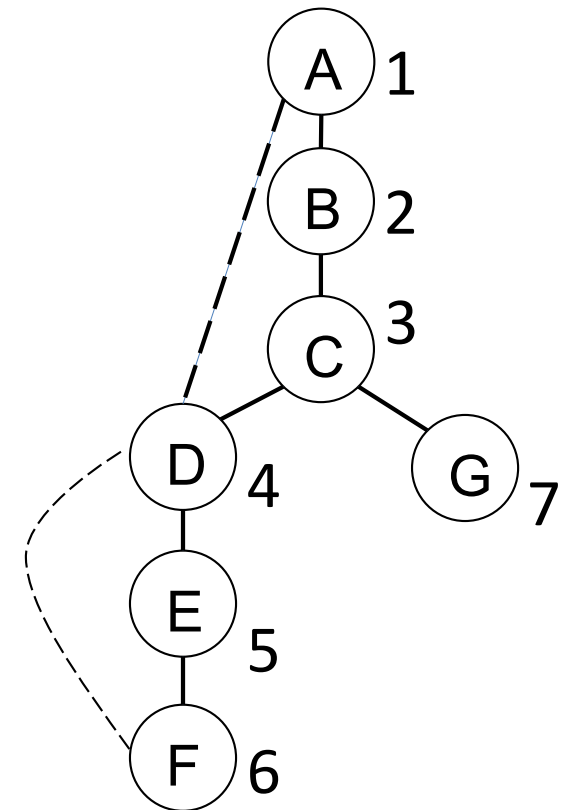
Finding All Articulation Points Efficiently

- **Idea:** an articulation point separates the graph into two groups, so that all paths from nodes in one group to nodes in the other group **MUST** go through the node.
- **Example:**
 - node U is an articulation point, since all paths between any node in group 1 and any node in group 2 must go through U.



Two Sets of a Node

- In the search tree, each node separates the nodes into two subsets
 - **Children set:** Nodes **in its subtree**
 - **Parents set:** Nodes **not in its subtree**
- Check if **Children** and **Parents** are separated after removing the node
 - The node is an articulation point if at least one **child** node and **parents** are separated after removing it
 - There is no alternative path



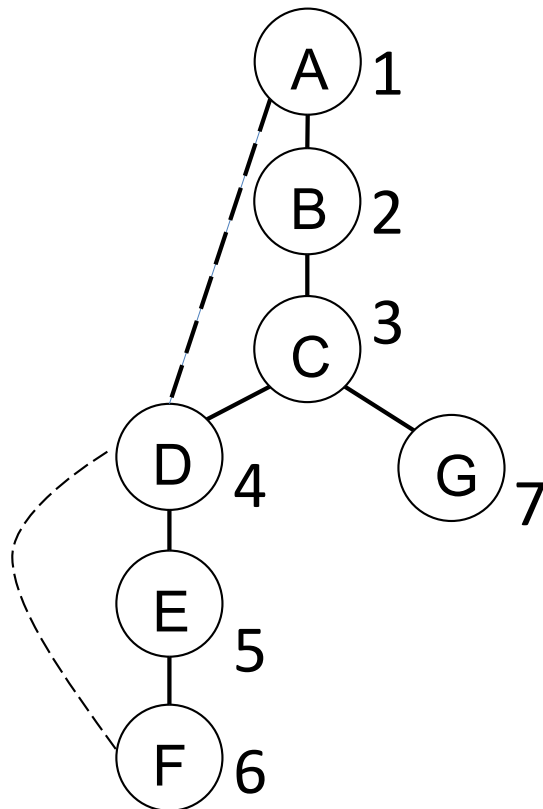
Articulation Points Algorithm

- **Theorem:** a non-root node A is an articulation point, if and only if there exists a child node B, for which there is no alternative path from B to any of the parents
 - Removing node A will separate B and the parents
 - This is independent of the root node of the DFS and order that the neighbours are checked
- Checking alternative paths for a child node B of node A
 - An edge in the graph, but not in the DFS tree
 - Directly link node B to a parent node
 - Link another child node C in the same subtree of B to a parent node
 - All the child node in the same subtree are connected without node A

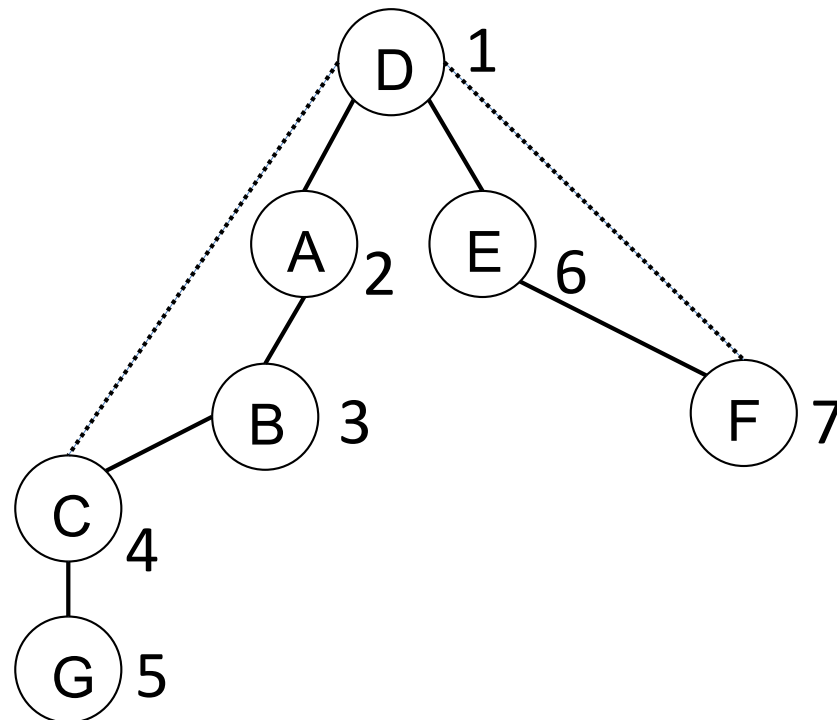
Articulation Points Algorithm

- **Theorem:** a **root node** is an **articulation point** if and only if it has **multiple sub-trees** in the DFS
 - Proof is easy (no alternative path between sub-trees)

A is not an articulation point
(one sub-tree)



D is an articulation point
(two sub-trees)



Recursive AP Algorithm

Initialise count number of all nodes as $\text{count}(\text{node}) = \infty$, meaning all nodes are unvisited;
Initially, $\text{APs} = \{\}$, that is, no articulation point is found;
Randomly select a node as the root node, set $\text{count}(\text{root}) = 0$, $\text{numSubTrees} = 0$;

```
for (each neighbour of root) {  
    if (count(neighbour) =  $\infty$ ) {  
        recArtPts(neighbour, 1, root); // recursive DFS for the neighbour  
        numSubTrees ++;  
    }  
  
    if (numSubTrees > 1) then add root into APs;  
}
```

```
recArtPts(node, count, parent) {  
    ...  
}
```

Recursive AP Algorithm

```
recArtPts(node, count, parent) {  
    count(node) = count;  
    // store the minimum count number the node can reach back via alternative path  
    reachBack = count;  
    for (each neighbour of node other than parent) {  
        // case 1: direct alternative path: neighbour is visited before  
        if (count(neighbour) <  $\infty$ )  
            reachBack = min(count(neighbour), reachBack);  
        // case 2: indirect alternative path: neighbour is an unvisited child in the same sub-tree  
        else {  
            // calculate alternative paths of the child, which can also be reached by itself  
            childReach = recArtPts(neighbour, count+1, node);  
            reachBack = min(childReach, reachBack);  
            // no alternative path from neighbour to any parent  
            if (childReach >= count) then add node into APs;  
        }  
    }  
    return reachBack;  
}
```


Iterative AP Algorithm

Initialise $\text{count}(\text{node}) = \infty$, $\text{APs} = \{\}$;

Randomly select a node as the root node, set $\text{count}(\text{root}) = 0$, $\text{numSubTrees} = 0$;

for (each neighbour of root) {

if ($\text{count}(\text{neighbour}) = \infty$) {

iterArtPts(neighbour, 1, root);

$\text{numSubTrees}++$;

 }

if ($\text{numSubTrees} > 1$) **then** add root into APs;

}

The only difference from
the recursive version

iterArtPts(node, count, parent) {

 ...

}

Iterative AP Algorithm

```
iterArtPts(firstNode, count, root) {  
  Initialise stack as a single element <firstNode, count, root>;  
  repeat until (stack is empty) {  
    peek <n*, count*, parent*> from stack;  
    if (count(n*) =  $\infty$ ) {  
      count(n*) = count, reachBack(n*) = count;  
      children(n*) = all the neighbours of n* except parent*;  
    }  
    else if (children(n*) is not empty) {  
      get a child from children(n*) and remove it from children(n*);  
      if (count(child) <  $\infty$ ) then reachBack(n*) = min(count(child), reachBack(n*));  
      else push <child, count+1, n*> into stack;  
    }  
    else {  
      if (n* is not firstNode) {  
        reachBack(parent*) = min(reachBack(n*), reachBack(parent*));  
        if (reachBack(n*) >= count(parent*)) then add parent* into APs;  
      }  
      remove <n*, count*, parent*> from stack;  
    }  
  }  
}
```

Example

